# Number Theory

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#### Goals:

- Basic Primality Testing
- Sieve of Eratosthenes
- Prime Factorization
- Binary Exponentiation
- GCD/LCM and their properties

### **Basic Primality Testing**

A primality test is an algorithm for determining whether a number is prime or composite.

There are many algorithms for primality testing such as:

- ✓ Brute-Force
- ✓ Square Root method
- ✓ Using Sieve (precomputation)

#### **Brute Force Approach:**

Check for every integer from 2 to n-1, if it divides n.

```
bool is_prime(int n)
{
    for (int i = 2; i < n; i++)
        if (n % i == 0) return false;
    return n > 1;
}
```

#### Some important observations:

- ✓ Factors always appear in pair.
- ✓ If number given n is not prime, then it must be possible to factor it into two values A and B, such that A.B = n.
- ✓ Hence with this we can see, either one of the two has to be less than root(n).
- ✓ Therefore, if the number is not a prime, it must have atleast one factor less than or equal to the square root of the number.

```
bool is_prime(int n)
{
    for (int i = 2; i*i <= n; i++)
        if (n % i == 0) return false;
    return n > 1;
}
```

#### Sieve of Eratosthenes

An algorithm for finding all the prime numbers in just O(n\*loglogn) operations.

#### Main Idea:

- ✓ Don't check for all numbers, multiples of prime numbers are always composite.
- ✓ Mark multiples of prime numbers as composite.

Time complexity: n\*log(log(n)) - Here

```
void sieve(int n)
     bool primes[n+1];
     fill(primes, primes+n+1, true);
     primes[0] = primes[1] = false;
     for (int i = 2; i*i <= n; i++) {
          if (primes[i])
               for (int j = i*i; j <= n; j += i)
                    primes[j] = false;
```

### **Prime Factorization**

Prime Factorization is finding all the prime factors of a given number.

There are multiple ways to factor primes, such as:

- ✓ Trial Division
- ✓ Sieve method (precomputation)

### Trial Division Method 👰

Here, we will use the fact that the smallest divisor of any number N is prime, and is less than square root of N.

N can be represented as product of powers of primes.

```
vector<int> factor(int n)
  vector<int> facts;
  for (long long d = 2; d * d <= n; d++) {
    while (n \% d == 0) \{
      facts.push back(d);
      n = d;
  if (n > 1)
    facts.push_back(n);
  return facts;
```

## Problem Problem

You are given an integer n.

Check if n has an **odd** divisor, greater than one (does there exist such a number x (x>1) that n is divisible by x and x is odd).

 $2 \le n \le 10^{14}$  - Here



Problem 1

Problem 2

Problem 3

### **Binary Exponentiation**



It is used to compute the value of  $A^B$  efficiently.

The naı̈ve approach takes O(B) whereas binary exponentiation has time complexity of  $O(\log_2 B)$ .

#### **Recursive Approach**

```
int binary_expo(int a, int b)
{
    if (b == 0) return 1;
    int c = binary_expo(a, b/2);
    if (b % 2 == 0) return c*c;
    return c*c*a;
}
```

#### Iterative Approach

```
int binary_expo(int a, int b) {
   int c = 1;
   while (b) {
     if (b % 2) c *= a;
     a *= a;
     b /= 2;
   }
   return c;
}
```

### **Greatest Common Divisor**

GCD(a,b) is the greatest common divisor of a and b.

LCM(a, b) is the least common multiple of a and b.

To calculate gcd efficiently, we use Euclidean Algorithm.

It states that GCD(a, b) = GCD(b, a%b). when b = 0, the solution is a.

Time Complexity: O(log(min(a,b))) - Proof

```
int gcd_(int a, int b)
{
    if (b == 0) return a;
    return gcd(a, a%b);
}
```

#### **Least Common Divisor**

$$lcm(a,b) = \frac{a*b}{\gcd(a,b)}$$

### Properties 🍥

- $\checkmark \gcd(a,b)$  can be represented as product of  $\min(p_i^{a_i},p_i^{b_i})$  for each prime factor.
- $\checkmark lcm(a,b)$  can be represented as product of  $max(p_i^{a_i},p_i^{b_i})$  for each prime factor.
- $\checkmark$  gcd(a, b, c, ...) is same as gcd(gcd(gcd(a, b), c), ...).
- $\checkmark lcm(a, b, c,...)$  is same as lcm(lcm(a, b), c),...).
- $\checkmark \gcd(a, a + 1) = 1.$

Problem: Here