Laboratory 9: Graphs

CSC205A Data structures and Algorithms Laboratory B. Tech. 2015

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Introduction and Purpose of Experiment

- Graphs are important data structures used in many applications like network modelling applications, routing, traffic, water supply etc.
- This experiment introduces graphs and its representation.



Aim and objectives

Aim:

 To design and develop C program to represent a Graph using adjacency list and adjacency matrix

Objectives:

At the end of this lab, the student will be able to

- Represent a graph data structure using adjacent list and adjacency list
- Compare and contrast between both the approaches.



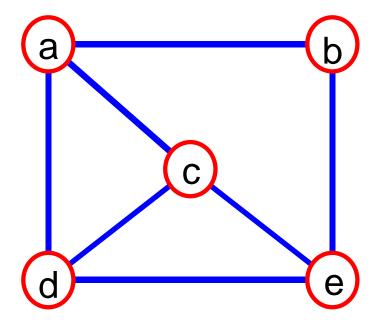
What is a Graph?

A graph G = (V,E) is composed of:

V: set of vertices

E: set of edges connecting the vertices in V

- An edge e = (u,v) is a pair of vertices
- Example:



$$V = \{a,b,c,d,e\}$$

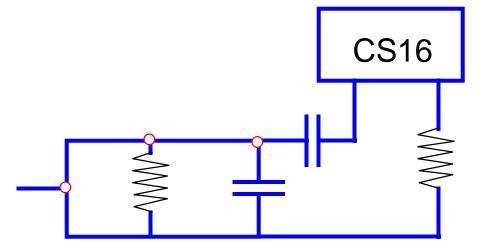


Some real life applications



Applications

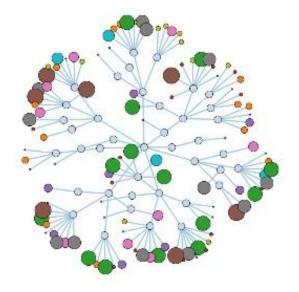
• Electronic circuits





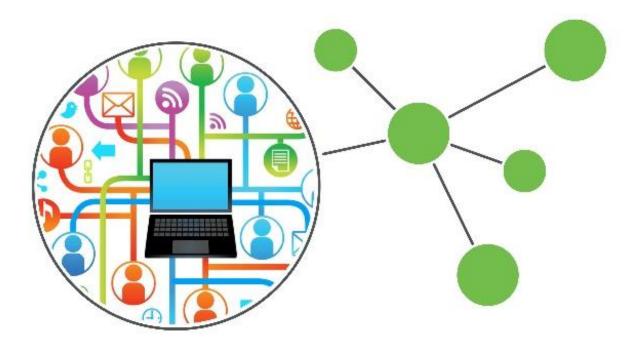
Graph Databases

- Store
- Manage
- Query



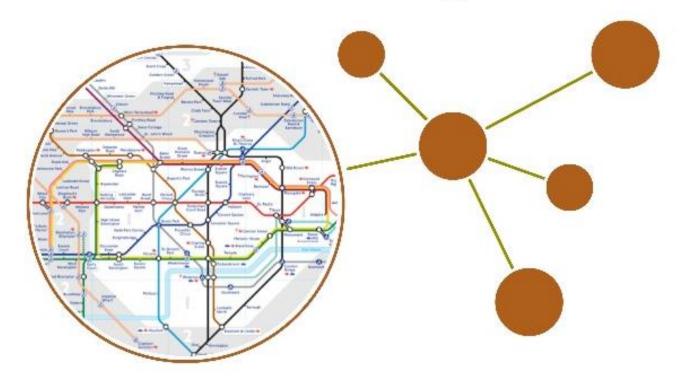


Social Network



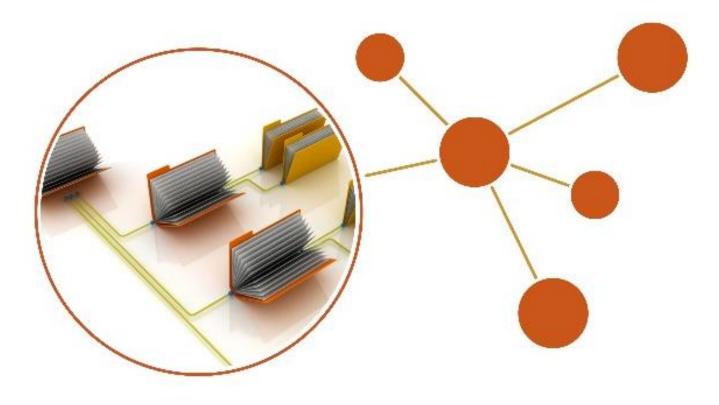


Route Finding





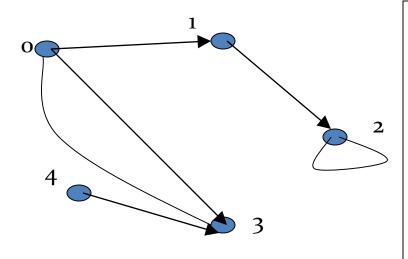
Access Control





Examples of Graphs

- V={0,1,2,3,4}
- $E=\{(0,1), (1,2), (0,3), (3,0), (2,2), (4,3)\}$



When (x,y) is an edge, we say that x is *adjacent to* y, and y is *adjacent from* x.

o is adjacent to 1.

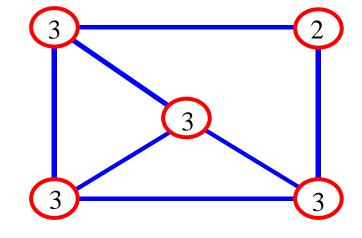
1 is not adjacent to o.

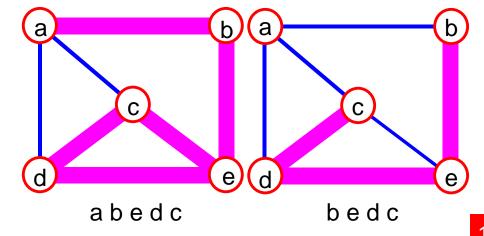
2 is adjacent from 1.



Terminology: Path

 Path: sequence of vertices v₁,v₂,...v_k such that consecutive vertices v_i and v_{i+1} are adjacent.

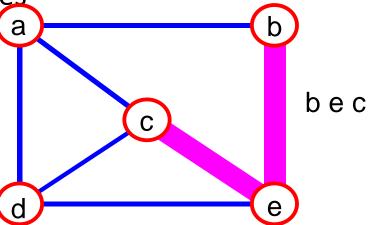




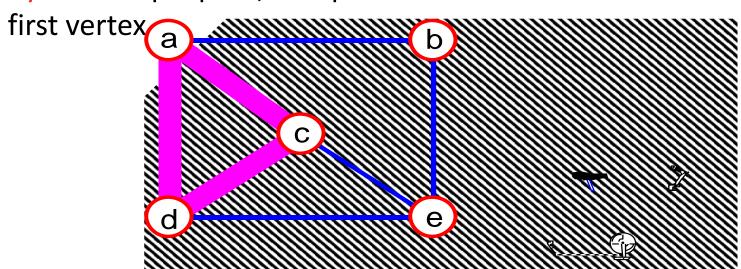


More Terminology

Simple path: no repeated vertices



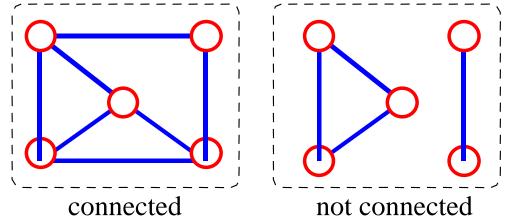
• Cycle: simple path, except that the last vertex is the same as the



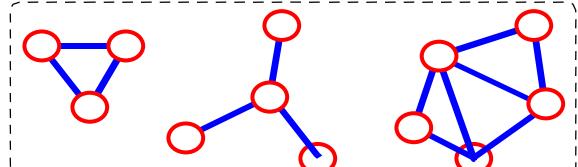


Even More Terminology

•Connected graph: any two vertices are connected by some path

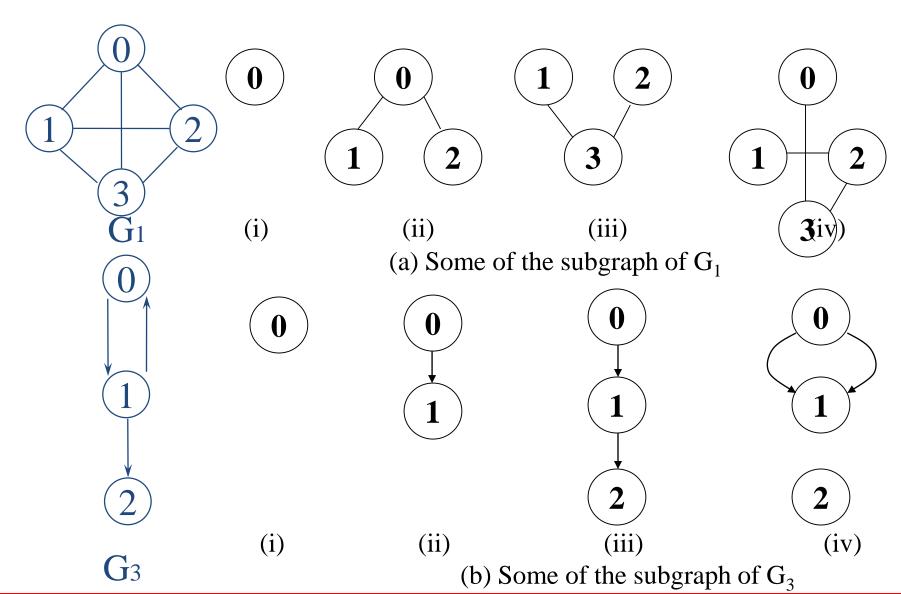


- Subgraph: subset of vertices and edges forming a graph
- Connected component: maximal connected subgraph. E.g., the graph below has 3 connected components.





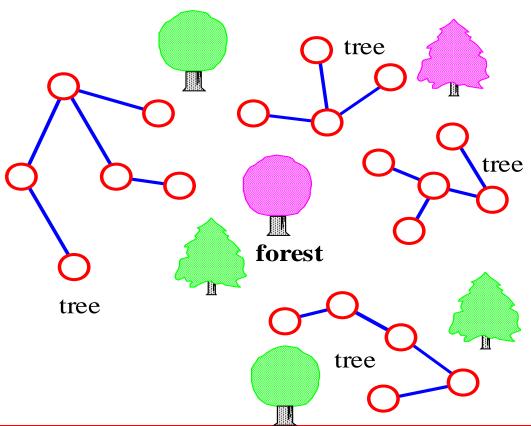
Subgraphs Examples





More...

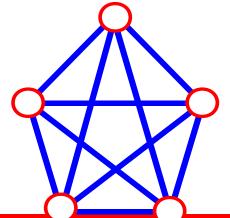
- tree connected graph without cycles
- forest collection of trees





Connectivity

- Let n = #vertices, and m = #edges
- A complete graph: one in which all pairs of vertices are adjacent
- How many total edges in a complete graph?
 - Each of the n vertices is incident to n-1 edges, however, we would have counted each edge twice! Therefore, intuitively, m = n(n-1)/2.
- Therefore, if a graph is not complete, m < n(n -1)/2



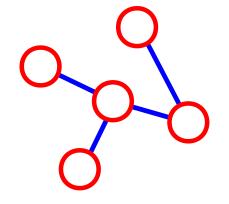
$$n = 5$$

 $m = (5 * 4)/2 = 10$



More Connectivity

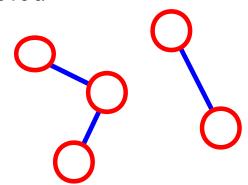
• For a tree m = n - 1



$$\mathbf{n} = 5$$

$$\mathbf{m} = 4$$

If
$$m < n - 1$$
, G is not connected



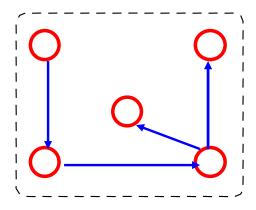
$$\mathbf{n} = 5$$

$$\mathbf{m} = 3$$



Oriented (Directed) Graph

A graph where edges are directed





Directed vs. Undirected Graph

- An undirected graph is one in which the pair of vertices in a edge is unordered, $(v_0, v_1) = (v_1, v_0)$
- A directed graph is one in which each edge is a directed pair of vertices, <v₀, v₁> != <v₁,v₀>

tail head

Graph Representation

- For graphs to be computationally useful, they have to be conveniently represented in programs
- There are two computer representations of graphs:
 - Adjacency matrix representation
 - Adjacency lists representation



Graph Representations

- For graphs to be computationally useful, they have to be conveniently represented in programs
- There are two computer representations of graphs:
 - Adjacency Matrix
 - Adjacency Lists

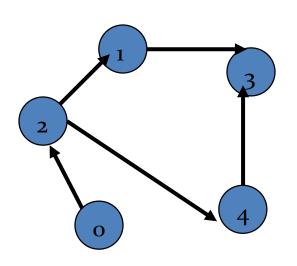


Adjacency Matrix

- A square grid of boolean values
- If the graph contains N vertices, then the grid contains N rows and N columns
- For two vertices numbered I and J, the element at row I and column J is true if there is an edge from I to J, otherwise false

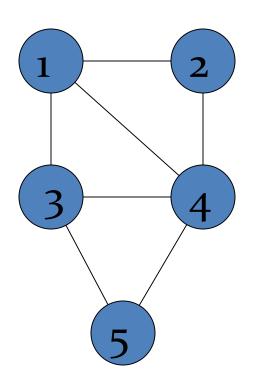


Adjacency Matrix



0	1	2	3	4
false	false	true	false	false
false	false	false	true	false
false	true	false	false	true
false	false	false	false	false
false	false	false	true	false
\				

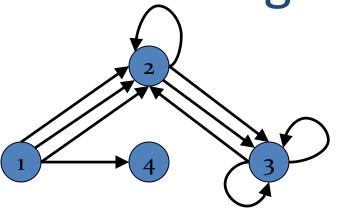
Adjacency Matrix



	1	2	3	4	5
1	0	1	1	1	0
2	1	0	0	1	0
3	1	0	0	1	1
4	1	1	1	0	1
5	0	0	1	1	0



Adjacency Matrix -Directed Multigraphs



A:

$$\begin{pmatrix} 0 & 3 & 0 & 1 \\ 0 & 1 & 2 & 0 \\ 0 & 1 & 2 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$



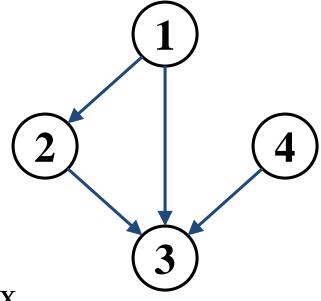
Adjacency Lists Representation

- A graph of n nodes is represented by a onedimensional array L of linked lists, where
 - L[i] is the linked list containing all the nodes adjacent from node i.
 - The nodes in the list L[i] are in no particular order



Graphs: Adjacency List

- Adjacency list: for each vertex v ∈ V, store a list of vertices adjacent to v
- Example:
 - $Adj[1] = \{2,3\}$
 - $Adj[2] = \{3\}$
 - $Adj[3] = \{\}$
 - $Adj[4] = \{3\}$
- Variation: can also keep
 a list of edges coming *into* vertex

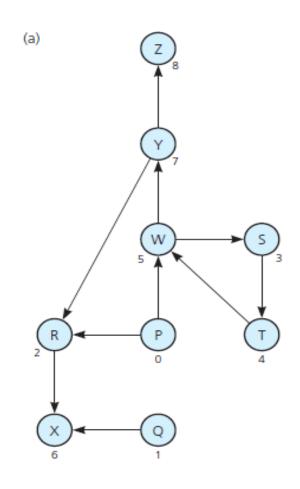


Graphs: Adjacency List

- How much storage is required?
 - The *degree* of a vertex v = # incident edges
 - Directed graphs have in-degree, out-degree
 - For directed graphs, # of items in adjacency lists is Σ out-degree(v) = |E| For undirected graphs, # items in adjacency lists is Σ degree(v) = 2 |E|
- So: Adjacency lists take O(V+E) storage



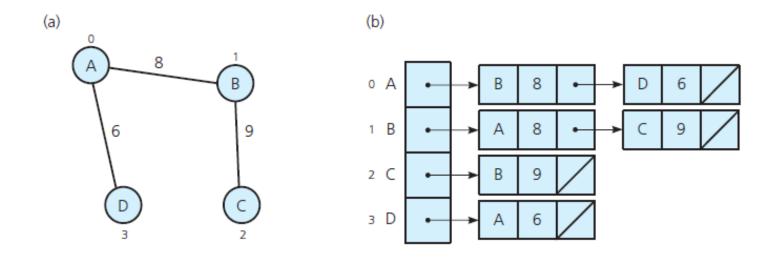
Implementing Graphs



(b)		0	1	2	3	4	5	6	7	8
		Р	Q	R	S	Т	W	X	Υ	Z
0	Р	0	0	1	0	0	1	0	0	0
1	Q	0	0	0	0	0	0	1	0	0
2	R	0	0	0	0	0	0	1	0	0
3	S	0	0	0	0	1	0	0	0	0
4	Т	0	0	0	0	0	1	0	0	0
5	W	0	0	0	1	0	0	0	1	0
6	Х	0	0	0	0	0	0	0	0	0
7	Υ	0	0	1	0	0	0	0	0	1
8	Z	0	0	0	0	0	0	0	0	0



Implementing Graphs



- (a) A weighted undirected graph and
 - (b) its adjacency list



Data structure for Adjacency list

```
// A structure to represent an adjacency list node
struct AdjListNode
    int dest:
    struct AdjListNode* next;
1:
// A structure to represent an adjacency list
struct AdjList
    struct AdjListNode *head; // pointer to head node of list
1:
// A structure to represent a graph. A graph is an array of adjacency lists.
// Size of array will be V (number of vertices in graph)
struct Graph
    int V:
    struct AdjList* array;
};
```



```
// A utility function to create a new adjacency list node
struct AdjListNode* newAdjListNode(int dest)
    struct AdiListNode* newNode =
             (struct AdjListNode*) malloc(sizeof(struct AdjListNode));
    newNode->dest = dest:
    newNode->next = NULL:
    return newNode:
// A utility function that creates a graph of V vertices
struct Graph* createGraph(int V)
    struct Graph* graph = (struct Graph*) malloc(sizeof(struct Graph));
   qraph->V = V;
    // Create an array of adjacency lists. Size of array will be V
   graph->array = (struct AdjList*) malloc(V * sizeof(struct AdjList));
     // Initialize each adjacency list as empty by making head as NULL
   int i:
    for (i = 0; i < V; ++i)
       graph->array[i].head = NULL;
     return graph;
```



Complete the Code for addEdge Function

```
// Adds an edge to an undirected graph
void addEdge(struct Graph* graph, int src, int dest)
{
    1. Add an edge from src to dest. A new node is added to the adjacency list of src. The node is added at the begining
    write the code necessary for adding an edge from src to dest
    2. Since graph is undirected, add an edge from dest to src also
    write the necessary code to connect dest to src
}
```



Experimental Procedure

- Analyse the problem statement
- Design an algorithm for the given problem statement and develop a flowchart/pseudo-code
- Implement the algorithm in C language
- Compile the C program
- Design test cases and test the implemented program
- Document the Results
- Analyse and discuss the outcomes of your experiment



Exercise

Design an algorithm and write a program to create a hash table and use suitable
hash function and demonstrate search operation for a given data. Tabulate the
output for various inputs and verify against expected values. Analyse the efficiency
of hash function used. Describe your learning along with the limitations of both, if
any. Suggest how these can be overcome.



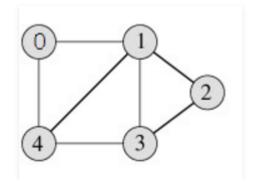
Key factors and discussion

- Implement adjacency matrix representation of graph
- Implement adjacency list representation of graph
 - Complete the code for addEdge function and write print graph function
 - Compare the efficiency of adjacency matrix and list representation of graph
 Collison resolution



Expected output for Adjacency Matrix

Input Graph



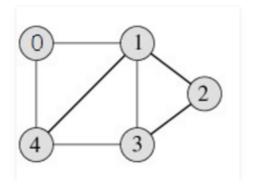
Output

	0	1	2	3	4
0	0	1	0	0	1
1	1	0	1	1	1
2	0	1	0	1	0
3	0	1	1	0	1
4	1	1	0	1	0



Expected output for Adjacency List

Input Graph



Output

```
Adjacency list of vertex 0
head -> 4-> 1
Adjacency list of vertex 1
head -> 4-> 3-> 2-> 0
Adjacency list of vertex 2
head -> 3-> 1
Adjacency list of vertex 3
head -> 4-> 2-> 1
Adjacency list of vertex 4
head -> 3-> 1-> 0
```



Results and Presentations

- Calculations/Computations/Algorithms
 The calculations/computations/algorithms involved in each program has to be presented
- Presentation of Results
 The results for all the valid and invalid cases have to be presented
- Analysis and Discussions
 how the data is manipulated or transformed, what are the
 key operations involved. Errors encounters and how they are
 resolved.
- Conclusions



Comments

- Limitations of Experiments
 Outline the loopholes in the program, data structures or solution approach.
- Limitations of Results
 Present the test cases; justify if the program is tested correctly considering all the outcomes. Mention what is not tested, if any.
- Learning happened
 What is the overall learning happened
- ConclusionsSummary



References

• Gilberg, R. F., and Forouzan, B. A. (2007): A Pseudocode Approach With C, 2nd edn. Cengage Learning

