

# Overview

- Register Transfer Language
- Register Transfer
- Bus and Memory Transfers
- Arithmetic Micro-operations
- **Logic Micro-operations**
- **Shift Micro-operations**
- **Arithmetic Logic Shift Unit**

# Logic Micro operations

## ◆ Logic microoperation

- Logic microoperations consider *each bit of the register separately* and treat them as binary variables

» *exam)*

$$P: R1 \leftarrow R1 \oplus R2$$

$$\begin{array}{r} 1010 \text{ Content of R1} \\ + 1100 \text{ Content of R2} \\ \hline 0110 \text{ Content of R1 after P=1} \end{array}$$

- Special Symbols

» Special symbols will be adopted for the logic microoperations **OR**( $\vee$ ), **AND**( $\wedge$ ), and **complement**(*a bar on top*), to distinguish them from the corresponding symbols used to express Boolean functions

» *exam)*

$$P+Q: R1 \leftarrow R2 + R3, R4 \leftarrow R5 \vee R6$$

Logic OR

Arithmetic ADD

## ◆ List of Logic Microoperation

- Truth Table for 16 functions for 2 variables : **Tab. 4-5**
- 16 Logic Microoperation : **Tab. 4-6**

∴ All other Operation can be derived

## ◆ Hardware Implementation

- 16 microoperation → Use only 4(AND, OR, XOR, Complement)
- One stage of logic circuit

# Logic Microoperations

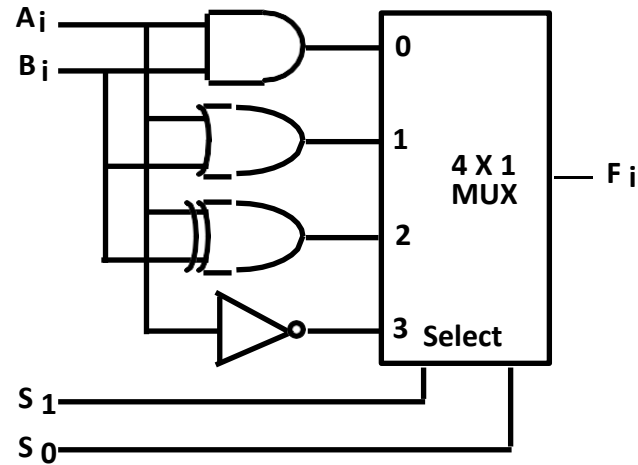
| X | Y | F <sub>0</sub> | F <sub>1</sub> | F <sub>2</sub> | F <sub>3</sub> | F <sub>4</sub> | F <sub>5</sub> | F <sub>6</sub> | F <sub>7</sub> | F <sub>8</sub> | F <sub>9</sub> | F <sub>10</sub> | F <sub>11</sub> | F <sub>12</sub> | F <sub>13</sub> | F <sub>14</sub> | F <sub>15</sub> |
|---|---|----------------|----------------|----------------|----------------|----------------|----------------|----------------|----------------|----------------|----------------|-----------------|-----------------|-----------------|-----------------|-----------------|-----------------|
| 0 | 0 | 0              | 0              | 0              | 0              | 0              | 0              | 0              | 0              | 1              | 1              | 1               | 1               | 1               | 1               | 1               | 1               |
| 0 | 1 | 0              | 0              | 0              | 0              | 1              | 1              | 1              | 1              | 0              | 0              | 0               | 0               | 1               | 1               | 1               | 1               |
| 1 | 0 | 0              | 0              | 1              | 1              | 0              | 0              | 1              | 1              | 0              | 0              | 1               | 1               | 0               | 0               | 1               | 1               |
| 1 | 1 | 0              | 1              | 0              | 1              | 0              | 1              | 0              | 1              | 0              | 1              | 0               | 1               | 0               | 1               | 0               | 1               |

TABLE 4-5. Truth Table for 16 Functions of Two Variables

| Boolean function   | Microoperation                       | Name       | Boolean function      | Microoperation                       | Name           |
|--------------------|--------------------------------------|------------|-----------------------|--------------------------------------|----------------|
| $F_0 = 0$          | $F \leftarrow 0$                     | Clear      | $F_8 = (x+y)'$        | $F \leftarrow \overline{A \vee B}$   | NOR            |
| $F_1 = xy$         | $F \leftarrow A \wedge B$            | AND        | $F_9 = (x \oplus y)'$ | $F \leftarrow \overline{A \oplus B}$ | Ex-NOR         |
| $F_2 = xy'$        | $F \leftarrow A \wedge \overline{B}$ |            | $F_{10} = y'$         | $F \leftarrow \overline{B}$          | Compl-B        |
| $F_3 = x$          | $F \leftarrow A$                     | Transfer A | $F_{11} = x+y'$       | $F \leftarrow A \vee \overline{B}$   |                |
| $F_4 = x'y$        | $F \leftarrow \overline{A} \wedge B$ |            | $F_{12} = x'$         | $F \leftarrow \overline{A}$          | Compl-A        |
| $F_5 = y$          | $F \leftarrow B$                     | Transfer B | $F_{13} = x'+y$       | $F \leftarrow \overline{A} \vee B$   |                |
| $F_6 = x \oplus y$ | $F \leftarrow A \oplus B$            | Ex-OR      | $F_{14} = (xy)'$      | $F \leftarrow \overline{A \wedge B}$ | NAND           |
| $F_7 = x+y$        | $F \leftarrow A \vee B$              | OR         | $F_{15} = 1$          | $F \leftarrow \text{all 1's}$        | set to all 1's |

TABLE 4-6. Sixteen Logic Microoperations

# Hardware Implementation



Function table

| $S_1$ | $S_0$ | Output           | $\mu$ -operation |
|-------|-------|------------------|------------------|
| 0     | 0     | $F = A \wedge B$ | AND              |
| 0     | 1     | $F = A \vee B$   | OR               |
| 1     | 0     | $F = A \oplus B$ | XOR              |
| 1     | 1     | $F = A'$         | Complement       |

# Applications of Logic Microoperations

- Logic microoperations can be used to manipulate individual bits or a portions of a word in a register
- Consider the data in a register A. In another register, B, is bit data that will be used to modify the contents of A

- |                        |                                  |
|------------------------|----------------------------------|
| ➤ Selective-set        | $A \leftarrow A + B$             |
| ➤ Selective-complement | $A \leftarrow A \div B$          |
| ➤ Selective-clear      | $A \leftarrow A \bullet B'$      |
| ➤ Mask (Delete)        | $A \leftarrow A \bullet B$       |
| ➤ Clear                | $A \leftarrow A \div B$          |
| ➤ Insert               | $A \leftarrow (A \bullet B) + C$ |
| ➤ Compare              | $A \leftarrow A \div B$          |

# Applications of Logic Microoperations

1. In a selective set operation, the bit pattern in B is used to *set* certain bits in A

$$\begin{array}{rcl}
 1100 & A_t & \\
 1010 & B & \\
 1110 & A_{t+1} & (A \leftarrow A + B)
 \end{array}$$

If a bit in B is set to 1, that same position in A gets set to 1, otherwise that bit in A keeps its previous value.

2. In a selective complement operation, the bit pattern in B is used to *complement* certain bits in A

$$\begin{array}{rcl}
 1100 & A_t & \\
 1010 & B & \\
 0110 & A_{t+1} & (A \leftarrow A \div B)
 \end{array}$$

If a bit in B is set to 1, that same position in A gets complemented from its original value, otherwise it is unchanged

# Applications of Logic Microoperations

3. In a selective clear operation, the bit pattern in B is used to *clear* certain bits in A

1 1 0 0  $A_t$

1 0 1 0 B

0 1 0 0  $A_{t+1}$  ( $A \leftarrow A \cdot B'$ )

If a bit in B is set to 1, that same position in A gets set to 0, otherwise it is unchanged

4. In a mask operation, the bit pattern in B is used to *clear* certain bits in A

1 1 0 0  $A_t$

1 0 1 0 B

1 0 0 0  $A_{t+1}$  ( $A \leftarrow A \cdot B$ )

If a bit in B is set to 0, that same position in A gets set to 0, otherwise it is unchanged

# Applications of Logic Microoperations

5. In a clear operation, if the bits in the same position in A and B are the same, they are cleared in A, otherwise they are set in A

1 1 0 0  $A_t$

1 0 1 0 B

0 1 1 0  $A_{t+1}$        $(A \leftarrow A \div B)$



# Applications of Logic Microoperations

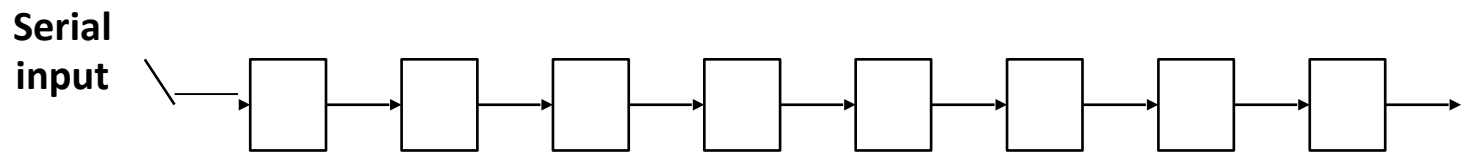
6. An insert operation is used to introduce a specific bit pattern into A register, leaving the other bit positions unchanged

This is done as

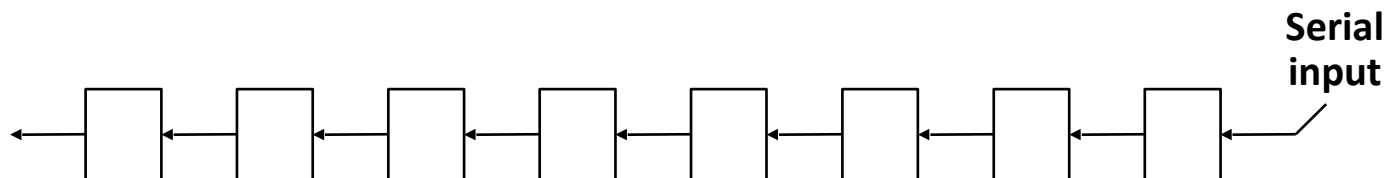
- A mask operation to clear the desired bit positions, followed by
- An OR operation to introduce the new bits into the desired positions
- Example
  - Suppose you wanted to introduce 1010 into the low order four bits of A:
  - 1101 1000 1011 0001       A (Original)
  - 1101 1000 1011 **1010**       A (Desired)
  
  - 1101 1000 1011 0001                                   A (Original)
  - 1111 1111 1111 0000                               Mask
  - 1101 1000 1011 0000                               A (Intermediate)
  - 0000 0000 0000 **1010**                             Added bits
  - 1101 1000 1011 **1010**                             A (Desired)

# Shift Microoperations

- There are three types of shifts
  - *Logical shift*
  - *Circular shift*
  - *Arithmetic shift*
- What differentiates them is the information that goes into the serial input
- A right shift operation

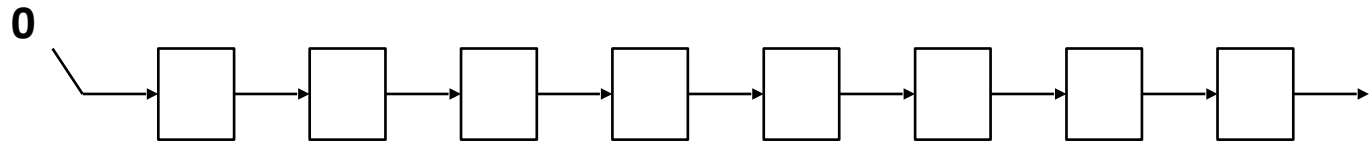


- A left shift operation

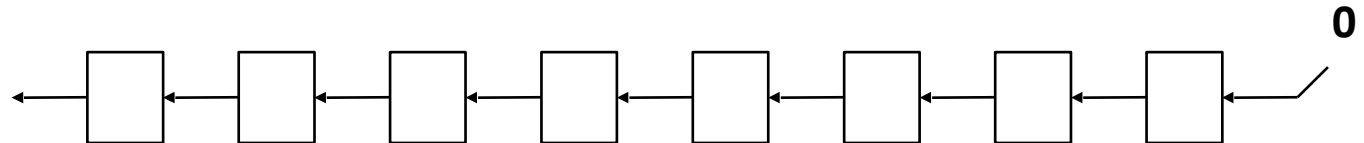


# Logical Shift

- In a logical shift the serial input to the shift is a 0.
- A right logical shift operation:



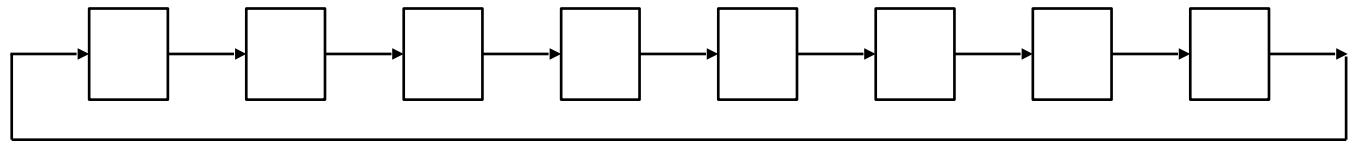
- A left logical shift operation:



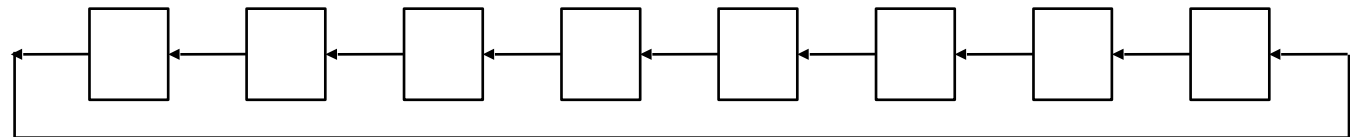
- In a Register Transfer Language, the following notation is used
  - *shl* for a logical shift left
  - *shr* for a logical shift right
  - Examples:
    - $R2 \leftarrow shr\ R2$
    - $R3 \leftarrow shl\ R3$

# Circular Shift

- In a circular shift the serial input is the bit that is shifted out of the other end of the register.
- A right circular shift operation:



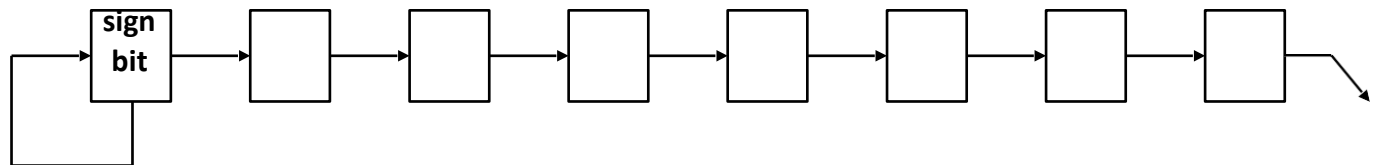
- A left circular shift operation:



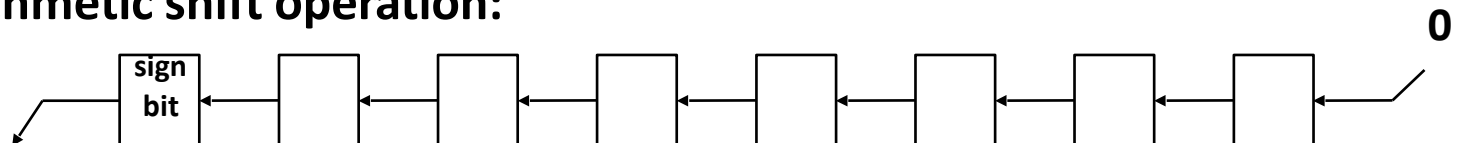
- In a RTL, the following notation is used
  - *cil* for a circular shift left
  - *cir* for a circular shift right
  - Examples:
    - $R2 \leftarrow cir\ R2$
    - $R3 \leftarrow cil\ R3$

# Arithmetic Shift

- An arithmetic shift is meant for signed binary numbers (integer)
- An arithmetic left shift **multiplies** a signed number **by two**
- An arithmetic right shift **divides** a signed number **by two**
- Sign bit : 0 for positive and 1 for negative
- The main distinction of an arithmetic shift is that it must keep the sign of the number the same as it performs the multiplication or division
- A right arithmetic shift operation:

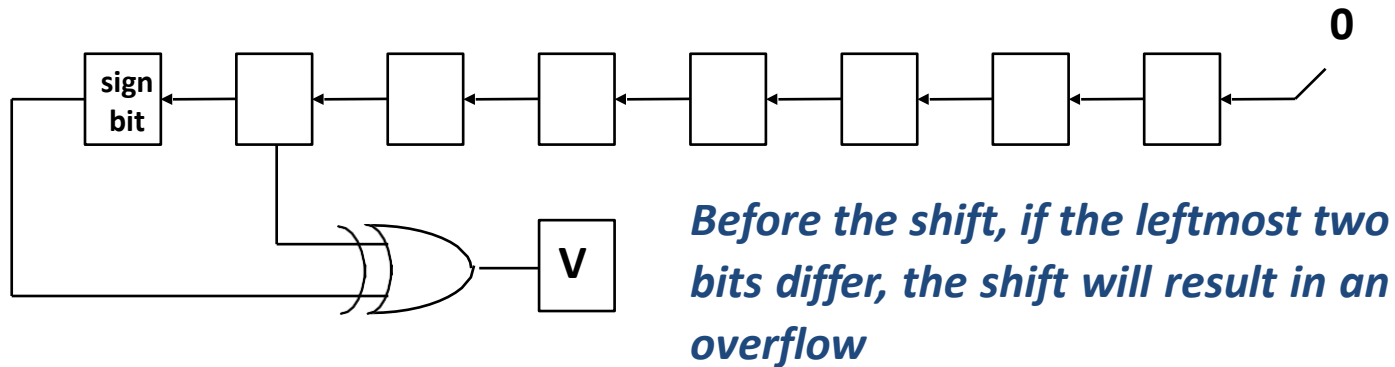


- A left arithmetic shift operation:



# Arithmetic Shift

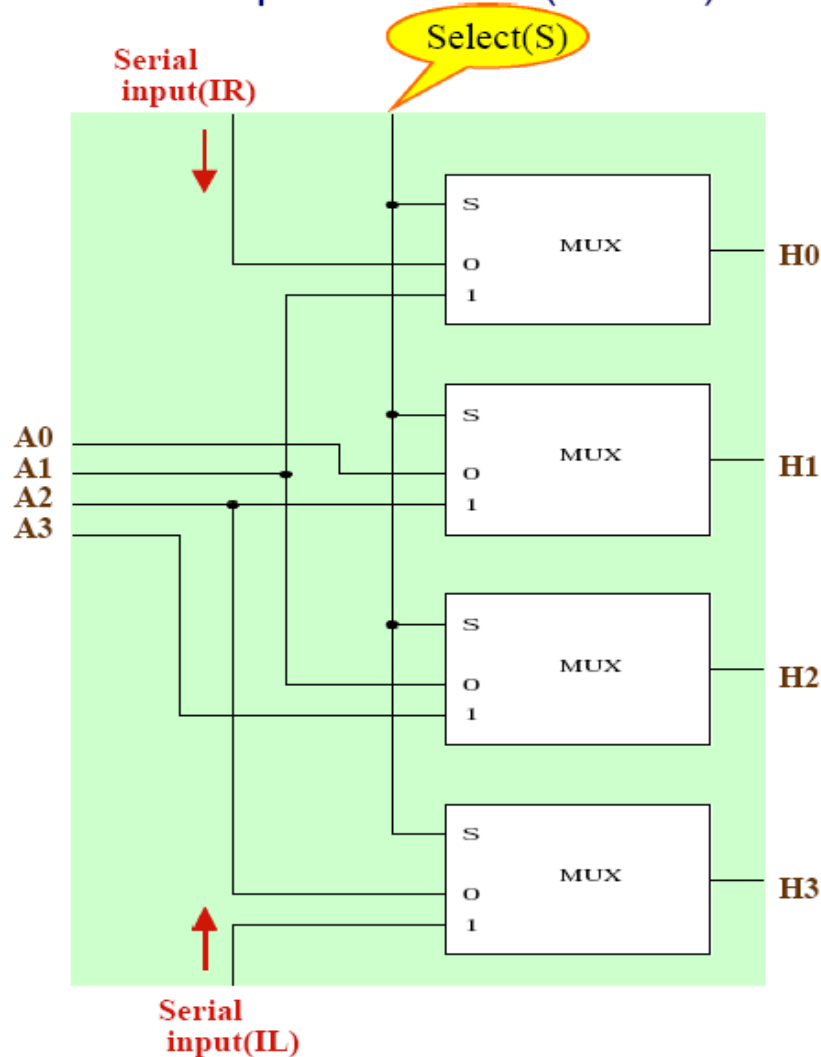
- An left arithmetic shift operation must be checked for the overflow



- In a RTL, the following notation is used
  - *ashl* for an arithmetic shift left
  - *ashr* for an arithmetic shift right
  - Examples:
    - »  $R2 \leftarrow ashr R2$
    - »  $R3 \leftarrow ashl R3$

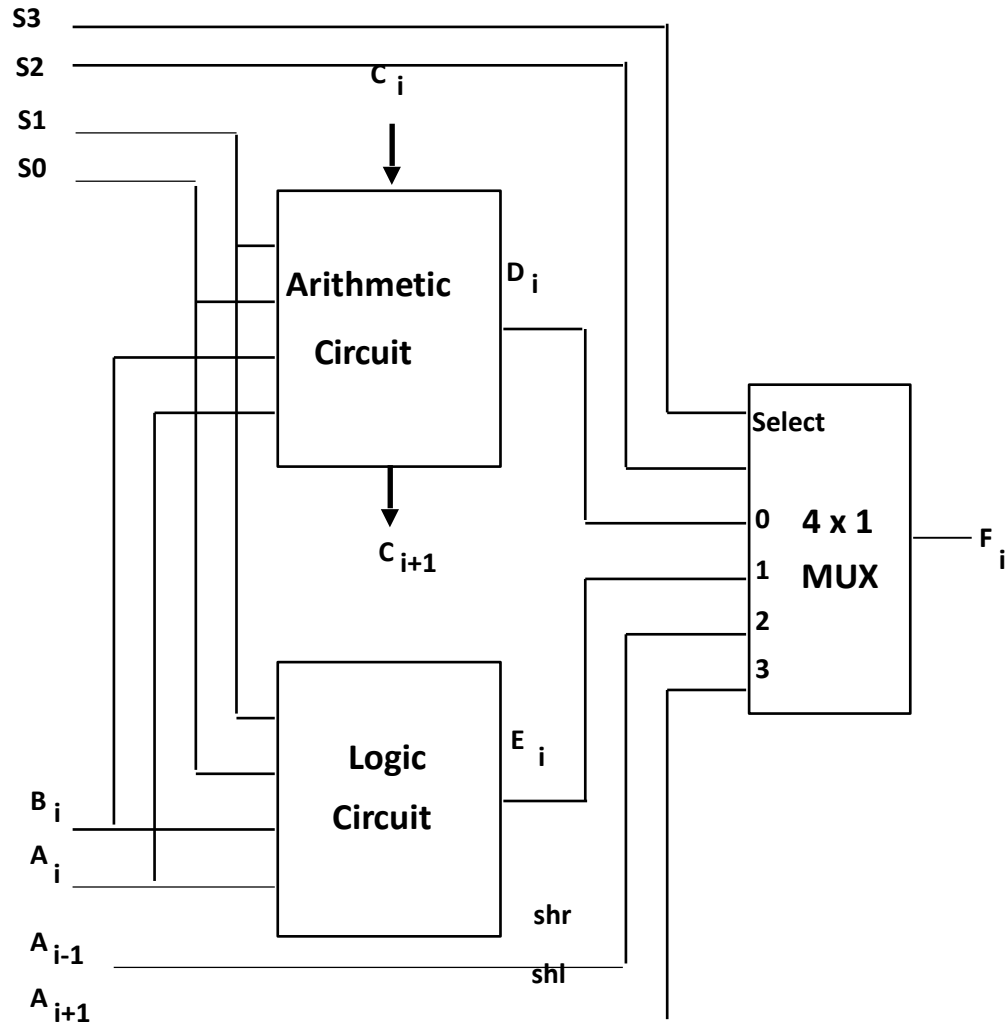
# Hardware Implementation of Shift Microoperation

## ◆ Hardware Implementation(Shifter) :



| Function Table |        |    |    |    |
|----------------|--------|----|----|----|
| Select         | output |    |    |    |
| S              | H0     | H1 | H2 | H3 |
| 0              | IR     | A0 | A1 | A2 |
| 1              | A1     | A2 | A3 | IL |

# Arithmetic Logic and Shift Unit



| s2 | s3 | Operation  |
|----|----|------------|
| 0  | 0  | Arithmetic |
| 0  | 1  | Logical    |
| 1  | 0  | Shr        |
| 1  | 1  | shl        |



**TABLE 4-8** Function Table for Arithmetic Logic Shift Unit

| Operation select |       |          |          |          | Operation                  | Function                 |
|------------------|-------|----------|----------|----------|----------------------------|--------------------------|
| $S_3$            | $S_2$ | $S_1$    | $S_0$    | $C_{in}$ |                            |                          |
| 0                | 0     | 0        | 0        | 0        | $F = A$                    | Transfer $A$             |
| 0                | 0     | 0        | 0        | 1        | $F = A + 1$                | Increment $A$            |
| 0                | 0     | 0        | 1        | 0        | $F = A + B$                | Addition                 |
| 0                | 0     | 0        | 1        | 1        | $F = A + B + 1$            | Add with carry           |
| 0                | 0     | 1        | 0        | 0        | $F = A + \overline{B}$     | Subtract with borrow     |
| 0                | 0     | 1        | 0        | 1        | $F = A + \overline{B} + 1$ | Subtraction              |
| 0                | 0     | 1        | 1        | 0        | $F = A - 1$                | Decrement $A$            |
| 0                | 0     | 1        | 1        | 1        | $F = A$                    | Transfer $A$             |
| 0                | 1     | 0        | 0        | $\times$ | $F = A \wedge B$           | AND                      |
| 0                | 1     | 0        | 1        | $\times$ | $F = A \vee B$             | OR                       |
| 0                | 1     | 1        | 0        | $\times$ | $F = A \oplus B$           | XOR                      |
| 0                | 1     | 1        | 1        | $\times$ | $F = \overline{A}$         | Complement $A$           |
| 1                | 0     | $\times$ | $\times$ | $\times$ | $F = \text{shr } A$        | Shift right $A$ into $F$ |
| 1                | 1     | $\times$ | $\times$ | $\times$ | $F = \text{shl } A$        | Shift left $A$ into $F$  |