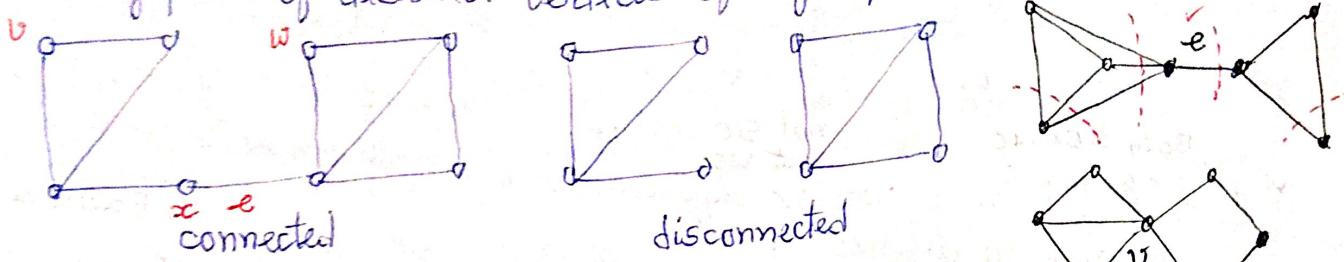


Connectedness in Undirected Graphs:

DEF. An undirected graph is called connected if there is a path between every pair of distinct vertices of the graph. Otherwise disconnected.



A disconnected graph consists of two or more connected graphs.

Each of these connected subgraphs is called a component.

Theorem: A graph is disconnected if and only if its vertex set V can be partitioned into two nonempty subsets V_1 and V_2 such that there exists no edge in G whose one end vertex is in V_1 and the other in V_2 .

Theorem: If a graph (connected or disconnected) has exactly two vertices of odd degree, there must be a path joining these two vertices.

Theorem: If G is a bipartite graph, then each cycle of G has even length.

Theorem: Let G be a simple graph on n vertices. If G has k components, then the number m of edges of satisfies

$$(n-k) \leq m \leq (n-k)(n-k+1)/2$$

Corollary: Any simple graph with n vertices and more than $(n-1)(n-2)/2$ edges is connected.

Q. Determine $k(G)$ and $\lambda(G)$ for each of : (i) C_6 , (ii) W_6 , (iii) $K_{4,7}$ (iv) Q_4 .

A disconnecting set in a connected graph is a set of edges whose removal disconnects G .

A disconnecting set, no proper subset of which is a disconnecting set is called a cutset. A cutset with only one edge is called a bridge.

Edge connectivity, $\lambda(G) = \text{size of the smallest cutset in } G$. If $\lambda(G) \geq k$, G is k -edge-connected.

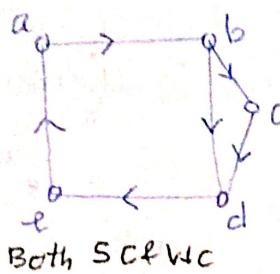
A separating set in a connected graph G is a set of vertices whose deletion disconnects G . A separating set with only one vertex is called a CUT-VERTEX. (Vertex) connectivity $k(G) = \text{size of smallest separating set in } G$. G is k -connected if $k(G) \geq k$.

Connectedness in Directed Graphs:

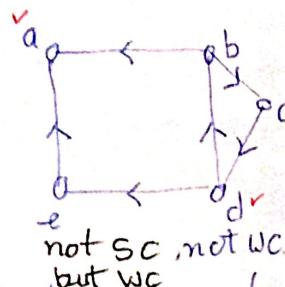
DEF. A directed graph is strongly connected if there is a path from a to b and from b to a whenever a and b are vertices in the graph.

DEF. A directed graph is weakly connected if there is a path between every two vertices in the underlying undirected graph.

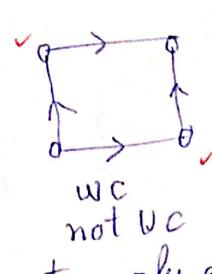
DEF. A digraph is unilaterally connected if for every pair of vertices one is reachable from the other.



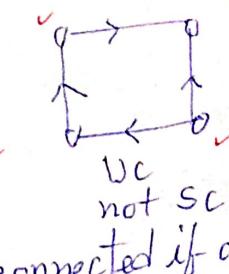
Both SC & WC



not SC, not WC
but WC



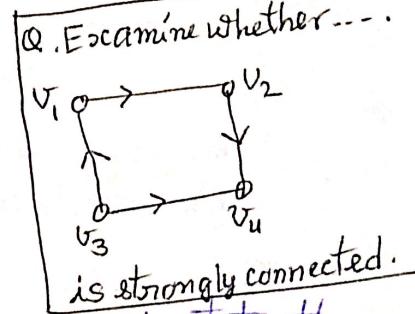
WC
not UC



UC
not SC

✓ THEOREM: An n -vertex digraph is strongly connected if and only if the matrix m defined by $m = A + A^2 + A^3 + \dots + A^n$

has no zero entry, A is the adjacency matrix.

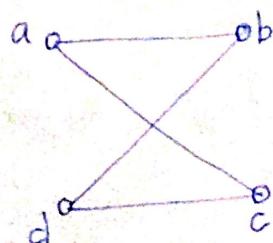


is strongly connected.

Counting Paths Between Vertices:

Theorem: Let G be a graph with adjacency matrix A with respect to the ordering $1, 2, 3, \dots, n$ (with directed or undirected edges, with multiple edges and loops allowed). The number of different paths of length r from i to j , where r is a positive integer, equals the (i, j) th entry of A^r .

Ex. Determine the number of paths of length four from a to d in the simple graph:



$$A = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}; \quad A^2 = \begin{bmatrix} 0 & 1 & 2 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 2 \\ 1 & 1 & 1 & 0 \end{bmatrix}; \quad A^3 = \begin{bmatrix} 0 & 1 & 3 & 2 \\ 1 & 0 & 2 & 2 \\ 1 & 1 & 1 & 1 \\ 2 & 2 & 1 & 0 \end{bmatrix}; \quad A^4 = \begin{bmatrix} 0 & 1 & 4 & 3 \\ 1 & 0 & 3 & 3 \\ 2 & 1 & 2 & 1 \\ 3 & 2 & 1 & 0 \end{bmatrix}$$

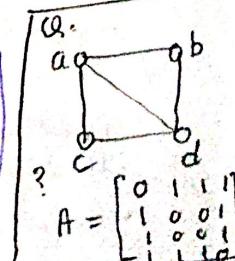
The number of paths of length four from a to d is the $(1, 4)$ th entry of A^4 :

$$? X = A + A^2 + A^3 = \begin{bmatrix} 0 & 1 & 2 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 2 \\ 1 & 2 & 1 & 0 \end{bmatrix}$$

✓ THEOREM: If A is the adjacency matrix of an undirected graph G with n vertices, and

$$X = A + A^2 + A^3 + \dots + A^{n-1}$$

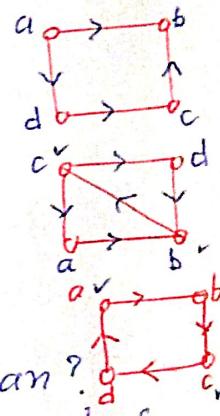
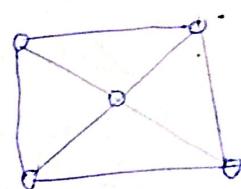
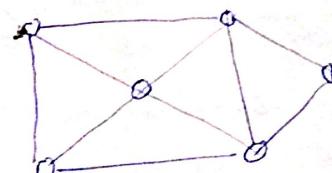
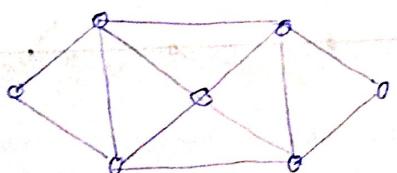
Then, G is disconnected if and only if there exists at least one entry in matrix X that is zero.



$$A = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

EULER GRAPHS

- A closed walk in a graph that contains every edge of the graph exactly once is called an Euler line/Euler circuit, and a graph that consists of an Euler is called an Euler graph.
- A open walk in a graph that includes (or traces or covers) all edges of the graph without retracing any edge is called a unicursal line or an open Euler or an Euler path.
- A (connected) graph that has a unicursal line will be termed/called a unicursal graph or semi-Euler graph.
- A graph that has neither Euler line nor unicursal line is called non-Euler graph.



Q. Discuss Königsberg bridges problem.

Q. Which of the following graphs are Eulerian? semi-Eulerian?
 (i) K_5 (ii) $K_{2,3}$ (iii) the graph of the cube (iv) the graph of the octahedron (v) the Petersen graph.

Q. Examine each of the following for an Euler graph.

(i) K_n (ii) $K_{m,n}$ (iii) λn (iv) Q_k (v) Platonic graphs.

Theorem: A connected graph G is an Euler graph if and only if all vertices of G are of even degree.

Theorem: A connected graph G is an Euler graph if and only if it can be decomposed into circuits/its set of edges can be split up into disjoint cycles.

Corollary: A connected graph is semi-Eulerian if and only if it has exactly two vertices of odd degree.

Theorem: In a connected graph G with exactly $2k$ odd vertices, there exist k edge-disjoint subgraphs such that they together contain all edges of G and that each is a unicursal graph.

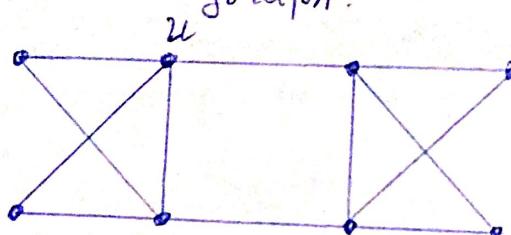
Fleury's Algorithm :

Theorem: Let G be an Eulerian graph. Then the following construction is always possible, and produces an Eulerian line of G .

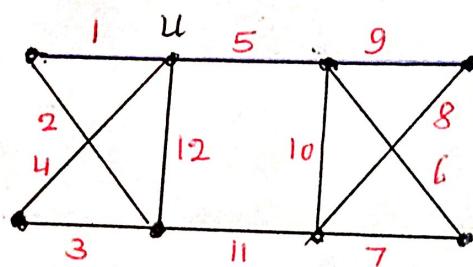
Start at any vertex u and traverse the edges in an arbitrary manner, subject only to the following rules:

- erase the edges as they are traversed, and if any isolated vertices result, erase them too;
- at each stage, use a bridge only if there is no alternative.

EX: Use Fleury's algorithm to produce an Eulerian trail/line/circuit for the graph.

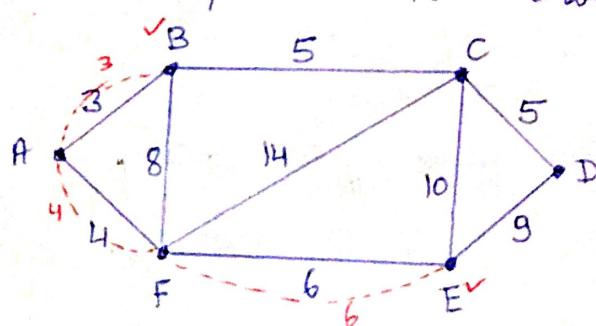


solution



Solution: There are many possible solutions; for example, traverse the edges in the order indicated by the adjoining diagram.

EX: Solve the Chinese postman for the weighted graph:



E.X. Find a closed walk in a graph G that passes through every vertex exactly once.
Find an open walk in a graph G that passes through every vertex exactly once.

