Java Transaction API (JTA)

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Database Transaction

Atomicity

Each transaction be "aller nothing"

Consistency

Bringing the database from one valid state to another

Isolation

 The concurrent execution of transactions results in a system state that would be obtained if transactions were executed serially

Durability

 Once a transaction has been committed, it will remain so, even in the event of power loss, crashes, or errors

Distributed Transaction



- A transaction in which two or more network hosts are involved
 - Hosts provide transactional resources
 - Transaction manager responsible for creating and managing a global transaction
- Two-phase commit (2PC): usually applied for updates able to commit in a short period of time (from milliseconds to minutes)
- Long-lived distributed transactions: more sophisticated techniques that involve multiple undo levels are used, practically implemented in systems based on Web Services

Two-Phase Commit Protocol (2PC)

- A distributed algorithm that coordinates all the processes that participate in a distributed atomic transaction on whether to commit or abort (roll back) the transaction
- Assumptions
 - One node is designated the coordinator, while the rest of nodes are designated the cohorts
 - There is stable storage at each node with a writeahead log
 - No node crashes forever
 - The data in the write-ahead log is never lost or corrupted in a crash
 - Any two nodes can communicate with each other

2PC's Basic Algorithm



- Commit request phase (voting phase)
 - 1. The coordinator sends a query to commit message to all cohorts and waits until it has received a reply from all cohorts
 - 2. Cohorts execute the transaction up to the point where they will be asked to commit; they each write an entry to their undo log and an entry to their redo log
 - 3. Each cohort replies with an AGREEMENT message (voting YES, to commit) if the cohort's actions succeeded, or an ABORT message (voting NO, not to commit) if the cohort experiences a failure that will make it impossible to commit

2PC's Basic Algorithm (cont.)

- Commit phase (completion phase)
 - Success: if the coordinator received an AGREEMENT message from all cohorts during the commit-request phase
 - 1. The coordinator sends a COMMIT message to all the cohorts
 - Each cohort completes the operation, and releases all the locks and resources held during the transaction
 - 3. Each cohort sends an ACKNOWLEDGMENT to the coordinator
 - The coordinator completes the transaction when all acknowledgments have been received
 - Failure: if any cohort votes NO during the commit-request phase (or the coordinator's timeout expires):
 - 1. The coordinator sends a ROLLBACK message to all the cohorts
 - 2. Each cohort undoes the transaction using the undo log, and releases the resources and locks held during the transaction
 - 3. Each cohort sends an ACKNOWLEDGMENT to the coordinator
 - 4. The coordinator undoes the transaction when all acknowledgements have been received

X/Open XA



- A specification by the Open Group for distributed transaction processing
- Specifying how a transaction manager will roll up the transactions against the different data-stores into an "atomic" transaction and execute this with the 2PC protocol for the transaction
- A type of transaction coordination, allowing many resources to participate in a single, coordinated, atomic update operation
- The XA specification describes what a resource manager must do to support transactional access, and resource managers that follow this specification are said to be XA-compliant

XA Transactions

- Applications that use global transactions involve one or more Resource Managers and a Transaction Manager
 - A Resource Manager (RM) provides access to transactional resources.
 - A Transaction Manager (TM) coordinates the transactions that are part of a global transaction
- The individual transactions within a global transaction are "branches" of the global transaction

XA Transactions (cont.)

- To carry out a global transaction, it is necessary to bring each component to a point when it can be committed or rolled back
- Depending on what each component reports about its ability to succeed, they must all commit or roll back as an atomic group
- The process for executing a global transaction
 - In the first phase, all branches are prepared
 - In the second phase, the TM tells the RMs whether to commit or roll back
- A global transaction might use one-phase commit (1PC), when a TM finds that a global transaction consists of only one transactional resource

Java Transaction API (JTA)

- Enabling distributed transactions to be done across multiple X/Open XA resources in a Java environment
- Architecture
 - A transaction manager or transaction processing monitor (TP monitor) coordinates the transactions across multiple resources such as databases and message queues
 - Each resource has its own resource manager, which typically has its own API for manipulating the resource
 - The resource manager allows a TP monitor to coordinate a distributed transaction between its own and other resource managers
 - The application communicates with the TP monitor to begin, commit or rollback the transactions
- The JTA architecture requires that each resource manager must implement the javax.transaction.xa.XAResource interface in order to be managed by the TP monitor

Lab

 Using JTA, MySQL and ActiveMQ to implement a distributed transaction processing application

Programming Practice

- Design and implement a Java application with distributed transaction
 - The system maintains a student list (matriculation numbers and names), where the matriculation number must be unique
 - The user inputs the matriculation number and name of a new student, to insert a record
 - Whenever a new record is inserted, a message is sent to another application via the broker