Elf Formats in Linux

ELF: Executable and Linking Format

 ELF is a portable object file format defining the composition and organization of the object file.
 Kernel and binary loader looks at this format to know how to load this file and find various pieces of information like code, initialized data, dependencies on shared libraries etc.

Elf File Types

- Elf defines the format of executable binary files. There are four different types:
 - Relocatable
 - •Created by compilers or assemblers. Need to be processed by the linker before running.
 - •holds code and data suitable to link with other object files
 - Executable
 - •Have all relocation done and all symbol resolved except perhaps shared library symbols that must be resolved at run time.
 - suitable for execution
 - Shared object
 - •Shared library containing both symbol information for the linker and directly runnable code for run time.
 - holds code and data suitable to link with other relocatable object or shared objectsObject files are created by the assembler and link editor, object files are binary representations of programs
 - Core file
 - •A core dump file.

ELF Structure

- Elf files have a dual nature:
 - Compilers, assemblers, and linkers treat the file as a set of logical sections described by a section header table.
 - The system loader treats the file as a set of segments described by a program header table.

ELF Structure

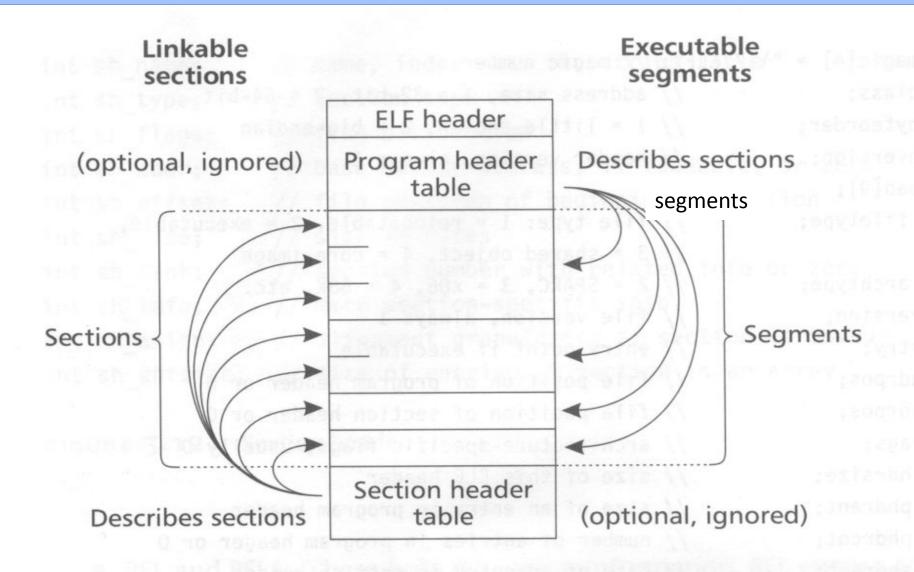


FIGURE 3.10 • Two views of an ELF file.

ELF File Format

- An ELF file has:
 - A header
 - A program header
 - ☐ A section header
 - Sections
 - More stuff which matters less (string table, debug symbols, ...)

Sections are used by the linker (compile time)

Segments are used by the loader (runtime)

ELF Structure

- A single segment usually consist of several sections. E.g., a loadable read-only segment could contain sections for executable code, read-only data, and symbols for the dynamic linker.
- Relocatable files have section header tables. Executable files have program header tables. Shared object files have both.
- Sections are intended for further processing by a linker, while the segments are intended to be mapped into memory.

An Example C Program

```
#include <stdio.h>
  int main() {
    printf("hello world!n");
   Save the code in a file named test.c, and then compile the
   source code into a *.o file and an executable with the
   commands below.
$ gcc test.c -c
$ gcc test.o -o test
   We can check if a file is ELF file by the file command as
   shown below.
$ file test.o
```

Program segments/sections

\$ size /usr/bin/test.o

```
text data bss dec hex 245138 3696 2832 251666 3d712 filename
```

/usr/bin/cc

\$ \$ readelf -h test

The -h option means to display the ELF file header. The Magic number is used to indicate the file is ELF file.

ELF Header

- The Elf header is always at offset zero of the file.
- The program header table and the section header table's offset in the file are defined in the ELF header.
- The header is decodable even on machines with a different byte order from the file's target architecture.
- \$ readelf -h test.o
- The -h option means to display the ELF file header. The Magic number is used to indicate the file is ELF file.

```
char magic[4] = "\177ELF"; // magic number
                        // address size, 1 = 32-bit, 2 = 64-bit
char class;
                        // 1 = little-endian, 2 = big-endian
char byteorder;
                     // header version, always 1
char hversion:
char pad[9];
                        // file type: 1 = relocatable, 2 = executable,
short filetype;
                         // 3 = shared object, 4 = core image
                        // 2 = SPARC, 3 = x86, 4 = 68K, etc.
short archtype;
int fversion;
                         // file version, always 1
                         // entry point if executable
int entry;
                         // file position of program header or 0
int phdrpos;
                         // file position of section header or 0
int shdrpos;
                         // architecture-specific flags, usually 0
int flags:
                         // size of this ELF header
short hdrsize:
                         // size of an entry in program header
short phdrent;
                         // number of entries in program header or 0
short phdrcnt;
                         // size of an entry in section header
short shdrent;
                         // number of entries in section header or 0
short shdrcnt;
                         // section number that contains section name strings
short strsec:
```

Display ELF header: readelf -h

```
ELF Header:
  Magic: 7f 45 4c 46 02 01 01 00 00 00 00 00 00 00 00
  Class:
                                       ELF64
  Data:
                                       2's complement, little
  endian
  Version:
                                       1 (current)
  OS/ABI:
                                       UNIX - System V
  ABI Version:
                                       0
                                       EXEC (Executable file)
  Type:
  Machine:
                                       Advanced Micro Devices
  X86-64
  Version:
                                       0x1
                                       0x400440
  Entry point address:
                                       64 (bytes into file)
  Start of program headers:
  Start of section headers:
                                       4424 (bytes into file)
  Flags:
                                       0x0
  Size of this header:
                                       64 (bytes)
                                       56 (bytes)
  Size of program headers:
  Number of program headers:
  Size of section headers:
                                       64 (bytes)
  Number of section headers:
                                       31
  Coation booden attained table indeed of
```

ELF File Format

Linking View ELF Header Program header table (optional) Section 1 Section 2

Section n

Section header table

Execution View

ELF Header

Program header table

Segment 1

Segment n

Section header table (optional)

ELF Files: Dual Views

Linking View **Execution View ELF Header ELF Header** Program header Program header table (required) table (optional) Section 1 Segment 1 Section 2 Segment 2 Segment n Section n Section header Section header table (required) table (optional)

Linker Overview



Headers

.text

.rodata

.data

.bss

ELF object file b

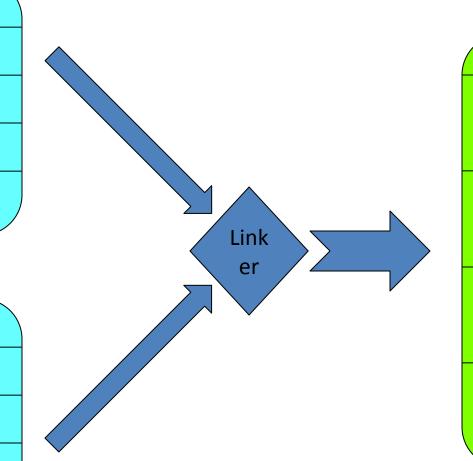
Headers

.text

.rodata

.data

.bss



Executable or library

Headers

.text a & b

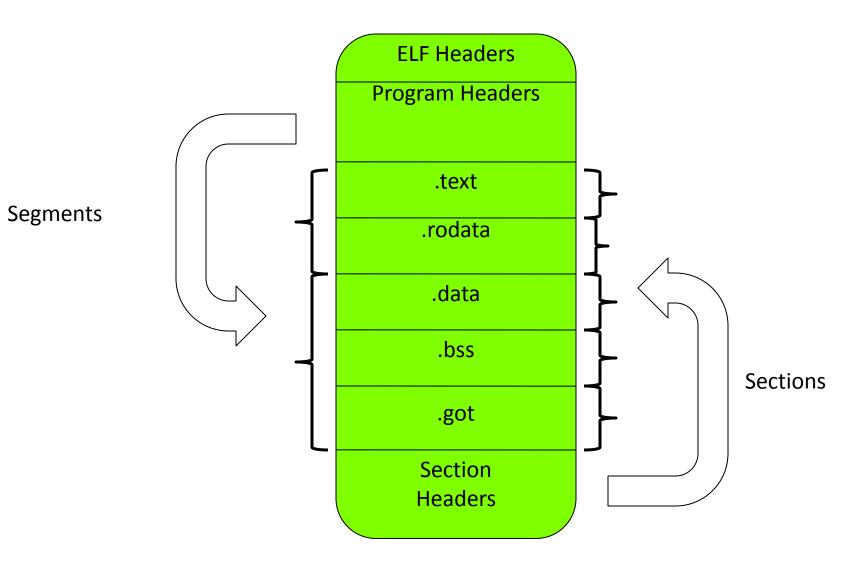
.rodata a & b

.data a & b

.bss a & b

Headers Overview

Executable or library



Relocatable Files

- A relocatable or shared object file is a collection of sections.
- Each section contains a single type of information, such as program code, read-only data, or read/write data, relocation entries, or symbols.
- Every symbol's address is defined relative to a section.
 - Therefore, a procedure's entry point is relative to the program code section that contains that procedure's code.

Section Header

FIGURE 3.12 • Section header.

Types in Section Header

- PROGBITS: This holds program contents including code, data, and debugger information.
- NOBITS: Like PROGBITS. However, it occupies no space.
- SYMTAB: These hold symbol table.
- STRTAB: This is a string table, like the one used in a.out.
- REL and RELA: These hold relocation information.
- DYNAMIC and HASH: This holds information related to dynamic linking.

Flags in Section Header

- WRITE: This section contains data that is writable during process execution.
- ALLOC: This section occupies memory during process execution.
- EXECINSTR: This section contains executable machine instructions.

- .text:
 - This section holds executable instructions of a program.
 - Type: PROGBITS
 - Flags: ALLOC + EXECINSTR
 - \$ readelf -S vipi | grep '.text'
- .data:
 - This section holds initialized data that contributes to the program's image.
 - Type: PROGBITS
 - Flags: ALLOC + WRITE
- .rodata:
 - This section holds read-only data.
 - Type: PROGBITS
 - Flags: ALLOC

\$objdump -s test.o

- .bss :
 - This section holds uninitialized data that contributed to the program's image. By definition, the system will initialize the data with zero when the program begins to run.
 - Type: NOBITS
 - Flags: ALLOC + WRITE
- .rel.text, .rel.data, and .rel.rodata:
 - These contain the relocation information for the corresponding text or data sections.
 - Type: REL
 - Flags: ALLOC is turned on if the file has a loadable segment that includes relocation.
- .symtab:
 - This section hold a symbol table.
- .strtab:
 - This section holds strings.

• .init:

- This section holds executable instructions that contribute to the process initialization code.
- Type: PROGBITS
- Flags: ALLOC + EXECINSTR

• .fini:

- This section hold executable instructions that contribute to the process termination code.
- Type: PROGBITS
- Flags: ALLOC + EXECINSTR
- C does not need these two sections. However, C++ needs them.

• .interp:

- This section holds the pathname of a program interpreter.
- Type: ALLOC
- Flags: PROGBITS
- If this section is present, rather than running the program directly, the system runs the interpreter and passes it the elf file as an argument.
- For many years (used in a.out), UNIX has had self-running interpreted text files, using
- #! /bin/csh as the first line of the file.
- Elf extends this facility to interpreters that run nontext programs.
- In practice, this is used to run the run-time dynamic linker to load the program and to link in any required shared libraries.

- .debug:
 - This section holds symbolic debugging information.
 - Type: PROGBIT
- .line:
 - This section holds line number information for symbolic debugging, which describes the correspondence between the program source and the machine code (ever used gdb?)
 - Type: PROGBIT
- .comment
 - This section may store extra information.
- .got:
 - This section holds the global offset table.
 - •Related toshared library.
 - Type: PROGBIT
- .plt:
 - This section holds the procedure linkage table.
 - Type: PROGBIT
- .note:
 - This section contains some extra information.

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ELF header	blosist care of the state of
(segment table)	(not considered sections)
.text	TE SEE WITCHES SO
.data	object fortuna section of the control of the contro
.rodata	A typical relocatable file.
.bss	relidata and entra a sendo
.sym	ection.
.rel.text	ileasille edende (Sie fadilie to la
.rel.data	will benedia beningan una mistell ou b
.rel.rodata	ing signing to our creating set as
.line	a rells how large a data object is (pa
.debug	musule globul vasione everywher
.strtab	Level per litter in som the restau
Section table	(not considered a section)

FIGURE 3.14 • Sample relocatable ELF file.

String Table

- Like the format used in a.out.
- String table sections hold null-terminated character sequences, commonly called strings.
- The object file uses these strings to represent symbol and section names.
- We use an index into the string table section to reference a string.
- The reason why we separate symbol names from symbol tables is that in C or C++, there is no limitation on the length of a symbol.

String Table

Index	+0	+1	+2	+3	+4	+5	+6	+7	+8	+9
0	\0	n	a	m	е		\0	V	a	r
10	i	a	b	1	е	\0	a	b	1	е
20	\0	\0	х	х	\0					

\$ readelf -p '.symtab' test.o

Index	String		
0	none		
1	name.		
7	Variable		
11	able		
16	able		
24	null string		

Symbol Table

- An object file's symbol table holds information needed to locate and relocate a program's symbolic definition and references.
- A symbol table index is a subscript into this array.

\$ readelf -s test.o

Symbol Table

```
// position of name string in string table
int name;
              // symbol value, section relative in reloc,
int value;
                 absolute in executable
                                             e.g., int, double
             // object or function size
int size;
char type:4; // data object, function, section, or special-case file
char bind:4; // local, global, or weak
                                                    If a definition is available
char other; // spare
                                                    for an undefined weak
             // section number, ABS, COMMON, or Userbol, the linker will use
short sect:
                                                    it. Otherwise, the value
FIGURE 3.13 • ELF symbol table.
                                                    defaults to 0.
```

The section relative to which the symbol is defined. (e.g., the function entry points are defined relative to .text)

Relocation Table

- Relocation is the process of connecting symbolic references with symbolic definitions.
- Relocatable files must have information that describes how to modify their section contents.
- A relocation table consists on many relocation structures.

- \$ readelf -I test.o to read program Headers
- *\$ readelf -S test.o* view the section headers and program headers with readelf command.

Executable Files

- An executable file usually has only a few segments.
 E.g.,
 - A read-only one for the code.
 - A read-only one for read-only data.
 - A read/write one for read/write data.
- All of the loadable sections are packed into the appropriate segments so that the system can map the file with just one or two operations.
 - E.g., If there is a .init and .fini sections, those sections will be put into the read-only text segment.
 - Segment are Page Aligned

Segments

- Executable and shared object files contain segments. A segment is a grouping of one or more sections in the object file. For example, the *code segment* includes the .text and .rodata sections while the *data segment* contains the .bss and .data sections.
- It Defines:
- Permission applied in segments
- Size & Section Information
- Sections with identical permission can mapped to one segments
- \$ readelf —I test.o

Dynamic linking

- Dynamic linking allows to resolve library functions at runtime
- Relocatable File: Holds code and data suitable for linking with other object files to create an executable or shared object file
 - Executable File: Holds a program suitable for execution
 - Shared Object File: Holds code and data suitable for linking in two contexts:
 - The **Link Editor** may process it with other relocatable and shared object files to create another object file
 - The **Dynamic Linker** combines it with an executable file and other shared objects to create a process image

Program Header

```
int type;  // loadable code or data, dynamic linking info, etc.
int offset;  // file offset of segment
int virtaddr;  // virtual address to map segment
int physaddr;  // physical address, not used
int filesize;  // size of segment in file
int memsize;  // size of segment in memory (bigger if contains bss)
int flags;  // Read, Write, Execute bits
int align;  // required alignment, invariably hardware page size
```

FIGURE 3.15 • ELF program header.

The Types in Program Header

- This field tells what kind of segment this array element describes:
 - PT_LOAD: This segment is a loadable segment.
 - PT_DYNAMIC: This array element specifies dynamic linking information.
 - PT_INTERP: This element specified the location and size of a null-terminated path name to invoke as an interpreter.

Executable File Example

	File offset	Load address	Type header
ELF header	0	0x8000000	
Program header	0x40	0x8000040	
Read-only text (size 0x4500)	0x100	0x8000100	LOAD, read/execute
Read/write data (file size 0x2200, memory size 0x3500)	0x4600	0x8005600	LOAD, read/write/ execute

Nonloadable information and optional section headers

FIGURE 3.16 • ELF loadable segments.

Elf Linking

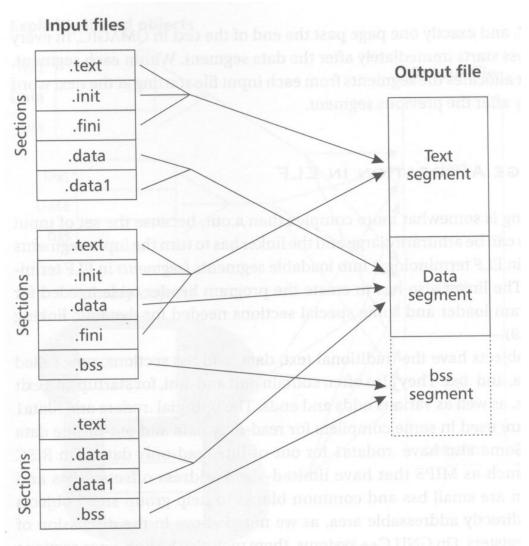


FIGURE 4.9 • ELF linking.

Elf File Trace

(We can use the objdump or nm command)

Program Header

```
Program Header:
   PHDR off 0x00000034 vaddr 0x08048034 paddr 0x08048034 align 2**2
        filesz 0x000000c0 memsz 0x000000c0 flags r-x
  INTERP off 0x000000f4 vaddr 0x080480f4 paddr 0x080480f4 align 2**0
        filesz 0x00000019 memsz 0x00000019 flags r--
   LOAD off
               0x00000000 vaddr 0x08048000 paddr 0x08048000 align 2**12
        filesz 0x00000564 memsz 0x00000564 flags r-x
   LOAD off
               0x00000564 vaddr 0x08049564 paddr 0x08049564 align 2**12
        filesz 0x000000a8 memsz 0x000000cc flags rw-
               0x0000059c vaddr 0x0804959c paddr 0x0804959c align 2**2
DYNAMIC off
        filesz 0x00000070 memsz 0x00000070 flags rw-
               0x00000110 vaddr 0x08048110 paddr 0x08048110 align 2**2
   NOTE off
```

filesz 0x00000018 memsz 0x00000018 flags r--

Dynamic Section

Dynamic Section:

NEEDED libe.so.4

INIT 0×8048390

FINI 0x8048550

HASH 0x8048128

STRTAB 0x80482c8

SYMTAB 0x80481b8

STRSZ 0xad

SYMENT 0×10

DEBUG 0×0

PLTGOT 0x8049584

PLTRELSZ 0x18

PLTREL 0x11

JMPREL 0x8048378

Need to link this shared library for printf()

Section Header

Sect	cions:								
Idx	Name	Size	VMA	LMA	File off	Algn			
0	.interp	00000019	080480f4	080480f4	000000f4	2**0			
		CONTENTS,	ALLOC, LOA	AD, READONI	LY, DATA				
1	.note.ABI-tag	00000018	08048110	08048110	00000110	2**2			
		CONTENTS,	ALLOC, LOA	AD, READONI	LY, DATA				
2	.hash	00000090	08048128	08048128	00000128	2**2			
		CONTENTS,	ALLOC, LOA	AD, READONI	LY, DATA				
3	.dynsym	00000110	080481b8	080481b8	000001b8	2**2			
		CONTENTS,	ALLOC, LOA	AD, READONI	LY, DATA				
4	.dynstr	000000ad	080482c8	080482c8	000002c8	2**0			
		CONTENTS,	ALLOC, LOA	AD, READONI	LY, DATA				
5	.rel.plt	00000018	08048378	08048378	00000378	2**2			
		CONTENTS,	ALLOC, LOA	AD, READONI	LY, DATA				
6	.init	0000000b	08048390	08048390	00000390	2**2			
		CONTENTS,	ALLOC, LOA	AD, READONI	LY, CODE				
7	.plt	00000040	0804839c	0804839c	0000039c	2**2			
		CONTENTS,	ALLOC, LOA	AD, READONI	LY, CODE				
8	.text	00000174	080483dc	080483dc	000003dc	2**2			
CONTENTS, ALLOC, LOAD, READONLY, CODE									

Section Header (cont'd)

9 .fini	00000006	08048550	08048550	00000550	2**2
	CONTENTS,	ALLOC, LO	AD, READON	LY, CODE	
10 .rodata	0000000e	08048556	08048556	00000556	2**0
	CONTENTS,	ALLOC, LO	AD, READON	LY, DATA	
11 .data	0000000c	08049564	08049564	00000564	2**2
	CONTENTS,	ALLOC, LO	AD, DATA		
12 .eh_frame	00000004	08049570	08049570	00000570	2**2
	CONTENTS,	ALLOC, LO	AD, DATA		
13 .ctors	80000008	08049574	08049574	00000574	2**2
	CONTENTS,	ALLOC, LO	AD, DATA		
14 .dtors	0000008	0804957c	0804957c	0000057c	2**2
	CONTENTS,	ALLOC, LO	AD, DATA		
15 .got	00000018	08049584	08049584	00000584	2**2
	CONTENTS,	ALLOC, LOZ	AD, DATA		
16 .dynamic	00000070	0804959c	0804959c	0000059c	2**2
	CONTENTS,	ALLOC, LOZ	AD, DATA		
17 .bss	00000024	0804960c	0804960c	0000060c	2**2
	ALLOC				
18 .stab	000001bc	0000000	00000000	0000060c	2**2
	CONTENTS,	READONLY,	DEBUGGING		
19 .stabstr	00000388	0000000	00000000	000007c8	2**0
	CONTENTS,	READONLY,	DEBUGGING		
20 .comment	000000c8	00000000	00000000	00000b50	2**0

Symbol Table

```
SYMBOL TABLE:
080480f4 1
                              0000000
             d
                .interp
08048110 1
                .note.ABI-tag 0000000
08048128 1
                .hash 00000000
080481b8 l
             d .dynsym
                               0000000
080482c8 1
                               0000000
                .dynstr
             d
08048378 1
             d .rel.plt
                               0000000
08048390 1
             d .init 00000000
             d .plt 00000000
0804839c l
080483dc 1
             d .text 00000000
             d .fini 00000000
08048550 1
08048556 1
             d .rodata
                               0000000
08049564 1
             d .data 00000000
08049570 1
                .eh frame
                              0000000
08049574 1
             d
                .ctors 00000000
0804957c l
             d .dtors 00000000
                       0000000
08049584 1
             d
                .got
0804959c 1
                               0000000
             d
                .dynamic
```

Symbol Table (cont'd)

```
• 0804960c l d .bss 00000000
• 00000000 l d .stab 00000000
• 00000000 l d .stabstr 00000000
• 00000000 l d .comment 00000000
• 00000000 l d .note 00000000
• 00000000 l d *ABS* 00000000
• 00000000 l d *ABS* 00000000
• 00000000 l d *ABS* 00000000
• 00000000 l df *ABS* 00000000 crtstuff.c
• 08048460 1
                .text 00000000 gcc2 compiled.
• 08049568 1 O .data 00000000 p.3
• 0804957c l O .dtors 0000000 DTOR LIST
• 0804956c l O .data 00000000 completed.4
• 08048460 l F .text 00000000 do global dtors aux
• 08049570 l O .eh frame 00000000 EH FRAME BEGIN
```

Symbol Table (cont'd)

```
080484b4 l F .text 00000000 fini dummy
0804960c l 0 .bss 00000018 object.11
080484bc l F .text 00000000 frame dummy
080484e0 l F .text 00000000 init dummy
08049570 l O .data 00000000 force to data
08049574 l O .ctors 0000000 CTOR LIST
00000000 l df *ABS* 00000000 crtstuff.c
08048520 l .text 00000000 gcc2 compiled.
08049578 1 O .ctors 0000000 CTOR END
08048548 l F .text 00000000 init dummy
08049570 1 O .data 00000000 force_to_data
08049580 l O .dtors 0000000 DTOR END
08049570 l O .eh frame 00000000 FRAME END
00000000 l df *ABS* 00000000 p10.c
080483ac F *UND* 00000031 printf
0804959c g O *ABS* 0000000 DYNAMIC
08048550 g O *ABS* 0000000 etext
08048390 g F .init 00000000 init
08049624 q 0 .bss 00000004 environ
00000000 w *UND* 00000000 deregister frame info
08049630 g O *ABS* 00000000 end
08049628 g O.bss 00000004 xx
```

Symbol Table (cont'd)

```
080483dc g F .text 00000083 start
080484e8 g F .text 00000038 main
08048550 g F .fini 0000000 fini
0804962c g O .bss 00000004 yy
080483bc
       F *UND* 00000070 atexit
0804960c g O *ABS* 0000000 edata
08049584 g O *ABS* 0000000 GLOBAL_OFFSET_TABLE_
08049630 g O *ABS* 00000000 end
080483cc
       F *UND* 0000005b exit
00000000 w *UND* 00000000 register frame info
```

Dynamic Symbol Table

```
DYNAMIC SYMBOL TABLE:
080483ac DF *UND* 00000031 printf
0804959c g DO *ABS* 0000000 DYNAMIC
08048550 g DO *ABS* 00000000 etext
08048390 g DF .init 00000000 init
08049624 g DO .bss 00000004 environ
00000000 w D *UND* 00000000 deregister frame info
08049630 g DO *ABS* 00000000 end
08049564 g DO .data 00000004 __progname
0804960c q DO *ABS* 00000000 bss start
08048550 g DF .fini 00000000 fini
080483bc DF *UND* 00000070 atexit
0804960c g DO *ABS* 00000000 edata
08049584 q DO *ABS*
                    0000000 GLOBAL OFFSET_TABLE_
                    00000000 end
08049630 g DO *ABS*
080483cc DF *UND* 0000005b exit
00000000 w D *UND* 00000000 register_frame_info
```

Dynamic Relocation Table

• DYNAMIC RELOCATION RECORDS

```
    OFFSET TYPE VALUE
    08049590 R_386_JUMP_SLOT printf
    08049594 R_386_JUMP_SLOT atexit
    08049598 R 386 JUMP SLOT exit
```

Library

Library is a collection of resources used to develop software.

These may include pre-written code and subroutines, classes, values or type specifications.

Libraries contain code and data that provide services to independent programs.

Adv- sharing and changing of code and data in a modular fashion, and eases the distribution of the code and data.

Some executables are both standalone programs and libraries, but most libraries are not executable. Executables and libraries make references known as links to each other through the process known as linking, which is typically done by a linker.

Shared Libraries

shared library or shared object is a file that is intended to be shared by executable files and further shared objects files.

Modules used by a program are loaded from individual shared objects into memory at load time or run time, rather than being copied by a linker when it creates a single monolithic executable file for the program.

Shared Libraries

Shared libraries remove the common library routines from the executable file

Maintaining a single copy of the library routine somewhere in memory that all processes reference.

This reduces the size of each executable file (But may add some runtime overhead, either when the program is first executed or the first time each shared library function is called)

Another advantage of shared libraries is that library functions can be replaced with new versions without having to relink edit every program that uses the library.

```
$ cc -static hello1.c prevent gcc from using shared libraries
```

```
$ Is -I a.out
```

-rwxrwxr-x 1 sar 475570 Feb 18 23:17 a.out

\$ size a.out

text data bss dec hex filename 375657 3780 3220 382657 5d6c1 a.out

If we compile this program to use shared libraries, the text and data sizes of the executable file are greatly decreased:

\$ cc hello1.c gcc defaults to use shared libraries

\$ Is -I a.out

-rwxrwxr-x 1 sar 11410 Feb 18 23:19 a.out

\$ size a.out

text data bss dec hex filename 872 256 4 1132 46c a.out

http://tldp.org/HOWTO/Program-Library-HOWT O/shared-libraries.html