Compiler Design

Report Of Phase - 2

Parser / Syntax Checker

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Introduction

Yacc provides a general tool for imposing structure on the input to a computer program. The Yacc user prepares a specification of the input process; this includes rules describing the input structure, code to be invoked when these rules are recognized, and a low-level routine to do the basic input. Yacc then generates a function to control the input process. This function, called a parser, calls the user-supplied low-level input routine (the lexical analyzer) to pick up the basic items (called tokens) from the input stream. These tokens are organized according to the input structure rules, called grammar rules; when one of these rules has been recognized, then user code supplied for this rule, an action, is invoked; actions have the ability to return values and make use of the values of other actions.

An important part of the input process is carried out by the lexical analyzer. This user routine reads the input stream, recognizing the lower level structures, and communicates these tokens to the parser. For historical reasons, a structure recognized by the lexical analyzer is called a terminal symbol, while the structure recognized by the parser is called a nonterminal symbol. To avoid confusion, terminal symbols will usually be referred to as tokens.

Specification Of Yacc

Like lex, yacc has its own specification language. A yacc specification is structured along the same lines as a Lex specification.

```
%{
    /* C declarations and includes */
%}
    /* Yacc token and type declarations */
%%
    /* Yacc Specification
    in the form of grammar rules like this:
    */
    symbol : symbols tokens
        { $$ = my_c_code($1); }
    ;
%%
    /* C language program (the rest) */
```

The Yacc Specification rules are the place where you "glue" the various tokens together that lex has conveniently provided to you.

Each grammar rule defines a symbol in terms of:

- 1) other symbols
- 2) tokens (or terminal symbols) which come from the lexer.

Each rule can have an associated action, which is executed after all the component symbols of the rule have been parsed. Actions are basically C-program statements surrounded by curly braces.

Lexer

The lexer will provide us with the following tokens:

- 1) Each possible keyword (TITLE, MENU, etc) as a separate token
- 2) A LABEL token for representing labels
- 3) A EXEC token for executable commands
- 4) Icon file names will be represented as 3 separate tokens: '<' LABEL '>'

To aid in error-recovery, newline characters will be considered significant, and will be passed up as a separate token (unless they are preceded by a '\').

The key aspects of the lexer are:

Each keyword returns a separate, unique token

A LABEL is either an identifier (loosely speaking) or an arbitrary string in quotes, or any alternating sequence of these two things

- 1) An EXEC token is identified by the fact that:
- 2) it is not a keyword.
- 3) it appears after we have scanned a LABEL token on the same line.
- 4) This is achieved by setting the start-condition ACT when we scan the LABEL token.

We use lex's "first match" rule to ensure that keywords get priority over the corresponding LABEL and EXEC interpretations, and that the EXEC interpretation gets priority over the LABEL interpretation in the state ACT.

It is thus essential that, where a keyword may appear on a line, the length of the other rules (for LABEL or EXEC) be no longer than the keyword rule. Otherwise, lex's "longest match" rule would override the "first match" rule.

Tokens

We have included a lot of return statements, but have not defined the tokens in the brackets. A token is simply a unique integer which yacc associates with an item that the lexer has found. Within yacc, all tokens must be declared in the "Yacc Declarations" section of our yacc specification, like this:

%token TITLE %token MENU

The yyerror() function

The yyerror() function is called when yacc encounters an invalid syntax. The yyerror() is passed a single string (char*) argument. Unfortunately, this string usually just says "parse error", so on its own, it's pretty useless. Error recovery is an important topic which we will cover in more detail later on. For now, we just want a basic yyerror() function like this:

```
yyerror(char *err) {
   fprintf(stderr, "%s\n",err);
}
```

Implementations Details

First we retrieved all the tokens from the lexical analyser. The strings which matches the the grammar of lex are known as tokens. Also in the lexer part the comments part has been taken care of. The comment is ignored whenever it is encountered. All the data types, arrays, functions, if else conditions and loops which are there in our abstract are written in lex language. Operators are also taken care. Now in the parser file we have called the tokens which are returned from the lexical phase. All the operators are written in the priority wise so that the problem of precedence is taken care off. In the operator's % left indicate that the precedence is from left to right and % right indicate right to left. Also % left means left recursive and % right means right recursive.

Now we have started our grammar by the start keyword, then followed by again start and function or declaration or start or it can can be empty also. The empty part is taken care by including the keyword empty which means there is nothing which matches the grammar. For structure the STRUCT keyword is defined followed by '{ 'start '} ', which means inside the structure we can define the start state. The function declaration takes type, ID then argument declaration in the parameter part then followed by the compound statements. statement can contain statement again or for statement or while statement or if statement or return statement in it followed by semicolon. Then we have defined the function by defining the grammar as first will be return type then variable name which is ID here the parenthesis then followed by argument declaration then statements. Then we have declaration statements which can be expression, structure or function call. We have also used error keyword which is for handling multi line errors.

Then we have defined the statement, it can be any compound statement or while loop statement, for loop statement, declaration, print, scan statement. Compound statement is a statement list enclose in curly braces. Declaration statement is defined by ID followed by assign which means assigning a value to the variable.

Then we have declared rules for while loop, for loop,do while cases separately in the grammar. Also we have declared rules for if and else if and else. Grammar for if statement is if statement declared then brackets then condition. Also dangling if-else part is taken care off in the yacc part.

We have also written grammar for array declaration by giving ID first followed by '[' then expression then ']'. The grammar for printf and scanf is also given in the code of the grammar. Also expr and rexpr are also defined so that it handle the cases with operators part. Also less than equal to , greater than equal to , not equal to cases are handled. Then there are operators such as '+', '-', '*', '/', etc are handled.

At the end if parsing is completed then we have shown that Parsing is complete. If parsing is not complete than we have shown that parsing error is there. Also we have declared the error line where the error will occur. We have also displayed the symbol table in which we have shown that symbol name, type, attribute and line number.

Code for Lexical Phase

```
%{
      int lineno=1
     char *tempid;
#include"hashtable.c"
#include<ctype.h>
#include<string.h>
char types[100];
                             %}
identifier
ndentifle
header
type
storage
qualifier
digits
decimal
lint
llint
double
float
scientific
scientificf
str_literal
character
space
next line
                              \n
"["|"]"|","|":"|"{"|"}"|"("|")"|"="|[-+*%/--&|^]
 s_operator
%x mlcomment
%x slcomment
%%
"/*"
                             BEGIN(mlcomment)
 -/*"
<mlcomment>[^*\n]*
<mlcomment>\n
<mlcomment>"*"+[^/]
<mlcomment>"*"+"/"
                                 BEGIN(INITIAL)
 "//"
<slcomment>[^\n]
                              BEGIN(slcomment)
```

Code For Yacc

```
STRUCT ID '{' start '}'
structure:
                                                                                                                                                                         TYPE ID '(' argdecls ')' compstmt
                      declstmt ';'
expr ';'
structure ';'
functcall ';'
error ';'
error '}'
declaration:
                       compstmt
forstmt
whilestmt
                       dowhilestmt
ifstmt
                       returnstmt
declaration
                       print
scan
';'
                       '{' stmtlist '}'
compstmt:
                       stmt stmtlist
%empty
stmtlist:
                       TYPE decllist
declstmt:
                       ID assign ',' decllist
```

```
ID assign ',' decllist
ID assign
arraydecl ',' decllist
arraydecl
arraydecl '=' '{' arrassign
decllist:
dowhilestmt:
                     DO compstmt WHILE '(' rexp ')' ';'
                     RETURN ';'
RETURN rexp ';'
returnstmt:
                     FOR '(' fexp ';' fexp ';' fexp ')' stmt
fexp:
                     rexp
%empty;
                      WHILE '(' rexp ')' stmt
whilestmt:
                     NUM ',' arrassign
NUM '}'
arrassign:
                     '=' rexp
%empty
assign:
                      ID '(' mul ')'
                       NUM | ID;
```

```
× scanner.y
106
107
108
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120
121
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                                          ID '(' mul ')'
          const:
                                        NUM | ID;
const ',' mul
const
         arraydecl:
                                               PRINTF '(' STR_LITERAL ')' ';'
PRINTF '(' STR_LITERAL numlist')' ';'
                                             SCANF '(' STR_LITERAL snumlist')' ';'
                                          TYPE ID argassign',' argdecls
TYPE ID argassign
%empty
         argdecls:
                                          '=' NUM
%empty
         argassign:
                                          ',' '&' ID snumlist
                                         ',' ID numlist
          numlist:
                                           expr
NA rexp
rexp EE rexp
          rexp :
```

```
| To Me expr | To DE expr | To
```

Code for Hash Table

Test Cases

Test case #1:

Input:

```
finclude<stdio.h>

int main()

{
    int a,b
    a = 100;
    b = 1000;
    sum=a+b;
    return 0;
}

// semicolon missing at line 5
```

Expected output: syntax error before ')'

Actual output :

Test case #2:

Test case #3

```
#Includesdtio.hp
#includesdtio.hp
#includesdtio.hp

#includesdtio.hp

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```

Test case #4

```
#include<stdio.h>
#include<stdio.h

#include<std
```

Test case #5:

```
# scanner | x scanner | x scanner | x scanner | x | scanne
```

Test case #6:

```
#includestdio.h>
#includestdib.h>

void find product(int a,int b)
{
    int c=a*b;
    printf("Product is %d",c);
}

int main()
{
    int a=10,b=10;
    find_product(a,);
    return 0;
}

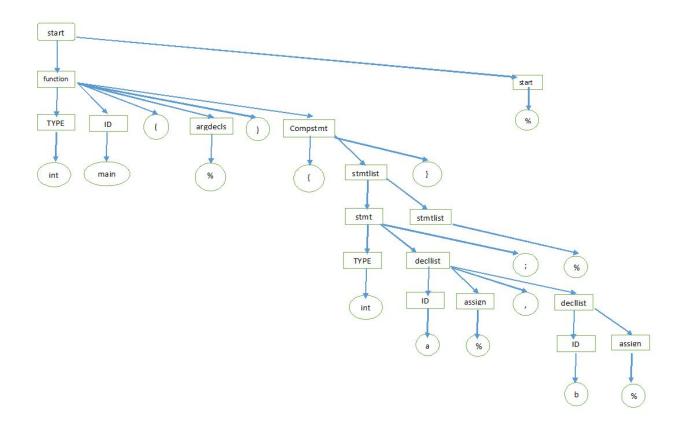
// function parameter given is wrong
```

Expected output: Syntax error at line number 10

Actual Output:

```
Parse Tree:
```

```
#include<stdio.h>
Int main(){
      int a,b;
}
```



Conclusion

So far we have finished the two stages of the C compiler lexical phase and syntax phase.Lexical phase will generate all tokens and show lexical errors whenever it is found.While the syntax phase will generate

Syntax related errors and the parse the Code.All the cases are handled very carefully.It will also handle multi-line errors and print on different lines.