Financial Engineering and Risk Management Review of vectors

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Reals numbers and vectors

- ullet We will denote the set of real numbers by ${\mathbb R}$
- Vectors are finite collections of real numbers
- Vectors come in two varieties
 - Row vectors: $\mathbf{v} = \begin{bmatrix} v_1 & v_2 & \dots & v_n \end{bmatrix}$

$$ullet$$
 Column vectors $oldsymbol{w} = egin{bmatrix} w_1 \\ w_2 \\ \vdots \\ w_n \end{bmatrix}$

- By default, vectors are column vectors
- ullet The set of all vectors with ${f n}$ components is denoted by ${\Bbb R}^{f n}$

Linear independence

• A vector \mathbf{w} is linearly dependent on $\mathbf{v}_1, \mathbf{v}_2$ if

$$\mathbf{w} = \alpha_1 \mathbf{v}_1 + \alpha_2 \mathbf{v}_2$$
 for some $\alpha_1, \alpha_2 \in \mathbb{R}$

Example:

$$\begin{bmatrix} 2 \\ 6 \\ 4 \end{bmatrix} = 2 \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} + 4 \begin{bmatrix} 0 \\ 1 \\ 1 \end{bmatrix}$$

- Other names: linear combination, linear span
- A set $\mathcal{V} = \{\mathbf{v}_1, \dots, \mathbf{v}_m\}$ are linearly independent if **no** \mathbf{v}_i is linearly dependent on the others, $\{\mathbf{v}_j : j \neq i\}$

Basis

ullet Every $old w \in \mathbb{R}^n$ is a linear combination of the linearly independent set

$$\mathcal{B} = \left\{ \begin{bmatrix} 1\\0\\\vdots\\0 \end{bmatrix}, \begin{bmatrix} 0\\1\\\vdots\\0 \end{bmatrix}, \dots \begin{bmatrix} 0\\0\\\vdots\\1 \end{bmatrix} \right\} \qquad \mathbf{w} = w_1 \underbrace{\begin{bmatrix} 1\\0\\\vdots\\0 \end{bmatrix}}_{2} + w_2 \underbrace{\begin{bmatrix} 0\\1\\\vdots\\0 \end{bmatrix}}_{2} + \dots + w_n \underbrace{\begin{bmatrix} 0\\0\\\vdots\\1 \end{bmatrix}}_{2}$$

- ullet Basis \equiv any linearly independent set that spans the entire space
- Any basis for \mathbb{R}^n has exactly n elements

Norms

- A function $\rho(\mathbf{v})$ of a vector \mathbf{v} is called a norm if
 - $\rho(\mathbf{v}) \geq 0$ and $\rho(\mathbf{v}) = 0$ implies $\mathbf{v} = \mathbf{0}$
 - $\rho(\alpha \mathbf{v}) = |\alpha| \, \rho(\mathbf{v})$ for all $\alpha \in \mathbb{R}$
 - $\rho(\mathbf{v}_1 + \mathbf{v}_2) \le \rho(\mathbf{v}_1) + \rho(\mathbf{v}_2)$ (triangle inequality)

ho generalizes the notion of "length"

• Examples:

- ℓ_2 norm: $\|\mathbf{x}\|_2 = \sqrt{\sum_{i=1}^n |x_i|^2}$... usual length
- ℓ_1 norm: $\|\mathbf{x}\|_1 = \sum_{i=1}^n |x_i|$
- $\bullet \ \ell_{\infty} \ \text{norm:} \ \left\| \mathbf{x} \right\|_{\infty} = \max_{1 \leq i \leq n} \left| x \right|_{i}$
- ℓ_p norm, $1 \le p < \infty$: $\|\mathbf{x}\|_p = \left(\sum_{i=1}^n |x|_i^p\right)^{\frac{1}{p}}$

Inner product

• The inner-product or dot-product of two vector $\mathbf{v}, \mathbf{w} \in \mathbb{R}^n$ is defined as

$$\mathbf{v} \cdot \mathbf{w} = \sum_{i=1}^{n} v_i w_i$$

- The ℓ_2 norm $\|\mathbf{v}\|_2 = \sqrt{\mathbf{v} \cdot \mathbf{v}}$
- The angle θ between two vectors \mathbf{v} and \mathbf{w} is given by

$$\cos(\theta) = \frac{\mathbf{v} \cdot \mathbf{w}}{\|\mathbf{v}\|_2 \|\mathbf{w}\|_2}$$

• Will show later: $\mathbf{v} \cdot \mathbf{w} = \mathbf{v}^{\top} \mathbf{w} = \text{product of } \mathbf{v} \text{ transpose and } \mathbf{w}$

Financial Engineering and Risk Management Review of matrices

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Matrices

- Matrices are rectangular arrays of real numbers
- Examples:

$$\bullet \ \, \mathbf{A} = \begin{bmatrix} 2 & 3 & 7 \\ 1 & 6 & 5 \end{bmatrix} \colon 2 \times 3 \ \text{matrix} \\ \bullet \ \, \mathbf{B} = \begin{bmatrix} 2 & 3 & 7 \end{bmatrix} \colon 1 \times 3 \ \text{matrix} \equiv \text{row vector} \\ \bullet \ \, \mathbf{A} = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{bmatrix} \colon \mathbf{m} \times \mathbf{n} \ \text{matrix} \dots \mathbb{R}^{\mathbf{m} \times \mathbf{n}} \\ \bullet \ \, \mathbf{I} = \begin{bmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & & & 1 \end{bmatrix} \dots \ n \times n \ \text{Identity matrix}$$

Vectors are clearly also matrices

Matrix Operations: Transpose

• Transpose: $\mathbf{A} \in \mathbb{R}^{m \times d}$

$$\mathbf{A}^{\top} = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1d} \\ a_{21} & a_{22} & \dots & a_{2d} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{md} \end{bmatrix}^{\top} = \begin{bmatrix} a_{11} & a_{21} & \dots & a_{m1} \\ a_{12} & a_{22} & \dots & a_{d2} \\ \vdots & \vdots & \ddots & \vdots \\ a_{1d} & a_{2d} & \dots & a_{md} \end{bmatrix} \in \mathbb{R}^{d \times m}$$

• Transpose of a row vector is a column vector

Example:

•
$$\mathbf{A} = \begin{bmatrix} 2 & 3 & 7 \\ 1 & 6 & 5 \end{bmatrix}$$
: 2×3 matrix ... $\mathbf{A}^{\top} = \begin{bmatrix} 2 & 1 \\ 3 & 6 \\ 7 & 5 \end{bmatrix}$: 3×2 matrix

•
$$\mathbf{v} = \begin{bmatrix} 2 \\ 6 \\ 4 \end{bmatrix}$$
: column vector ... $\mathbf{v}^{\top} = \begin{bmatrix} 2 & 6 & 4 \end{bmatrix}$: row vector

Matrix Operations: Multiplication

• Multiplication: $\mathbf{A} \in \mathbb{R}^{\mathbf{m} \times \mathbf{d}}$, $\mathbf{B} \in \mathbb{R}^{\mathbf{d} \times \mathbf{p}}$ then $\mathbf{C} = \mathbf{A} \mathbf{B} \in \mathbb{R}^{\mathbf{m} \times \mathbf{p}}$

$$c_{ij} = \left[egin{array}{cccc} a_{i1} & a_{i2} & \dots & a_{id} \end{array}
ight] \left[egin{array}{cccc} b_{1j} \ b_{2j} \ dots \ b_{dj} \end{array}
ight]$$

- row vector $\mathbf{v} \in \mathbb{R}^{1 \times d}$ times column vector $\mathbf{w} \in \mathbb{R}^{d \times 1}$ is a scalar.
- Identity times any matrix $\mathbf{AI}_n = \mathbf{I}_m \mathbf{A} = \mathbf{A}$

• Examples:

•
$$\begin{bmatrix} 2 & 3 & 7 \\ 1 & 6 & 5 \end{bmatrix} \begin{bmatrix} 2 \\ 6 \\ 4 \end{bmatrix} = \begin{bmatrix} 2(2) + 3(6) + 7(4) \\ 1(2) + 6(6) + 5(4) \end{bmatrix} = \begin{bmatrix} 50 \\ 58 \end{bmatrix}$$

$$\bullet \ \ell_2 \ \text{norm:} \ \left\| \begin{bmatrix} 1 \\ -2 \end{bmatrix} \right\|_2 = \sqrt{1^2 + (-2)^2} = \sqrt{\begin{bmatrix} 1 \\ -2 \end{bmatrix}} = \sqrt{\begin{bmatrix} 1 \\ -2 \end{bmatrix}}^\top \begin{bmatrix} 1 \\ -2 \end{bmatrix}$$

• inner product: $\mathbf{v} \cdot \mathbf{w} = \mathbf{v}^{\top} \mathbf{w}$

Linear functions

• A function $f: \mathbb{R}^d \mapsto \mathbb{R}^m$ is linear if

$$f(\alpha \mathbf{x} + \beta \mathbf{y}) = \alpha f(\mathbf{x}) + \beta f(\mathbf{y}), \quad \alpha, \beta \in \mathbb{R}, \mathbf{x}, \mathbf{y} \in \mathbb{R}^d$$

- ullet A function f is linear if and only if $f(\mathbf{x}) = \mathbf{A}\mathbf{x}$ for matrix $\mathbf{A} \in \mathbb{R}^{m \times d}$
- Examples

•
$$f(\mathbf{x}) : \mathbb{R}^3 \mapsto \mathbb{R}$$
: $f(\mathbf{x}) = \begin{bmatrix} 2 & 3 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = 2x_1 + 3x_2 + 4x_3$

•
$$f(\mathbf{x}) : \mathbb{R}^3 \mapsto \mathbb{R}^2$$
: $f(\mathbf{x}) = \begin{bmatrix} 2 & 3 & 4 \\ 1 & 0 & 2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 2x_1 + 3x_2 + 4x_3 \\ x_1 + 2x_3 \end{bmatrix}$

- Linear constraints define sets of vectors that satisfy linear relationships
 - Linear equality: $\{x : Ax = b\}$... line, plane, etc.
 - Linear inequality: $\{x: Ax \leq b\}$... half-space

Rank of a matrix

- ullet column rank of ${f A} \in \mathbb{R}^{m imes d} = {\sf number}$ of linearly independent columns
 - range(\mathbf{A}) = { \mathbf{y} : \mathbf{y} = $\mathbf{A}\mathbf{x}$ for some \mathbf{x} }
 - column rank of A = size of basis for range(A)
 - column rank of $\mathbf{A} = m \Rightarrow \operatorname{range}(\mathbf{A}) = \mathbb{R}^m$
- row rank of A = number of linearly independent rows
- Fact: row rank = column rank $\leq \min\{m, d\}$
- Example:

$$\mathbf{A} = \begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & 6 \end{bmatrix}, \quad \mathsf{rank} = 1, \quad \mathsf{range}(\mathbf{A}) = \left\{ \lambda \begin{bmatrix} 1 \\ 2 \end{bmatrix} : \lambda \in \mathbb{R} \right\}$$

• $\mathbf{A} \in \mathbb{R}^{\mathbf{n} \times \mathbf{n}}$ and $\mathbf{rank}(\mathbf{A}) = n \Rightarrow \mathbf{A}$ invertible, i.e. $\mathbf{A}^{-1} \in \mathbb{R}^{\mathbf{n} \times \mathbf{n}}$

$$A^{-1}A = AA^{-1} = I$$

Financial Engineering and Risk Management

Review of linear optimization

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Hedging problem

- d assets
- Prices at time t = 0: $\mathbf{p} \in \mathbb{R}^d$
- Market in m possible states at time t=1
- ullet Price of asset j in state $i=S_{ij}$

$$\mathbf{S}_{j} = \begin{bmatrix} S_{1j} \\ S_{2j} \\ \vdots \\ S_{mj} \end{bmatrix} \quad \mathbf{S} = \begin{bmatrix} \mathbf{S}_{1} & \mathbf{S}_{2} & \dots & \mathbf{S}_{d} \end{bmatrix} = \begin{bmatrix} S_{11} & S_{12} & \dots & S_{1d} \\ S_{21} & S_{22} & \dots & S_{2d} \\ \vdots & \vdots & \ddots & \vdots \\ S_{m1} & S_{m2} & \dots & S_{md} \end{bmatrix} \in \mathbb{R}^{m \times d}$$

- Hedge an obligation $\mathbf{X} \in \mathbb{R}^m$
 - Have to pay X_i if state i occurs
 - Buy/short sell $\boldsymbol{\theta} = (\theta_1, \dots, \theta_d)^{\top}$ shares to cover obligation

Hedging problem (contd)

- ullet Position $oldsymbol{ heta} \in \mathbb{R}^d$ purchased at time t=0
 - ullet $heta_j =$ number of shares of asset j purchased, $j=1,\ldots,d$
 - Cost of the position $oldsymbol{ heta} = \sum_{j=1}^d p_j heta_j = \mathbf{p}^ op oldsymbol{ heta}$
- Payoff from liquidating position at time t=1
 - payoff y_i in state i: $y_i = \sum_{j=1}^d S_{ij}\theta_j$
 - ullet Stacking payoffs for all states: ${f y}={f S}{m heta}$
 - Viewing the payoff vector \mathbf{y} : $\mathbf{y} \in \mathsf{range}(\mathbf{S})$

$$\mathbf{y} = egin{bmatrix} \mathbf{S}_1 & \mathbf{S}_2 & \dots & \mathbf{S}_d \end{bmatrix} egin{bmatrix} eta_1 \ eta_2 \ eta_d \ eta_d \end{bmatrix} = \sum_{j=1}^d heta_j \mathbf{S}_j$$

• Payoff **y** hedges **X** if $y \ge X$.

Hedging problem (contd)

Optimization problem:

$$\begin{array}{ll} \min & \sum_{j=1}^d p_j \theta_j & (\equiv \mathbf{p}^\top \boldsymbol{\theta}) \\ \text{subject to} & \sum_{j=1}^d S_{ij} \theta_j \geq X_i, \quad i=1,\ldots,m \quad (\equiv \mathbf{S} \boldsymbol{\theta} \geq \mathbf{X}) \end{array}$$

- Features of this optimization problem
 - Linear objective function: $\mathbf{p}^{\top} \theta$
 - Linear inequality constraints: $S\theta > X$
- Example of a linear program
 - Linear objective function: either a min/max
 - Linear inequality and equality constraints $\begin{aligned} &\max/\min_{\mathbf{x}} \quad \mathbf{c}^{\top}\mathbf{x} \\ &\text{subject to} \quad \mathbf{A}_{eq}\mathbf{x} = \mathbf{b}_{eq} \\ &\mathbf{A}_{in}\mathbf{x} < \mathbf{b}_{in} \end{aligned}$

Linear programming duality

Linear program

$$P = \min_{\mathbf{x}} \mathbf{c}^{\top} \mathbf{x}$$
 subject to $\mathbf{A} \mathbf{x} \ge \mathbf{b}$

Dual linear program

$$D = \max_{\mathbf{u}} \mathbf{b}^{\mathsf{T}} \mathbf{u}$$
subject to $\mathbf{A}^{\mathsf{T}} \mathbf{u} = \mathbf{c}$
 $\mathbf{u} \ge \mathbf{0}$

Theorem.

- Weak Duality: P > D
- Bound: \mathbf{x} feasible for P, \mathbf{u} feasible for D, $\mathbf{c}^{\top}\mathbf{x} \geq P \geq D \geq \mathbf{b}^{\top}\mathbf{u}$
- Strong Duality: Suppose P or D finite. Then P = D.
- Dual of the dual is the primal (original) problem

More duality results

Here is another primal-dual pair

• General idea for constructing duals

$$\begin{array}{ll} P & = & \min\{\mathbf{c}^{\top}\mathbf{x}: \mathbf{A}\mathbf{x} \geq \mathbf{b}\} \\ & \geq & \min\{\mathbf{c}^{\top}\mathbf{x} - \mathbf{u}^{\top}(\mathbf{A}\mathbf{x} - \mathbf{b}): \mathbf{A}\mathbf{x} \geq \mathbf{b}\} \text{ for all } \mathbf{u} \geq \mathbf{0} \\ & \geq & \mathbf{b}^{\top}\mathbf{u} + \min\{(\mathbf{c} - \mathbf{A}^{\top}\mathbf{u})^{\top}\mathbf{x}: \mathbf{x} \in \mathbb{R}^{n}\} \\ & = & \left\{ \begin{array}{ll} \mathbf{b}^{\top}\mathbf{u} & \mathbf{A}^{\top}\mathbf{u} = \mathbf{c} \\ -\infty & \text{otherwise} \end{array} \right. \\ & \geq & \max\{\mathbf{b}^{\top}\mathbf{u}: \mathbf{A}^{\top}\mathbf{u} = \mathbf{c}\} \end{array}$$

• Lagrangian relaxation: dualize constraints and relax them!

Financial Engineering and Risk Management

Review of nonlinear optimization

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Unconstrained nonlinear optimization

Optimization problem

$$\min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x})$$

- Categorization of minimum points
 - \mathbf{x}^* global minimum if $f(\mathbf{y}) \geq f(\mathbf{x}^*)$ for all \mathbf{y}
 - \mathbf{x}^*_{loc} local minimum if $f(\mathbf{y}) \geq f(\mathbf{x}^*_{loc})$ for all \mathbf{y} such that $\|\mathbf{y} \mathbf{x}^*_{loc}\| \leq r$
- Sufficient condition for local min

• gradient
$$\nabla f(\mathbf{x}) = \begin{bmatrix} \frac{\partial f}{\partial x_1} \\ \vdots \\ \frac{\partial f}{\partial x_n} \end{bmatrix} = \mathbf{0}$$
: local stationarity

$$\bullet \text{ Hessian } \pmb{\nabla}^2 f(\mathbf{x}) = \begin{bmatrix} \frac{\partial^2 f}{\partial x_1^2} & \frac{\partial^2 f}{\partial x_1 \partial x_2} & \cdots & \frac{\partial^2 f}{\partial x_1 \partial x_n} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{\partial^2 f}{\partial x_n \partial x_1} & \frac{\partial^2 f}{\partial x_n \partial x_2} & \cdots & \frac{\partial^2 f}{\partial x_n^2} \end{bmatrix} \text{ positive semidefinite }$$

• Gradient condition is sufficient if the function $f(\mathbf{x})$ is convex.

Unconstrained nonlinear optimization

Optimization problem

$$\min_{\mathbf{x} \in \mathbb{R}^2} x_1^2 + 3x_1x_2 + x_2^3$$

Gradient

$$\mathbf{\nabla} f(\mathbf{x}) = \begin{bmatrix} 2x_1 + 3x_2 \\ 3x_1 + 3x_2^2 \end{bmatrix} = \mathbf{0} \quad \Rightarrow \quad \mathbf{x} = \mathbf{0}, \quad \begin{bmatrix} -\frac{9}{4} \\ \frac{3}{2} \end{bmatrix}$$

- Hessian at **x**: $\mathbf{H} = \begin{bmatrix} 2 & 3 \\ 3 & 6x_2 \end{bmatrix}$
 - $\mathbf{x} = \mathbf{0}$: $\mathbf{H} = \begin{bmatrix} 2 & 3 \\ 3 & 0 \end{bmatrix}$. Not positive definite. Not local minimum.
 - $\mathbf{x} = \begin{bmatrix} -\frac{9}{4} \\ \frac{3}{2} \end{bmatrix}$: $\mathbf{H} = \begin{bmatrix} 2 & 3 \\ 3 & 9 \end{bmatrix}$. Positive semidefinite. Local minimum

Lagrangian method

• Constrained optimization problem

$$\begin{aligned} \max_{\mathbf{x} \in \mathbb{R}^2} & 2\ln(1+x_1) + 4\ln(1+x_2), \\ \text{s.t.} & x_1 + x_2 = 12 \end{aligned}$$

- Convex problem. But constraints make the problem hard to solve.
- Form a Lagrangian function

$$\mathcal{L}(\mathbf{x}, v) = 2\ln(1 + x_1) + 4\ln(1 + x_2) - v(x_1 + x_2 - 12)$$

 \bullet Compute the stationary points of the Lagrangian as a function of \emph{v}

$$\nabla \mathcal{L}(\mathbf{x}, v) = \begin{bmatrix} \frac{2}{1+x_1} - v \\ \frac{4}{1+x_2} - v \end{bmatrix} = \mathbf{0} \quad \Rightarrow \quad x_1 = \frac{2}{v} - 1, \quad x_2 = \frac{4}{v} - 1$$

• Substituting in the constraint $x_1 + x_2 = 12$, we get

$$\frac{6}{v} = 14 \quad \Rightarrow \quad v = \frac{3}{7} \quad \Rightarrow \quad \mathbf{x} = \frac{1}{3} \begin{bmatrix} 11\\25 \end{bmatrix}$$

Portfolio Selection

Optimization problem

$$\max_{\mathbf{x}} \quad \boldsymbol{\mu}^{\top} \mathbf{x} - \lambda \mathbf{x}^{\top} \mathbf{V} \mathbf{x}$$
 s.t.
$$\mathbf{1}^{\top} \mathbf{x} = 1$$

Constraints make the problem hard!

• Lagrangian function

$$\mathcal{L}(\mathbf{x}, v) = \boldsymbol{\mu}^{\top} \mathbf{x} - \lambda \mathbf{x}^{\top} \mathbf{V} \mathbf{x} - v(\mathbf{1}^{\top} \mathbf{x} - 1)$$

• Solve for the maximum value with no constraints

$$\nabla_x \mathcal{L}(\mathbf{x}, v) = \mu - 2\lambda \mathbf{V} \mathbf{x} - v \mathbf{1} = \mathbf{0} \quad \Rightarrow \quad \mathbf{x} = \frac{1}{2\lambda} \cdot \mathbf{V}^{-1} (\mu - v \mathbf{1})$$

ullet Solve for v from the constraint

$$\mathbf{1}^{\top} \mathbf{x} = 1 \quad \Rightarrow \quad \mathbf{1}^{\top} \mathbf{V}^{-1} (\boldsymbol{\mu} - v \mathbf{1}) = 2\lambda \quad \Rightarrow \quad v = \frac{\mathbf{1}^{\top} \mathbf{V}^{-1} \boldsymbol{\mu} - 2\lambda}{\mathbf{1}^{\top} \mathbf{V}^{-1} \mathbf{1}}$$

 \bullet Substitute back in the expression for x