

Q.1 Can you explain the logic and working of the Tower of Hanoi algorithm by writing a Java program?

How does the recursion work, and how are the movements of disks between rods accomplished?

The **Tower of Hanoi** is a classic problem in recursion. The problem involves three rods and n disks of different sizes, initially stacked in decreasing order on the first rod. The goal is to move all the disks to the third rod following these rules:

1. Only one disk can be moved at a time.
2. A disk can only be placed on top of a larger disk or on an empty rod.
3. Only the topmost disk of a stack can be moved.

Recursive Logic:

1. Move $n-1$ disks from **source** to **auxiliary** rod using **destination** as a helper.
2. Move the largest disk (n th disk) directly from **source** to **destination**.
3. Move the $n-1$ disks from **auxiliary** to **destination** using **source** as a helper.

This forms the recursive breakdown:

$$T(n) = 2T(n-1) + 1$$

Solving this recurrence gives:

$$T(n) = 2^n - 1$$

which represents the minimum number of moves needed.

Q.2 Given two strings word1 and word2, return the minimum number of operations required to convert word1 to word2.

Example 1:

Input: word1 = "horse", word2 = "ros"

Output: 3

Explanation:

horse -> rorse (replace 'h' with 'r')

rorse -> rose (remove 'r')

rose -> ros (remove 'e')

We define **edit distance** as the minimum number of operations to convert word1 to word2, using:

1. **Insertion** (Insert a character into word1).
 2. **Deletion** (Delete a character from word1).
 3. **Replacement** (Replace a character in word1).
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Dynamic Programming Approach:

1. **Define $dp[i][j]$** as the minimum operations required to convert the first i characters of word1 to the first j characters of word2.
2. **Base Case:**
 - If word1 is empty, we need j insertions.
 - If word2 is empty, we need i deletions.
3. **Recurrence Relation:**
 - If $word1[i-1] == word2[j-1]$, no operation is needed: $dp[i][j] = dp[i-1][j-1]$
 - Otherwise, take the minimum of:
 - **Insert** ($dp[i][j-1] + 1$)
 - **Delete** ($dp[i-1][j] + 1$)
 - **Replace** ($dp[i-1][j-1] + 1$)

Step-by-step conversion

1. **Replace** 'h' → 'r' → "rorse"
 2. **Delete** 'r' → "rose"
 3. **Delete** 'e' → "ros"
-

Time & Space Complexity

- **Time Complexity:** $O(m \times n)$, since we fill an $m \times n$ table.
- **Space Complexity:** $O(m \times n)$, as we use a 2D DP table.

Space Optimization

We can optimize space to $O(n)$ by using only two rows (prev and curr), since each row only depends on the previous row.

Example 2:

Input: word1 = "intention", word2 = "execution"

Output: 5

Explanation:

intention -> inention (remove 't')

inention -> enention (replace 'i' with 'e')

enention -> exention (replace 'n' with 'x')

exention -> exection (replace 'n' with 'c')

exection -> execution (insert 'u')

Input:

word1 = "intention"

word2 = "execution"

Step-by-Step Execution:

1. **Remove** 't': "intention" → "inention"
2. **Replace** 'i' with 'e': "inention" → "enention"
3. **Replace** 'n' with 'x': "enention" → "exention"
4. **Replace** 'n' with 'c': "exention" → "exection"
5. **Insert** 'u': "exection" → "execution"

Total operations: 5

Q. 3 Print the max value of the array [13, 1, -3, 22, 5].

Method 1: Using max()

arr = [13, 1, -3, 22, 5]

print("Max value:", max(arr))

Method 2: Using a Loop

arr = [13, 1, -3, 22, 5]

```
max_val = arr[0] # Initialize with the first element
```

```
for num in arr:
```

```
    if num > max_val:
```

```
        max_val = num
```

```
print("Max value:", max_val)
```

Q.4 Find the sum of the values of the array [92, 23, 15, -20, 10].

Method 1: Using sum()

python

CopyEdit

```
arr = [92, 23, 15, -20, 10]
```

```
print("Sum of array values:", sum(arr))
```

Method 2: Using a Loop

python

CopyEdit

```
arr = [92, 23, 15, -20, 10]
```

```
total_sum = 0 # Initialize sum
```

```
for num in arr:
```

```
    total_sum += num # Add each element to total_sum
```

```
print("Sum of array values:", total_sum)
```

Q.5 Given a number n. Print if it is an armstrong number or not. An armstrong number is a number if the sum of every digit in that number raised to the power of total digits in that number is equal to the number.

Example : $153 = 1^3 + 5^3 + 3^3 = 1 + 125 + 27 = 153$ hence 153 is an armstrong number. (Easy)

Input1 : 153

Output1 : Yes

Input 2 : 134

Output2 : No

```
def is_armstrong_number(n: int) -> str:
```

```
    # Convert number to string to easily iterate through digits
```

```
    digits = str(n)
```

```
    num_digits = len(digits)
```

```
    # Sum of digits raised to the power of the number of digits
```

```
    sum_of_powers = sum(int(digit) ** num_digits for digit in digits)
```

```
    # Check if the sum equals the original number
```

```
    if sum_of_powers == n:
```

```
        return "Yes"
```

```
    else:
```

```
        return "No"
```

```
# Example Usage
```

```
input1 = 153
```

```
input2 = 134
```

```
print(is_armstrong_number(input1)) # Output: Yes
```

```
print(is_armstrong_number(input2)) # Output: No
```