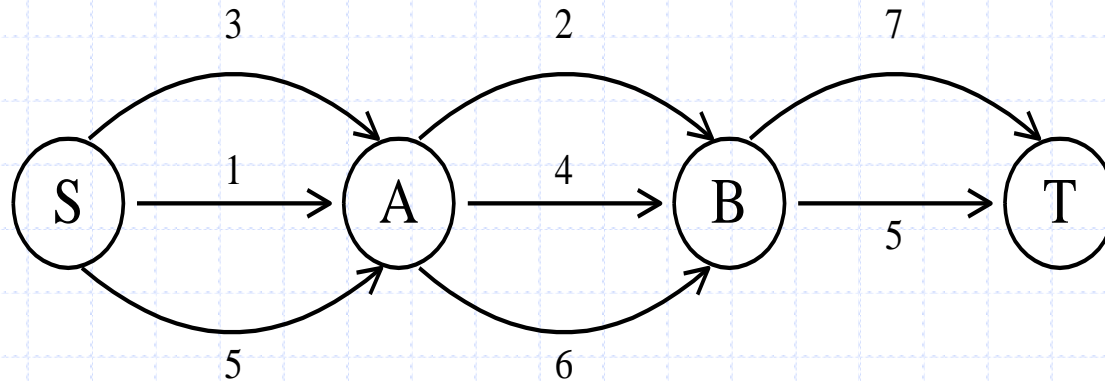




Multistage graph

The shortest path

◆ To find a shortest path in a multi-stage graph

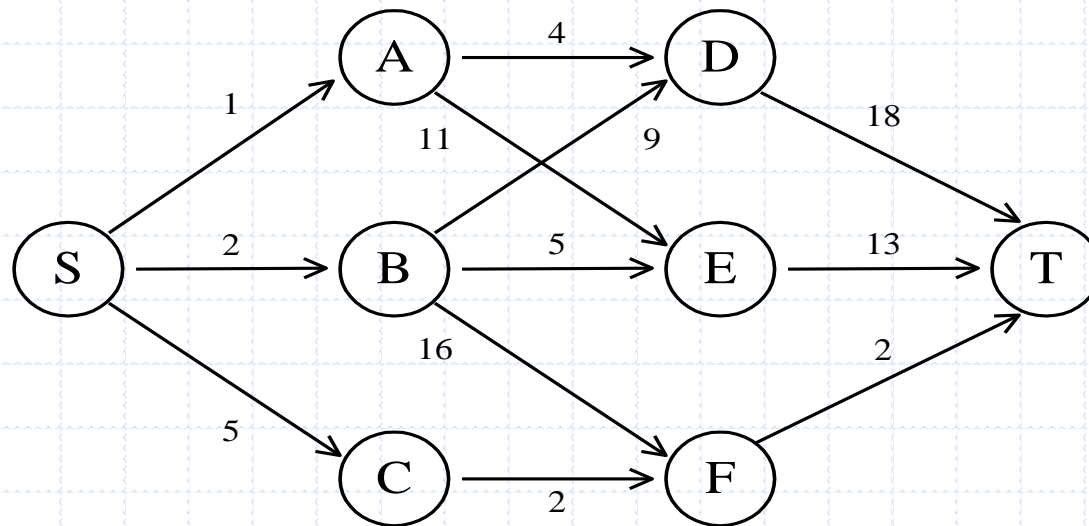


◆ Apply the greedy method :
the shortest path from S to T :

$$1 + 2 + 5 = 8$$

The shortest path in multistage graphs

◆ e.g.



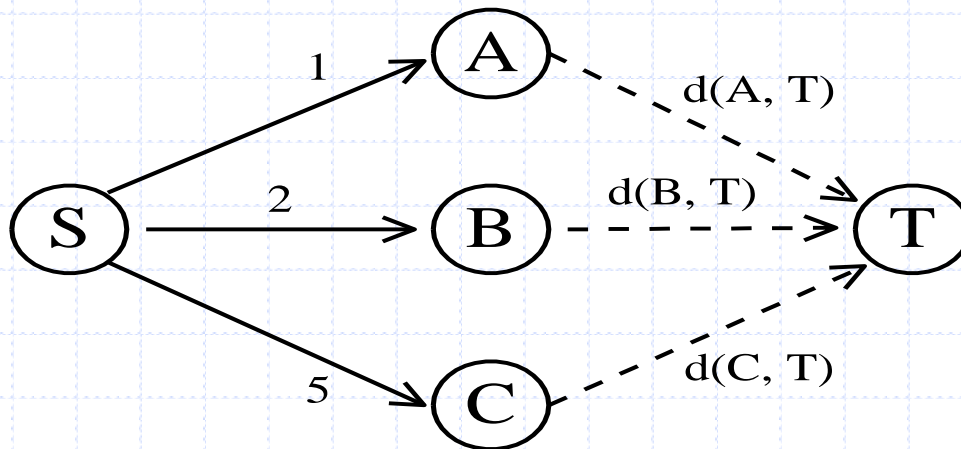
◆ The greedy method can not be applied to this case: (S, A, D, T) $1+4+18 = 23$.

◆ The real shortest path is:

(S, C, F, T) $5+2+2 = 9$.

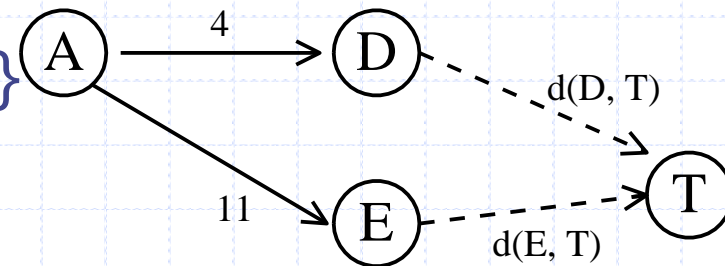
Dynamic programming approach

- ◆ Dynamic programming approach (forward approach):



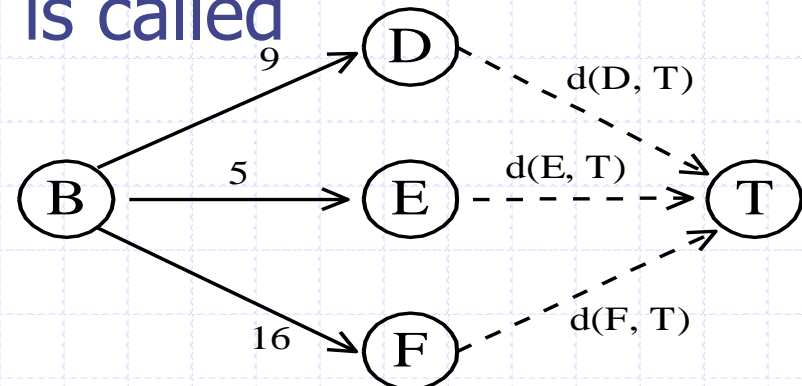
- ◆ $d(S, T) = \min\{1+d(A, T), 2+d(B, T), 5+d(C, T)\}$

- $d(A, T) = \min\{4+d(D, T), 11+d(E, T)\}$
 $= \min\{4+18, 11+13\} = 22.$



Dynamic programming

- ◆ $d(B, T) = \min\{9+d(D, T), 5+d(E, T), 16+d(F, T)\}$
 $= \min\{9+18, 5+13, 16+2\} = 18.$
- ◆ $d(C, T) = \min\{2+d(F, T)\} = 2+2 = 4$
- ◆ $d(S, T) = \min\{1+d(A, T), 2+d(B, T), 5+d(C, T)\}$
 $= \min\{1+22, 2+18, 5+4\} = 9.$
- ◆ The above way of reasoning is called backward reasoning.



Backward approach (forward reasoning)

◆ $d(S, A) = 1$


$d(S, B) = 2$

$d(S, C) = 5$

◆ $d(S, D) = \min\{d(S, A) + d(A, D), d(S, B) + d(B, D)\}$
 $= \min\{1 + 4, 2 + 9\} = 5$

$d(S, E) = \min\{d(S, A) + d(A, E), d(S, B) + d(B, E)\}$
 $= \min\{1 + 11, 2 + 5\} = 7$

$d(S, F) = \min\{d(S, A) + d(A, F), d(S, B) + d(B, F)\}$
 $= \min\{2 + 16, 5 + 2\} = 7$


$$\begin{aligned}\diamond d(S,T) &= \min\{d(S, D)+d(D, T), d(S,E)+ \\ &\quad d(E,T), d(S, F)+d(F, T)\} \\ &= \min\{ 5+18, 7+13, 7+2 \} \\ &= 9\end{aligned}$$

Principle of optimality

- ◆ Principle of optimality: Suppose that in solving a problem, we have to make a sequence of decisions D_1, D_2, \dots, D_n . If this sequence is optimal, then the last k decisions, $1 < k < n$ must be optimal.
- ◆ e.g. the shortest path problem
If i, i_1, i_2, \dots, j is a shortest path from i to j , then i_1, i_2, \dots, j must be a shortest path from i_1 to j
- ◆ In summary, if a problem can be described by a multistage graph, then it can be solved by dynamic programming.

Dynamic programming

- ◆ Forward approach and backward approach:
 - Note that if the recurrence relations are formulated using the forward approach then the relations are solved backwards . i.e., beginning with the last decision
 - On the other hand if the relations are formulated using the backward approach, they are solved forwards.
- ◆ To solve a problem by using dynamic programming:
 - Find out the recurrence relations.
 - Represent the problem by a multistage graph.

The resource allocation problem

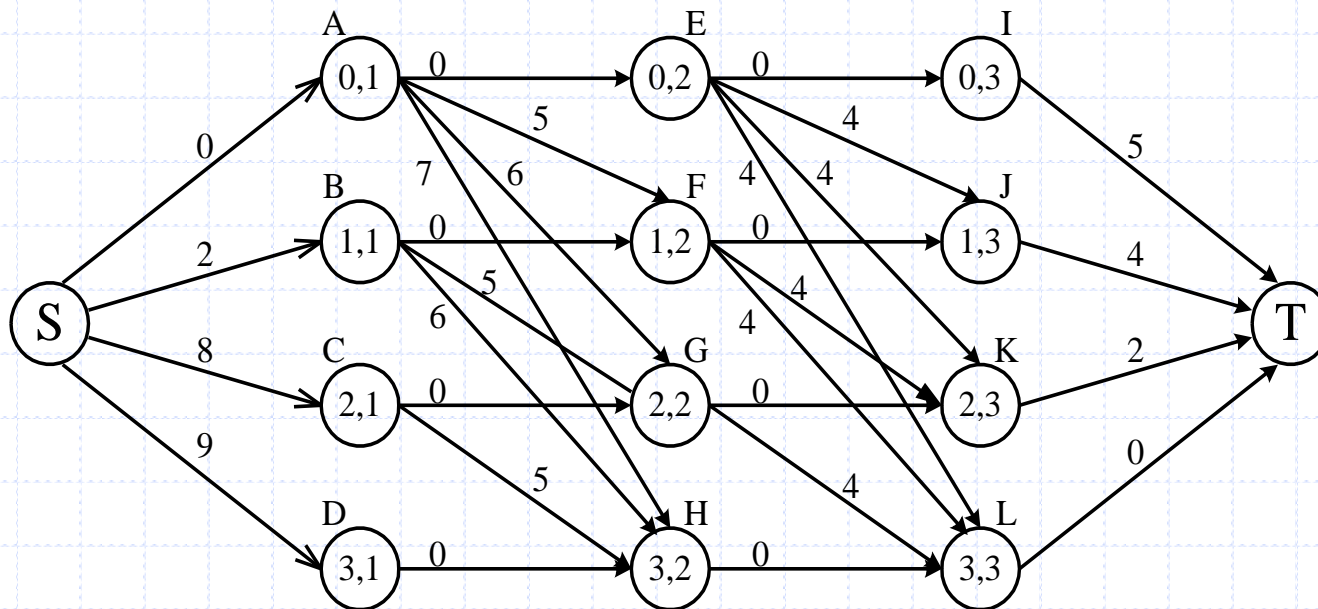
◆ m resources, n projects

profit $p(i, j)$: j resources are allocated to project i .

maximize the total profit.

Resource Project	1	2	3
1	2	8	9
2	5	6	7
3	4	4	4
4	2	4	5

The multistage graph solution



- ◆ The resource allocation problem can be described as a multistage graph.
- ◆ (i, j) : i resources allocated to projects 1, 2, ..., j
e.g. node $H=(3, 2)$: 3 resources allocated to projects 1, 2.

◆ Find the longest path from S to T :

(S, C, H, L, T), $8+5+0+0=13$

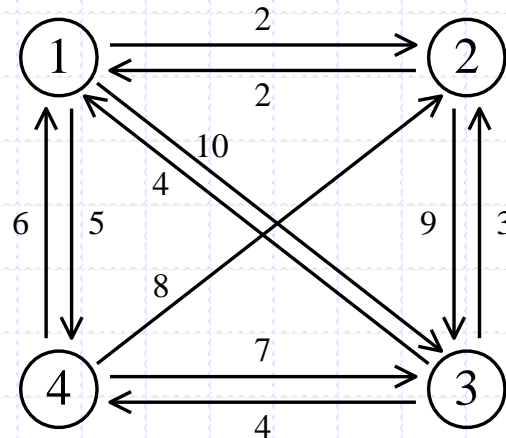
2 resources allocated to project 1.

1 resource allocated to project 2.

0 resource allocated to projects 3, 4.

The traveling salesperson (TSP) problem

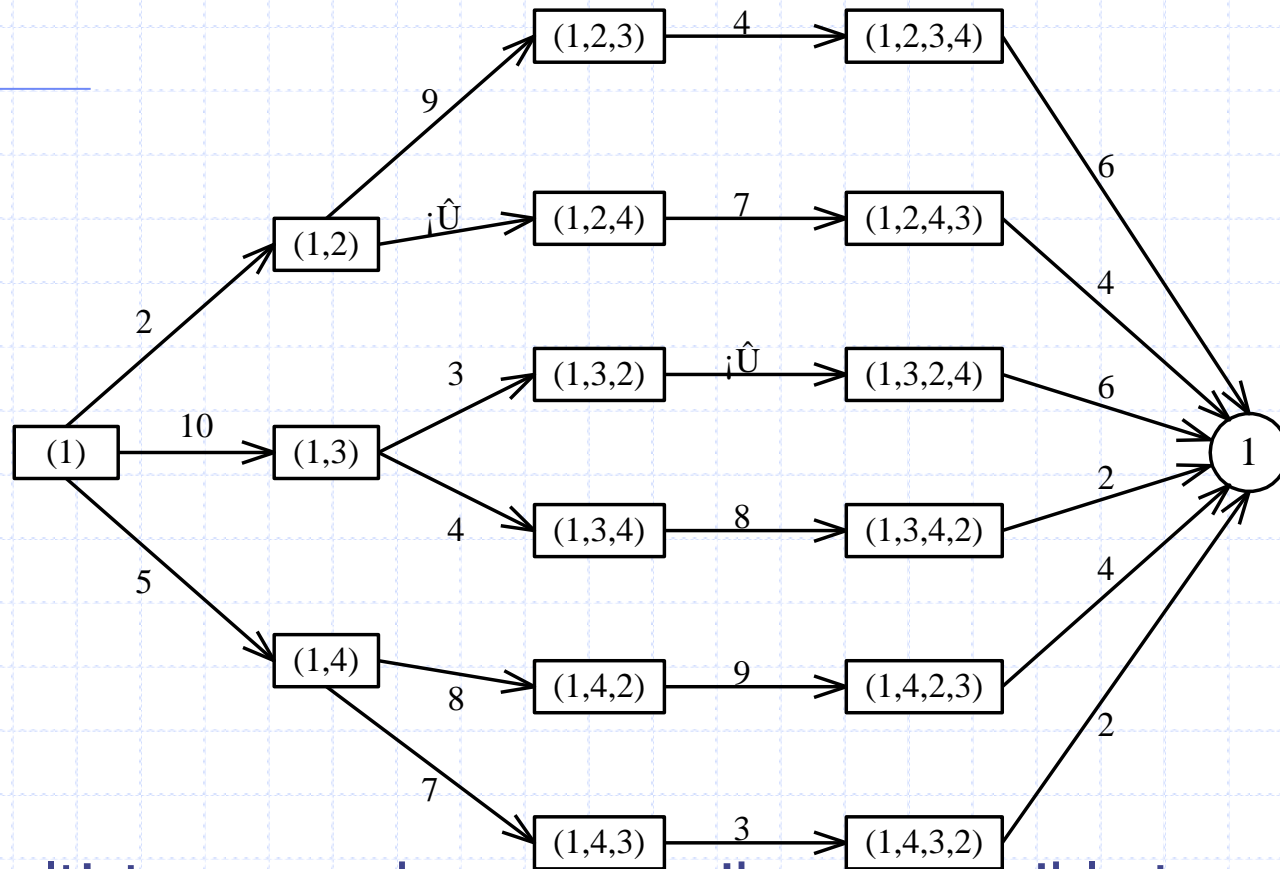
◆ e.g. a directed graph :



◆ Cost matrix:

	1	2	3	4
1	∞	2	10	5
2	2	∞	9	∞
3	4	3	∞	4
4	6	8	7	∞

The multistage graph solution



- ◆ A multistage graph can describe all possible tours of a directed graph.
- ◆ Find the shortest path: $(1, 4, 3, 2, 1)$ $5+7+3+2=17$