

HW2

1. a.) $T(n) = 3T(n/27) + 1$

$a=3, b=27, f(n)=1, \log_{27}(3) = 1/3$

$f(n) \leq O(n^{1/3-\epsilon}), \epsilon = 1/3$

$= O(n^0)$

$= O(1)$

$\therefore T(n) = \Theta(n^{1/3})$

b.) $T(n) = 7T(n/8) + \lg n$

$a=7, b=8, f(n) = \lg n, \lg_8(7) \approx 0.94$

$f(n) \leq O(n^{0.94-\epsilon}), \epsilon = 0.1$

$\leq O(n^{0.84})$

$\therefore T(n) = \Theta(n^{0.94})$

c.) $T(n) = 2T(n/4) + n$

$a=2, b=4, f(n)=n, \lg_4(2) = 1/2$

$f(n) \geq \Omega(n^{1/2+\epsilon}), \epsilon = 1/2$

$= \Omega(n)$

$2(\frac{n}{4}) \leq cn \Rightarrow \frac{n}{2} \leq cn, c = 1/2$

$= \frac{n}{2}$

$\therefore T(n) = \Theta(n)$

Hw 2

1. d.) $T(n) = 2T(n/4) + n^2$

$$a=2, b=4, f(n)=n^2, \lg_4(2) = 1/2$$

$$f(n) \geq \Omega(n^{1/2+\epsilon}), \epsilon = 3/2$$

$$= \Omega(n^2)$$

$$2\left(\frac{n}{4}\right)^2 \leq cn^2 \Rightarrow \frac{n^2}{8} \leq cn^2, c = 1/8$$
$$= \frac{n^2}{8}$$

$$\therefore T(n) = \Theta(n^2)$$

e.) $T(n) = 2T(n/4) + n^{1/2} \lg n$

$$a=2, b=4, f(n) = \sqrt{n} \lg n, \lg_4(2) = 1/2$$

$$f(n) = \Theta(\sqrt{n})$$

$$\therefore T(n) = \Theta(\sqrt{n} \lg^2 n)$$

2. There is no change, $27 > 17 \wedge 3$, no swaps

HWZ

3. Build Maxheap(A) = [25, 13, 20, 8, 7, 17, 2, 5, 4]

Then we ~~are~~ swap last and first
& run max heapify

$$A = [2, 4, 5, 7, 8, 13, 17, 20, 25]$$

which is sorted array

4. $T(n) = 2T(n/2 - 3) + n$

$$a=2, b=2, f(n)=n, \lg_2(2)=1, k=0$$

$$f(n) = \Theta(n^1)$$

$$\therefore T(n) = \Theta(n \lg n) \text{ which is in } \Omega(n \lg n)$$

5. a) there are 23 swaps, 5 in build max heap
and 18 are from heapsort re-orders. I got
these numbers from programming it.

b) There will be more comparisons but same
number of swaps.

c) It would increase to $O(n^2)$ because of the
extra comparisons.

HWZ

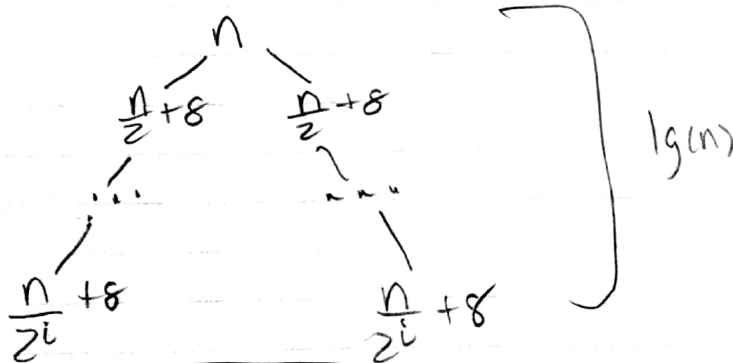
6. a) $T(n) = 2T(n/2) + \sin(n)$? No

↳ Can't use, $\sin(n)$ range is $[-1, 1]$, and running time can't be negative.

b) $T(n) = T(n/2) + n\sin(n) + 2n$? Yes
 $= T(n/2) + n(2 + \sin(n))$

Yes because $n(2 + \sin(n))$ will always be a positive number, unlike a)

7. $T(n) = 2T(n/2 + 8) + n$



$\therefore T(n) = \Theta(n \lg n)$, upper and lower bound

HW2

8. Loop invariant: The start of every iter of the for loop inside the Build Max Heap is a root of Max Heap

Any thing below $\lfloor \frac{n}{2} \rfloor$ has no leaf and doesn't need max heapify

When i is decremented after each iter down to 0 establishes loop invariant.

\therefore When $i=0$ we know the loop terminates and know there is a max heap