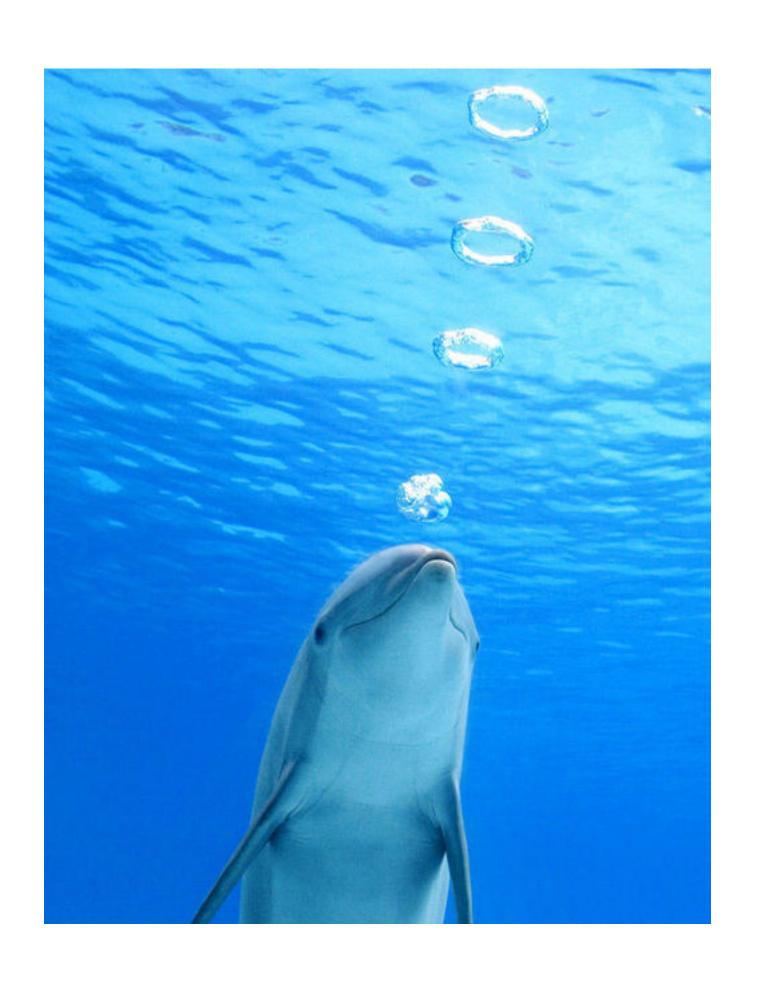
本节内容

知识总览 冒泡排序 交换排序 (快速排序)

基于"交换"的排序:根据序列中两个元素关键字的比较结果来对换这两个记录在序列中的位置

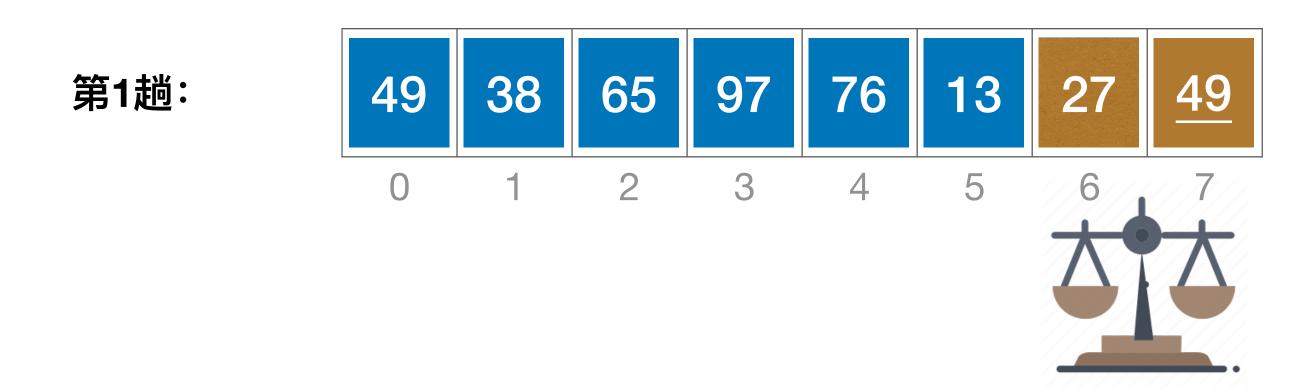
冒泡?

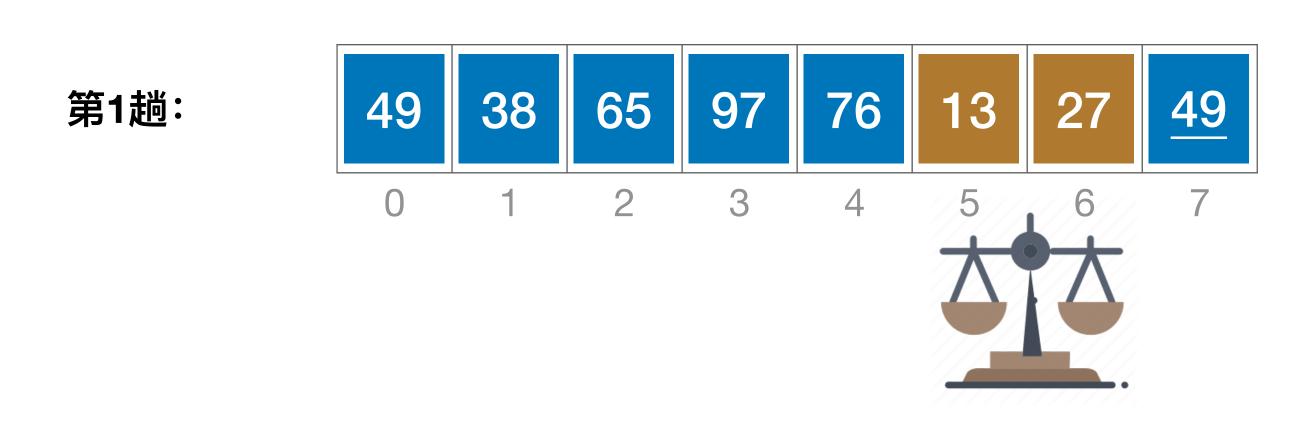


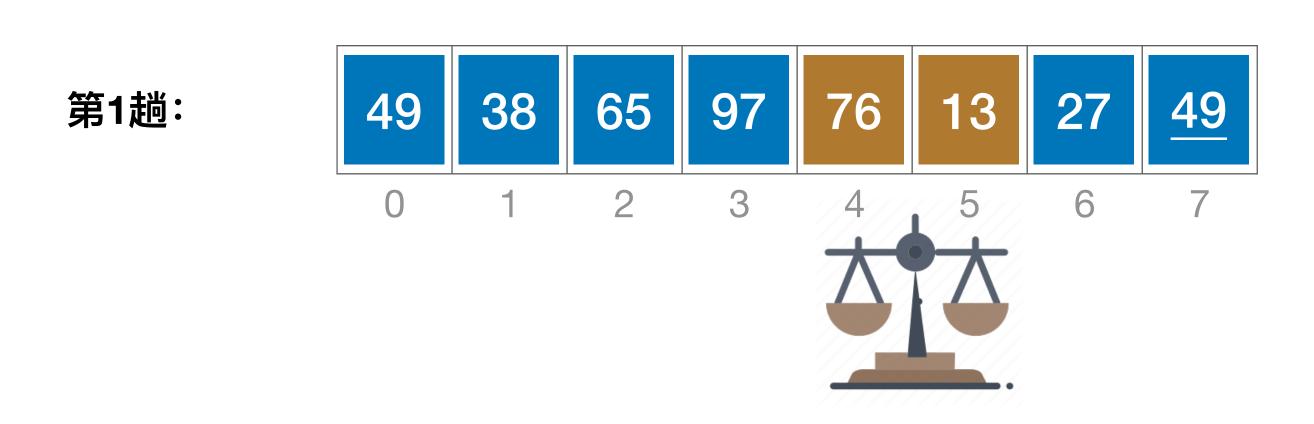
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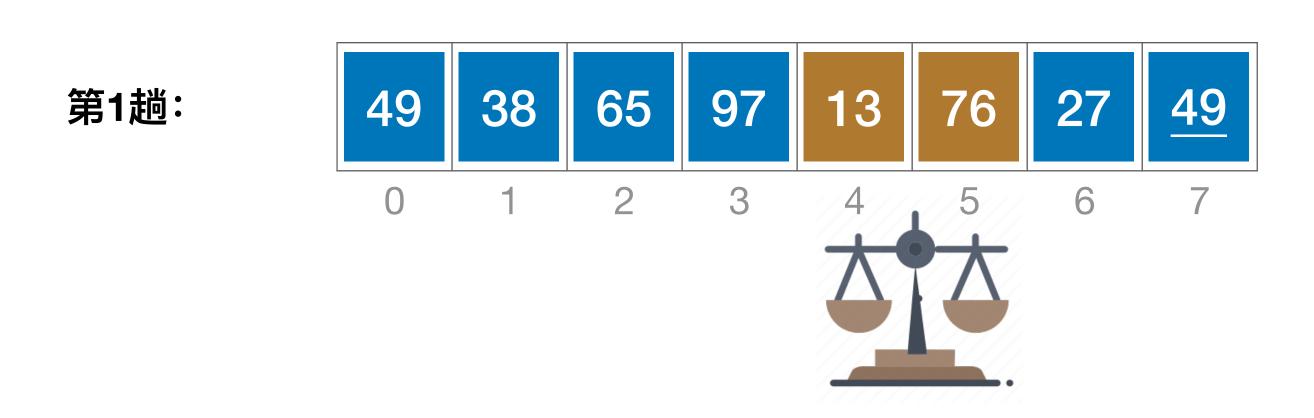


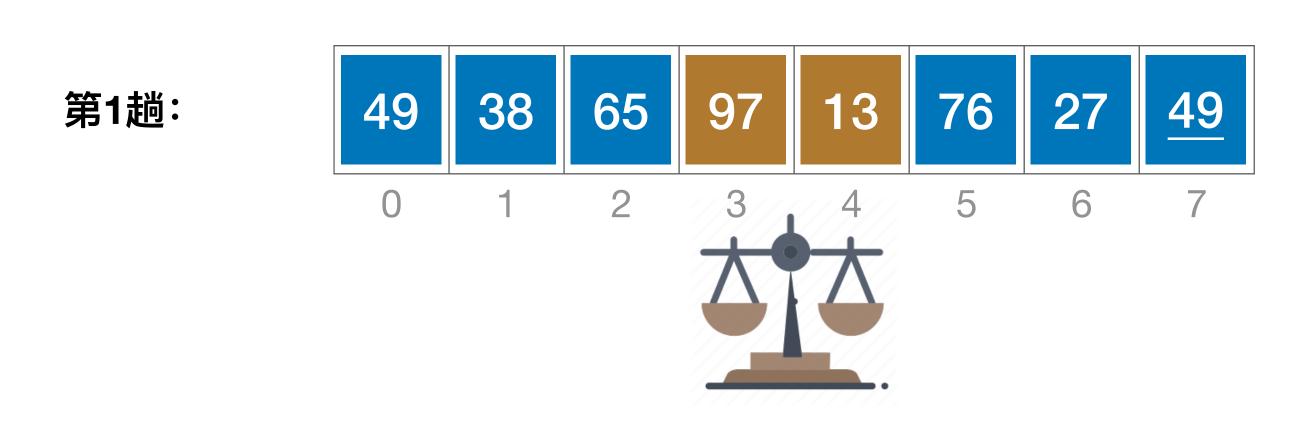
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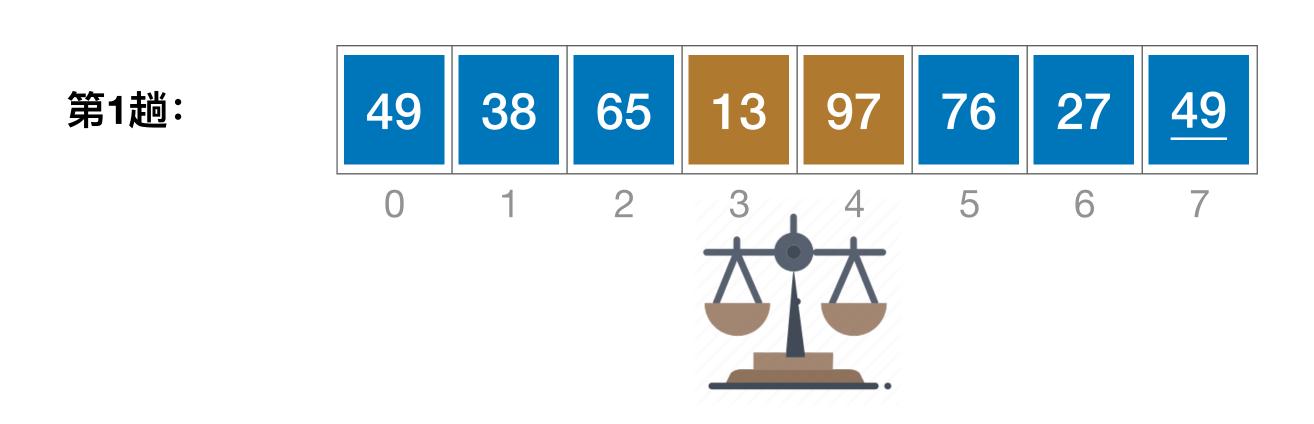


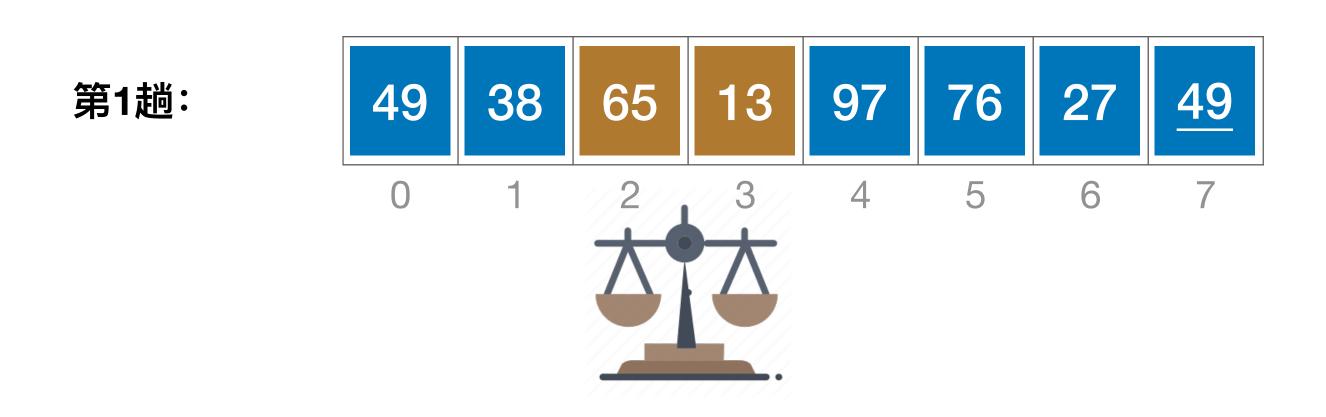


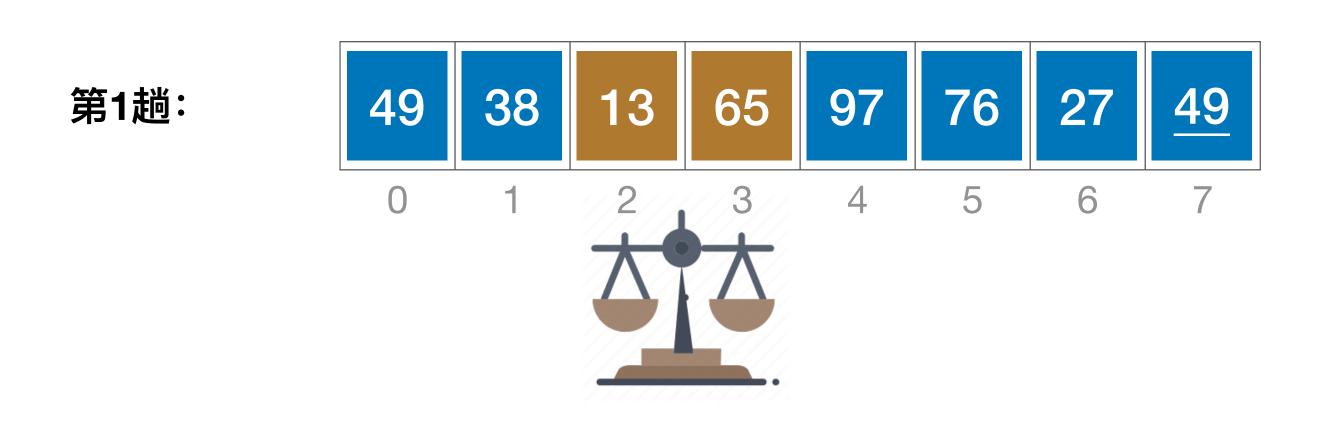


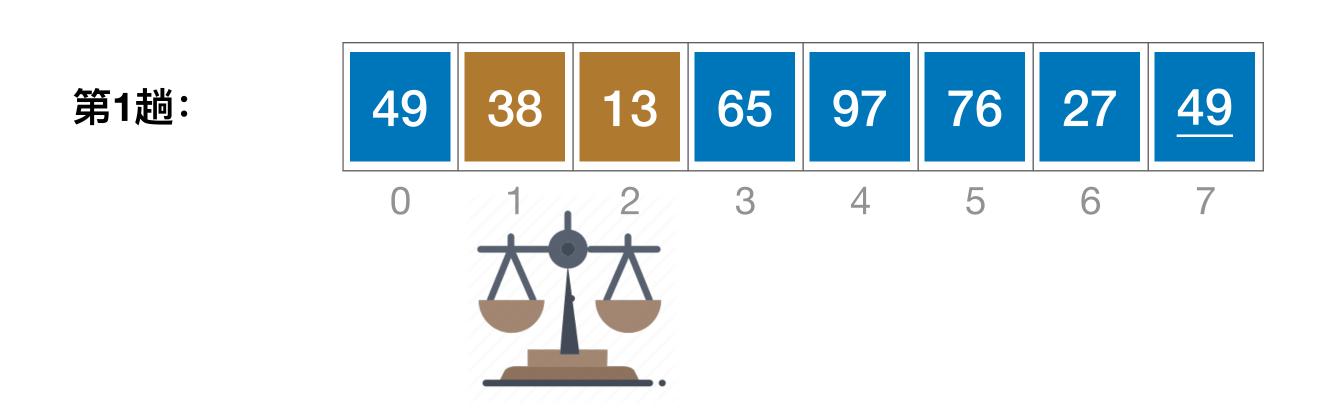


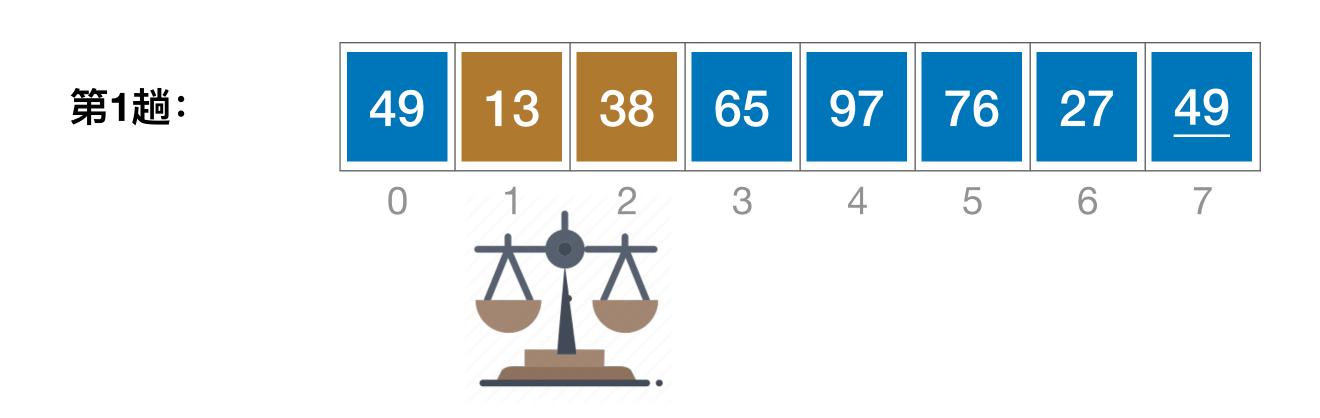


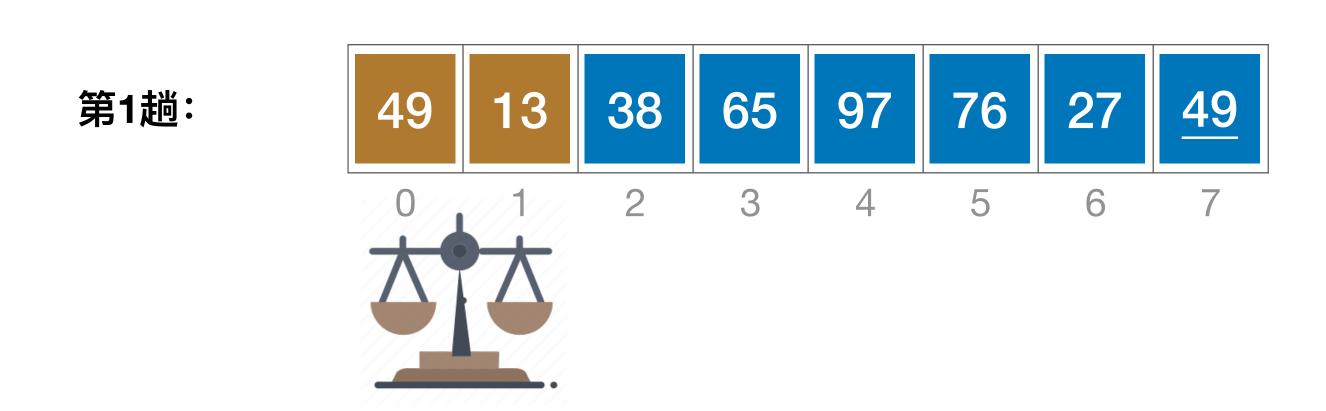


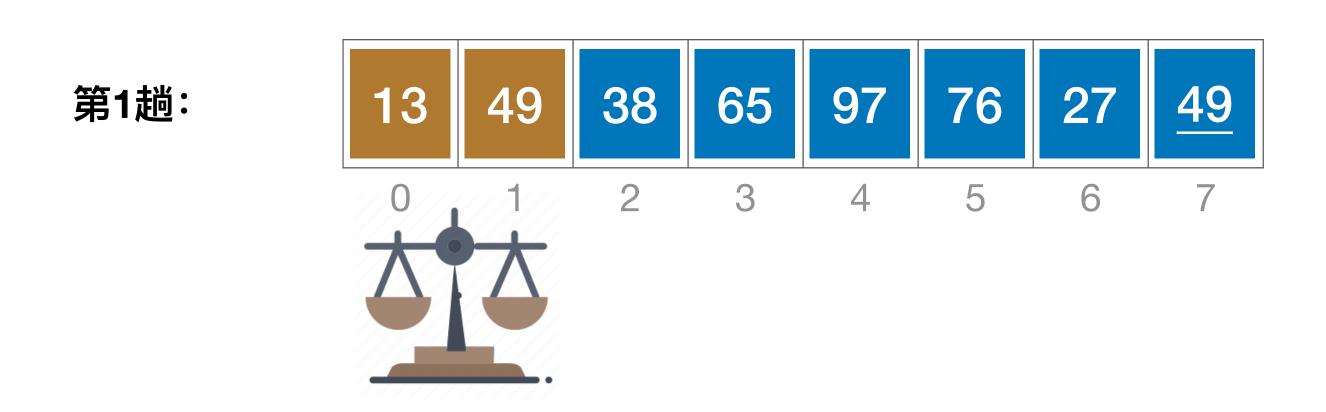






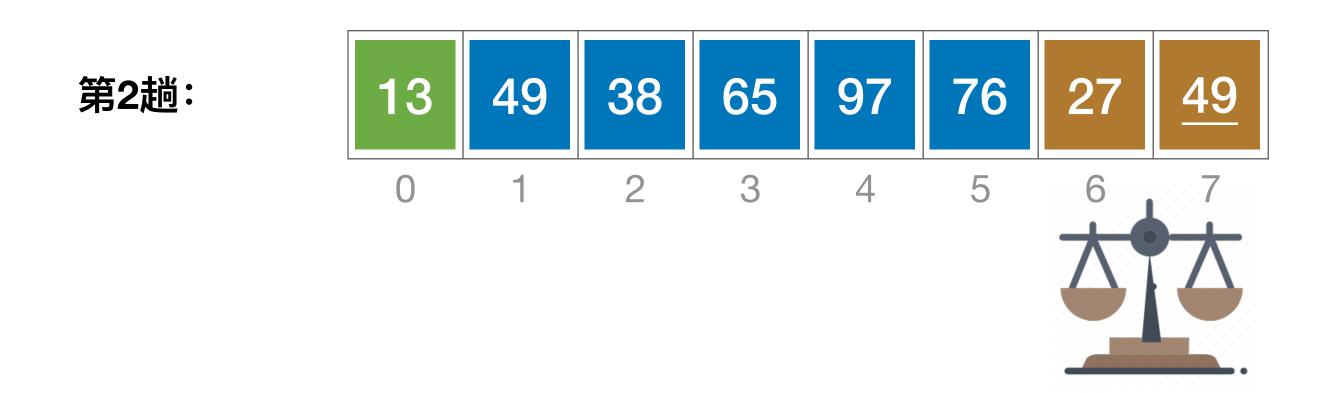


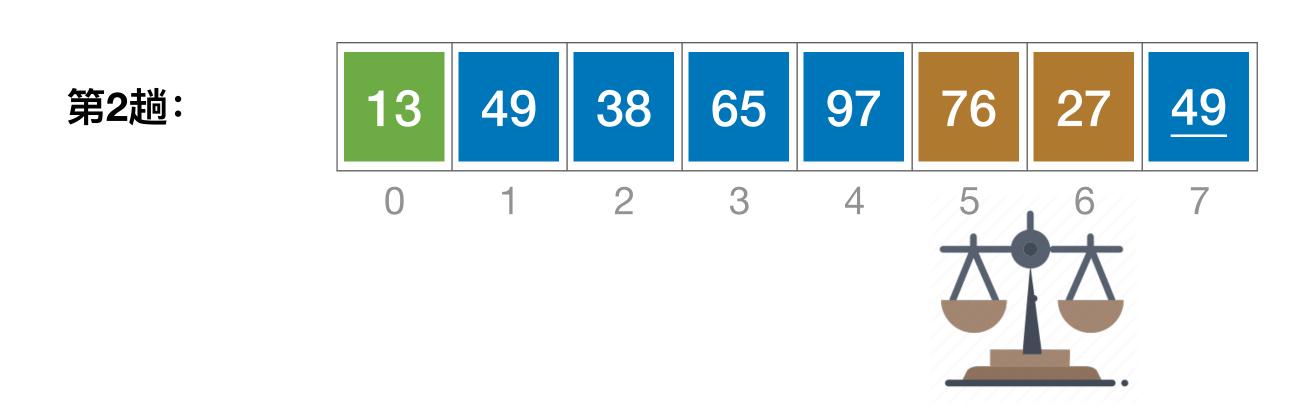


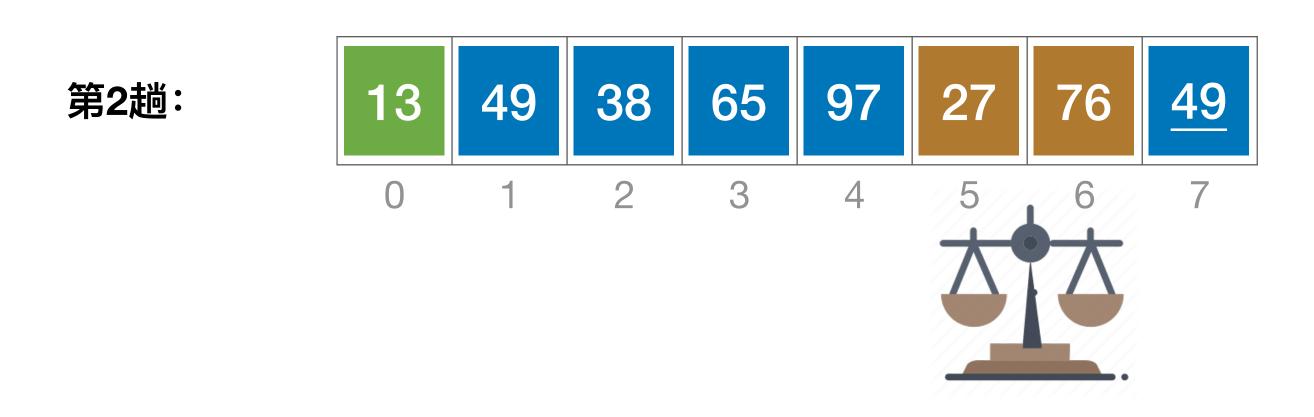


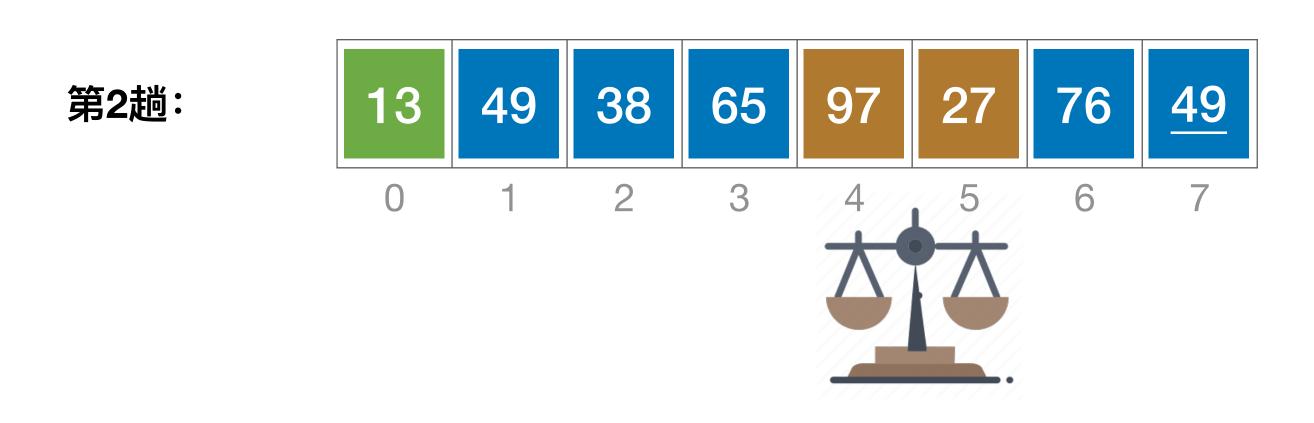


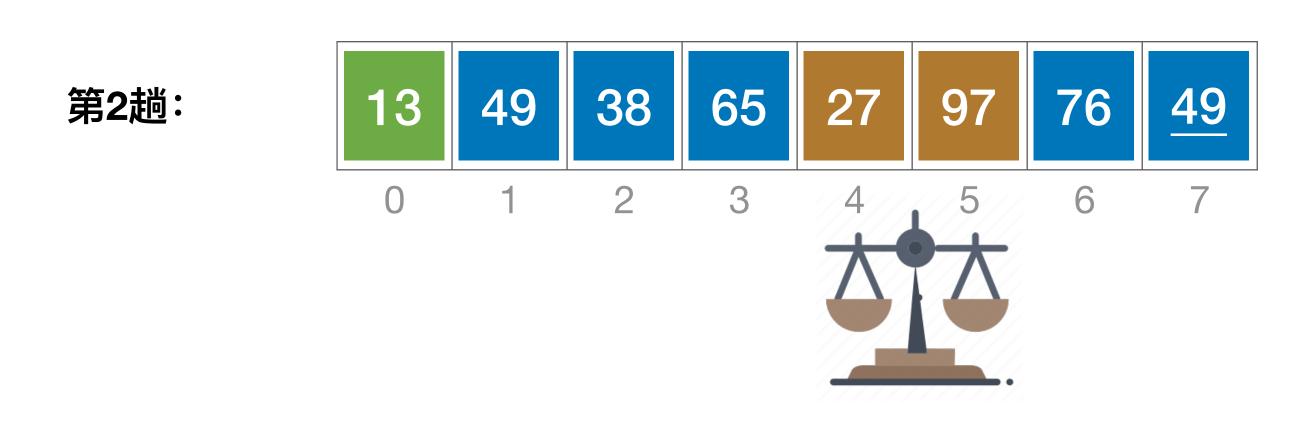
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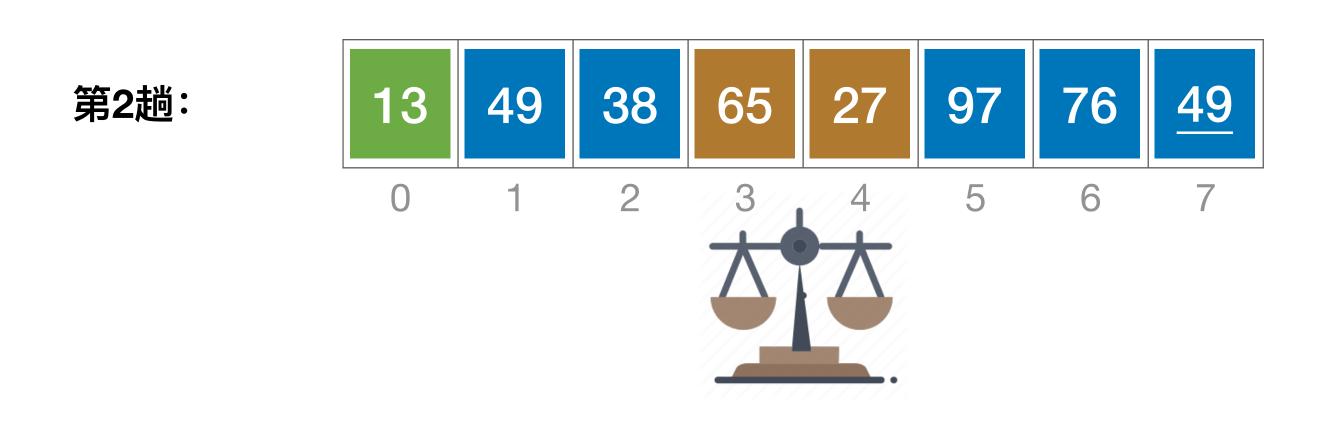


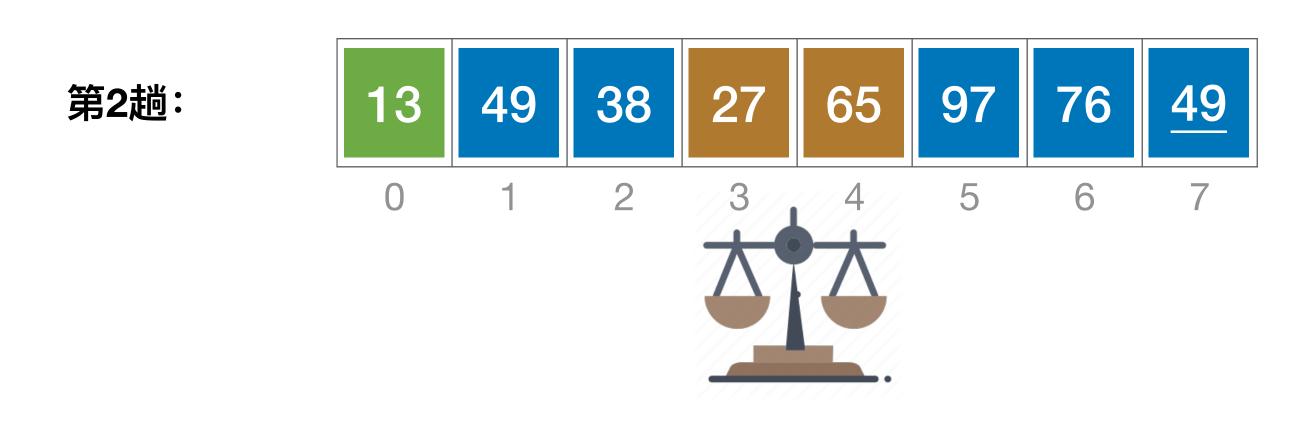


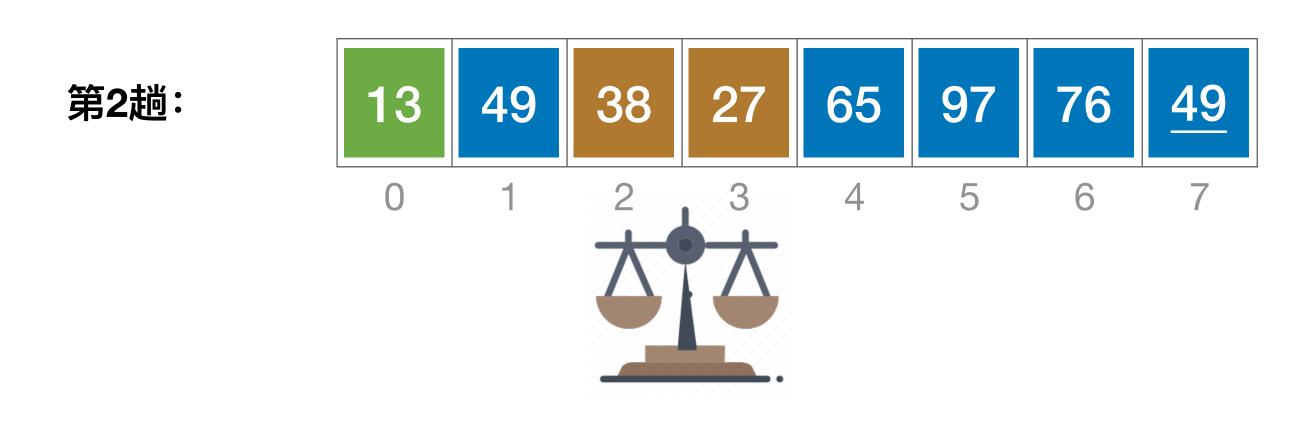


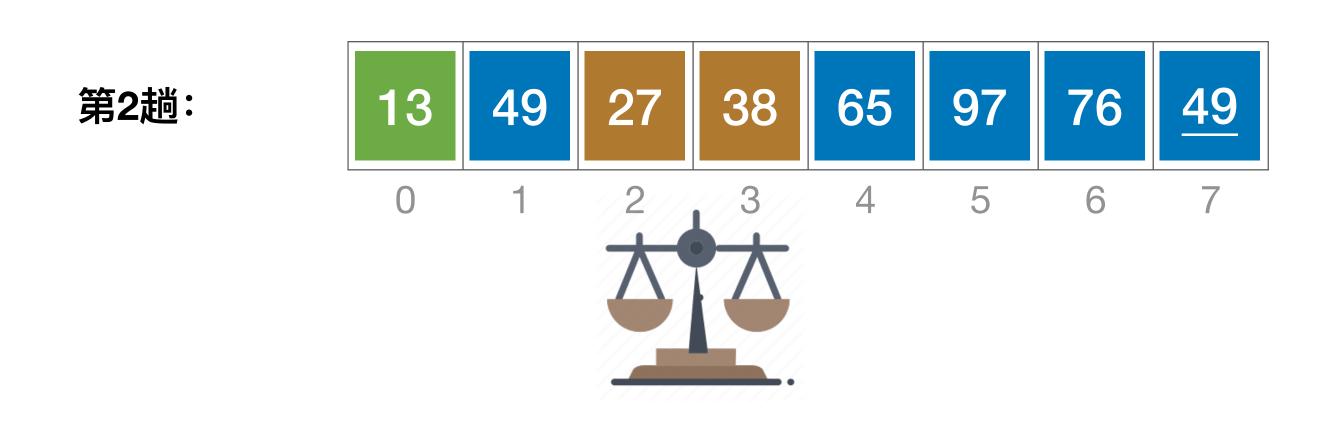


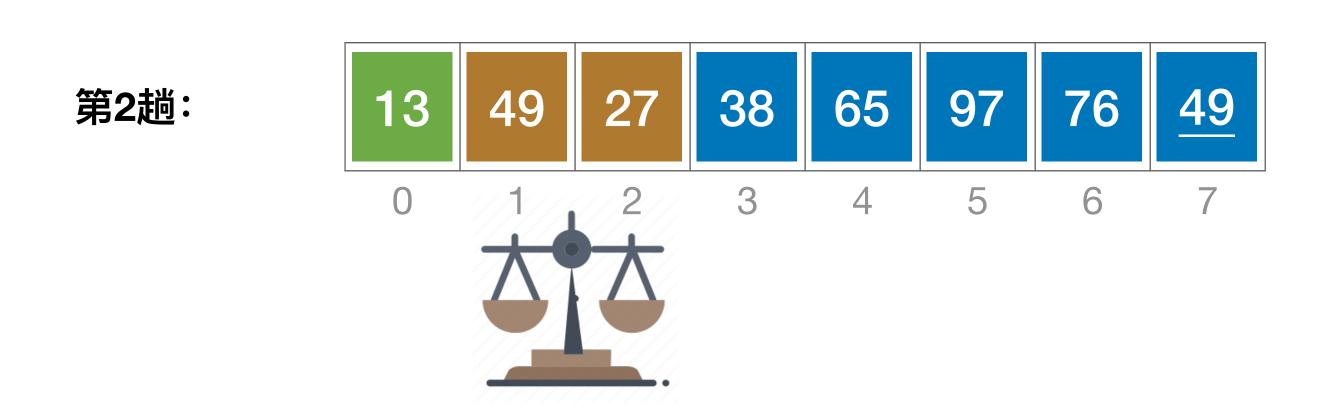


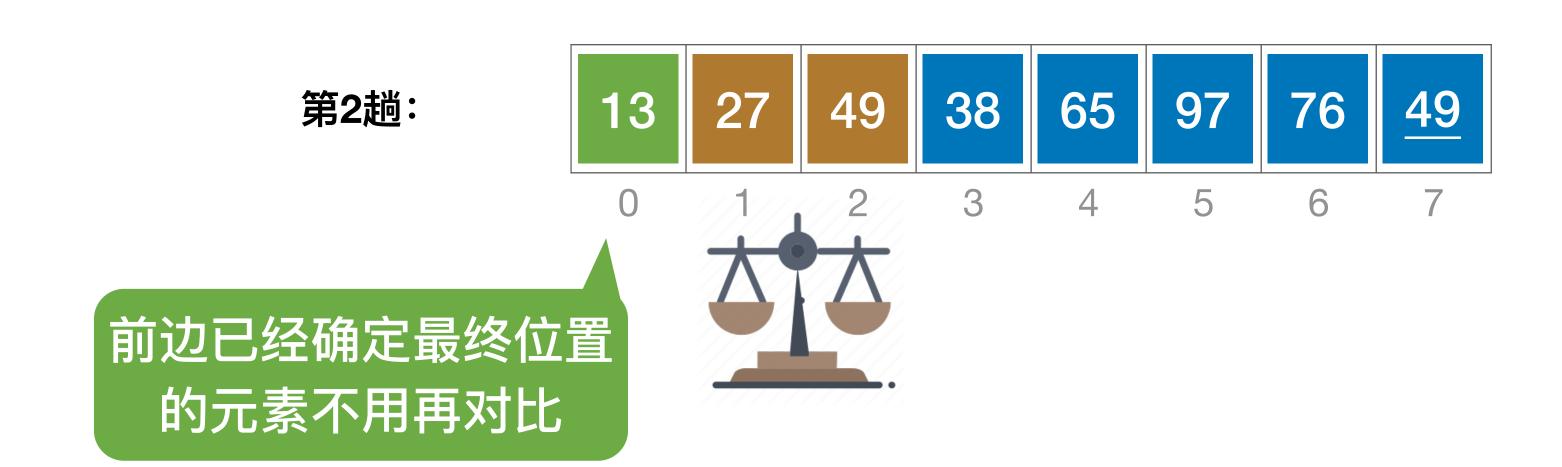








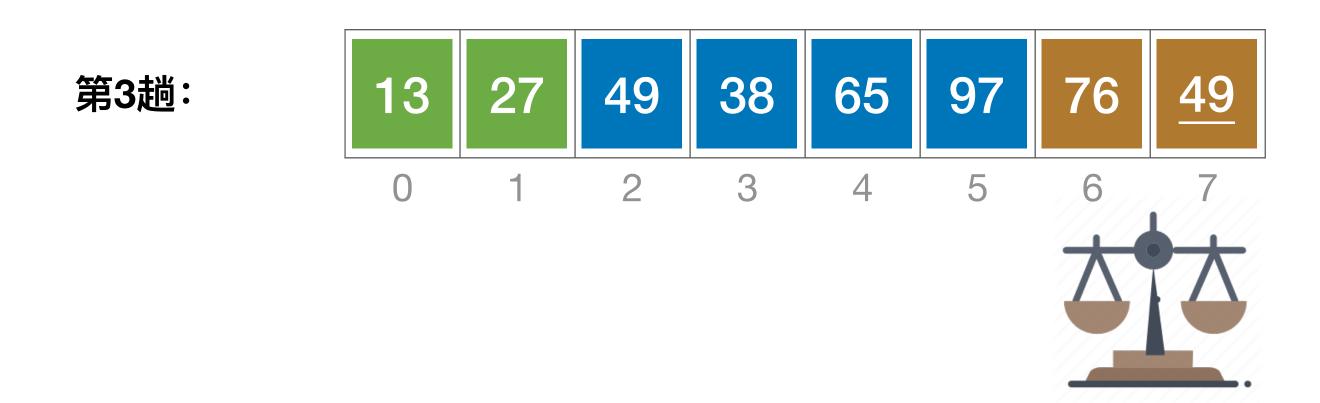


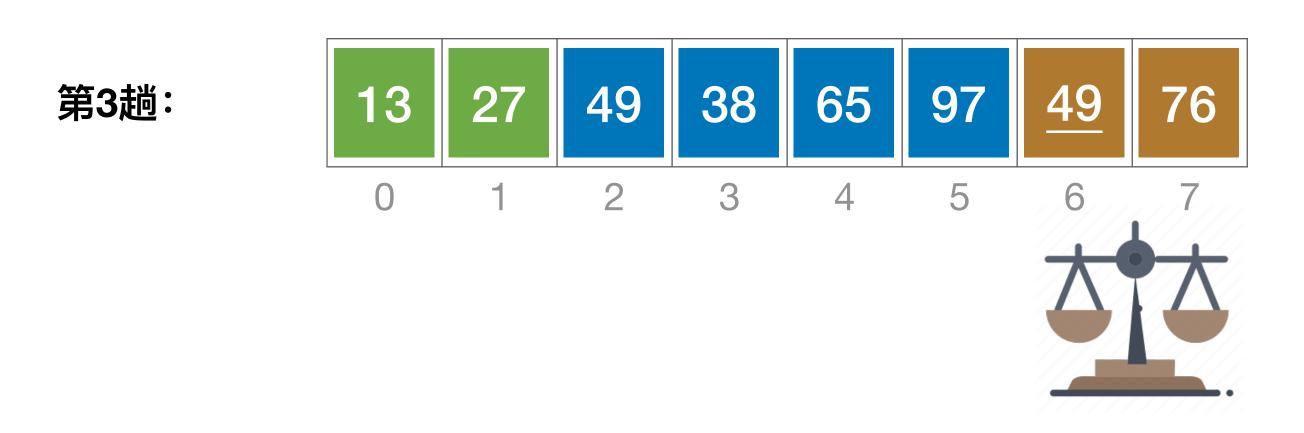


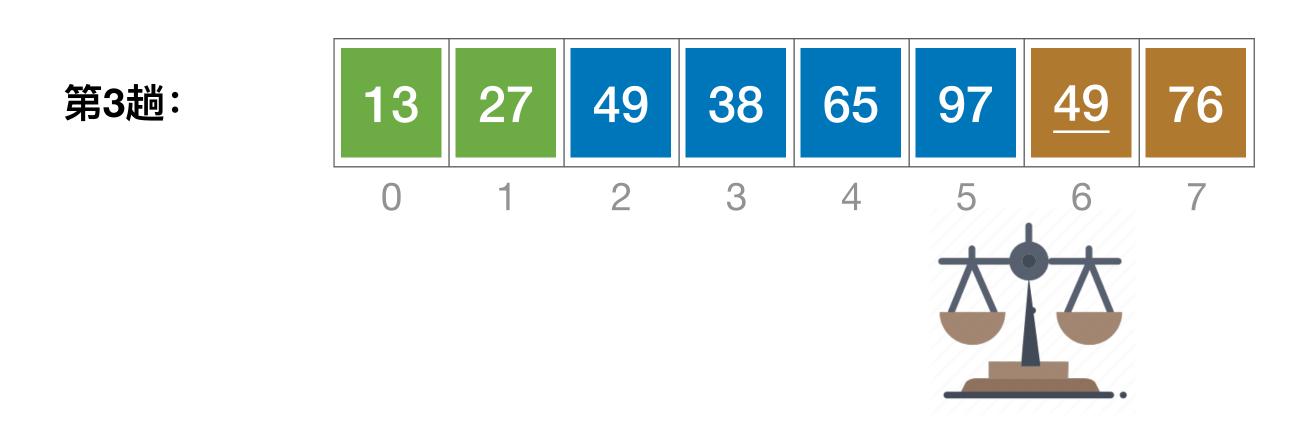


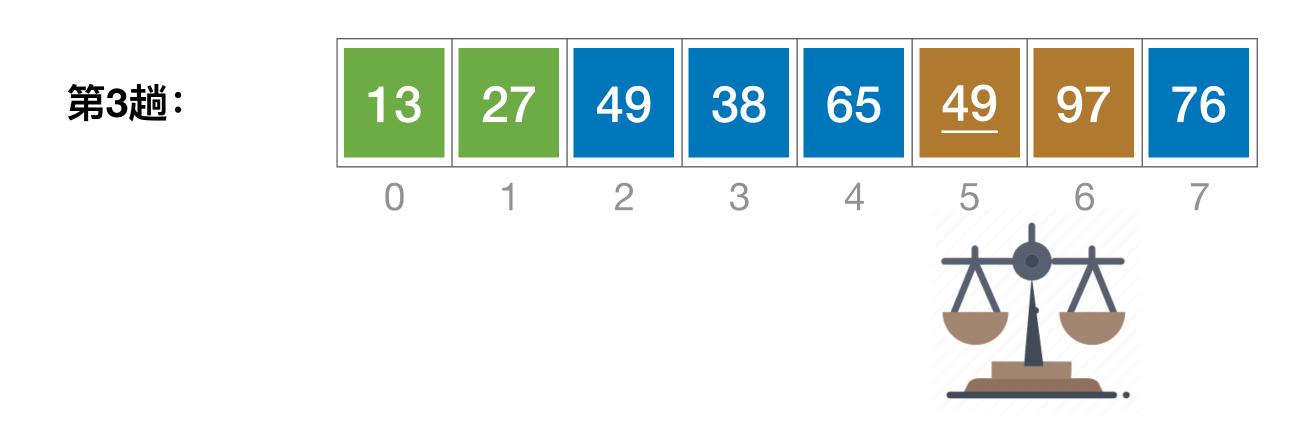
第2趟结束后,最小的两个 元素会"冒"到最前边

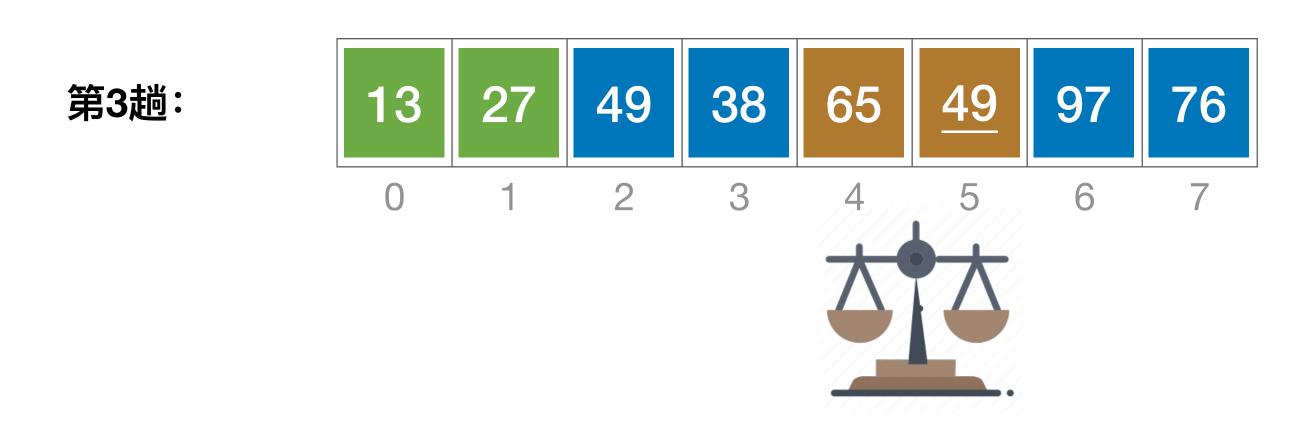
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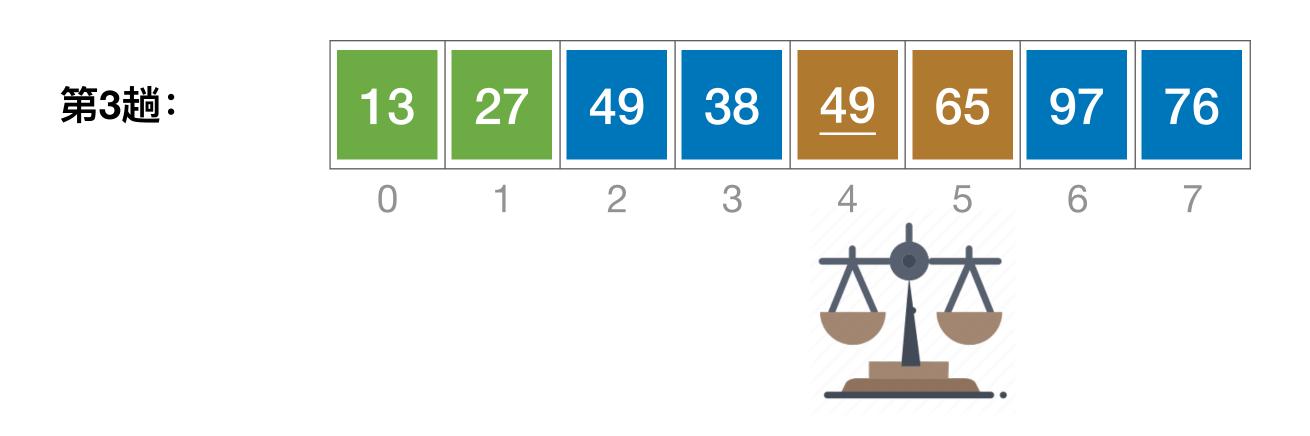


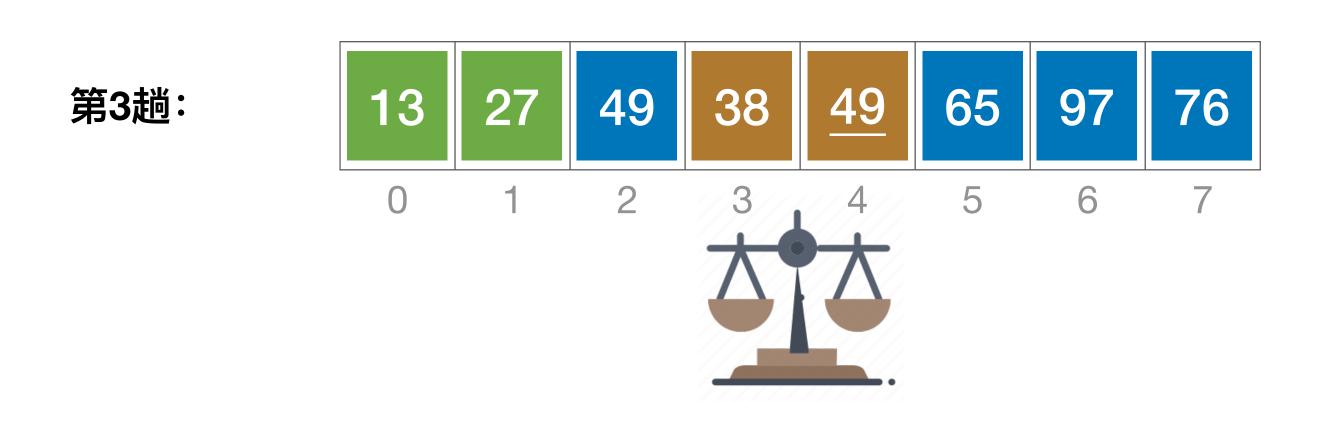


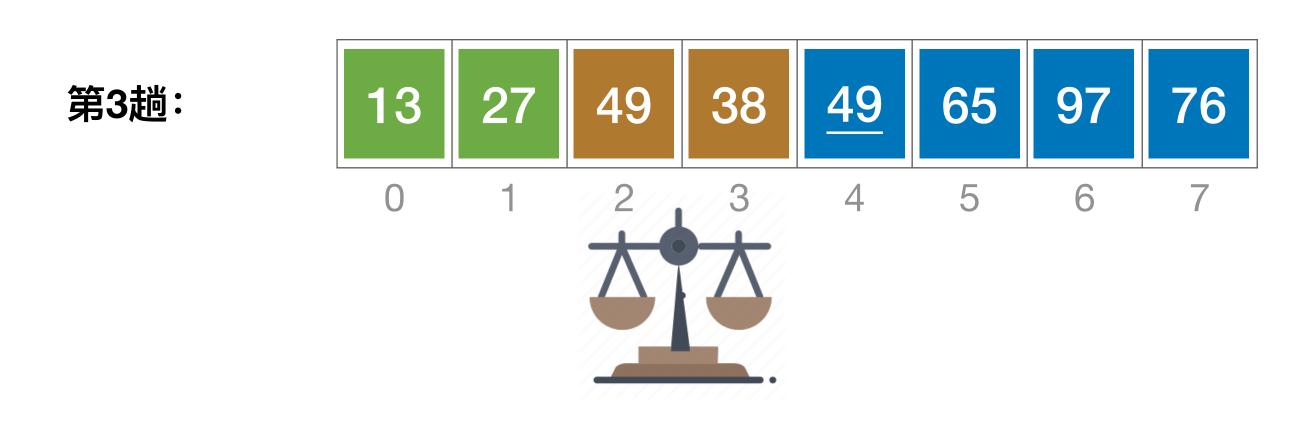










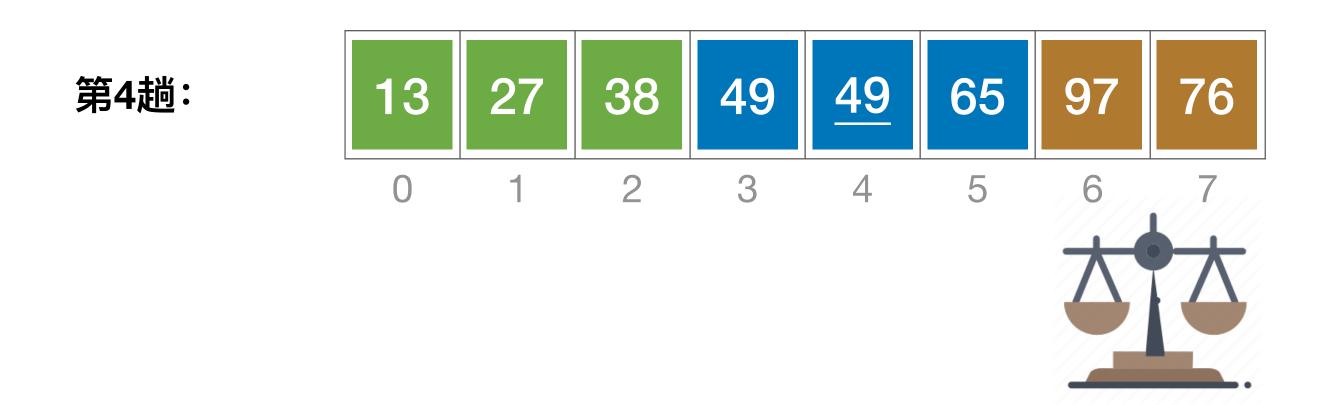


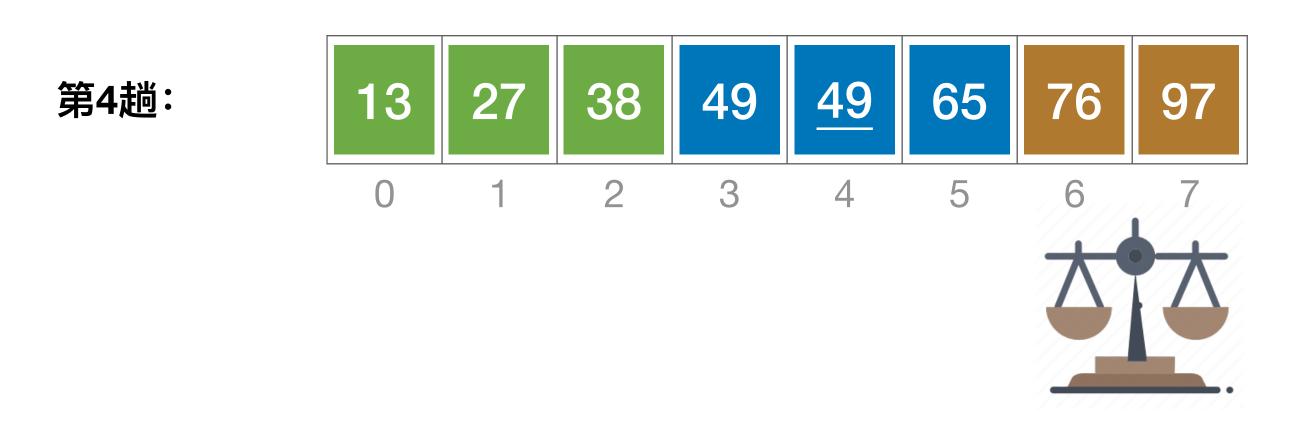


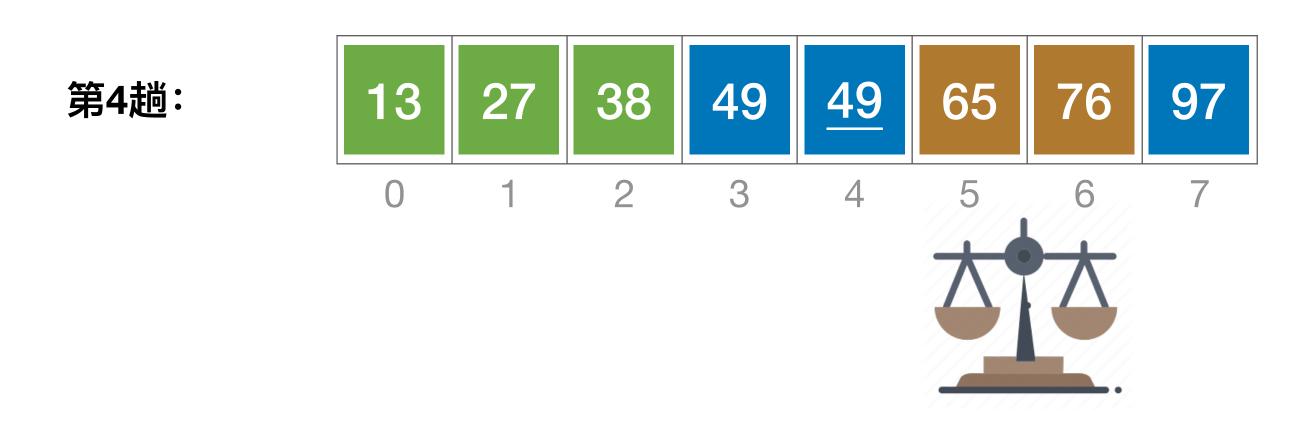


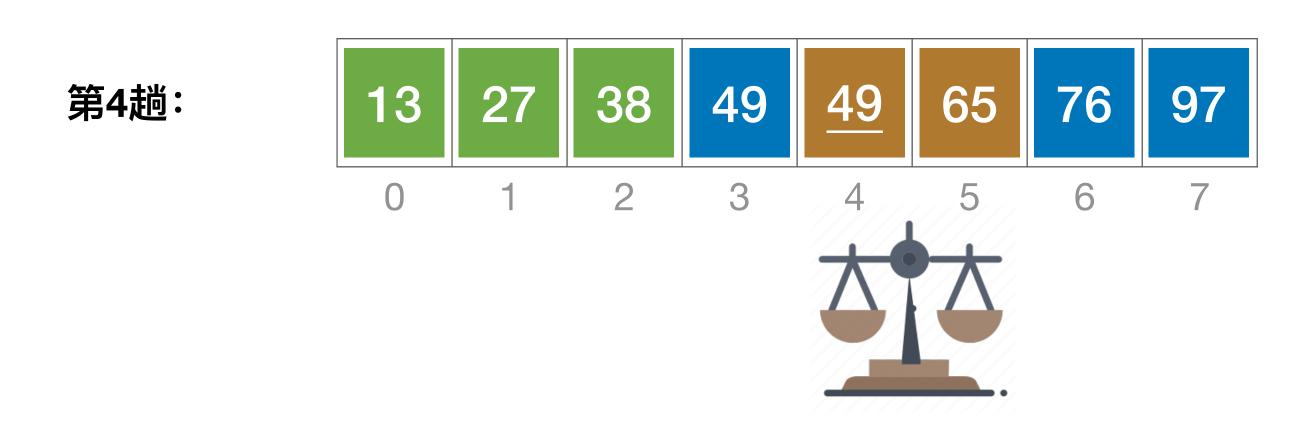
第3趟结束后,最小的3个元 素会"冒"到最前边

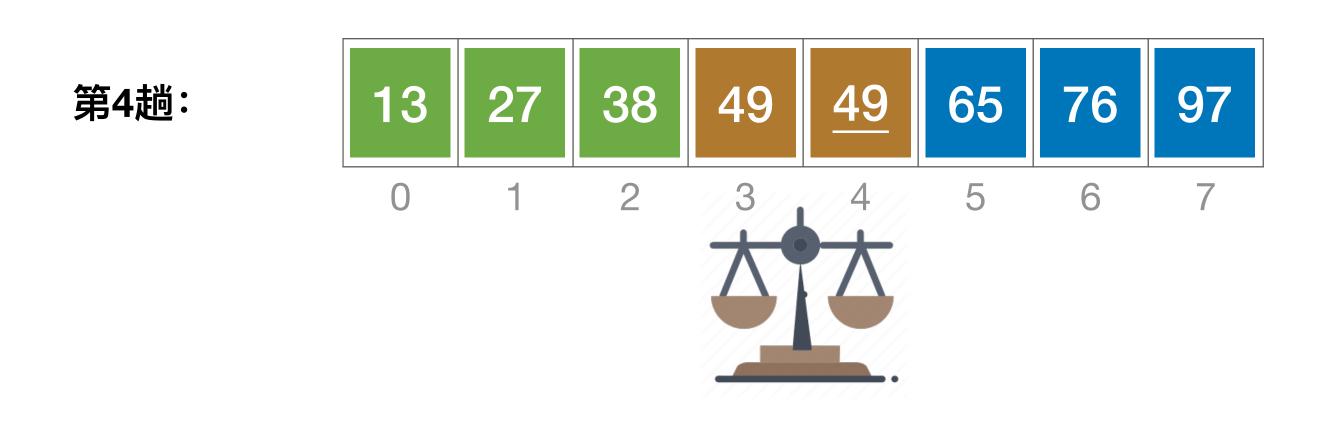
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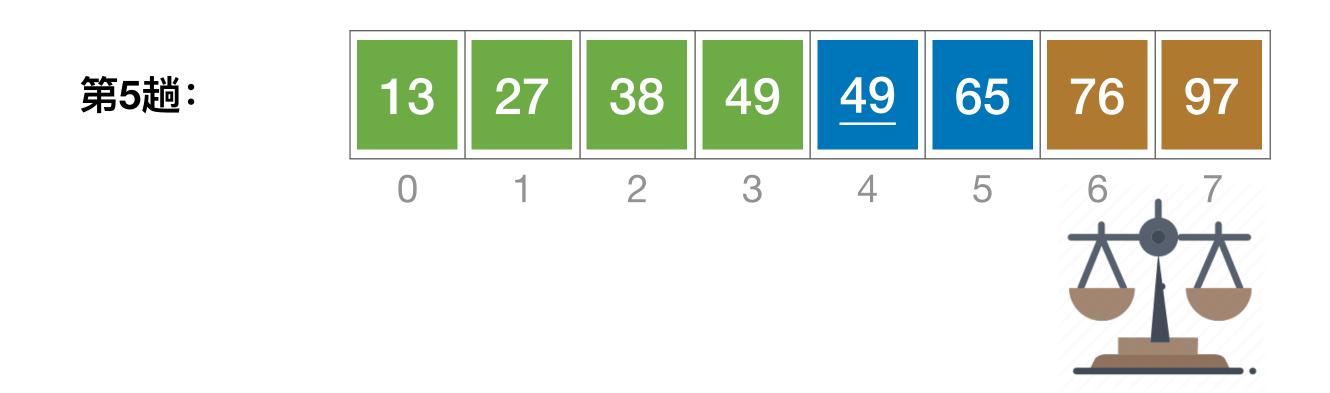


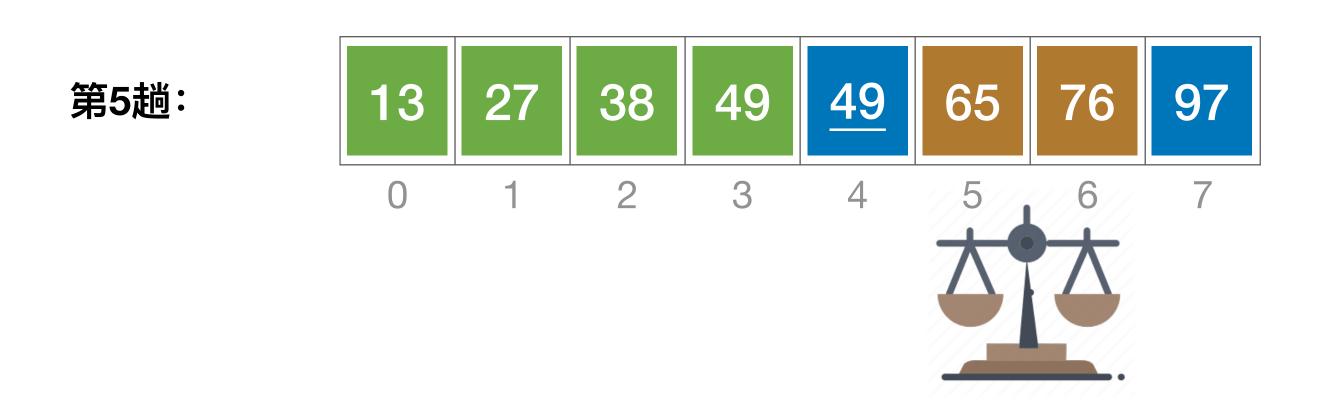






从后往前(或从前往后)两两比较相邻元素的值,若为逆序(即A[i-1]>A[i]),则交换它们,直到序列比较完。称这样过程为"一趟"冒泡排序。

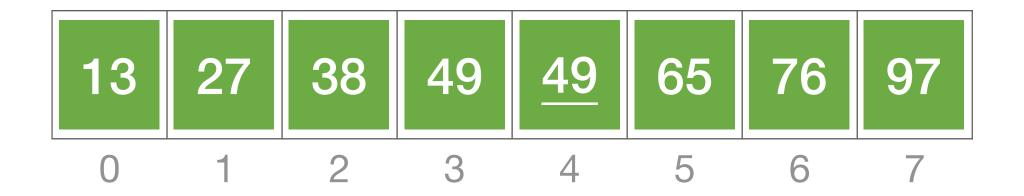




从后往前(或从前往后)两两比较相邻元素的值,若为逆序(即A[i-1]>A[i]),则交换它们,直到序列比较完。称这样过程为"一趟"冒泡排序。总共需进行 n-1 趟冒泡。







算法实现

```
//冒泡排序
//交换
                           void BubbleSort(int A[],int n){
void swap(int &a, int &b){
                              for(int i=0;i<n-1;i++){</pre>
   int temp = a;
                                  bool flag=false;
                                                //表示本趟冒泡是否发生交换的标志
   a = b;
                                  for(int j=n-1; j>i; j--) //一趟冒泡过程
   b = temp;
                                    - if(A[j-1]>A[j]){ //若为逆序
}
              只有A[j-1]>A[j]时才交换,
                                        swap(A[j-1],A[j]); //交换
                  因此算法是稳定的
                                        flag=true;
                                  if(flag==false)
                                     return;
                                                  //本趟遍历后没有发生交换,说明表已经有序
                 老哥, 稳。
                                         97
                                 38
                                     65
          i所指位置之前的元素
              都已"有序"
```

算法性能分析

空间复杂度: O(1)

最好情况(有序): 1 2 3 4 5 6 7 8

比较次数=n-1;交换次数=0 最好时间复杂度=O(n)

最坏情况(逆序): 8 7 6 5 4 3 2 1

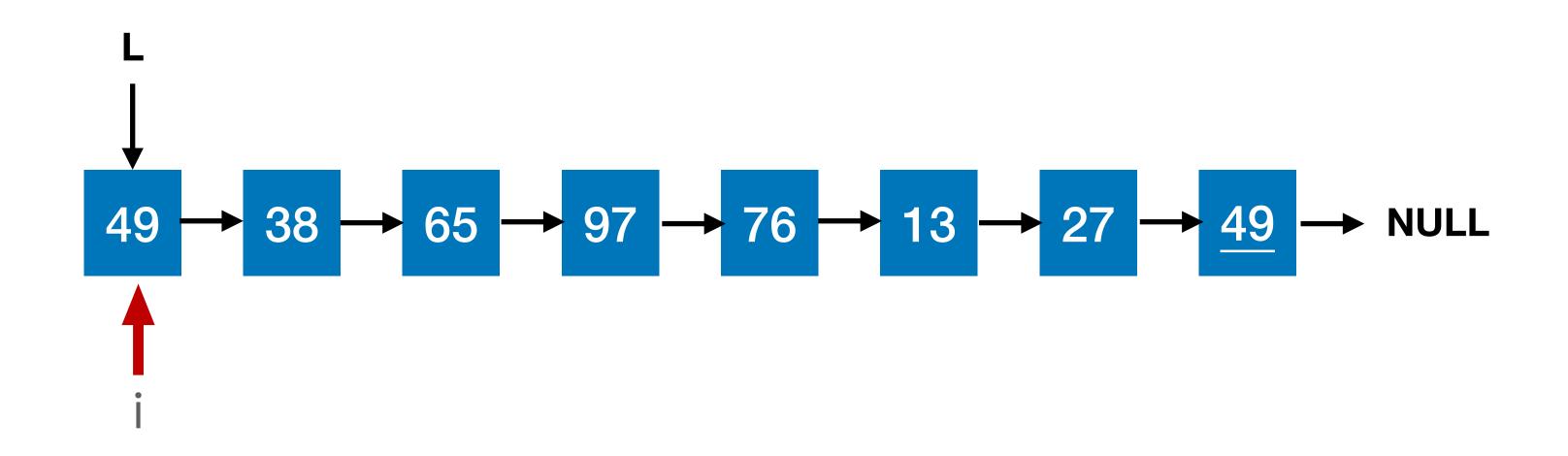
比较次数= $(n-1)+(n-2)+...+1 = \frac{n(n-1)}{2} = 交换次数$ 最坏时间复杂度= $O(n^2)$

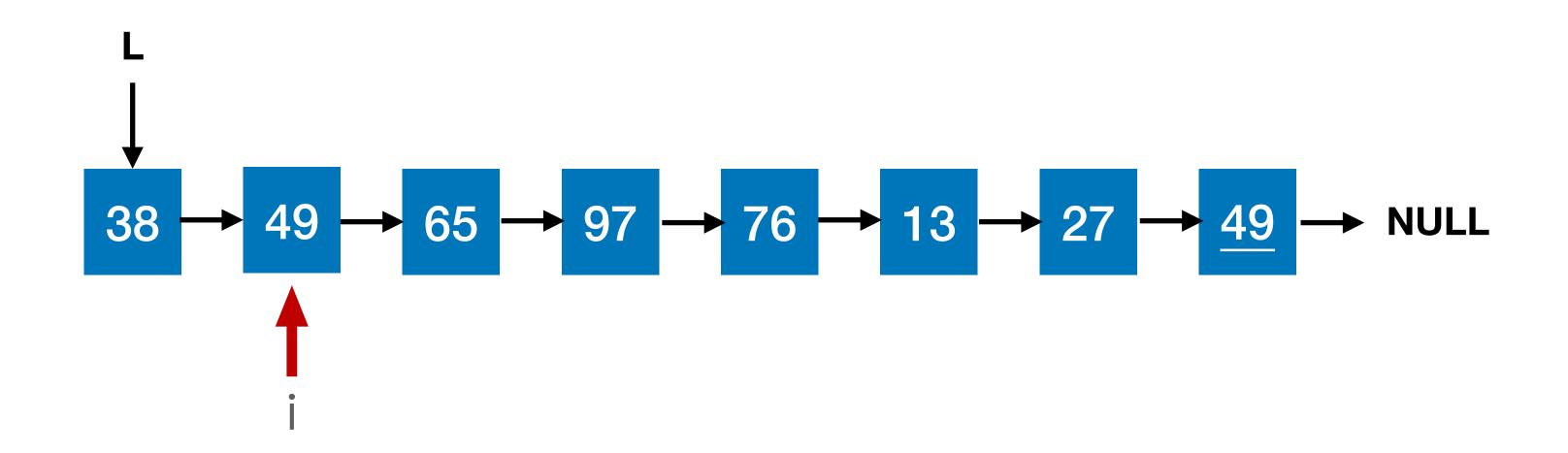
平均时间复杂度=O(n2)

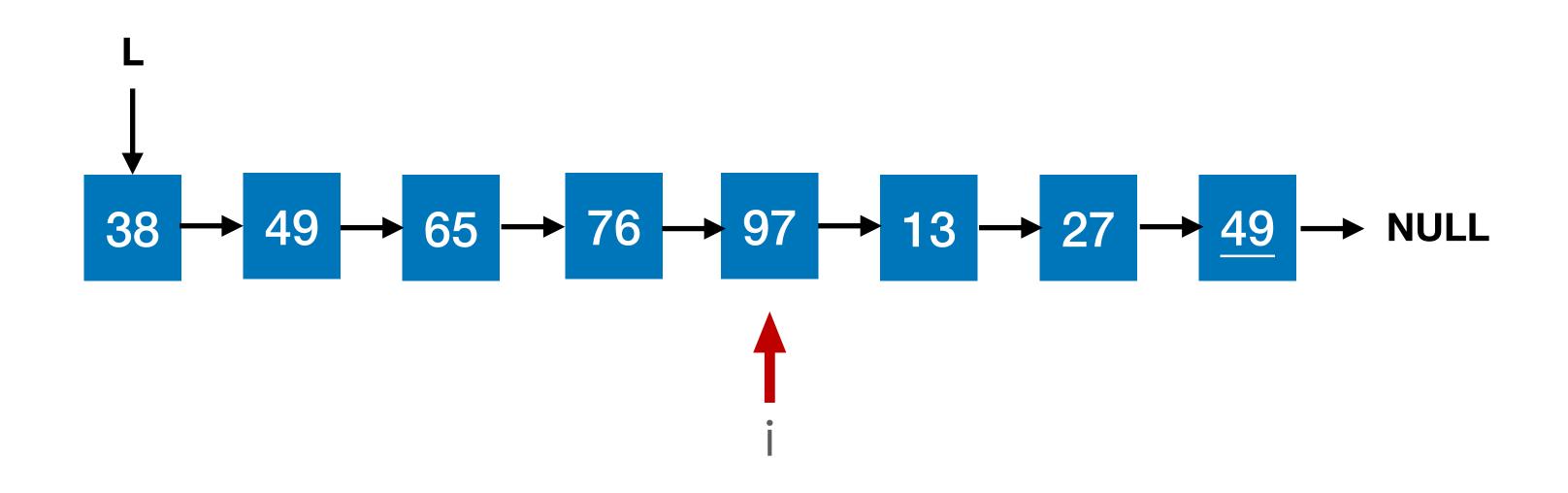
每次交换都需要 移动元素3次

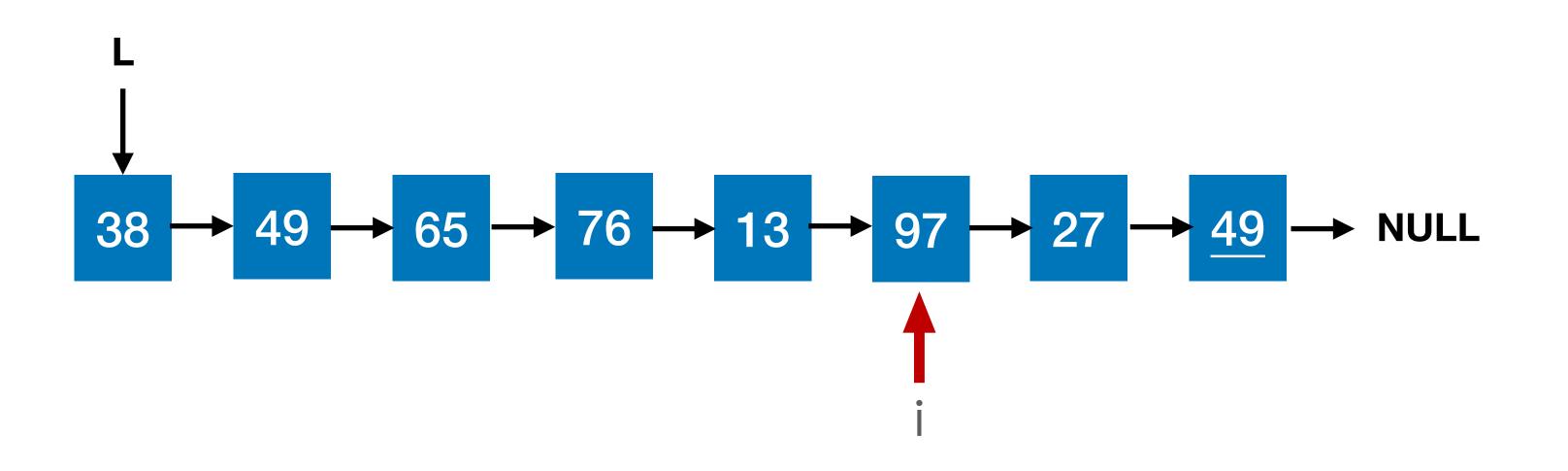
//交换

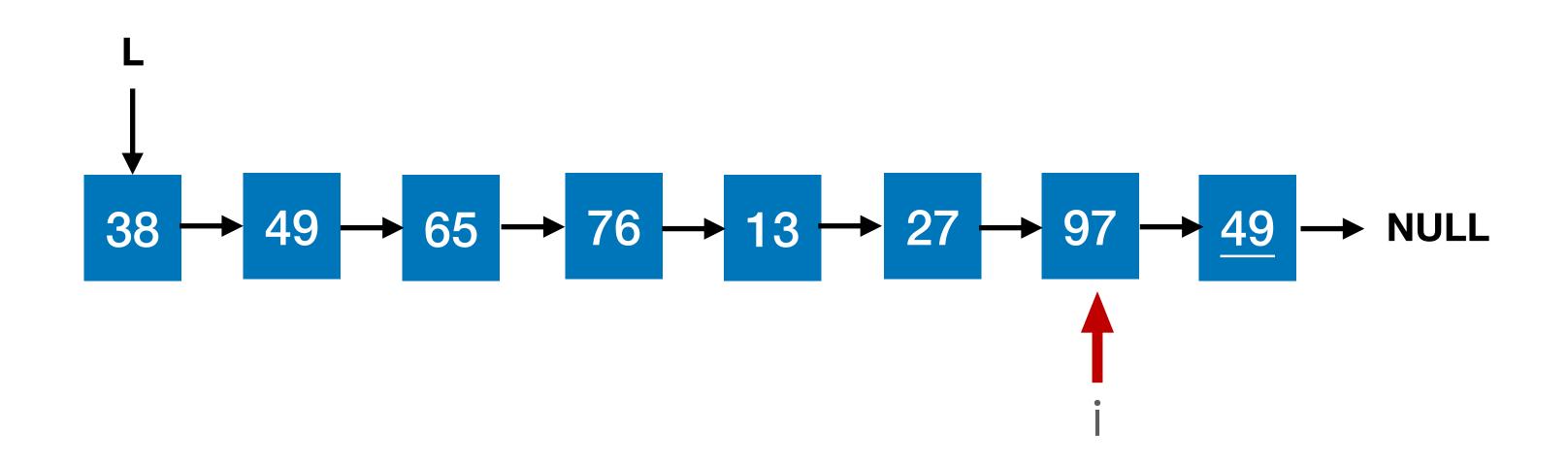
```
void swap(int &a, int &b){
   int temp = a;
   a = b;
   b = temp;
}
```

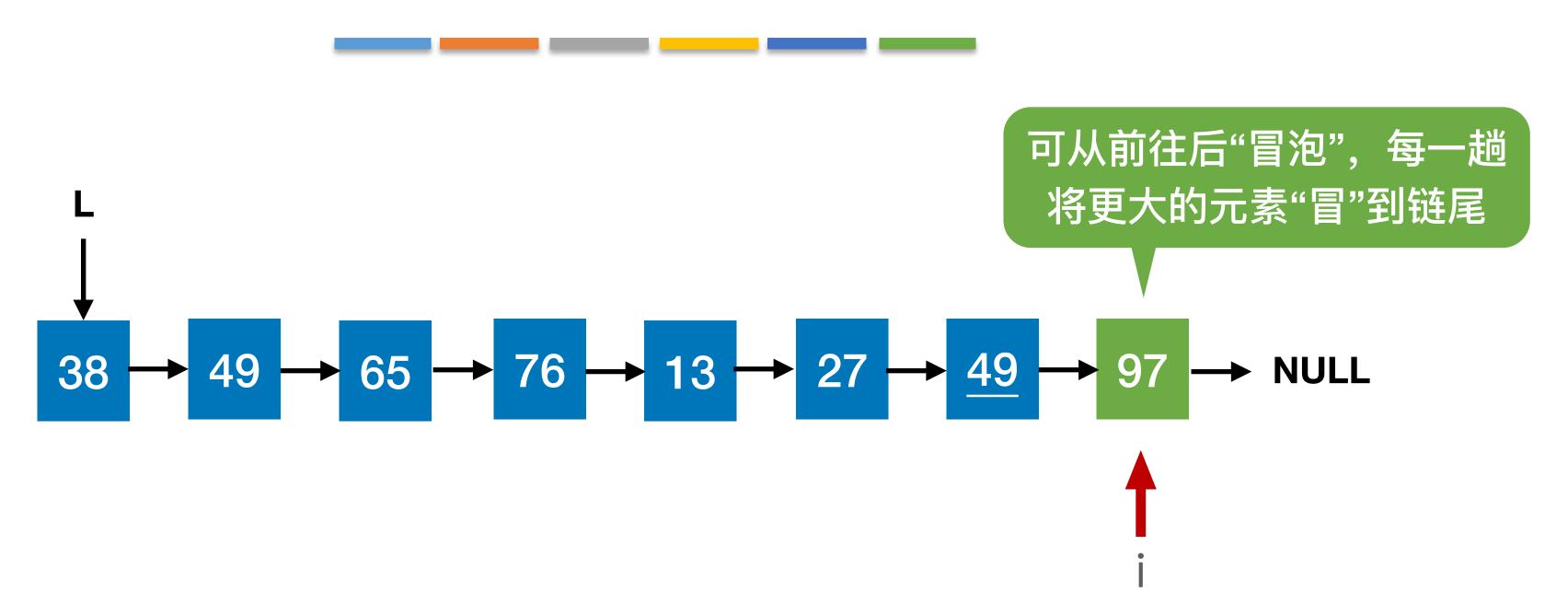




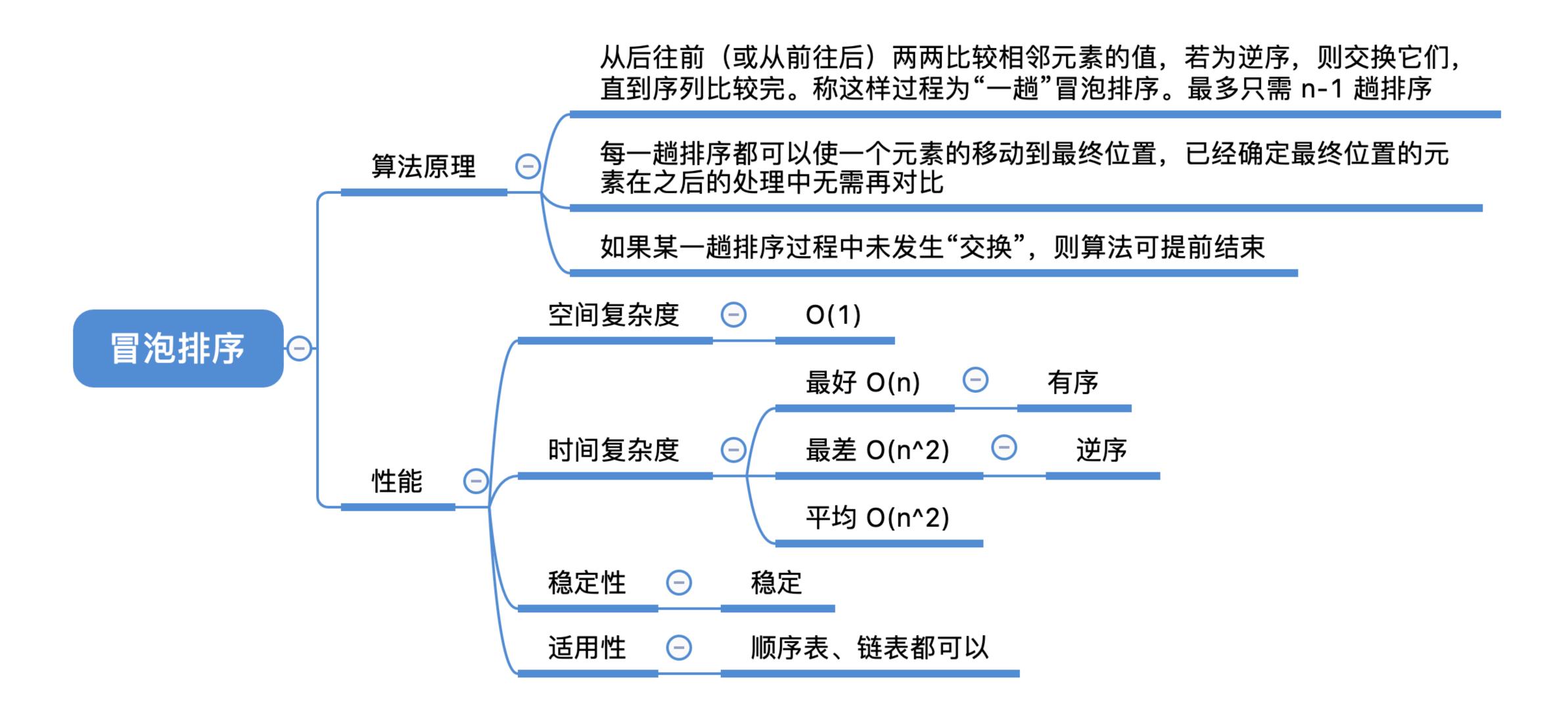








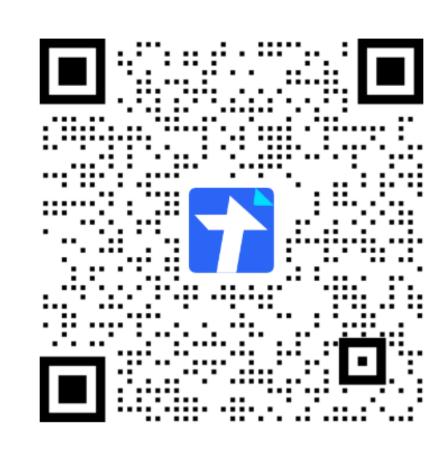
知识回顾与重要考点



欢迎大家对本节视频进行评价~



学员评分: 8.3.1 冒泡排序



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