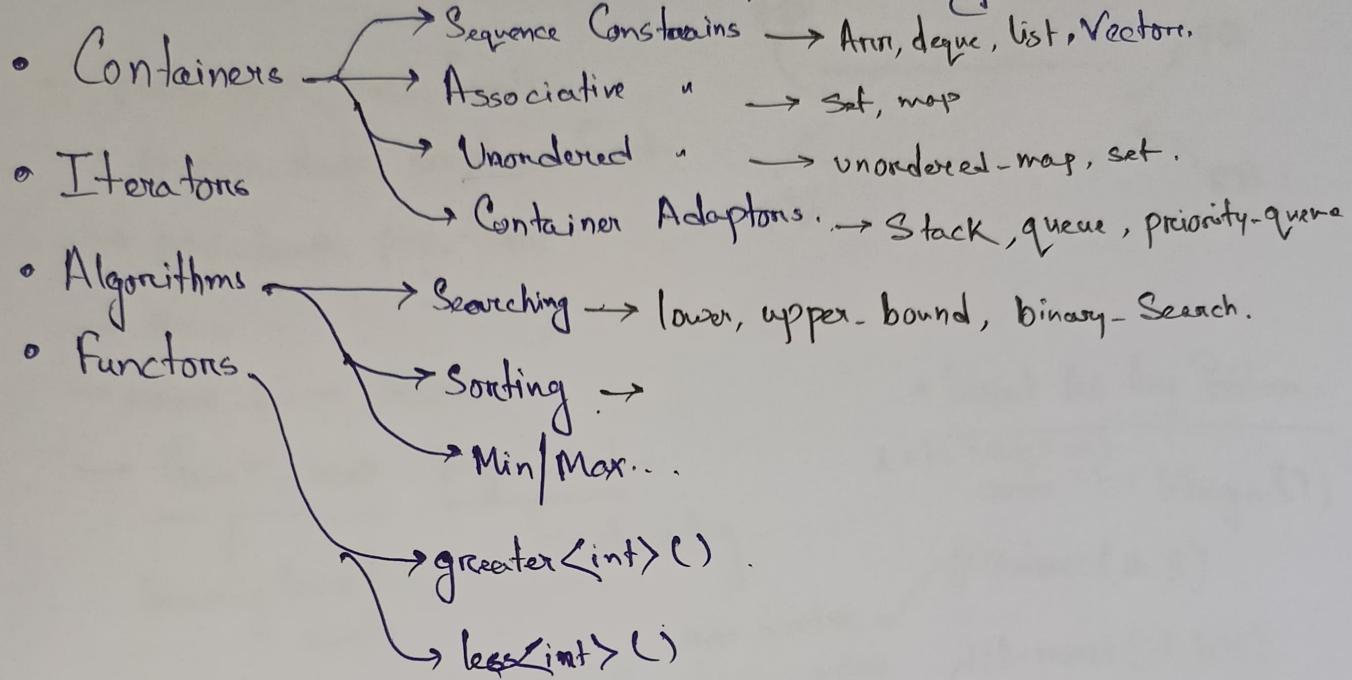


C++ STL

Standard Template Library.



Vector := (Dynamic / Resize)

`Vector<int> Vec;` //empty Vector

`Vector<int> Vec1 = {1,2};` // initialize.

`Vector<int> Vec2(3,10);` // size=3, value=10;

`Vector<int> Vec3(Vec2);` // copy Constructor.

[1 | 2 | 3 | 4 | 5]
↑
it

`Cout << * (it)`
↑
print Value

`Cout << it`
↑
print address
not Value

{ 1, 2, 3, 4, 5 }

→ Size & Capacity → if no space left it double's itself.

[1]

[1 | 2]

[1 | 2 | 3 | 1]

[1 | 2 | 3 | 4 | 1]

[1 | 2 | 3 | 4 | 5 | 1 | 1]

→ push_back / pop_back.

→ emplace_back.

→ at() and [i] → index Value.

{ 1, 2, 3, 4, 5 }

→ front() / back() .

→ erase, insert, clear, empty.

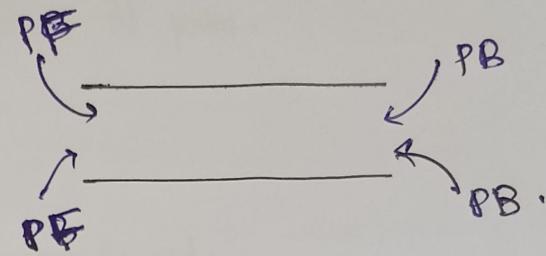
`Vec.insert (vec.begin() + 1, 100);`
`Vec = {1, 100, 2, 3, 4, 5}`

Allow insertion & deletion from both side -
 List := (Doubly linked list) → No Index Access ($\underline{\underline{\text{no } \&[0]}}$)

`list<int> l = {1, 2, 3}`

- push-back, push-front.
- emplace-back, emplace-front.
- pop-back, pop-front.
- size
- erase, clear, insert.

front, back, begin, end.
↓ Accessing front & back → iterating



• Insert At Any Position
 $\underline{l = \{1, 2, 3, 4, 5, 6\}}$
 $\text{auto it} = l.begin();$

$// \text{advance(it, 3)}$
 $l.l.insert(it, 100)$
 $\underline{l = \{1, 2, 3, 100, 4, 5, 6\}}$
 insert
 100 before 4

Deque := (Double Ended Queue).

→ Allow insertion & deletion from both side.

→ Index Access. (unlike List).

`deque<int> d = {1, 2, 3}`

- push-front, push-back | emplace
- pop-back, pop-front.
- front, back.
- erase, insert, begin, end.
- size, empty, clear.
- [] and at().

→ Same As List

Pair

→ Storing as Key-Value.

→ When using vector / list / deque of pairs.

pair <int, int> p1 = {3, 5};

pair <char, int> p2 = {'A', 1},
 ↑ → p2.Second,
 p2.First

vector<pair<int, int>> vec = {{1, 2}, {2, 3}, {3, 4}};

→ push-back() → vec.push-back({4, 5});

→ emplace-back() → vec.emplace-back(4, 5)

NON-SEQUENCE CONSTRAINTS

Stack := (LIFO)

→ Reversing things (String, array).

→ Valid Parentheses Problem.

→ DFS

→ Stack is not Iterable
(No begin(), end()).

→ push() → adds elem to TOP

→ emplace() → in-place.

→ top() → return the elem at TOP.

→ pop() → removes the TOP elem

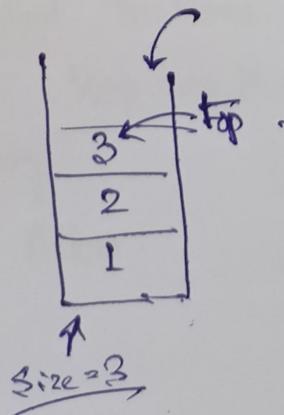
→ size()
→ empty() → swap()

Stack<int> s;

s.push(1) // 1

s.push(2) // 2 1

s.push(3) // 3 2 1



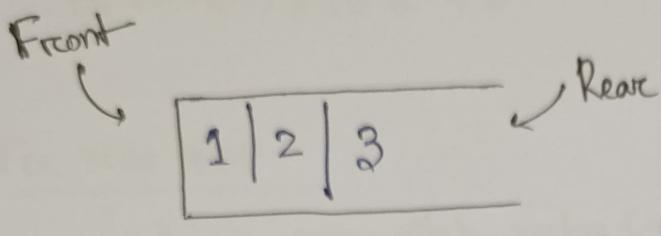
Queue := (FIFO).

→ BFS (Lvl order) in Trees/Graphs

→ Task Scheduling, CPU Scheduling.

→ Producer - Consumer Problems.

→ No Iteration (No begin(), end())



- push() → adds elem at the BACK
- emplace() → elems in-place at BACK
- front() → elems at the FRONT (first inserted).
- pop() → removes FRONT elem.
- size()
- empty()
- swap().

```
❸ queue<int> q;
q.push(1); // 1
q.push(2); // 1 2
q.push(3); // 1 2 3
```

q.front() // 1

q.pop() // 2 3

Priority Queue := (MAX HEAP By Default).

→ stores elems such that the largest elems always stays at the top.

→ When u always need the largest/smallest elems. quickly.

→ Dijkstra's Algo. (shortest Path).

→ K largest /smallest Problems.

→ Streaming data (keep top K elem)

→ Task Scheduling

→ Duplicate Allow.

Priority-queue <int> pq; // MAX HEAP
Default

pq.push(5); // 5

pq.push(3); // 5 3

pq.push(10); // 10 5 3

pq.push(4); // 10 5 4 3

→ push() → insert elems.

→ emplace() → faster insertion

→ top() → returns current maxi

→ pop() → removes maximum

→ size()

→ empty()

For MIN HEAP,

Priority-queue <int, vector<int>, greater<int>> minPq; // MIN HEAP

↳ makes smallest elems highest priority.

if we use
Pair

Priority-queue <pair<int,int>> pairPq;

{ pairPq.push({2,100});

{ pairPq.emplace(1,100); // in-place.

here max-heap

Based on first elem

→ output: (2,200), (1,100)

Map :

- map $\langle K, V \rangle$ stores key-value pairs in Sorted Order (ascending)
- Keys are Unique.
- Balanced BST.

map $\langle \text{String, int} \rangle m$; → Key wise sorting.

// Operator [] inserts OR updates a key.

$m["tv"] = 100;$
 $m["laptop"] = 100;$
 $m["headphones"] = 50;$
 $m["tablet"] = 120;$
 $m["watch"] = 50;$

Output:- (Keys Sorted Alphabetically)

camera → 25
headphones → 50
laptop → 100
tablet → 120
watch → 50.

- $m.\text{insert}(\{\text{"camera"}, 25\});$
- $m.\text{emplace}(\text{"mic"}, 25);$
- $m.\text{size}()$ → $m.\text{erase}(\text{"watch"});$ // remove by key.
- $m.\text{clear}()$ → $m.\text{find}(\text{"tv"});$
- $m.\text{empty}()$ → $m.\text{count}(\text{"camera"});$

↳ if found return 1.

otherwise 0.

→ $m["tv"]$

↳ returns the

Value = 100.

Unordered Map :=

- unordered_map <K, V>
- no specific order (random order)
- Keys must be unique.
- Fastest avg time $O(1)$

unordered_map <String, int> um;

Same As Map

But Random Order

Multi Map := (Sorted but allows duplicate keys).

- multimap <K, V>
- Stores key in sorted order.
- Allow multiple values for the same key.
- insert : $O(\log n)$.
- find() returns iterator to the first matching key.

multimap <String, int> mm;

~~mm.push("tv", 100);~~
~~mm["tv"] = 200;~~
~~mm["tv"] = 300;~~

mm.insert(mm.find("tv")); Increasing only the first occurrence of "tv".

↳ Output

tv → 200

tv → 300

Set := (Balanced BST)

→ Stores ~~gives~~ Unique Values in Sorted Order (ascending).

- insert, emplace
- count
- erase
- find
- size, empty, erase

→ you need fast search (insert / delete)
→ No Indexing $O(\log n)$.

Set<int> s;

s.insert(1);

s.insert(2);

s.insert(1); // duplicates ignored.

s.insert(3);

s.insert(4);

Output :- 1 2 3 4

→ s.size() // 4 (Coz duplicate ignored).

→ s.count(3); // checks if exists 1 else 0.

Multi Set :=

multiset<int> s;

→ Allows duplicate Value.

→ print all duplicate Values also.

→ Also give total size including duplicates.

s = 1 2 2 3 3 3 4 5

s.size() // 8

Unordered Set :=

→ No upper & lower-bound Concept

→ random order like unordered-map.

→ Duplicates ignored.

Lower Bound :- Set stores unique & sorted values(BST).

- Returns iterator to the FIRST elem $> x$.
- works in $O(\log n)$
- if no elem exists → returns s.end().

Set = 1, 2, 3, 5, 6, 7.

s.lower_bound(4); // not found so first elem greater than 4 is 5, so output is 5. If 4 was present then Output 4.

s.lower_bound(10); // == s.end().

Upper Bound :-

a a [b b b] c c d
0 1 2 3 4 5 6 7

upper_bound('b') → greater than Key. Means ans is 5 (c).
it can't be 'b'.

Set = 1, 2, 3, 5, 6

* (s.upper_bound(4)) == 5. → first elem $> x$.

Algorithms

Sorting :-
~~int arr[] = {1, 8, 5, 3};
 int n = 4;~~

→ Sort(arr, arr + n) // Sort in ascending Order
 || 1, 3, 5, 8

→ Sort(arr, arr + n, greater<int>()) // Sort in descending order.
 using greater<int>().
 || 8, 5, 3, 1.

Vector<pair<int, int>> rec = {{3, 1}, {2, 1}, {7, 1}, {5, 2}};

→ Sort(v.begin(), v.end()); || 2 1 (Sorted by first)
 3 1
 5 2
 7 1

Custom Operator For PAIR

Sort by Second elem
 // [] it's a lambda (like mini function without a name)
 auto comparator = [] (const pair<int, int> &p1, const pair<int, int> &p2)

{
 return p1.second < p2.second; // Sort by Second Value.

→ Sort(rec.begin(), rec.end(), comparator);

Output :-
 3 1
 2 1
 7 1
 5 2

Other Algos:-

• reverse() → When u need to flip array / Vector quickly.

• next_permutation() → Solving permutation-based Prob (Strings, numbers).

• max_element() | min_element() → To Find global max/min.

• binary_search() → binary search

• __builtin_popcount() → Bit manipulation, DP, set-bit problems.