

## Data Systems Group Assignment 3 – Due Monday 5<sup>th</sup> October 11 PM

1. Consider the following information

Relation/number of rows	Column/Attribute	Length (in bytes)	Number of distinct Values/Distribution
R/ $2^{20}$ rows.	KR	32	key
	A	24	10 - uniform
	B	28	$2^{10}$ - uniform
	C	4	2 values - uniform
	KS	16	Foreign Key – (see relation S) uniform
	Other columns	192	-
S/ $2^8$ rows	KS	16	key
	A	24	10 - uniform
	B	28	$2^8$ - uniform
	C	4	2 values - uniform
	Other columns	188	-

Consider the

- a) Disk block size (B) is 512 bytes, 4096 bytes, and 32,768 bytes. Three possibilities.
- b) Block Pointer length is 8 bytes.
- c) The record pointer length is 16 bytes.

Now compute the number of block accesses needed **(for each of the above block sizes)** to execute the following **retrievals**. Provide detailed answers with all intermediate calculations, formulae used, and the rationale. Use floor and ceil appropriately for the computation of the number of block accesses.

- a) R and S are unordered files, determine the number of block accesses for the retrieval
  - i. KR=val
  - ii. KR>val
  - iii. S.A=val
  - iv. S.A>val
- b) R and S are ordered files, determine the number of block accesses for the retrieval
  - i. S.KS=val
  - ii. S.KS>val
  - iii. R.B=val
  - iv. R.B>val
- c) Size in the number of blocks for Primary Index
  - i. On R.KR and the number of block accesses for R.KR=val.
  - ii. On S.KS and the number of block accesses for S.KS =val.

- d) Size in the number of blocks for the Clustered Index
  - i. On R.B and the number of block accesses for R.B <val.
  - ii. On S.A and the number of block accesses for R.A=val
- e) Size in the number of blocks for the secondary index on key
  - i. On S.KS and the number of block accesses for S.KS=val
- f) Size in the number of blocks for the secondary index on non-key
  - i. On R.C and the number of block accesses for R.C =val
  - ii. On S.B and the number of block accesses for S.B<val
  - iii. On R.KS and the number of block accesses for R.KS =val
- g) Consider a B-Tree on R.KR with each block  $\frac{2}{3}$ <sup>rd</sup> full. Determine the number of block accesses for R.KR=val retrieval.
- h) Consider a B+ Tree on R.KS with each block  $\frac{2}{3}$ <sup>rd</sup> full. Determine the number of block accesses for R.KS=val retrieval.
- i) Consider the main memory buffer of  $(2^{10}+2)$  blocks. Compute the number of block accesses for the block nested loop equijoin between R and S on attributes R.KS=S.KS.
- j) Consider an equijoin of R and S, on attributes R.KS=S.KS, compute the number of block accesses for nested loop join using B+tree index on S.KS.
- k) Consider B+tree secondary indices on R.C and S.C are available. Which join algorithm will you use, and how many block accesses will you need to perform R and S's equijoin on attributes R.C=S.C.
- l) There is a retrieval on a single relation using an indexing method that will take the maximum number of block accesses. What is that retrieval, and how many block accesses will it take?
- m) Change the block size so that the complete primary index on S.KS will fit in a single block. Recompute the number of block accesses for the above query (l).