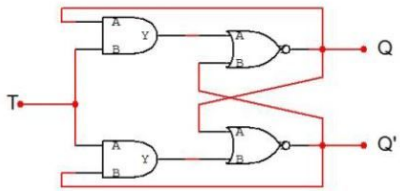
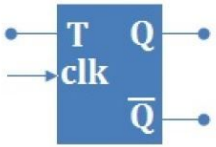
Sıralı Mantık Tasarımı

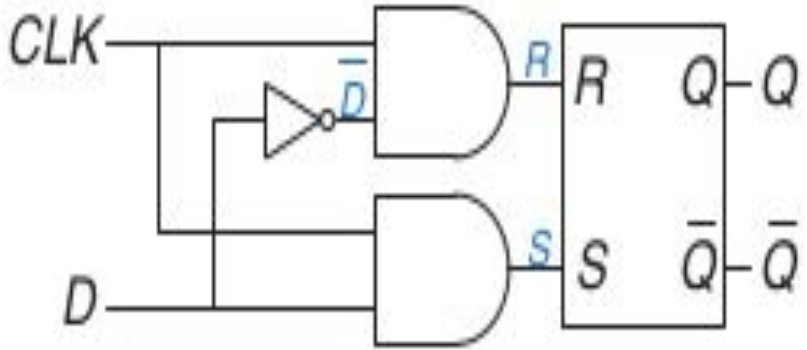




Suhap SAHIN

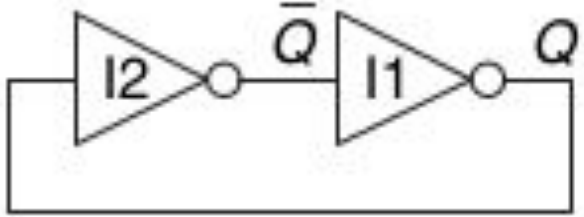
LATCHES & FLIP-FLOPS

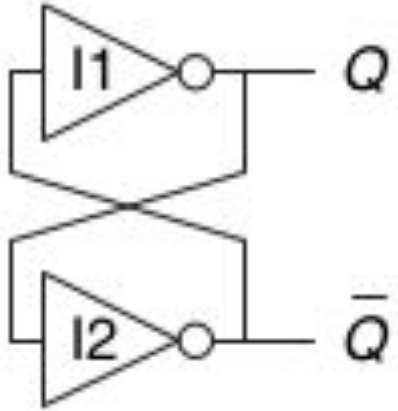
- Adısık devrelerin analizi zor oldugundan; bu bölümde çokca kullanılan senkron akdısık devrelerden bahsedilecektir.
    - Senkron ardısık devreler, birlesik lojik devrelerden ve devrenin bir önceki durumunu tutan flip-floplar olusmaktadır.
- Bu bölümde, ardısık devreleri tasarlamanın kolay bir yolu olan sonlu durum makineleri açıklanmaktadır.
- Son olarak, sıralı devrelerin hızını analiz edilecek ve paralellik kavramı hızı arttırmanın bir yolu olarak tartısılacaktır.



LATCHES & FLIP-FLOPS

En temel hafıza birimi iki durumlu çıkısa sahip elemanlardan olusur.



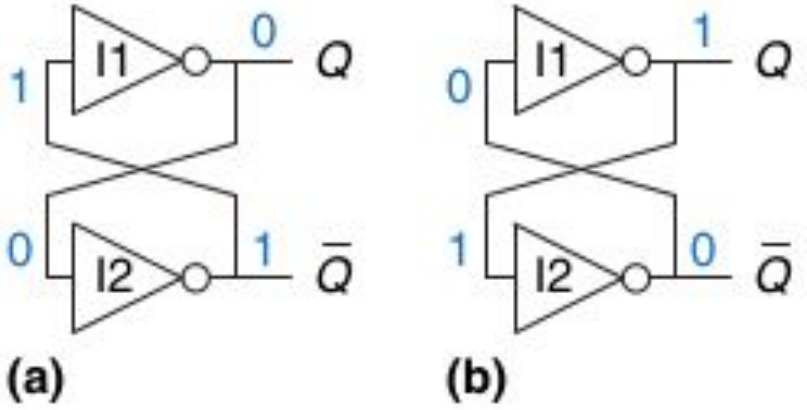


cross-coupled inverter

Q degeri, devrenin gelecekteki davranısını açıklamak için geçmis bilgileri içerir

Baslangıc durumuna ait gecmis bilgi olmadıgında, baslangıc degerleri kestirilemez.

LATCHES & FLIP-FLOPS



Case I: Q=0

I2 => TRUE (Q')

I1 => FALSE (Q)

DEVRE KARARLIDIR

Case I: Q=1

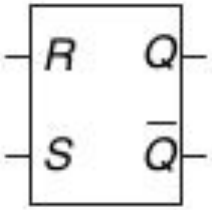
I2 => FALSE (Q')

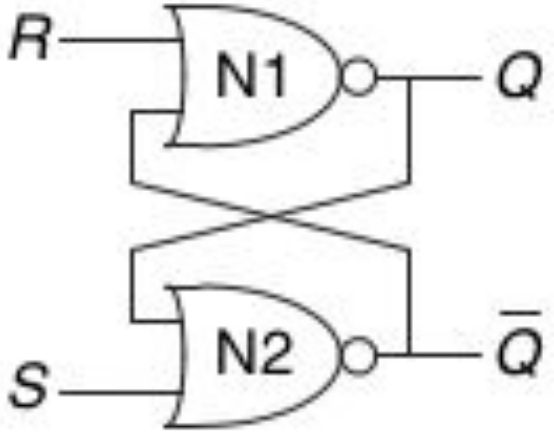
I1 => TRUE (Q)

DEVRE KARARLIDIR

Girise sahip degil

SR Latch





Case I: R=1, S=0

R=1 Q'=? --> N1 => FALSE

S=0 Q =0 --> N2 => TRUE

Case II: R=0, S=1

S=1 Q' =? --> N2 => FALSE

R=0 Q = 1 --> N1 => TRUE

Case III: R=1, S=1

R=1 Q'=? --> N1 => FALSE

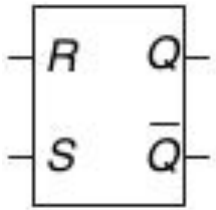
S=1 Q =? --> N2 => FALSE

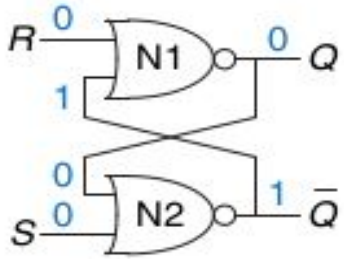
Case IV: R=0, S=0

R=0 Q'=? --> N1 => ?

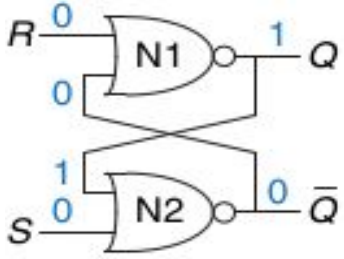
S=0 Q =? --> N2 => ?

SR Latch





(a)



(b)

Case IV: R=0, S=0

R=0 Q'=? --> N1 => ?

S=0 Q =? --> N2 => ?

Case IVa: Q=0

S=0 Q =0 --> N2 => TRUE

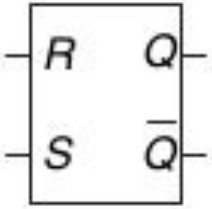
R=0 Q'=1 --> N1 => FALSE

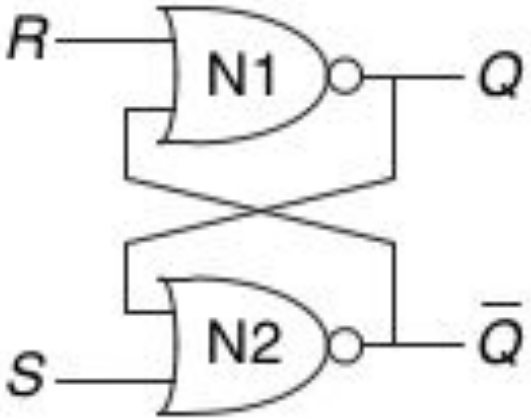
Case IVb: Q=1

S=0 Q =1 --> N2 => FALSE

R=0 Q'=0 --> N1 => TRUE

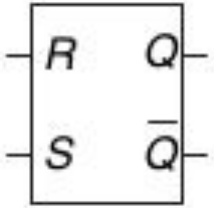
SR Latch



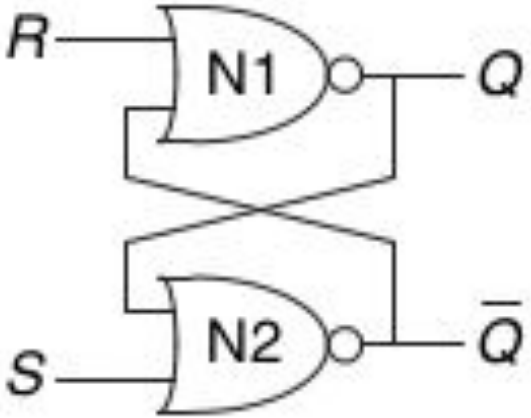


Case	S	R	Q	$\overline{Q}$
IV	0	0	Q_{prev}	$\overline{Q}_{prev}$
I	0	1	0	1
II	1	0	1	0
III	1	1	0	0

SR Latch



R=1 ve S=1 oldugu durumda devre davranısı kestirilemez

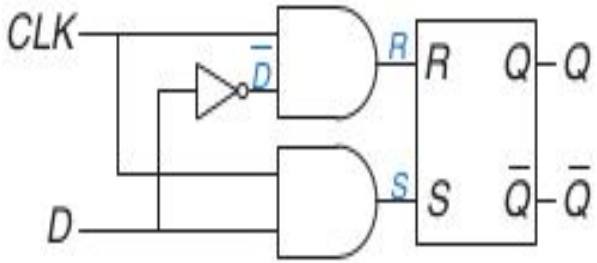


Case	S	R	Q	$\overline{Q}$
IV	0	0	Q_{prev}	$\overline{Q}_{prev}$
I	0	1	0	1
II	1	0	1	0
III	1	1	0	0

D Latch

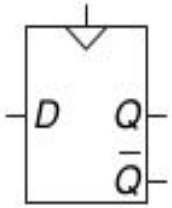
CLK = 1 ⇒ Latch is transparent

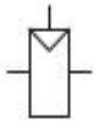
CLK = 0 ⇒ Latch is opaque



CLK	D	$\bar{D}$	S	R	Q	$\bar{Q}$
0	X	$\bar{X}$	0	0	Q_{prev}	$\bar{Q}_{prev}$
1	0	1	0	1	0	1
1	1	0	1	0	1	0

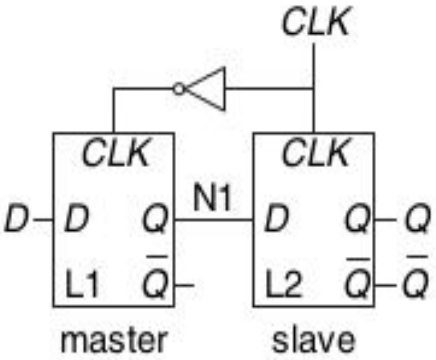
D Flip-Flop

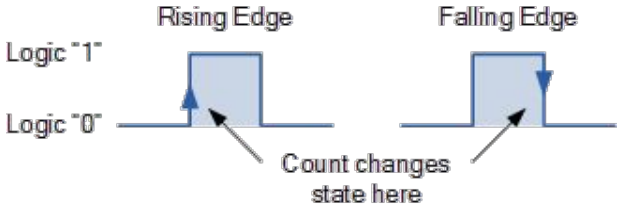




CLK = 0 ⇒ Master = transparent Slave = opaque

CLK = 1 ⇒ Master = opaque Slave = transparent





D tipi flip flop;

Yükselen kenarda D verisini çıkısa (Q) aktarır.

Aksi taktirde durumunu korur(hatırlar).

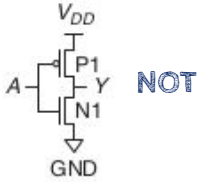
D tipi flip flop;

master-slave flip-flop

edge-triggered flip-flop

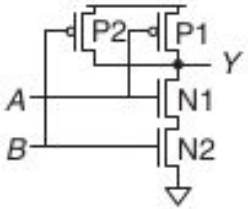
positive edge-triggered flip-flop

Örnek: D Flip-Flop Transistör sayısı

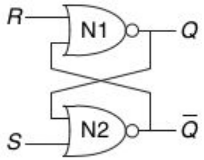


NOT

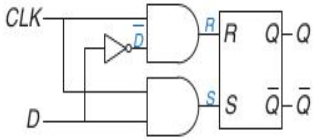
NAND



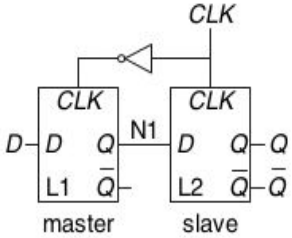
SR LATCH



D LATCH



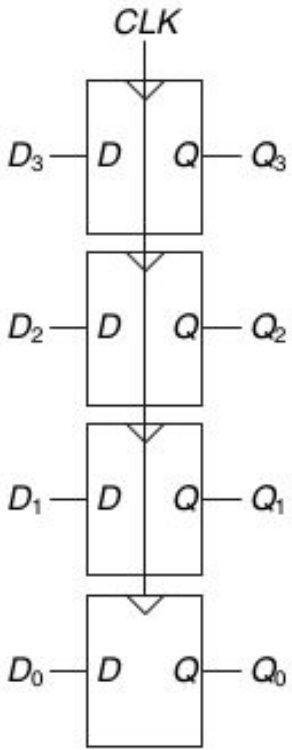
D Flip-Flop

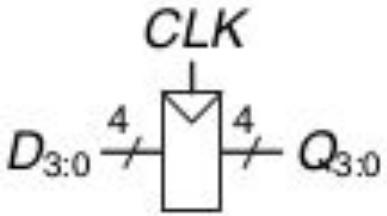


D tipi flip flop'taki transistör sayısı kaçtır?

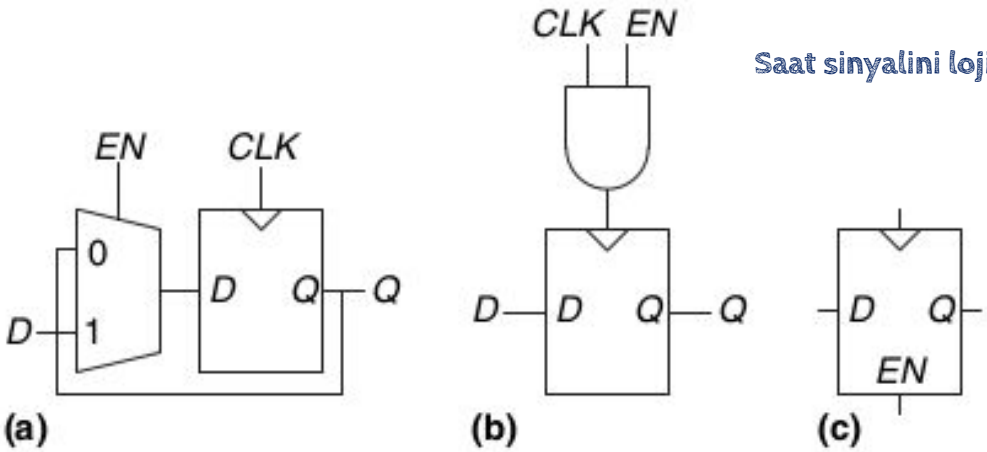
NAND / NOR	= 4 transistör		
NOT	= 2 transistör		
AND	= NAND + NOT	= 6 transistör	
SR LATCH	= 2xNOR	= 8 transistör	
D LATCH	= SR +2xAND + NOT	= 8+2x6+2	= 22 transistör
D flip-flop	= 2xD LATCH + NOT	= 2X22 + 2	= 46 transistör

Register / Saklayıcı





Enabled Flip-Flop



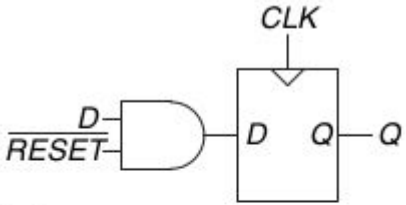
Saat sinyalini lojik bir islemden gecirmek kötü bir fikir

Verinin ne zaman kabul edilecegini belirleyen bir EN/ENABLE girisine sahiptir.

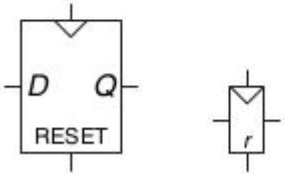
EN = TRUE ⇒ enable flip-flop = normal flip-flop

EN = FALSE ⇒ eski durum korunur

Resetlenebilir Flip-Flop



(a)



(b)

(c)

RESET = FALSE ⇒ enable flip-flop = normal flip-flop

RESET = TRUE ⇒ D = 0

Sistemde ilk calısmaya basladıgında;

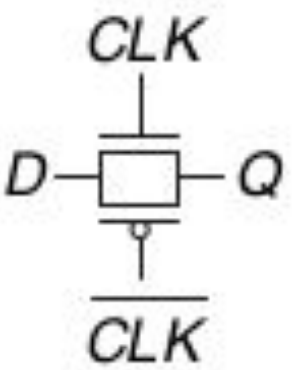
bütün flip-flop ları 0(baslangıc) durumuna getirmek için kullanılır.

Senkron resetlenebilir flip-flop: RESET girisi CLK sinyaline bagli

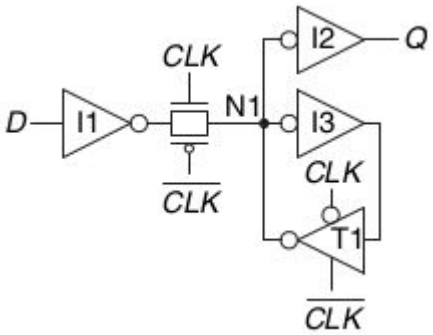
Asenkron resetlenebilir flip-flop: RESET girisi CLK sinyallinden bagımsız

Transistör Seviyesinde Latch ve Flip-Flop Tasarımı

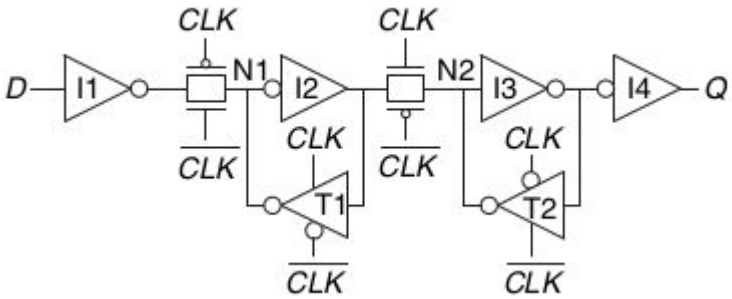
D latch



D latch

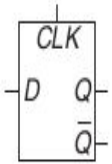


D flip flop



D Latch VS D Flip-Flop

D latch

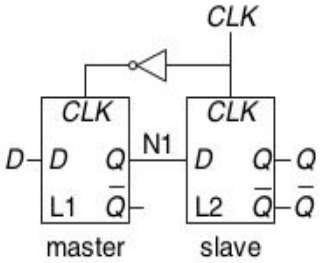


Level Triggering

CLK = 1 ⇒ Latch is transparent — D girisini Q çıkısına aktarır

CLK = 0 ⇒ Latch is opaque — Q önceki durumunu korur

D flip flop

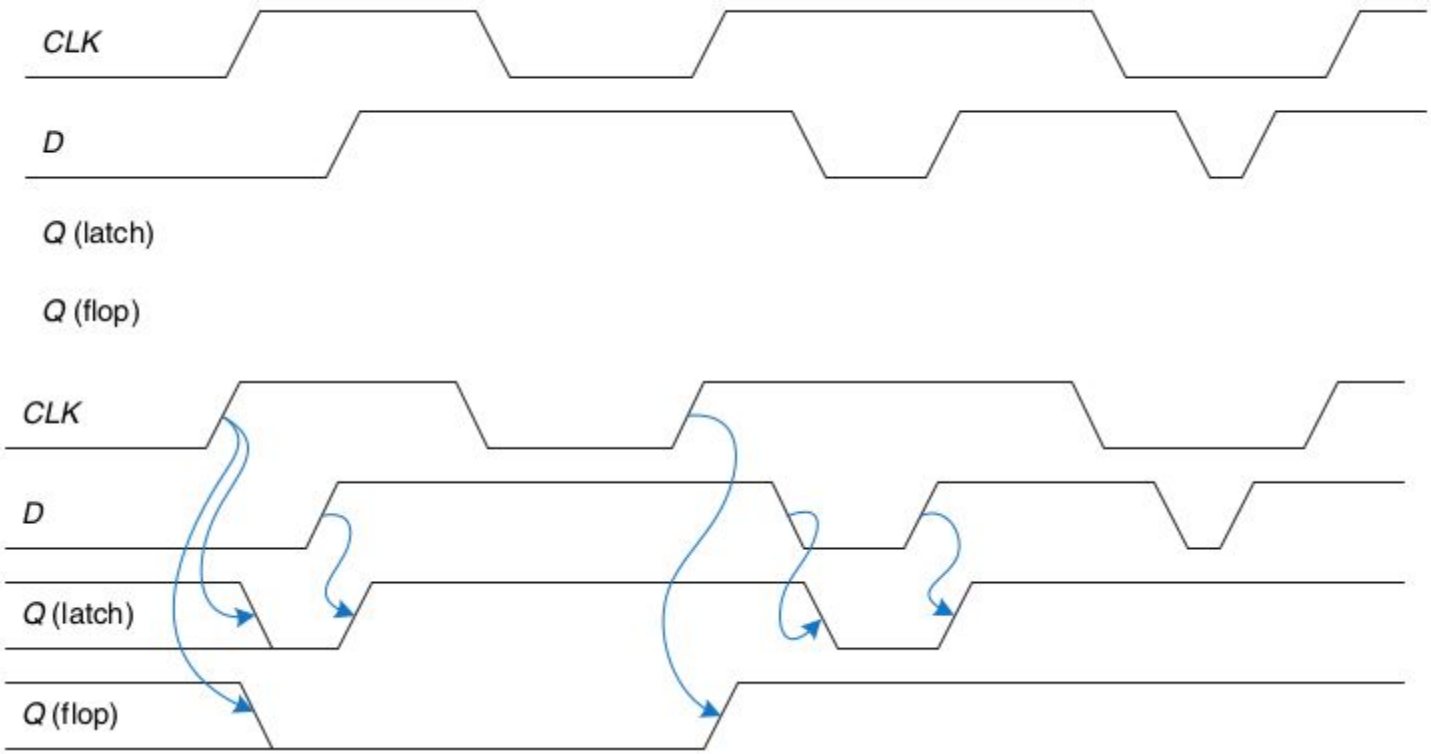


Positive Edge Triggering

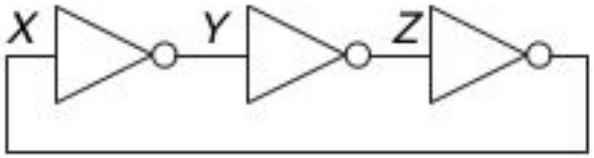
CLK = 0 ⇒ Master = transparent, Slave = opaque — D girisini Q çıkısına kopyalar

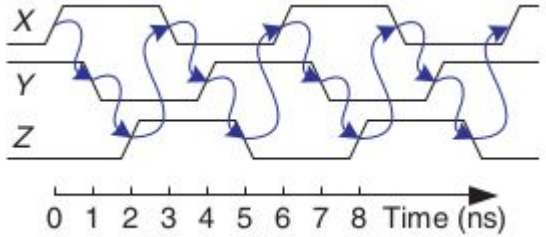
CLK = 1 ⇒ Master = opaque, Slave = transparent — Q önceki durumunu korur

D Latch VS D Flip-Flop



Senkron Lojik Tasarım - Halka Ösilatörü -





Inverter lerin yayılma gecikmesi 1ns

X=0 kabul edilirse;
Y=1 , Z=0 , X=1 (Kabul ile tutarsız olur)

0ns => X yükselir 1ns = > Y düser 2ns => Z yükselir ,
3ns => X düser 4ns => Y yükselir 5ns => Z düser
6ns => X yükselir

Yukarıdaki desen her 6 ns' de bir tekrar eder.
Bu devreye halka osilatörü denir.

Halka osilatörünün süresi = her invertörün yayılma gecikmesi
Gecikme = sürücünün üretimi, güç kaynagı voltajı

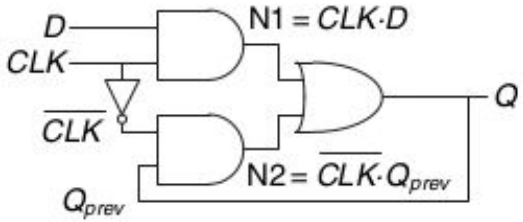
Halka osilatör periyodunun dogru bir sekilde tahmin edilmesi zordur.
Halka osilatörü, sıfır girisli ve periyodik olarak degisen bir çıkıslı sıralı bir devredir.

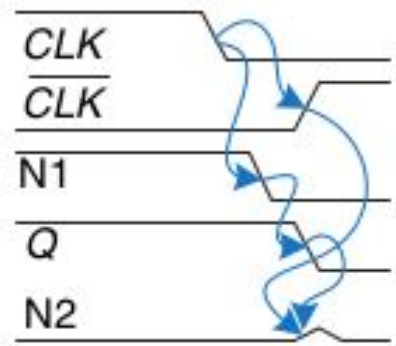
Senkron Lojik Tasarım - Yarış Durumu -

$$Q = CLK \cdot D + \overline{CLK} \cdot Q_{prev}$$

Ali az sayıda kapı kullanılarak, daha basarılı bir D Latch tasarladıgını idda ediyor.

CLK	D	Q_{prev}	Q
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1





CLK = D = 1 => Q=1 (transparent)

CLK = düser => Q=1(eski degerini hatırlar)

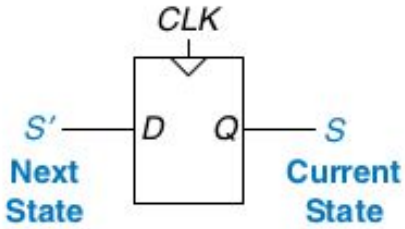
CLK dan CLK' gecis gecikmesinin göreceli olarak daha uzun oldugu varsayılsın

CLK' = yükselir => N1 ve Q; CLK' dan önce düser

Bu durumda;

N2 kesinlikle yükselemeyecektir.

Senkron Ardısık Devreler

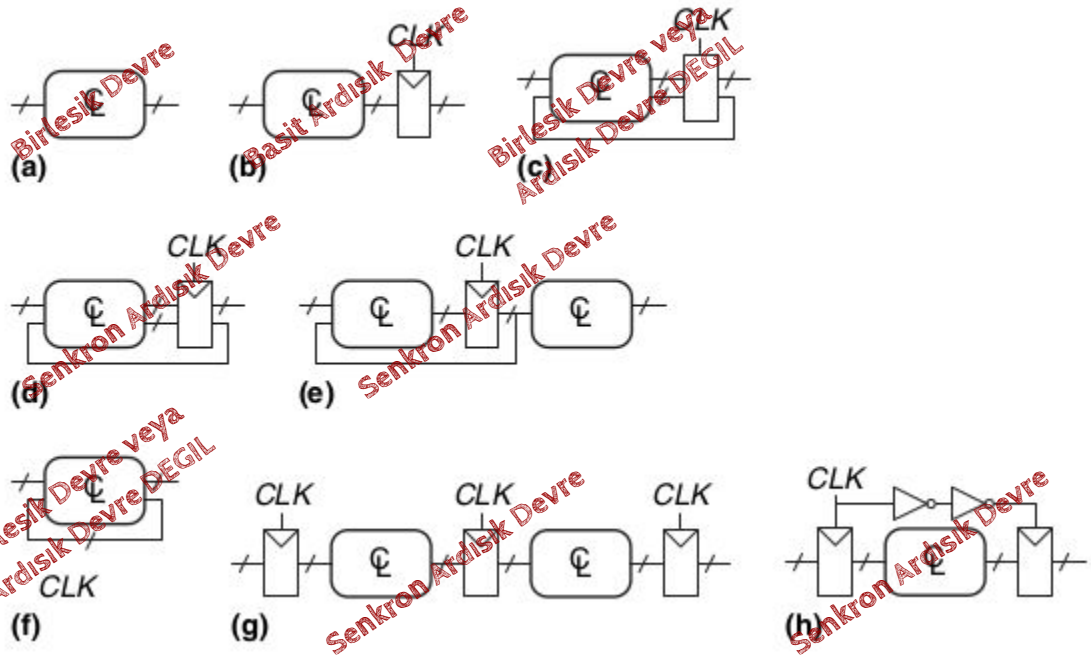


- Herbir elemanı saklayıcı veya birlesik devre
- En az bir elemanı saklayıcı
- Bütün saklayıcılar aynı saat sinyalini kullanılır
- Herbir döngüsel yol bir sakalayıcı(döngüyü bekletebilen) içerir

Bir girise(D) bir saat sinyaline(CLK), bir çıkısa (Q) ve iki duruma {0, 1} sahip; Bir flip-flop, en basit senkron ardısık devredir.
Mevcut durum S, sonraki durum ise S' ile ifade edilmistir.

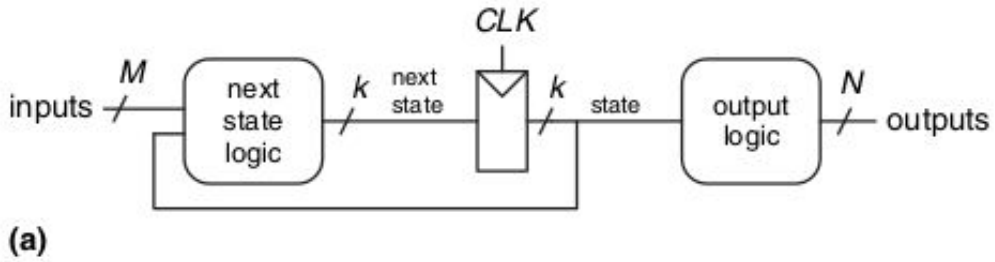
Senkron Ardısık Devreler

Asagıdalilerden hagisi senkron ardısık devredir.



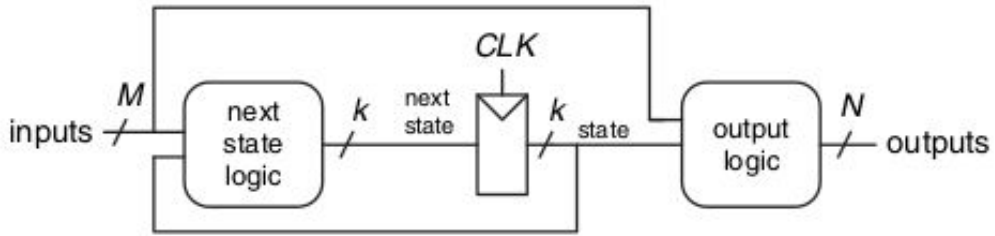
Sonlu Durum Makinaları

İstenilen fonksiyonu yerine getiren Senkron Ardısıl devrelerin tasarlanması için sematik bir yol sunar.



(a)

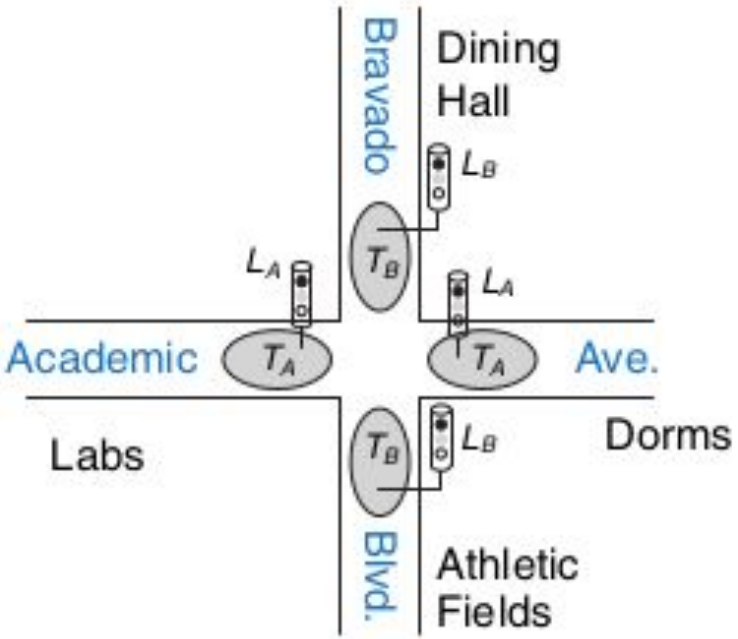
Moore machine

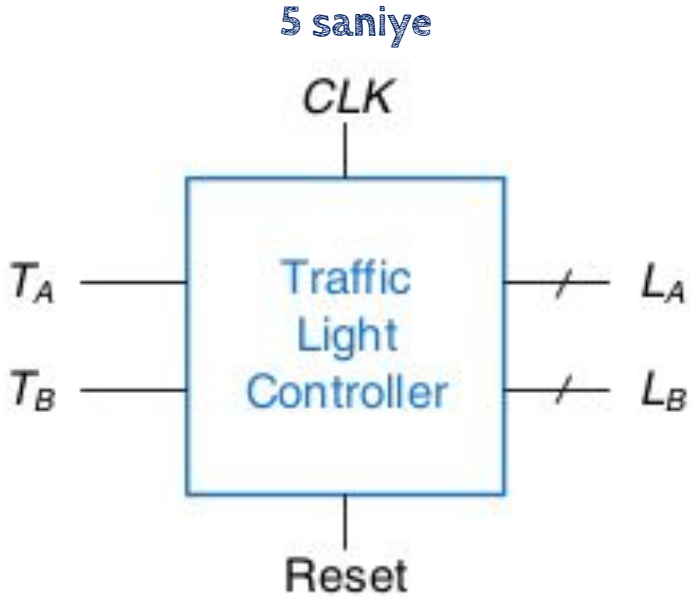


(b)

Mealy machine

Sonlu Durum Makinaları Örnegi





Sonlu Durum Makinaları Örnegi

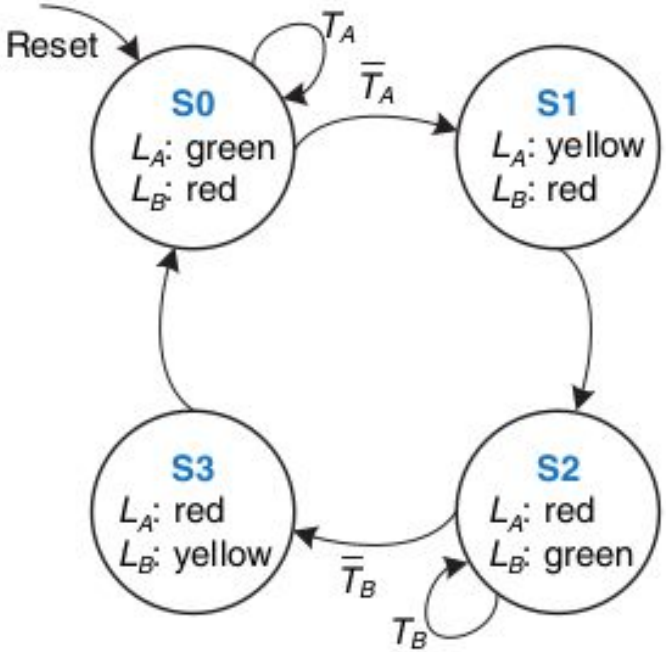


Table 3.1 State transition table

Current State S	Inputs T_A	Inputs T_B	Next State S'
S0	0	X	S1
S0	1	X	S0
S1	X	X	S2
S2	X	0	S3
S2	X	1	S2
S3	X	X	S0

Sonlu Durum Makinaları Örnegi

Table 3.1 State transition table

Current State S	Inputs T_A	Inputs T_B	Next State S'
S0	0	X	S1
S0	1	X	S0
S1	X	X	S2
S2	X	0	S3
S2	X	1	S2
S3	X	X	S0

Table 3.2 State encoding

State	Encoding $S_{1:0}$
S0	00
S1	01
S2	10
S3	11

Table 3.3 Output encoding

Output	Encoding $L_{1:0}$
green	00
yellow	01
red	10

Sonlu Durum Makinaları Örnegi

Table 3.1 State transition table

Current State S	Inputs T_A	Inputs T_B	Next State S'
S0	0	X	S1
S0	1	X	S0
S1	X	X	S2
S2	X	0	S3
S2	X	1	S2
S3	X	X	S0

Table 3.2 State encoding

State	Encoding $S_{1:0}$
S0	00
S1	01
S2	10
S3	11

Table 3.3 Output encoding

Output	Encoding $L_{1:0}$
green	00
yellow	01
red	10

Table 3.4 State transition table with binary encodings

Current State S_1	Current State S_0	Inputs T_A	Inputs T_B	Next State S'_1	Next State S'_0
0	0	0	X	0	1
0	0	1	X	0	0
0	1	X	X	1	0
1	0	X	0	1	1
1	0	X	1	1	0
1	1	X	X	0	0

$$S'_1 = \overline{S}_1 S_0 + S_1 \overline{S}_0 \overline{T}_B + S_1 \overline{S}_0 T_B \qquad S'_1 = S_1 \oplus S_0$$

$$S'_0 = \overline{S}_1 \overline{S}_0 \overline{T}_A + S_1 \overline{S}_0 \overline{T}_B \qquad S'_0 = \overline{S}_1 \overline{S}_0 \overline{T}_A + S_1 \overline{S}_0 \overline{T}_B$$

Sonlu Durum Makinaları Örnegi

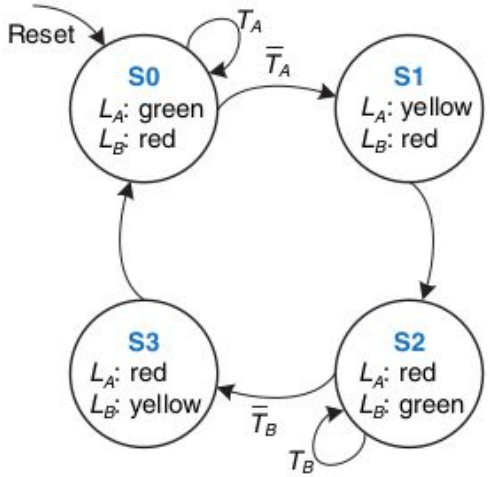


Table 3.5 Output table

Current State		Outputs			
S_1	S_0	L_{A1}	L_{A0}	L_{B1}	L_{B0}
0	0	0	0	1	0
0	1	0	1	1	0
1	0	1	0	0	0
1	1	1	0	0	1

Table 3.2 State encoding

State	Encoding $S_{1:0}$
S0	00
S1	01
S2	10
S3	11

Table 3.3 Output encoding

Output	Encoding $L_{1:0}$
green	00
yellow	01
red	10

$$L_{A1} = S_1$$

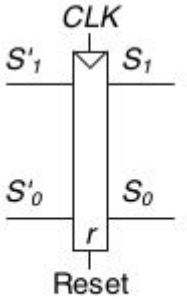
$$L_{A0} = \overline{S}_1 S_0$$

$$L_{B1} = \overline{S}_1$$

$$L_{B0} = S_1 S_0$$

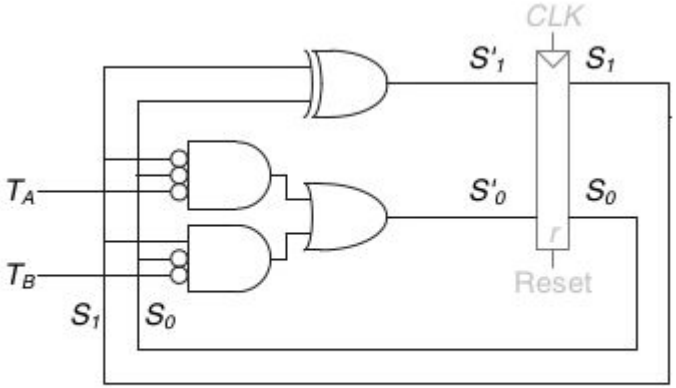
Sonlu Durum Makinaları Örnegi

Moore FSM



state register

(a)



inputs next state logic state register

(b)

$$S_1' = S_1 \oplus S_0$$

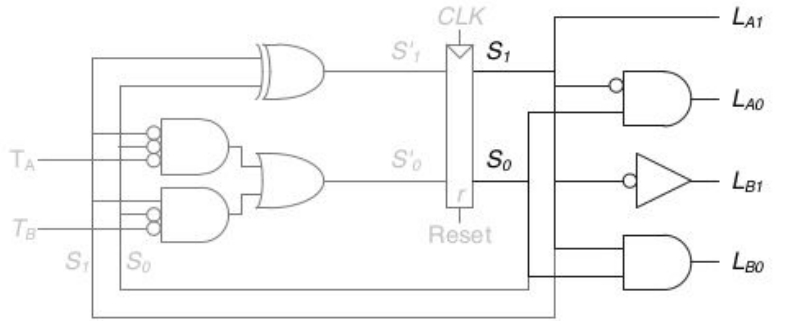
$$S_0' = \overline{S}_1 \overline{S}_0 \overline{T}_A + S_1 \overline{S}_0 \overline{T}_B$$

$$L_{A1} = S_1$$

$$L_{A0} = \overline{S}_1 S_0$$

$$L_{B1} = \overline{S}_1$$

$$L_{B0} = S_1 S_0$$



inputs next state logic state register output logic outputs

(c)

Sonlu Durum Makinaları Örnegi

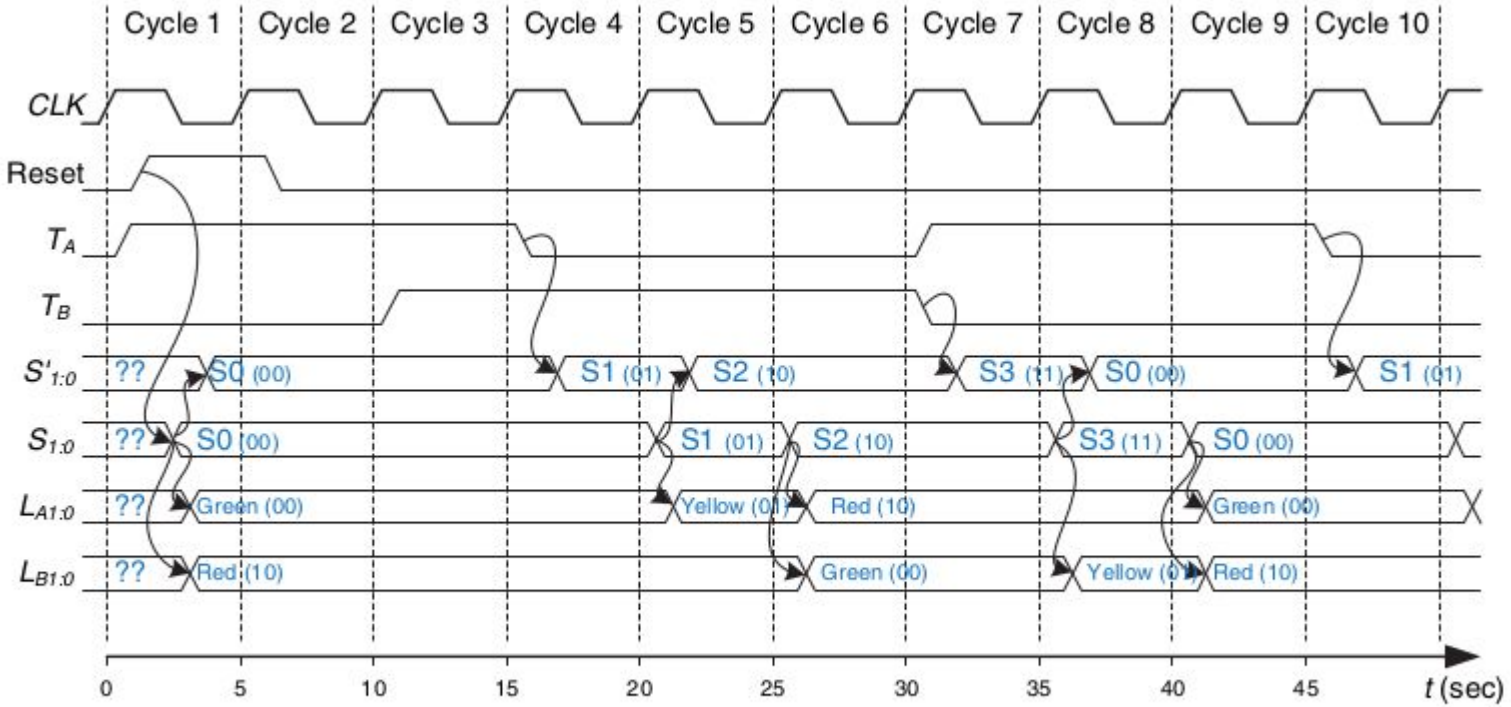


Figure 3.27 Timing diagram for traffic light controller

Durum Kodlama

ikili durum kodlama: Bir önceki örnekte oldugu gibi her durum ikili kodlama ile kodlanmaktadır. Daha fazla lojik kapı kullanılır.

tek-bit kodlama: Her durumda tek bir bit 1 degerini alan kodlamadır. Üç durum için 001, 010, 100 Daha fazla flip-flop kullanılır.

Durum Kodlama Örnek (Divide-by-N counter)

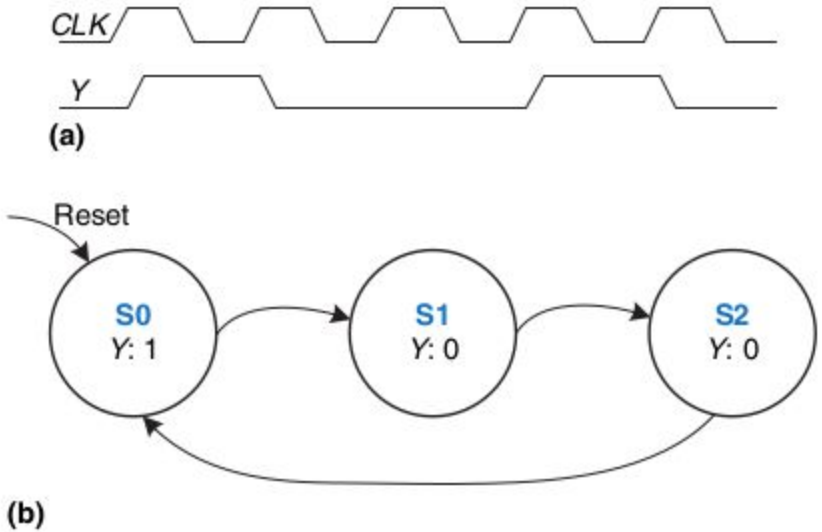


Figure 3.28 Divide-by-3 counter (a) waveform and (b) state transition diagram

Durum Kodlama Örnek (Divide-by-N counter)

Table 3.6 Divide-by-3 counter state transition table

Current State	Next State
S0	S1
S1	S2
S2	S0

Table 3.7 Divide-by-3 counter output table

Current State	Output
S0	1
S1	0
S2	0

Table 3.8 One-hot and binary encodings for divide-by-3 counter

State	One-Hot Encoding			Binary Encoding	
	S_2	S_1	S_0	S_1	S_0
S0	0	0	1	0	0
S1	0	1	0	0	1
S2	1	0	0	1	0

Durum Kodlama Örnek (Divide-by-N counter)

Table 3.6 Divide-by-3 counter state transition table

Current State	Next State
S0	S1
S1	S2
S2	S0

Table 3.7 Divide-by-3 counter output table

Current State	Output
S0	1
S1	0
S2	0

Table 3.9 State transition table with binary encoding

Current State		Next State	
S_1	S_0	S_1'	S_0'
0	0	0	1
0	1	1	0
1	0	0	0

$$S_1' = \bar{S}_1 S_0$$
$$S_0' = \bar{S}_1 \bar{S}_0$$

$$Y = \bar{S}_1 \bar{S}_0$$

Durum Kodlama Örnek (Divide-by-N counter)

Table 3.6 Divide-by-3 counter state transition table

Current State	Next State
S0	S1
S1	S2
S2	S0

Table 3.7 Divide-by-3 counter output table

Current State	Output
S0	1
S1	0
S2	0

Table 3.10 State transition table with one-hot encoding

Current State			Next State		
S_2	S_1	S_0	S_2'	S_1'	S_0'
0	0	1	0	1	0
0	1	0	1	0	0
1	0	0	0	0	1

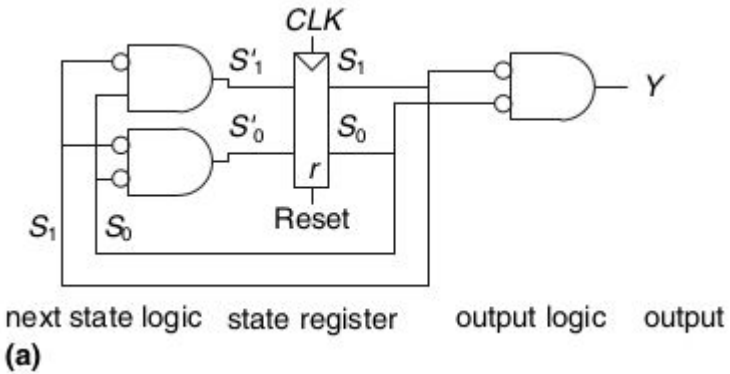
$$S_2' = S_1$$
$$S_1' = S_0$$
$$S_0' = S_2$$

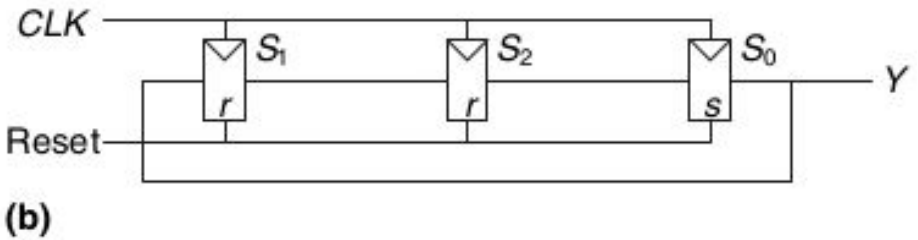
$$Y = S_0$$

Durum Kodlama Örnek (Divide-by-N counter)

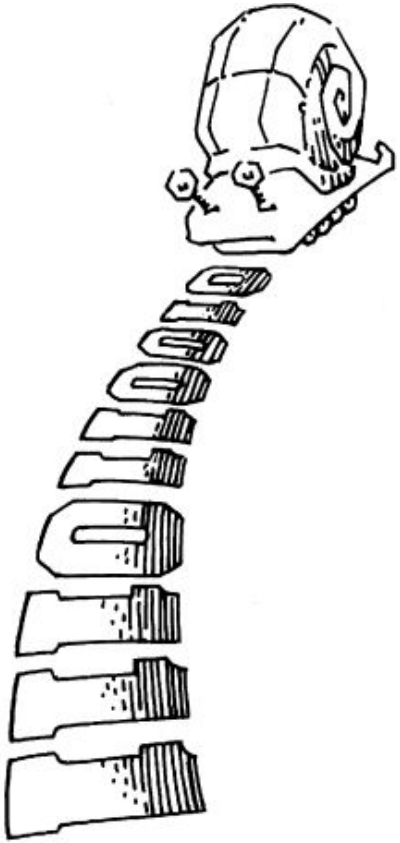
$$S_1' = \bar{S}_1 S_0$$
$$S_0' = \bar{S}_1 \bar{S}_0$$
$$Y = \bar{S}_1 \bar{S}_0$$

$$S_2' = S_1$$
$$S_1' = S_0$$
$$S_0' = S_2$$
$$Y = S_0$$



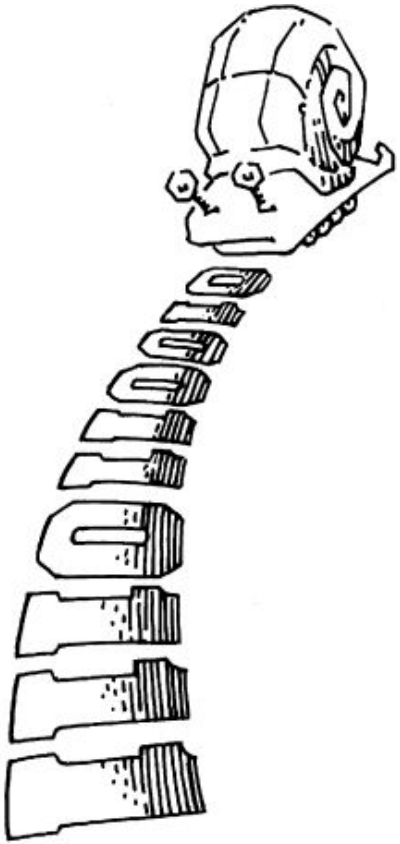


Moore and Mealy Makineleri

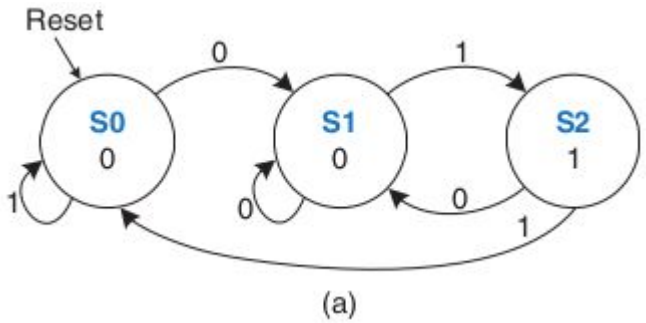


- Ali, FSM mantığı ile hareket eden robot bir salyangoz tasarladı.
- Salyangoz, 1 ve 0 dizisini içeren bir kağıt boyunca sağdan sola doğru ilerlemektedir.
- Salyangoz her saat dögüsünde bir bit geçmektedir.
- Geçtiği son iki bir 01 olursa salyangoz gülümsemektedir.
- Salyangozun ne zaman gülümseyeceğini hesaplayan FSM yi gelistiriniz.
- A girisi, salyangoz anteninin altındaki bittir.
- Salyangoz gülümsediğinde, Y çıkısı TRUE olur.
- Moore ve Mealy durum makinesi tasarımlarını karsılastırın.
- Alyssa'nın salyangozu 0100110111 dizisi boyunca gezinirken girisi, durumları ve çıktıyı gösteren her makine için bir zamanlama diyagramı çizin.

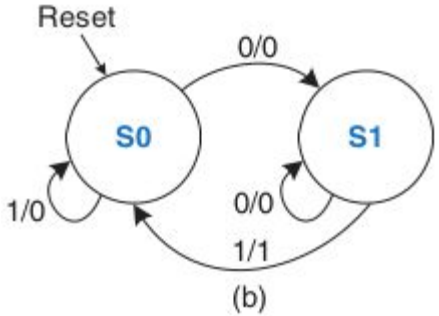
Moore and Mealy Makineleri



Moore Makinesi



Mealy Makinesi



A/Y

A: Gecisi tetikleyen girisler

Y: Gecise karsılık gelen cıkıs

Moore Makinesi

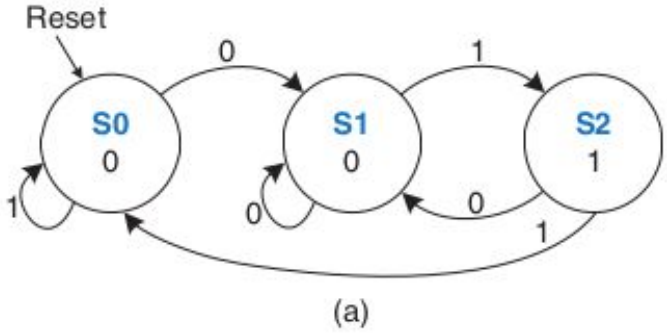


Table 3.11 Moore state transition table

Current State S	Input A	Next State S'
S0	0	S1
S0	1	S0
S1	0	S1
S1	1	S2
S2	0	S1
S2	1	S0

Table 3.12 Moore output table

Current State S	Output Y
S0	0
S1	0
S2	1

Moore Makinesi

ikilik durum kodlama:
S0=00, S1=01, S2=10

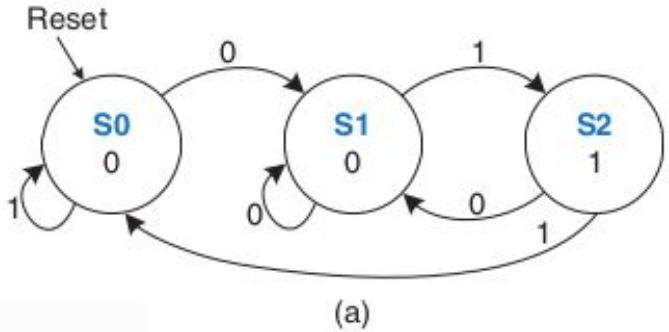


Table 3.13 Moore state transition table with state encodings

Current State S_1	Current State S_0	Input A	Next State S'_1	Next State S'_0
0	0	0	0	1
0	0	1	0	0
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	0	0

Table 3.14 Moore output table with state encodings

Current State S_1	Current State S_0	Output Y
0	0	0
0	1	0
1	0	1

$$S'_1 = S_0 A$$
$$S'_0 = \overline{A}$$

$$Y = S_1$$

Mealy Makinesi

ikilik durum kodlama:
S0=0, S1=1

Table 3.15 Mealy state transition and output table

Current State S	Input A	Next State S'	Output Y
S0	0	S1	0
S0	1	S0	0
S1	0	S1	0
S1	1	S0	1

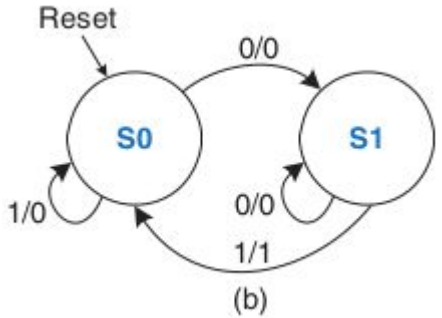


Table 3.16 Mealy state transition and output table with state encodings

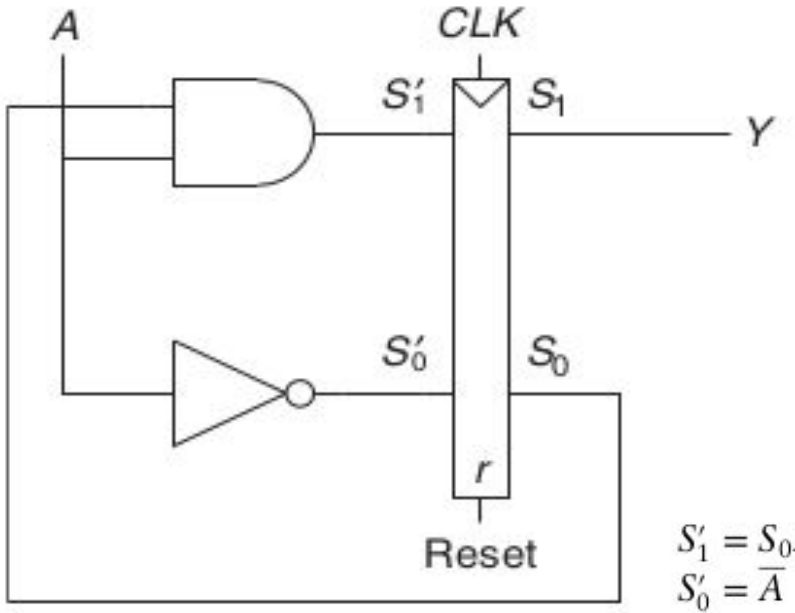
Current State S_0	Input A	Next State S'_0	Output Y
0	0	1	0
0	1	0	0
1	0	1	0
1	1	0	1

$$S'_0 = \overline{A}$$

$$Y = S_0 A$$

Moore and Mealy Makineleri

Moore Makinesi

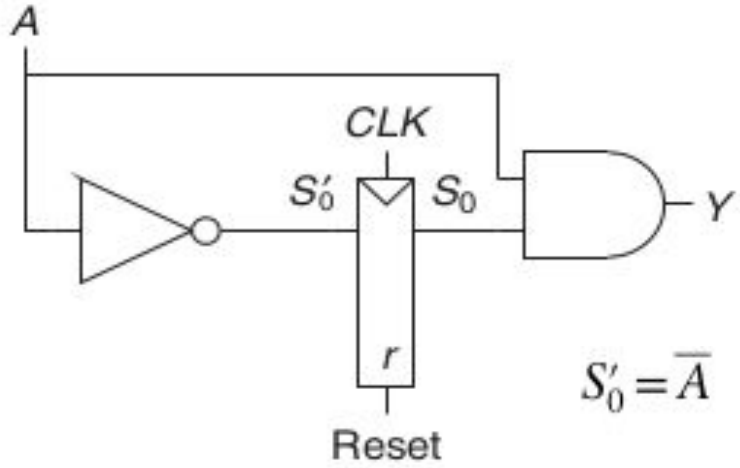


(a)

$$S_1' = S_0 A$$
$$S_0' = \overline{A}$$

$$Y = S_1$$

Mealy Makinesi

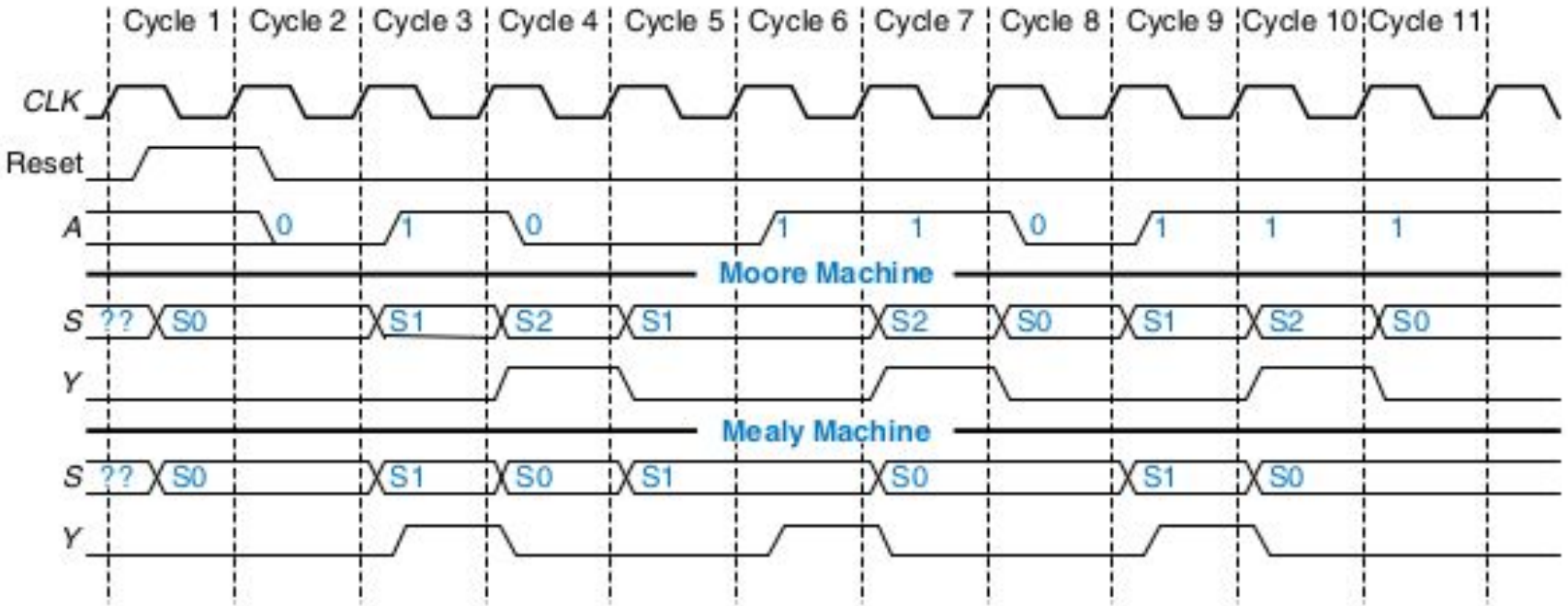


(b)

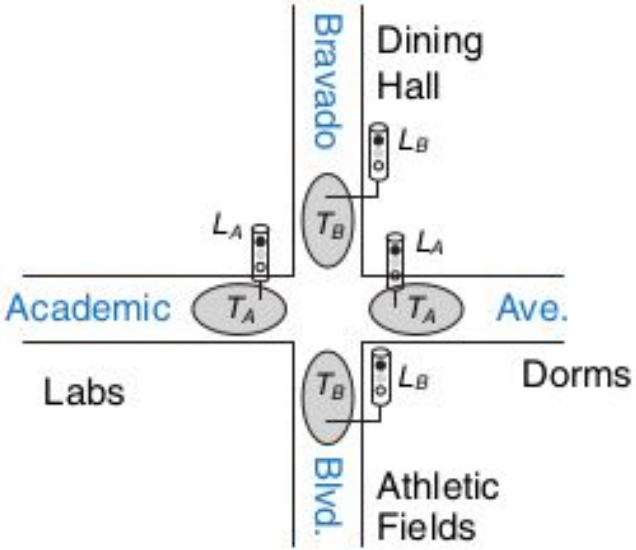
$$S_0' = \overline{A}$$

$$Y = S_0 A$$

Moore and Mealy Makineleri

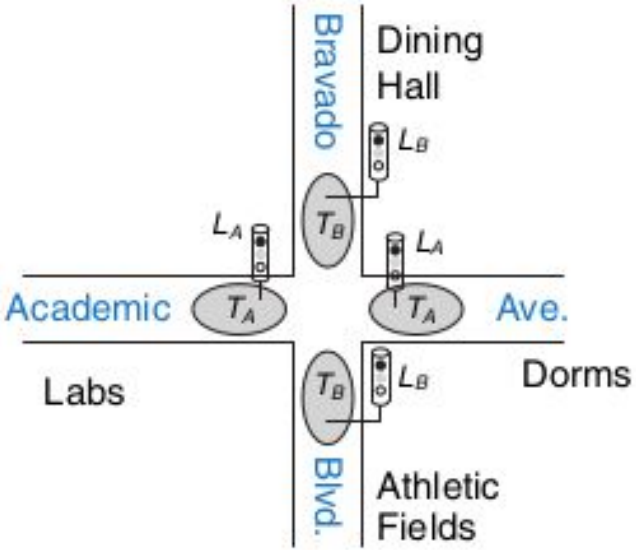


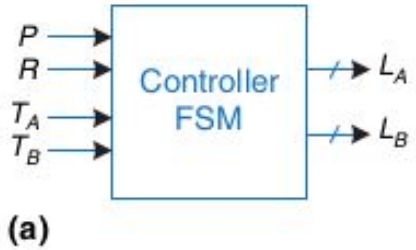
Durum Makinaları Üretimi



- Gecit Töreni Durum Makinası:
- Bravadı Bulvarında, futbol takımı ve bando takımı gecene kadar ısıgın yesilde kalması gerekmektedir.
- Dolayısıyla Kontrolör iki giris daha alır: P ve R
- Gecit törenini baslatan döngü: P
- Gecit töreninisonlandıran döngü: R
- Geçit töreni modundayken, LB yesil olana kadar normal sıralamasını takip eder ve geçit modu sona erene kadar LB yesil durumduna kalır.

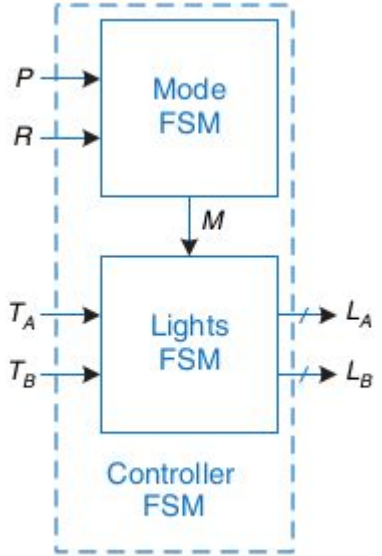
Durum Makinaları Üretimi





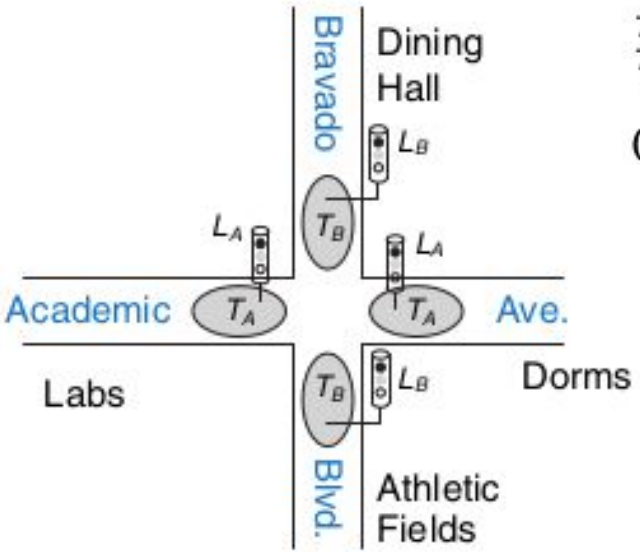
(a)

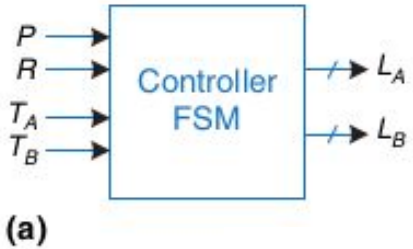
M: Gecit Töreni Modu



(b)

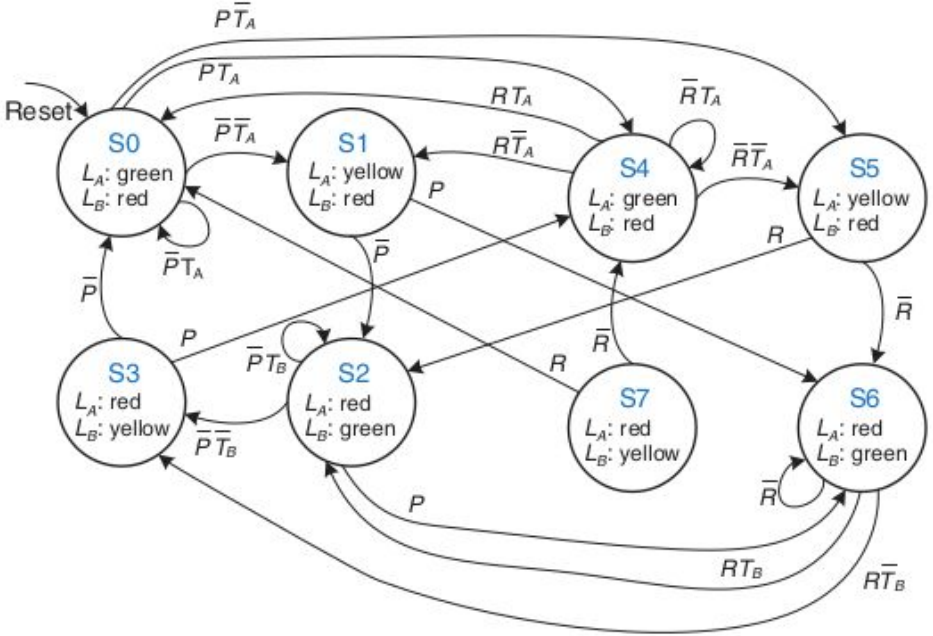
Durum Makinaları Üretimi





S0-S3: Normal Mod

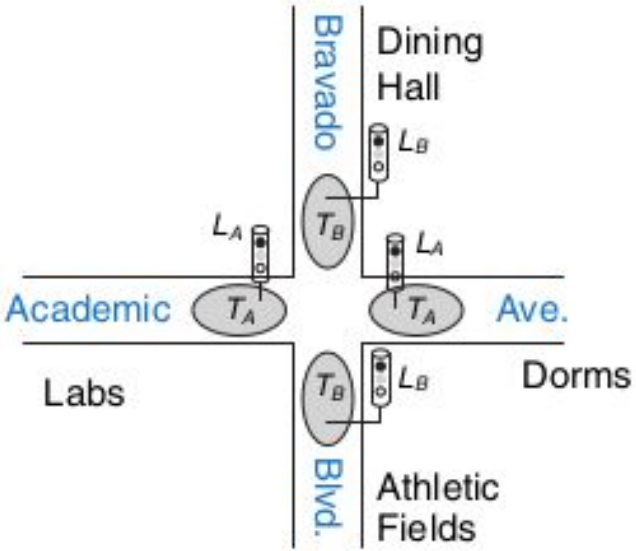
S4-S7: Gecit Modu

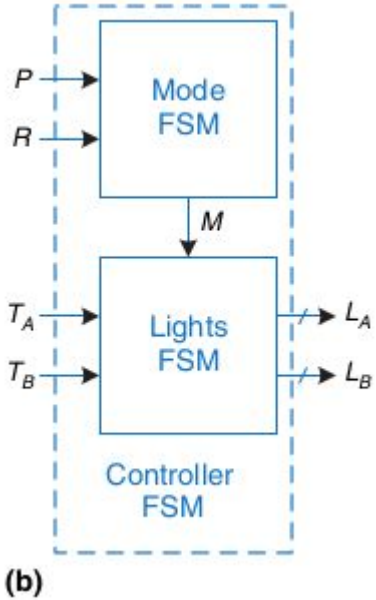


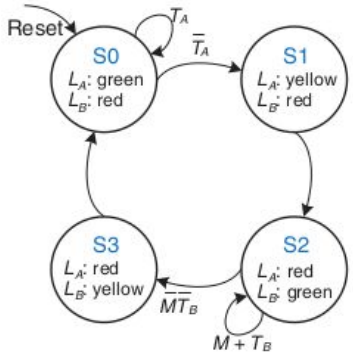
Durum Makinaları Üretimi

S0-S3: Normal Mod

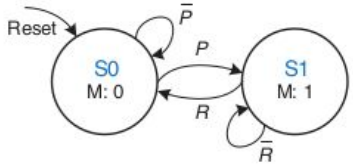
S4-S7: Gecit Modu





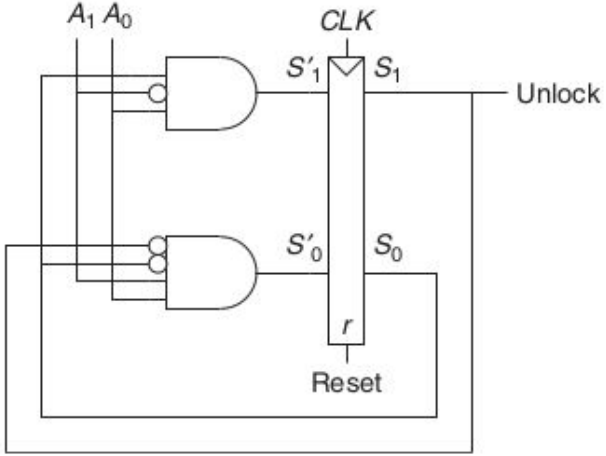


Lights FSM



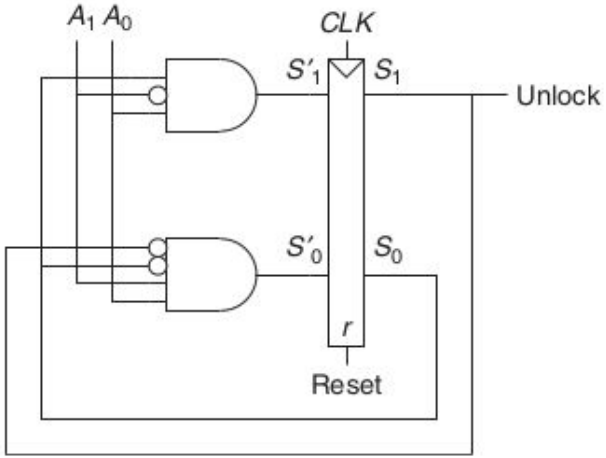
Mode FSM

FSM'yi Sematikten Çıkarma



- Ali eve geldi, ancak kapıdaki tus takımı kilidi yenilendigi icin eski kodu artık çalısmıyor.
- Ancak tus takımını yenileyen ustalar kapının üstünde yandaki devre seması bırakmıslardır.
- Ali, devrenin sonlu bir durum makinesi olabilecegini düsünüyor ve iceri girmek için durum geçis diyagramını çıkarmaya karar veriyor.

FSM'yi Sematikten Çıkarma



- Bu bir Moore makinesidir çünkü çıktı yalnızca durum bitlerine bağlıdır.
- Girişler : A0, A1
- Çıkış: Unlock

$$S_1' = S_0 \overline{A_1} A_0$$

$$S_0' = \overline{S_1}\ \overline{S_0} A_1 A_0$$

$$Unlock = S_1$$

FSM'yi Sematikten Çıkarma

$$S_1' = S_0\overline{A_1}A_0$$

$$S_0' = \overline{S_1}\,\overline{S_0}A_1A_0$$

$$Unlock = S_1$$

Table 3.17 Next state table derived from circuit in Figure 3.35

Current State S_1	Current State S_0	Input A_1	Input A_0	Next State S_1'	Next State S_0'
0	0	0	0	0	0
0	0	0	1	0	0
0	0	1	0	0	0
0	0	1	1	0	1
0	1	0	0	0	0
0	1	0	1	1	0
0	1	1	0	0	0
0	1	1	1	0	0
1	0	0	0	0	0
1	0	0	1	0	0
1	0	1	0	0	0
1	0	1	1	0	0
1	1	0	0	0	0
1	1	0	1	1	0
1	1	1	0	0	0
1	1	1	1	0	0

Table 3.18 Output table derived from circuit in Figure 3.35

Current State S_1	Current State S_0	Output *Unlock*
0	0	0
0	1	0
1	0	1
1	1	1

FSM'yi Sematikten Çıkarma

$$S_1' = S_0 \overline{A_1} A_0$$

$$S_0' = \overline{S_1}\ \overline{S_0} A_1 A_0$$

$$Unlock = S_1$$

Table 3.19 Reduced next state table

Current State		Input		Next State	
S_1	S_0	A_1	A_0	S_1'	S_0'
0	0	0	0	0	0
0	0	0	1	0	0
0	0	1	0	0	0
0	0	1	1	0	1
0	1	0	0	0	0
0	1	0	1	1	0
0	1	1	0	0	0
0	1	1	1	0	0
1	0	X	X	0	0

Table 3.18 Output table derived from circuit in **Figure 3.35**

Current State		Output
S_1	S_0	*Unlock*
0	0	0
0	1	0
1	0	1
1	1	1

FSM'yi Sematikten Çıkarma

$$S_1' = S_0 \overline{A_1} A_0$$

$$S_0' = \overline{S_1}\ \overline{S_0} A_1 A_0$$

$$Unlock = S_1$$

Table 3.21 Symbolic next state table

Current State S	Input A	Next State S'
S0	0	S0
S0	1	S0
S0	2	S0
S0	3	S1
S1	0	S0
S1	1	S2
S1	2	S0
S1	3	S0
S2	X	S0

Table 3.22 Symbolic output table

Current State S	Output *Unlock*
S0	0
S1	0
S2	1

FSM'yi Sematikten Çıkarma

$$S_1' = S_0\overline{A_1}A_0$$

$$S_0' = \overline{S_1}\,\overline{S_0}A_1A_0$$

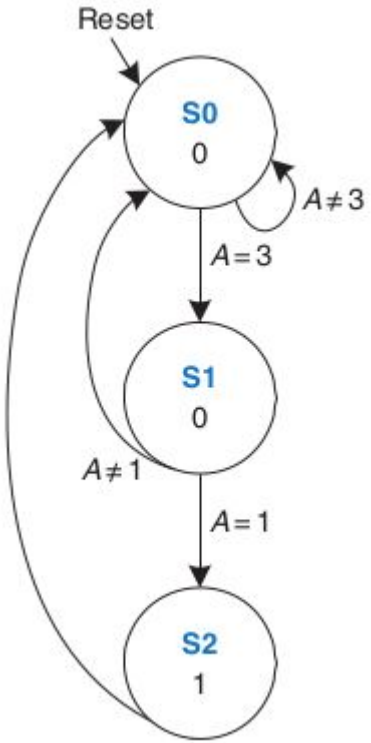
$$Unlock = S_1$$

Table 3.21 Symbolic next state table

Current State S	Input A	Next State S'
S0	0	S0
S0	1	S0
S0	2	S0
S0	3	S1
S1	0	S0
S1	1	S2
S1	2	S0
S1	3	S0
S2	X	S0

Table 3.22 Symbolic output table

Current State S	Output *Unlock*
S0	0
S1	0
S2	1



FSM Özeti

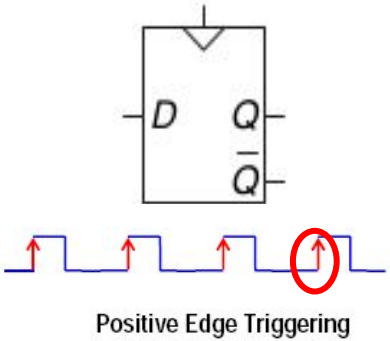
Sonlu durum makinaları, davranısı verilmis ardısıl devrelerin tasarımı için güclü bir yöntemdir.

Bir FSM tasarlamak için asagıdaki adımlar izlenmelidir:

- Devre giris ve çıkıslarını tanımlayın.
- Bir durum geçis diyagramı çizin.
- Moore makinesi için:
  - Bir durum geçis tablosu yazın.
  - Bir çıktı tablosu yazın.
- Mealy makinesi için:
  - Durum geçis ve çıktı tablosu beraber yazın.
- Kodlamayı secin-Kodlama tasarlanan donanımı etkiler-
- Sonraki durum ve çıkıs lojigi için boolean denklemlerini yazın
- Devre semasını çizin.

3.5 Ardışık Mantık Zamanlaması

D örneklenmesi



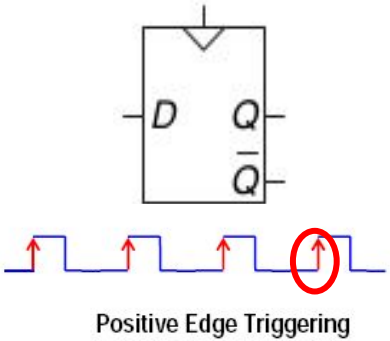
D = 1

D = 0

D 1 => 0 veya 1=>0

3.5 Ardışık Mantık Zamanlaması

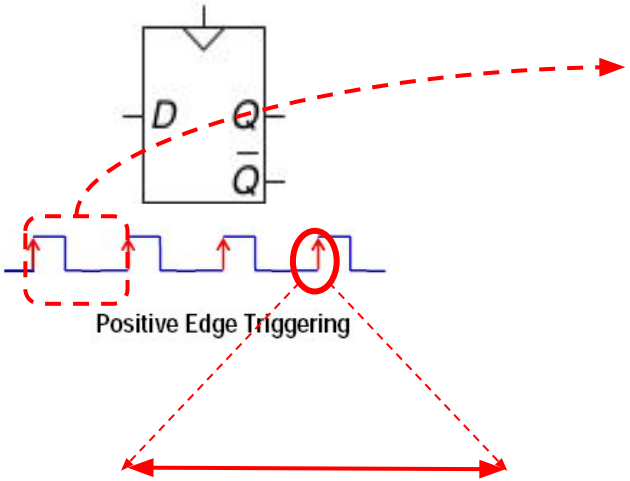
D örneklenmesi



Positive Edge Triggering

3.5 Ardısık Mantık Zamanlaması

D örneklenmesi



Açıklık Süresi

Kurulum Süresi + Tutma süresi

Dinamik Disiplin;

D sinyali saat sinyalin çıkma/inme islemi bittikten sonra kararlı bir hal aldılgında kullanılabilir.

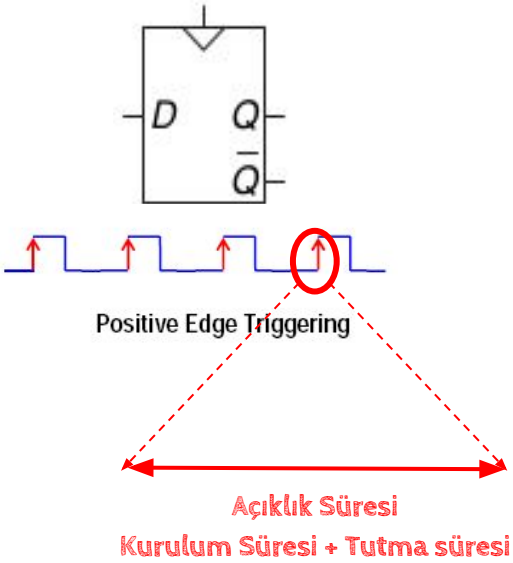
t dogal sayı ve n tam sayılı olmak kosulu ile;

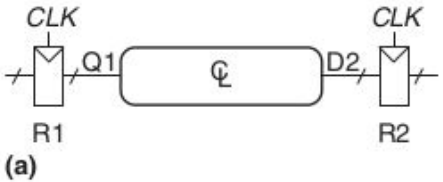
t. saat sinyali sonunda A[t] yazmak yerine

n.saat sinyali sonunda A[n] yazılabilir.

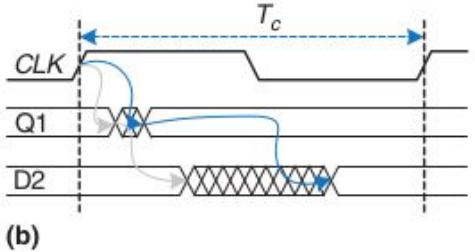
3.5 Ardışık Mantık Zamanlaması

D örneklenmesi





Tc Yeterince uzun olmalı

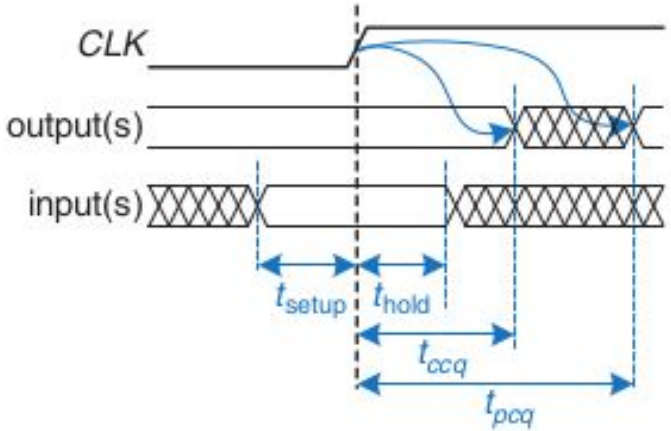


Dinamik Disiplin

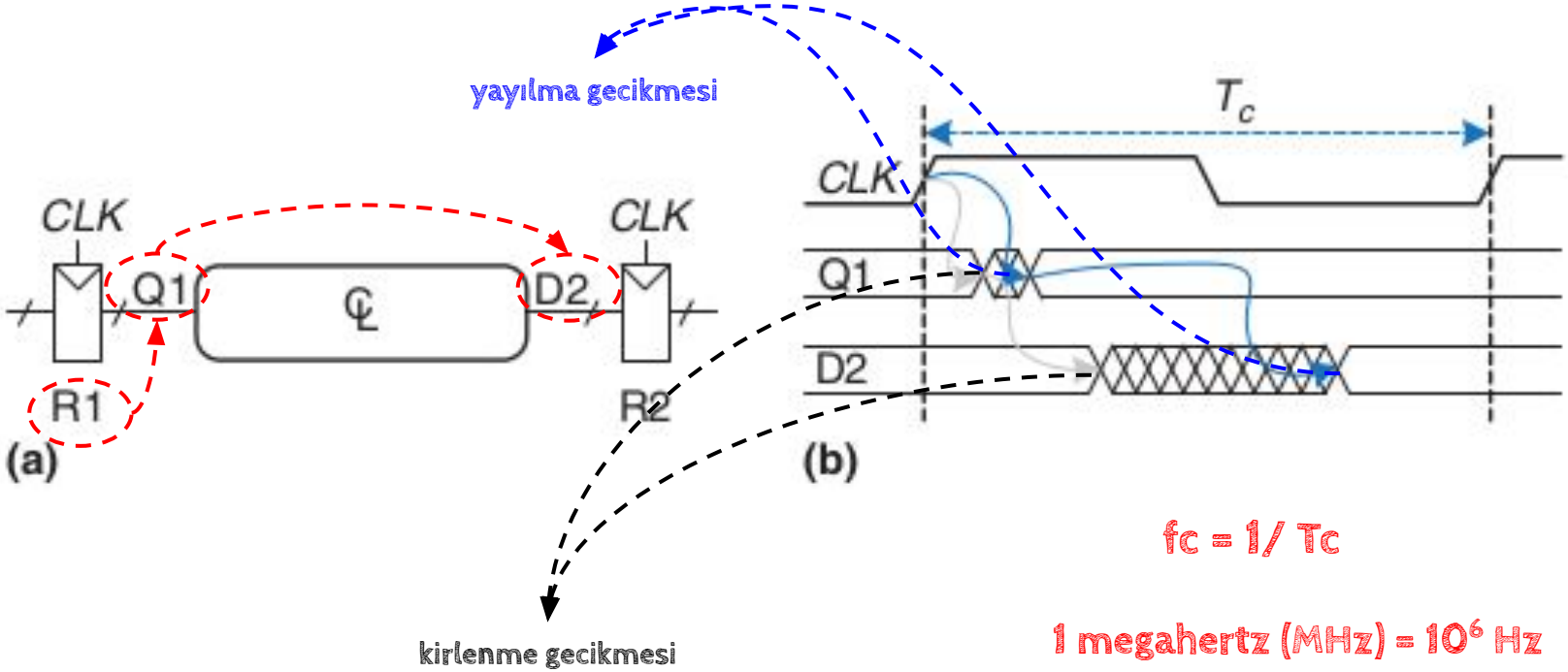
Cıkıs: t_{ccq} (clock-q contamination delay) gecikmesinden sonra, degismeye baslamalı
t_{pcq} (clock-q propagation delay) içinde son degerine ulasmalıdır.

Giris: yükselen kenarından önce kurulum süresi (t_{setup})
yükselen kenarından sonra tutma süresi(t_{hold}) toplamında sabit kalmalıdır.

Dinamik disiplin, bir senkron ardısık devrenin girislerinin açıklık süresince (aperture time = t_{setup} + t_{hold}) kararlı olmasını garanti eder.



Sistem Zamanlaması

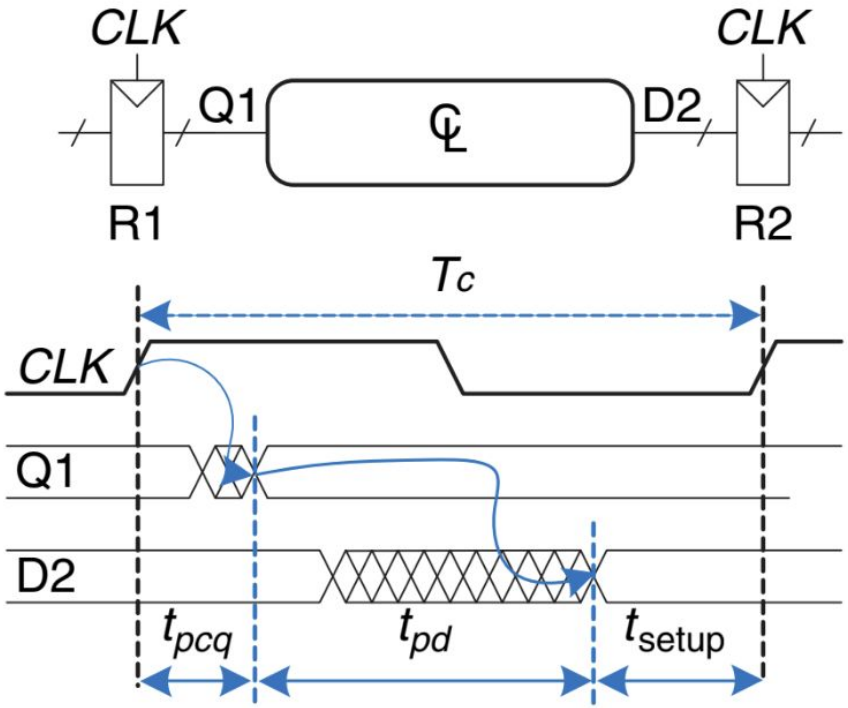


$f_c = 1 / T_c$

1 megahertz (MHz) = 10^6 Hz

1 gigahertz (GHz) = 10^9 Hz

Kurulum Süresi Kısıtlaması



$T_c \geq t_{pcq} + t_{pd} + t_{setup}$

Ticari tasarımlarda,

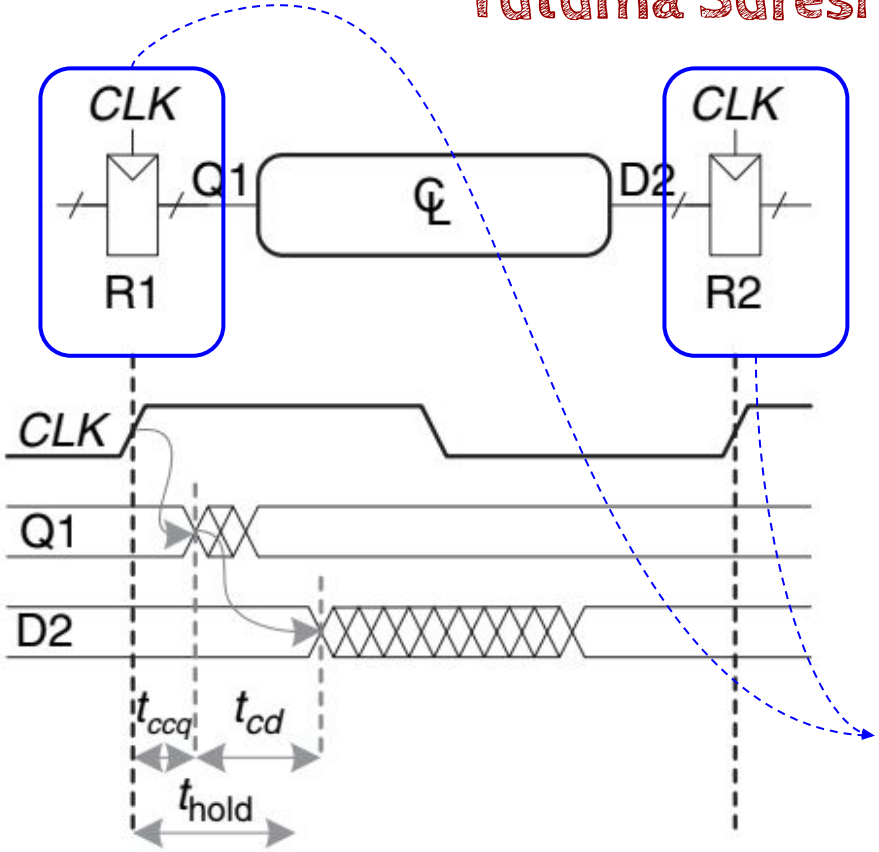
- saat periyodu : clock
- yayılma gecikmesi : tpcq
- kurulum süresi : tsetup

üretici tarafından belirlenir.

Tasarımcının birlesik mantık yayılma gecikmesini kontrol edebilir.

$t_{pd} \geq T_c - (t_{pcq} + t_{setup})$

Tutuma Süresi Kısıtlaması



D2, saatin yükselen kenarından sonra tccq + tcd oldugu anda degisebilir

tccq + tcd ≥ thold

tcd ≥ thold - tccq

tcd = 0

tccq ≥ thold

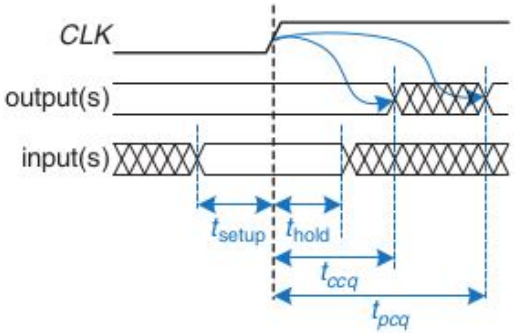
Dinamik Disiplin (Hatırlatma)

Cıkıs: t_{ccq} (clock-q contamination delay) : Çıktı saat yükseldikten, tccq sonra degismeye baslayabilir
(En hızlı gecikme)

t_{pcq} (clock-q propagation delay): saatten tpcq sonra nihai degere oturması gerekir.
(En yavas gecikme)

Giris: yükselen kenarından önce kurulum süresi (t_{setup})
yükselen kenarından sonra tutma süresi(t_{hold}) toplamında sabit kalmalıdır.

Dinamik disiplin, bir senkron ardısık devrenin girislerinin açıklık süresince (aperture time = t_{setup} + t_{hold}) kararlı olmasını garanti eder.



Hepsi Bir Arada

Flip Flop:

t_{ccq} -clock-to-Q kirlenme (contamination)- gecikmesi : 30ps

t_{pcq} -clock-to-Q yayılım(propagation)- gecikmesi: 80ps

t_{setup}: Kurulum süresi: 50ps

t_{hold}: Bekleme süresi: 60ps

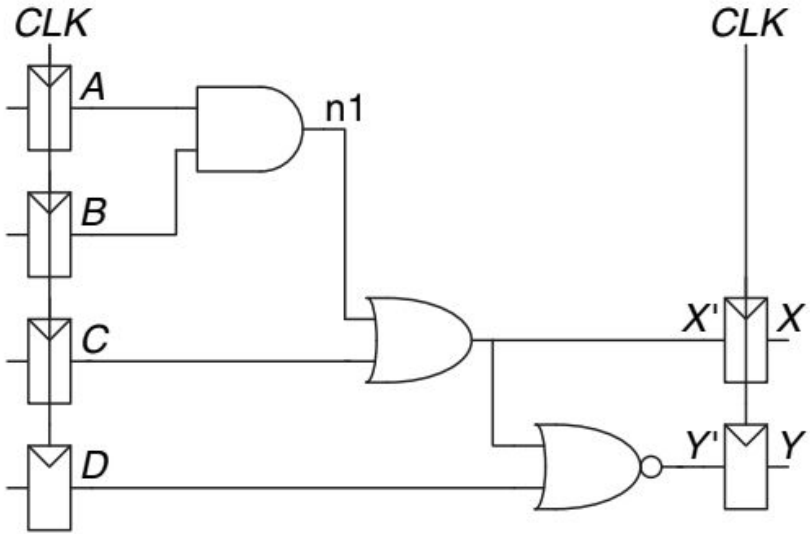
Lojik kapı:

t_{pd} :yayılım gecikmesi, 40ps

t_{cd} :kirleme gecikmesi: 25ps

Maksimum clock frekansı ?

Bekleme süresi ihlali ?

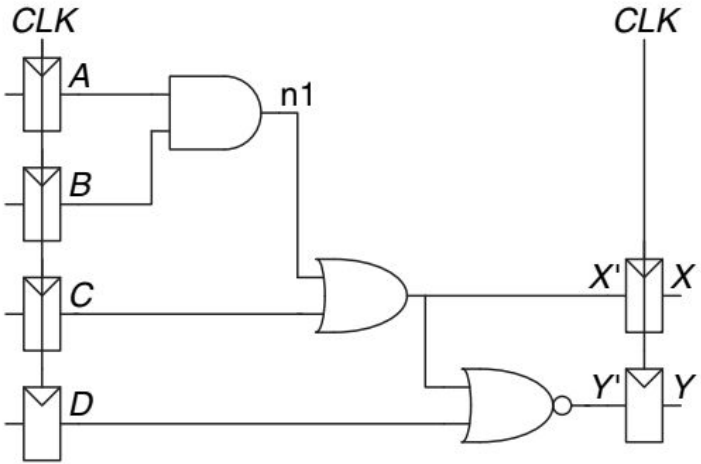


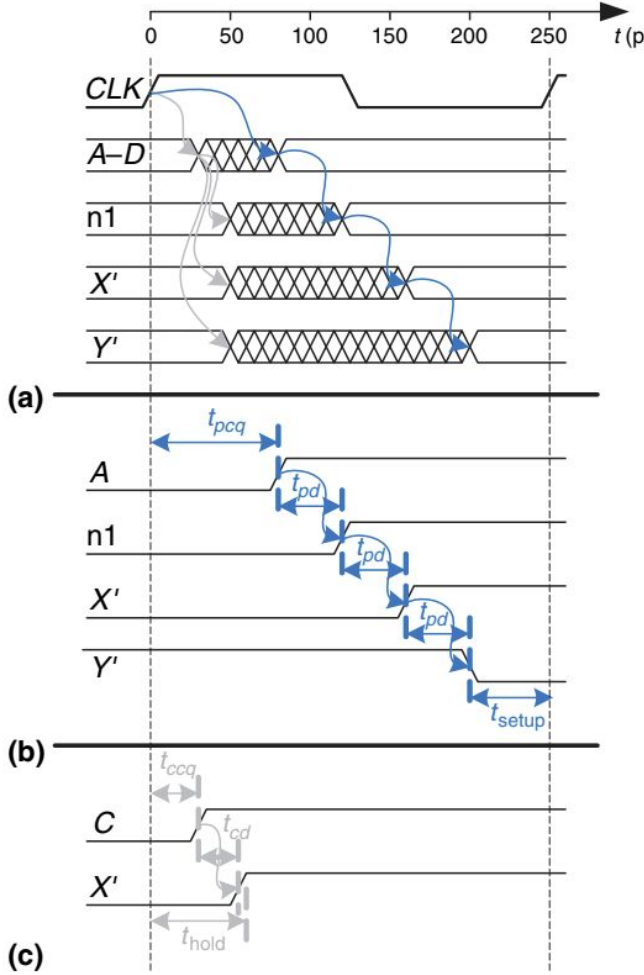
Hepsi Bir Arada

Zaman analizi:

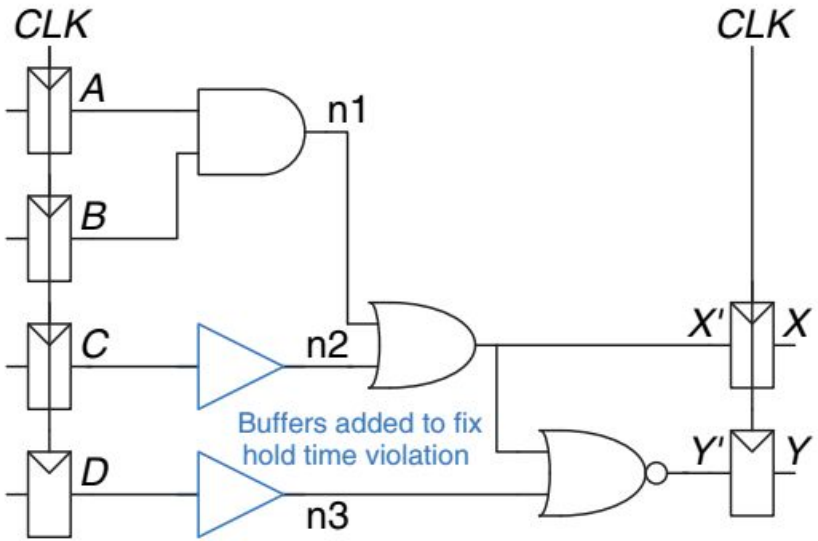
$$T_c \geq t_{pcq} + 3\,t_{pd} + t_{\text{setup}} = 80 + 3 \times 40 + 50 = 250\text{ps}$$

The maximum clock frequency is $f_c = 1/T_c = 4$ GHz.

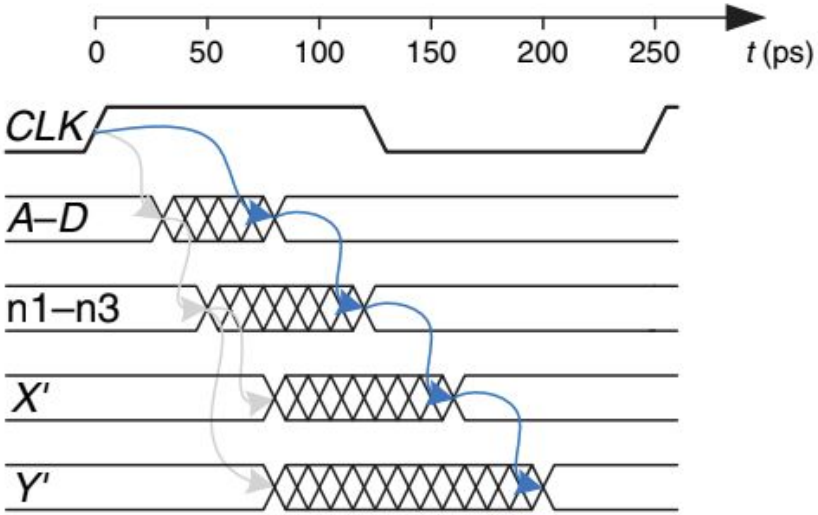




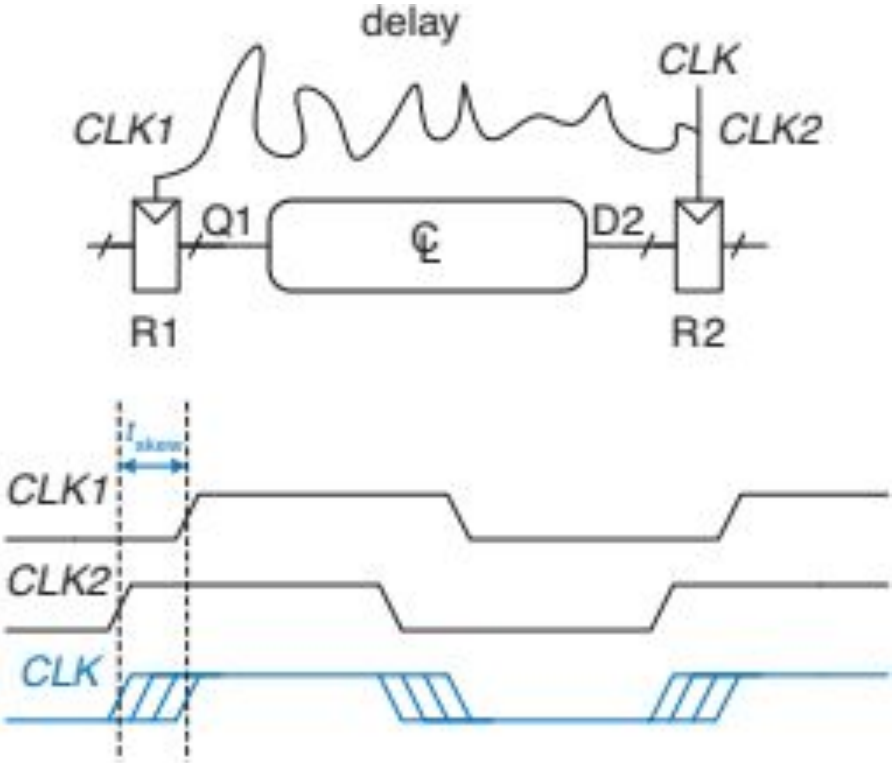
SABİT TUTMA SÜRESİ İHLALLERİ



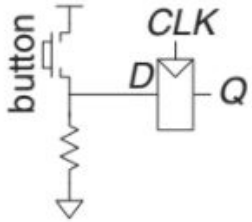
SABİT TUTMA SÜRESİ İHLALLERİ

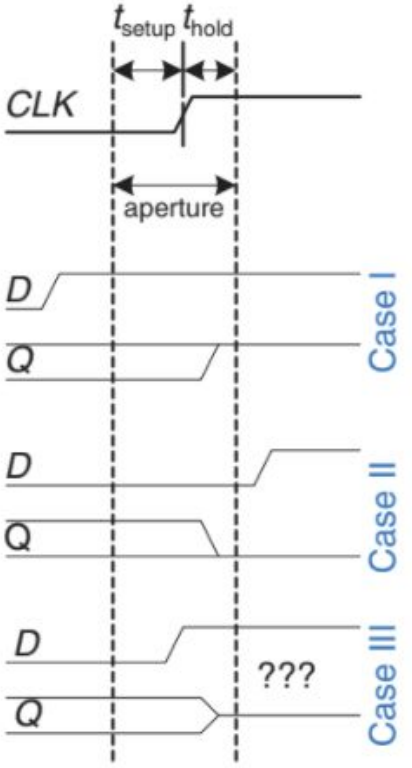


Saat Çarpıklığı*

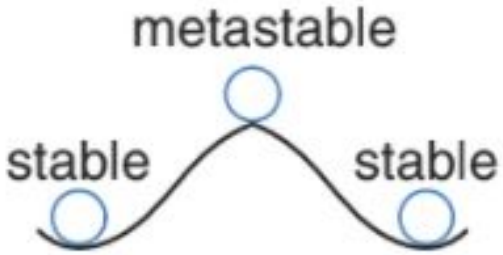


3.5.4 Meta kararlılık

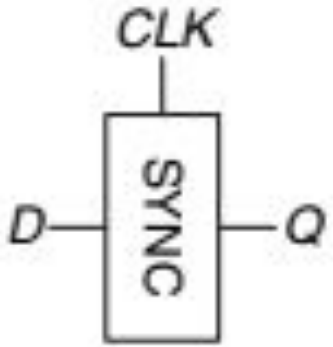




Meta kararlılık Durumu



Es Zamanlama



Paralellik

Paralellik

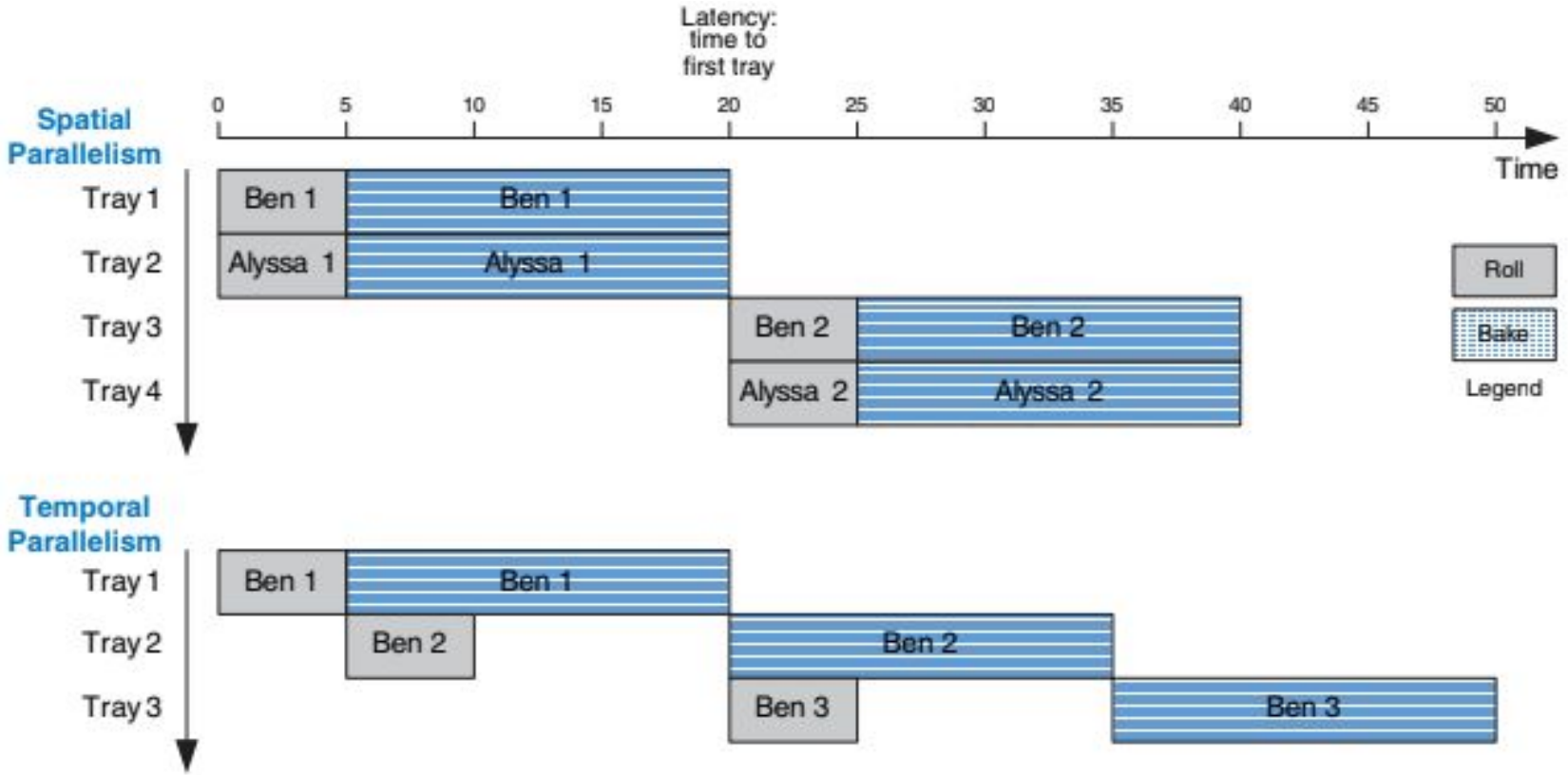
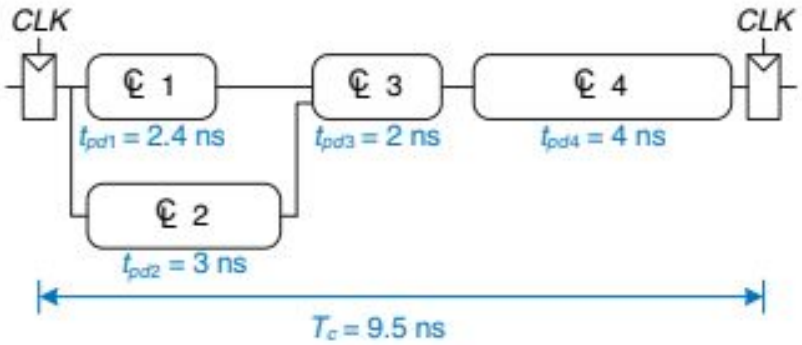
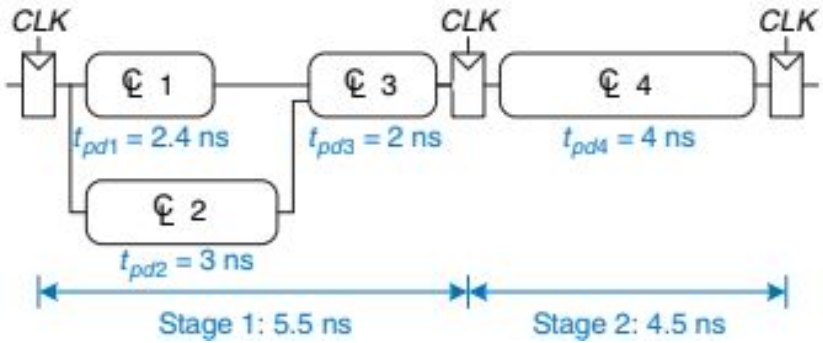


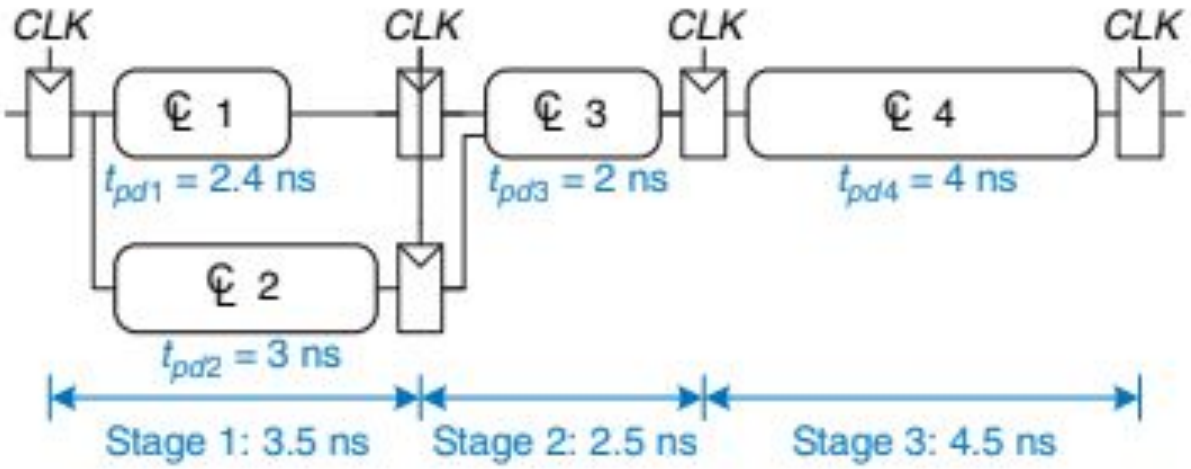
Figure 3.57 Spatial and temporal parallelism in the cookie kitchen

Paralellik


CLK
CLK
1
2
3
4
t_{pd1} = 2.4 ns
t_{pd2} = 3 ns
t_{pd3} = 2 ns
t_{pd4} = 4 ns
T_c = 9.5 ns


CLK
CLK
CLK
1
2
3
4
t_{pd1} = 2.4 ns
t_{pd2} = 3 ns
t_{pd3} = 2 ns
t_{pd4} = 4 ns
Stage 1: 5.5 ns
Stage 2: 4.5 ns

Paralellik


CLK
CLK
CLK
CLK
C 1
C 2
C 3
C 4
t_{pd1} = 2.4 ns
t_{pd2} = 3 ns
t_{pd3} = 2 ns
t_{pd4} = 4 ns
Stage 1: 3.5 ns
Stage 2: 2.5 ns
Stage 3: 4.5 ns

Sorular