

Bellman-Ford

1. Definition

The **Bellman–Ford algorithm** is a single-source shortest path algorithm that computes the minimum distance from a starting node to all other nodes in a weighted graph.

Unlike Dijkstra's algorithm, **Bellman–Ford can handle negative weight edges** and can also **detect negative weight cycles**.

It works on both:

- Directed graphs ✓
- Undirected graphs (if edges are added twice) ✓

2. Solution Approach (How Bellman–Ford Works)

Key Idea

Bellman–Ford performs *edge relaxation* repeatedly.

Relaxation means:

If the distance to vertex v can be improved using edge $(u \rightarrow v)$, update it.

Algorithm Steps

1. **Initialize distances**
 - $\text{dist}[\text{source}] = 0$
 - $\text{dist}[\text{others}] = +\infty$
2. **Relax all edges exactly $(V - 1)$ times**
 - For every edge (u, v, w) , check if
 $\text{dist}[v] > \text{dist}[u] + w$
→ if yes, update.

Why $V - 1$ times?

- The longest possible simple path in a graph has $(V - 1)$ edges.
3. **Check for Negative Weight Cycles**
 - Run one more relaxation round.
 - If any distance still improves → cycle exists.

3. Pseudocode (C++ Style)

```
// Bellman-Ford Algorithm
void bellmanFord(int n, int m, int src, vector<Edge>& edges) {
    const long long INF = 1e18;
    vector<long long> dist(n + 1, INF);

    dist[src] = 0;

    // Step 1: Relax all edges (n-1) times
    for (int i = 1; i <= n - 1; i++) {
        for (auto &e : edges) {
            int u = e.u;
            int v = e.v;
            long long w = e.weight;

            if (dist[u] != INF && dist[v] > dist[u] + w) {
                dist[v] = dist[u] + w;
            }
        }
    }

    // Step 2: Check for negative weight cycles
    for (int i = 1; i <= n - 1; i++) {
        for (auto &e : edges) {
            int u = e.u;
            int v = e.v;
            long long w = e.weight;

            if (dist[u] != INF && dist[v] > dist[u] + w) {
                cout << "Negative weight cycle detected!" << endl;
                return;
            }
        }
    }
}
```

```

    }
}

// Step 2: Check for negative weight cycle
for (auto &e : edges) {
    int u = e.u;
    int v = e.v;
    long long w = e.weight;

    if (dist[u] != INF && dist[v] > dist[u] + w) {
        cout << "Negative weight cycle detected!\n";
        return;
    }
}

// Print distances
for (int i = 1; i <= n; i++) {
    cout << "Distance to " << i << ": " << dist[i] << "\n";
}

```

Edge Structure

```

struct Edge {
    int u, v;
    long long weight;
};

```

4. Time Complexity

Time Complexity: $O(V \times E)$

Because:

- You relax all **E edges**,
- For **(V – 1)** iterations → approximately **$V \times E$** .

Space Complexity: $O(V)$

Only the distance array is stored.