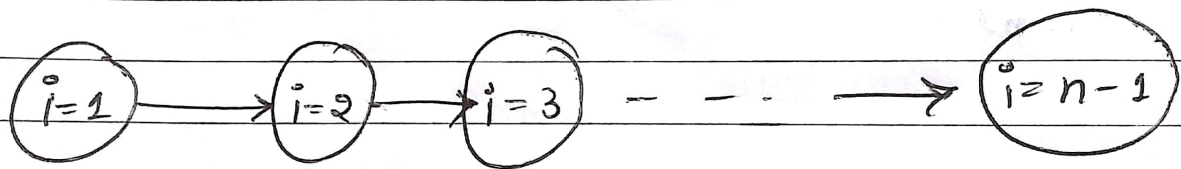


Question 3:



Width = 1

Work = $n-1$

CP = $n-1$.

Question 3.1.2:

template <typename T, typename op>

T reduce(T* array, size_t n)

{ int res = 0;

while (n > 1)

{

int j = 0;

for (int i = 0; i <= n-1; ~~i++~~ ⁱ⁺⁼²)

{

if (i == n-1)

{

arr[j] = op(arr[i], arr[i-1]);

}

else

arr[j] = arr[0];

j++;

}

reduce(*array, n/2);

}

res = arr[0];

return res;

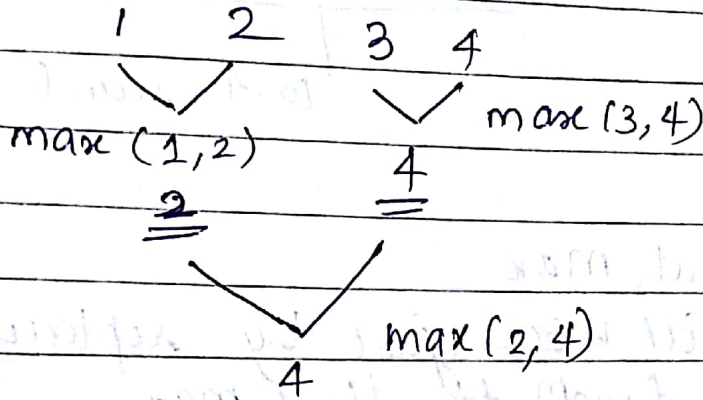
}

3.2

① Yes, the parallel version will work for max too. by substituting $op()$ by $\max()$

For max:

eg:-

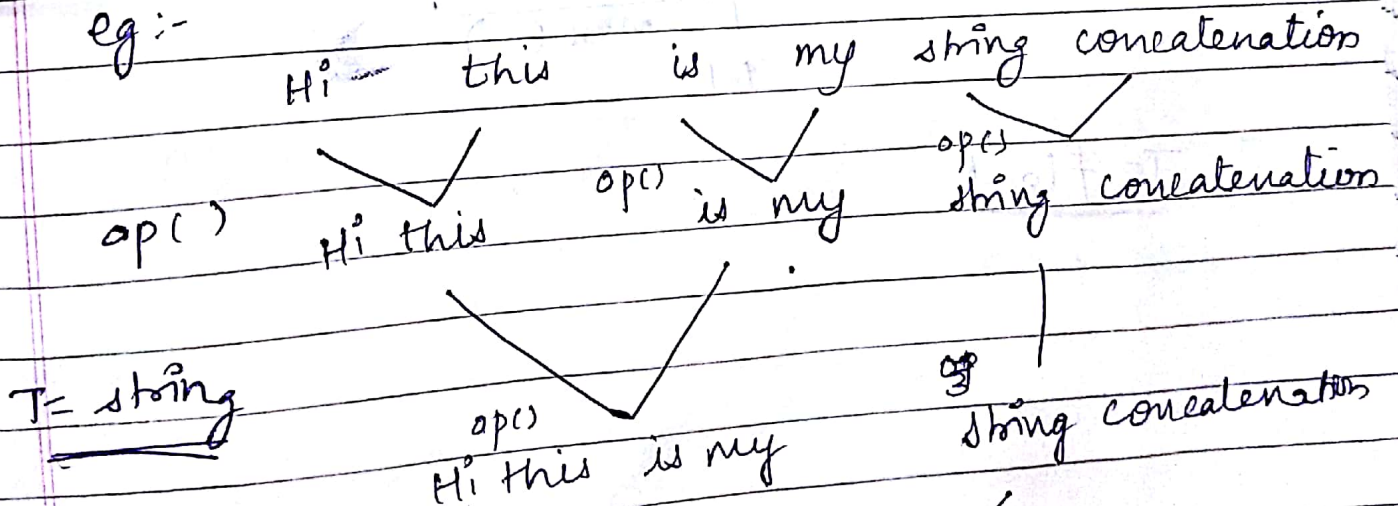


T = int

② for string concat:-

Yes, the code can also be used for achieving parallelism in concatenation. by substituting $op()$ by $\text{string concat}()$

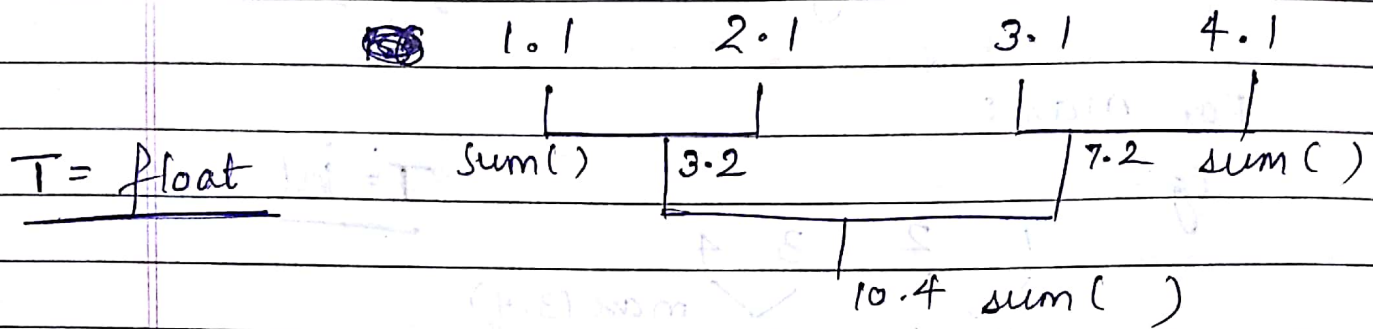
eg:-



T = string

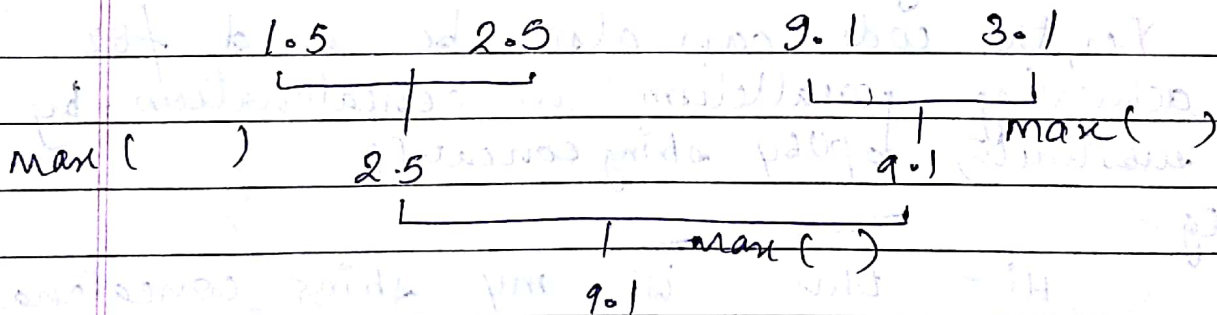
* op here is string concat fn.

Yes, It will work just by replacing op by float, sum
 3) for float, sum. functⁿ.



4) for float, max

Yes, it will work just by replacing op by max functⁿ for float max



T = float

$\Theta: 3.1$

Consider

$j=0$

$i=0$
 $i=1$

$j=1$

$i=2$
 $i=3$

$j=2$

$i=4$
 $i=5$

$j=3$

$i=6$
 $i=7$

j

$n=8$

$j=0$

$i=0$
 $i=1$

$j=1$

$i=2$
 $i=3$

$j=0$

$i=0$
 $i=1$

$width = n/2$

$work = (n-1)$

$CP = n/4$

$$\begin{matrix} i = 1 \\ j = 1 \end{matrix}$$

Question 3 :

