

3.1 Let  $A, B, C$  be objects in  $\mathcal{C}$  and hence objects in  $\mathcal{C}^{op}$ . Let  $f \in \text{Hom}_{\mathcal{C}^{op}}(A, B)$  and  $g \in \text{Hom}_{\mathcal{C}^{op}}(B, C)$ . Thus  $f \in \text{Hom}_{\mathcal{C}}(B, A)$  and  $g \in \text{Hom}_{\mathcal{C}}(C, B)$ . Since  $\mathcal{C}$  is a category, the morphism  $f \circ_{\mathcal{C}} g$  exists in  $\mathcal{C}$ . Define  $\circ_{\mathcal{C}^{op}}$  as  $g \circ_{\mathcal{C}^{op}} f := f \circ_{\mathcal{C}} g$ . We are guaranteed that the composition of morphisms exists since  $\mathcal{C}$  is a category. We gain associativity for the same reason. The composition of identity in  $\mathcal{C}^{op}$  becomes composition of identity on the other side in  $\mathcal{C}$ , which is still an identity.

3.2  $\text{End}_{\text{Set}}(A)$  is defined as  $\text{Hom}_{\text{Set}}(A, A)$ . This is the set of all set-functions  $A \rightarrow A$ ; in other words, the set  $A^A$ . Thus  $|\text{End}_{\text{Set}}(A)| = |A^A| = |A|^{|A|}$  (the last equality from a previous exercise.)

3.3 Let  $f : a \rightarrow b$  be some morphism. Want to show that  $f \circ 1_a =_{(1)} f =_{(2)} 1_b \circ f$ .

(1).  $1_a = (a, a) \implies a \sim a, f = (a, b) \implies a \sim b$ . The transitive property allows us to conclude that  $a \sim b$ , picking that as the composition, so  $1_a \circ f = f$ .

(2).  $f = (a, b) \implies a \sim b, 1_b = (b, b) \implies b \sim b$ . The transitive property gives us that  $a \sim b$ , which corresponds to  $f$ . So  $f \circ 1_b = f$ .

3.5 Example 3.3 describes a category of a set of objections with a reflexive, transitive relation between them.  $\hat{S}$  has a set of objects, the objects being subsets of  $S$ . The reflexive, transitive relation is  $\subseteq$ . In all aspects it takes the form of the abstract category described in Example 3.3.

3.6 Let  $\mathcal{V}$  be a category with  $\text{Obj}_{\mathcal{V}} = \mathbb{N}$ ,  $\text{Hom}_{\mathcal{V}}(n, m)$  being the set of  $m$  by  $n$  matrices. Let  $A : n \rightarrow m, B : m \rightarrow k$  be morphisms. Then define  $B \circ A := BA$ , the matrix product. This product is defined, since  $B$  has height the same as the width of  $A$ .

This composition is associative: consider  $A : n \rightarrow m, B : m \rightarrow k, C : k \rightarrow l$ . Then  $C \circ (B \circ A) = C(BA) = (CB)A = (C \circ B) \circ A$ .

The identity morphism  $1_n$  is the  $n \times n$  identity matrix  $I_n$ . Let  $A : n \rightarrow m$ . We have that  $A \circ 1_n = AI_n = A = I_m A = 1_m \circ A$ .

I suppose this category can be considered with the objects representing the sets of isomorphic finite-dimensional vector spaces of dimension  $n$ .

3.7 Objects in our coslice category are arrows in  $\mathcal{C}^A$  from some fixed object  $A$  to an arbitrary object  $Z$ :

$$\begin{array}{c} A \\ \downarrow f \\ Z \end{array}$$

A morphism  $f_1 \rightarrow f_2$  is a commutative diagram of the following form:

$$\begin{array}{ccc} A & & \\ \downarrow f_1 & \searrow f_2 & \\ Z_1 & \xrightarrow{\tau} & Z_2 \end{array}$$

We obtain composition of morphisms by observing that since the following diagram commutes:

$$\begin{array}{ccccc} & & A & & \\ & f_1 \swarrow & \downarrow f_2 & \searrow f_3 & \\ Z_1 & \xrightarrow{\tau} & Z_2 & \xrightarrow{\sigma} & Z_3 \end{array},$$

this diagram must also commute:

$$\begin{array}{ccc} & A & \\ f_1 \swarrow & & \searrow f_3 \\ Z_1 & \xrightarrow{\sigma\tau} & Z_3 \end{array}.$$