# Analysis of Code Running Time using Big-O Notation

# Part I: Analyzing Code for Running Time

# A.

```
int array_sum(int[] a, int n) {
    int i;
    int sum = 0;
    for (i = 0; i < n; i++) {
        sum = sum + a[i];
    }
    return sum;
}</pre>
```

## Analysis:

- The loop runs from i = 0 to i < n, iterating n times.
- Each iteration involves a constant-time operation.

Time Complexity: O(n) (Linear time)

## В.

#### **Analysis:**

- The first three loops iterate n times each, contributing  $O(n^3)$ .
- The condition i == j == k holds for O(n) cases.
- Inside the condition, the innermost loop runs  $O(n^3)$  times.
- Final complexity:  $O(n^3) \times O(n^3) = O(n^6)$ .

Time Complexity:  $O(n^6)$  (Polynomial time)

# $\mathbf{C}.$

```
int sum(int a, int b) {
    int c = a + b;
    return c;
}
```

#### **Analysis:**

• This function simply performs an addition and return, both of which take constant time.

Time Complexity: O(1) (Constant time)

#### D.

#### **Analysis:**

- The outer loop runs n times.
- The inner loop runs  $O(\log i)$  times for each i.
- Summing up over all i:

$$\sum_{i=1}^{n} O(\log i) = O(n \log n)$$

Time Complexity:  $O(n \log n)$ 

## $\mathbf{E}.$

```
int binarySearch(int[] arr, int target) {
    int left = 0;
    int right = arr.length - 1;

while (left <= right) {
        int mid = left + (right - left) / 2;

        if (arr[mid] == target) {
            return mid;
        } else if (arr[mid] < target) {
                left = mid + 1;
        } else {
                 right = mid - 1;
        }
    }
    return -1;
}</pre>
```

## Analysis:

- Binary search repeatedly halves the search space.
- $\bullet$  Each iteration reduces the problem size by a factor of 2.

Time Complexity:  $O(\log n)$  (Logarithmic time)