

# Announcement

- Assignment 3
  - Due Monday, October 20
  - Simulating an Accumulator Machine – Part II
  - Assignment 2 to be returned this afternoon.

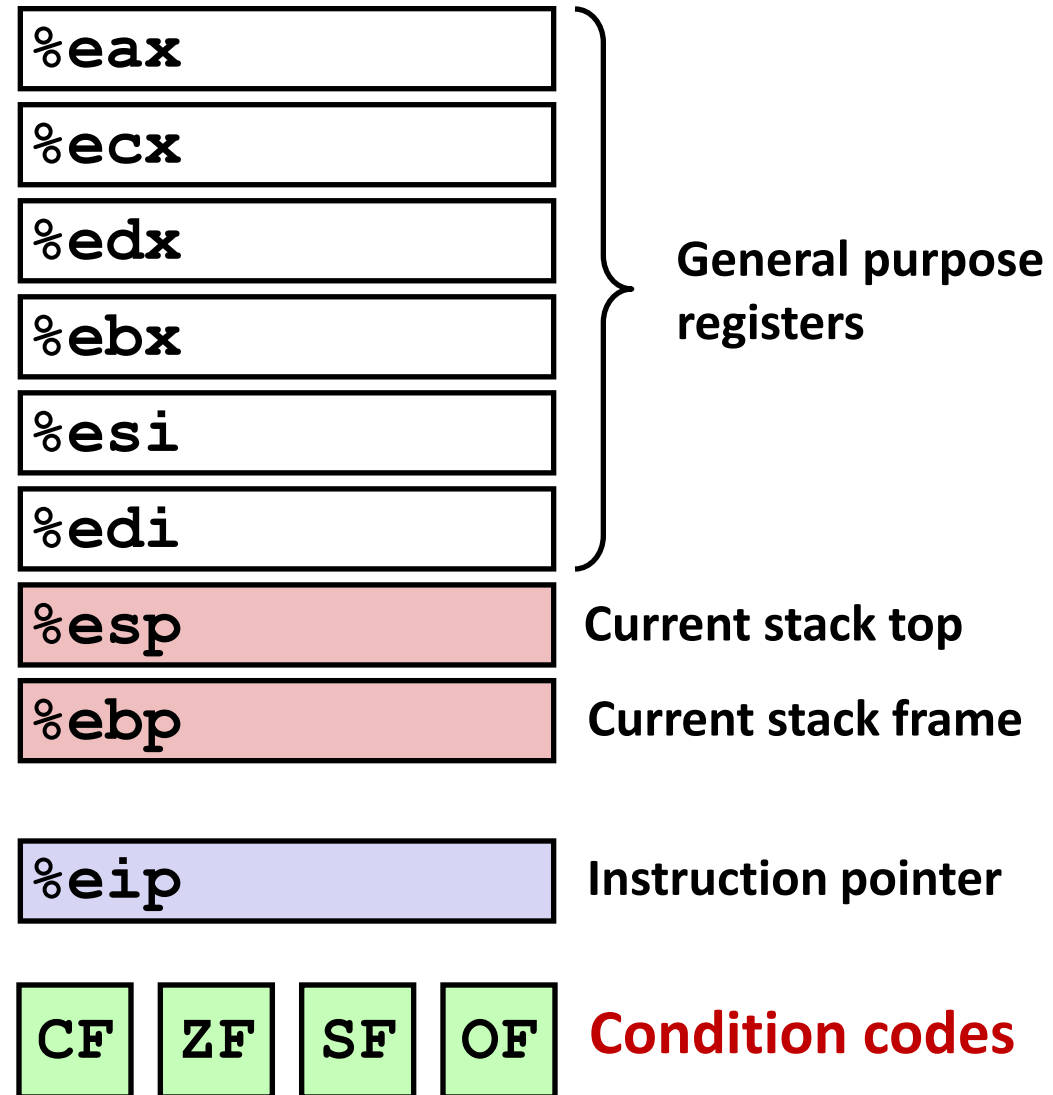
Lecture 17

# Control: Iteration

CPSC 275  
Introduction to Computer Systems

# Processor State

- Information about currently executing program
  - Temporary data ( **%eax**, ... )
  - Location of runtime stack ( **%ebp**, **%esp** )
  - Location of current code control point ( **%eip**, ... )
  - Status of recent tests ( **CF**, **ZF**, **SF**, **OF** )



# Setting Condition Codes

- Implicit setting by arithmetic instructions

```
    addl %ebx,%eax
    je .exit
    ...
.exit:
```

- Explicit setting by compare instruction

```
    cmpl    %eax,%edx
    jle     .L6
    ...
.L6:
```

# Jumping

- $jx$  Instructions
  - Jump to a different part of the code depending on the condition codes

| $jx$             | Description               | Condition                            |
|------------------|---------------------------|--------------------------------------|
| <code>jmp</code> | Unconditional             | 1                                    |
| <code>jz</code>  | Equal / Zero              | $ZF$                                 |
| <code>jnz</code> | Not Equal / Not Zero      | $\sim ZF$                            |
| <code>js</code>  | Negative                  | $SF$                                 |
| <code>jns</code> | Nonnegative               | $\sim SF$                            |
| <code>jg</code>  | Greater (Signed)          | $\sim (SF \wedge OF) \ \& \ \sim ZF$ |
| <code>jge</code> | Greater or Equal (Signed) | $\sim (SF \wedge OF)$                |
| <code>jl</code>  | Less (Signed)             | $(SF \wedge OF)$                     |
| <code>jle</code> | Less or Equal (Signed)    | $(SF \wedge OF) \   \ ZF$            |
| <code>ja</code>  | Above (unsigned)          | $\sim CF \ \& \ \sim ZF$             |
| <code>jb</code>  | Below (unsigned)          | $CF$                                 |

# “Do-While” Loop Example

## C Code

```
int pcount_do(unsigned x) {  
    int result = 0;  
    do {  
        result += x & 0x1;  
        x >>= 1;  
    } while (x);  
    return result;  
}
```

## Goto Version

```
int pcount_do(unsigned x) {  
    int result = 0;  
loop:  
    result += x & 0x1;  
    x >>= 1;  
    if (x)  
        goto loop;  
    return result;  
}
```

What does this function do?

# “Do-While” Loop Compilation

## Goto Version

```
int pcount_do(unsigned x) {  
    int result = 0;  
loop:  
    result += x & 0x1;  
    x >>= 1;  
    if (x)  
        goto loop;  
    return result;  
}
```

|      |        |
|------|--------|
| %edx | x      |
| %eax | result |

```
        movl    $0,%eax        # result = 0  
.L2:    # loop:  
        movl    %edx,%ecx  
        andl    $1,%ecx        # t = x & 1  
        addl    %ecx,%eax      # result += t  
        shrl    %edx           # x >>= 1  
        jne     .L2            # If !0, goto loop
```

# “While” Loop Example

## C Code

```
int pcount_while(unsigned x) {  
    int result = 0;  
    while (x) {  
        result += x & 0x1;  
        x >>= 1;  
    }  
    return result;  
}
```

## Goto Version

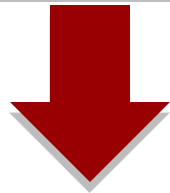
```
int pcount_do(unsigned x) {  
    int result = 0;  
    if (!x) goto done;  
loop:  
    result += x & 0x1;  
    x >>= 1;  
    if (x)  
        goto loop;  
done:  
    return result;  
}
```



# General “While” Translation

## While version

```
while (Test)  
    Body
```



## Do-While Version

```
if (!Test)  
    goto done;  
do  
    Body  
    while (Test) ;  
done:
```



## Goto Version

```
if (!Test)  
    goto done;  
loop:  
    Body  
    if (Test)  
        goto loop;  
done:
```

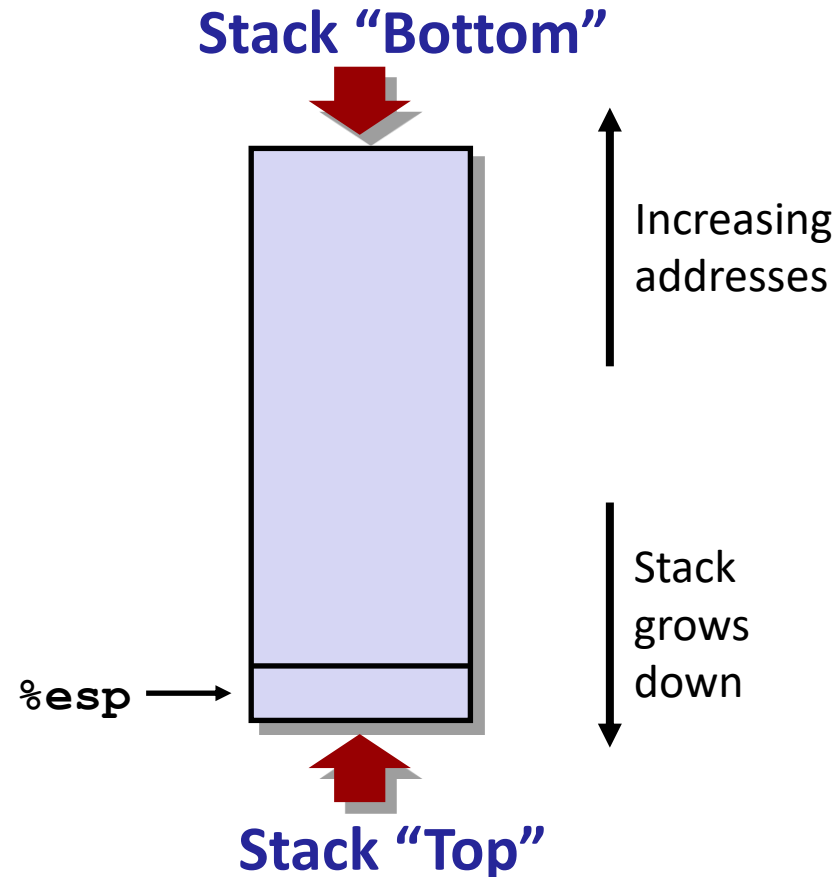


What happens when a function is called?

# Control: Procedures

# IA-32 Stack

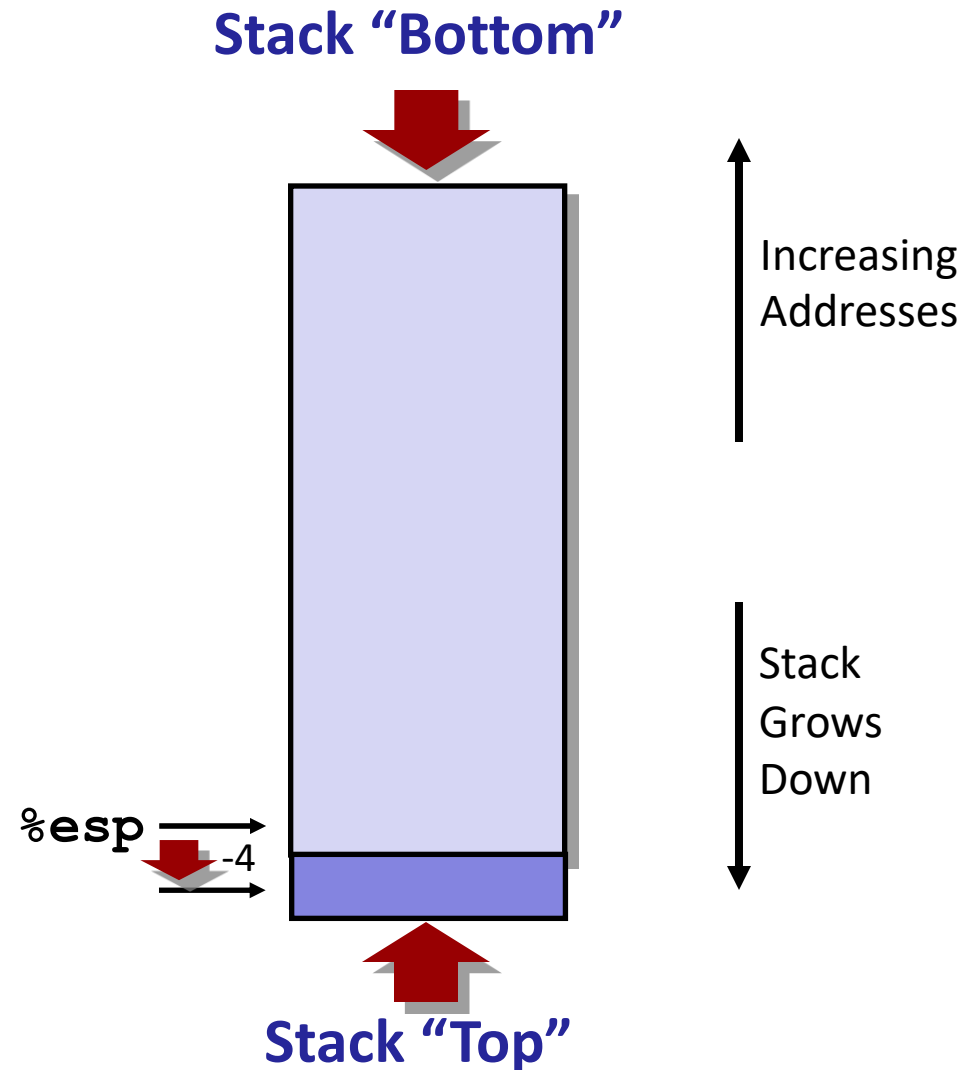
- Grows toward lower addresses
- Register `%esp` contains lowest stack address
  - address of “top” element
- Two operations on stack:
  - **push**: adding a new item on the stack
  - **pop**: removing the top element from the stack



# IA32 Stack: **push**

**pushl src**

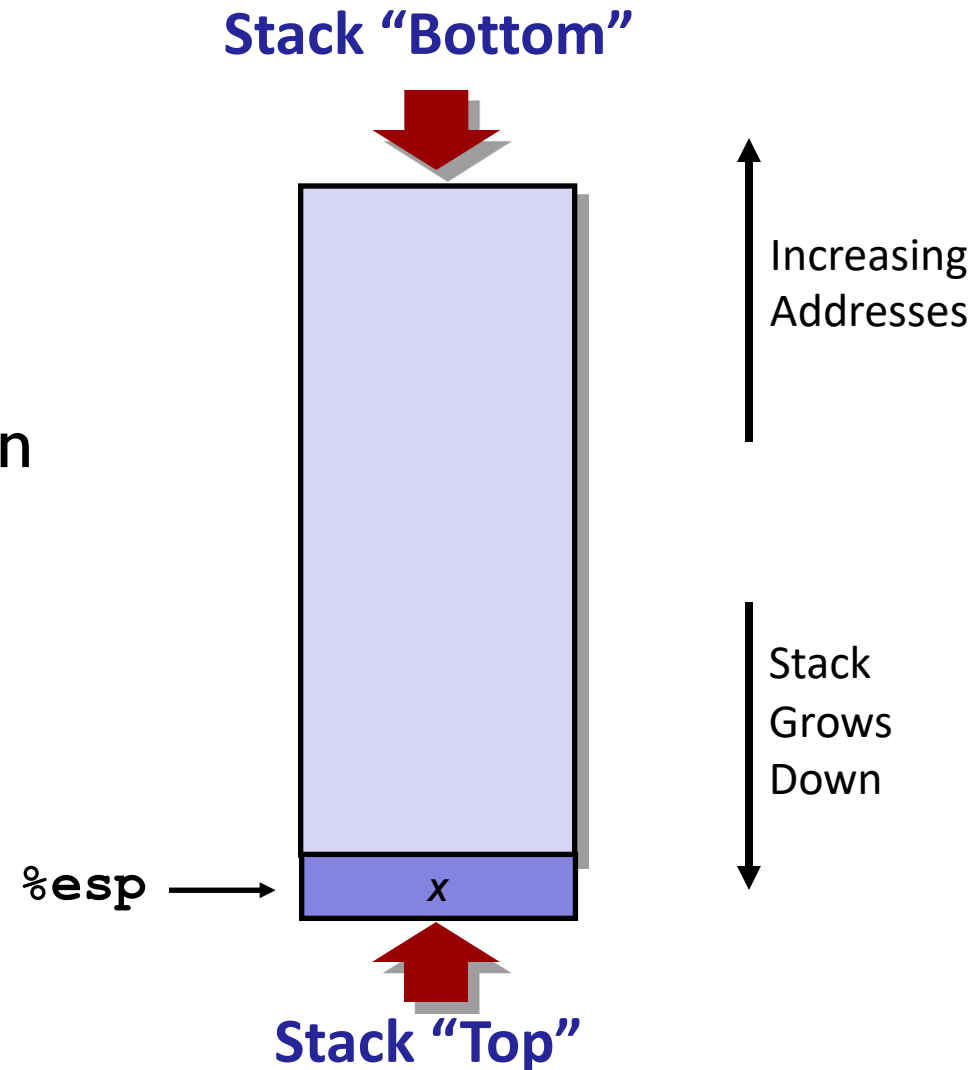
- Fetch operand at *src*
- Decrement **%esp** by 4



# IA32 Stack: **push**

## **pushl src**

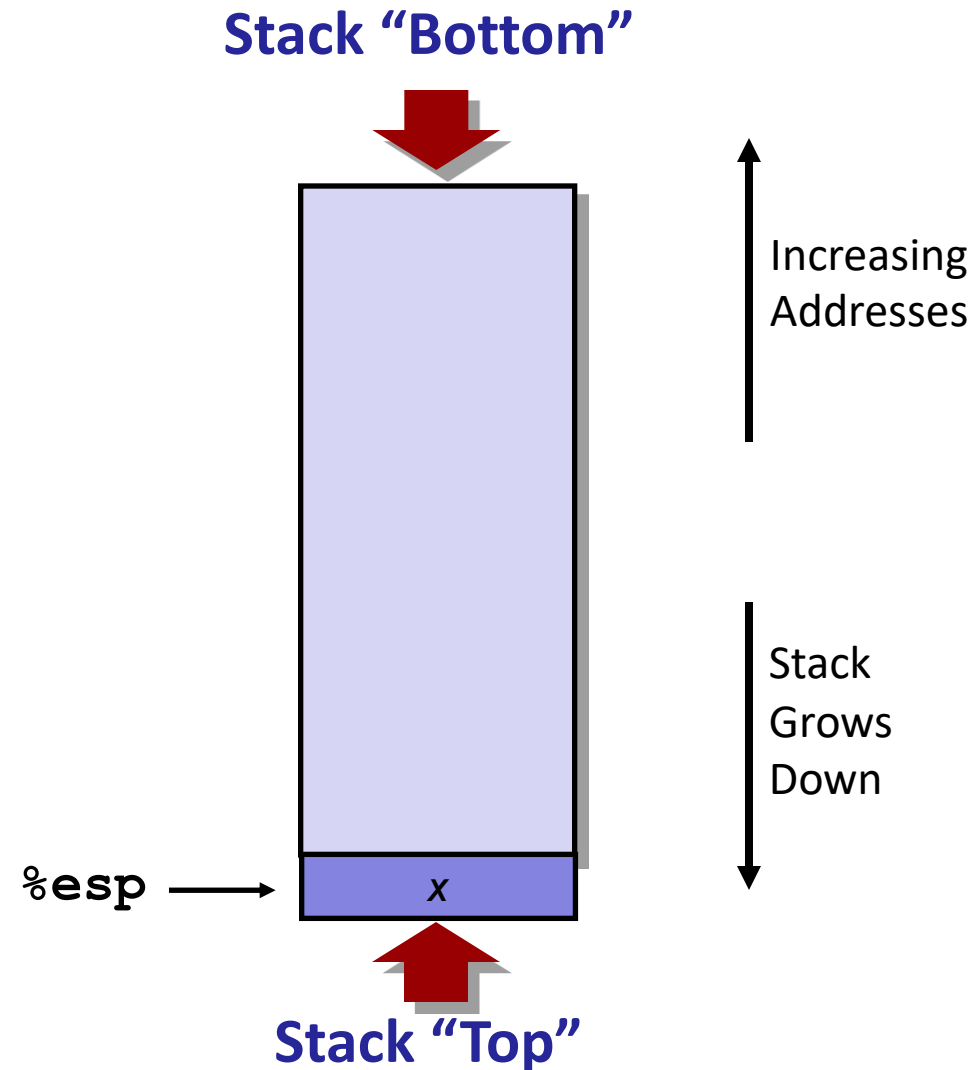
- Fetch operand at *src*
- Decrement **%esp** by 4
- Write operand at address given by **%esp**



# IA32 Stack: **pop**

**pop1** *dest*

- Copy the stack top item to *dest*.

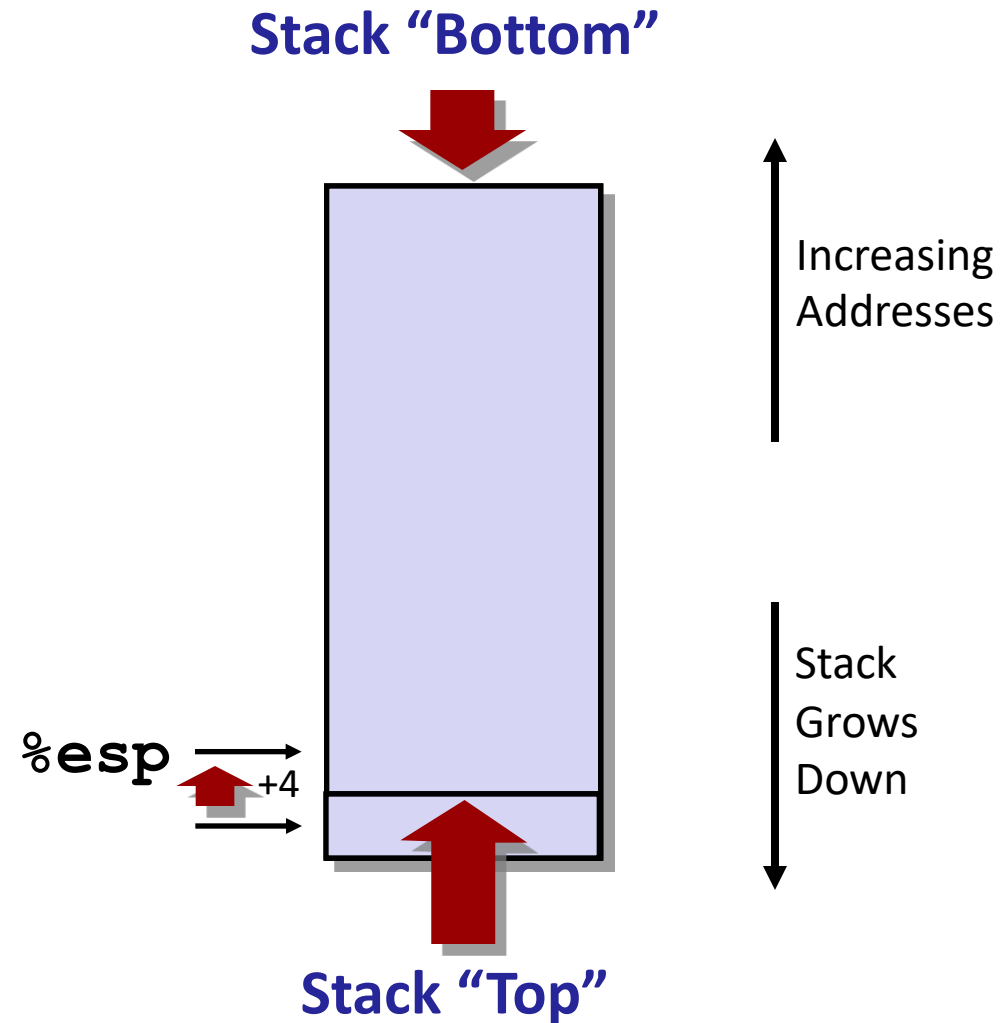




# IA32 Stack: **pop**, cont'd

**pop1** *dest*

- Copy the stack top item to *dest*.
- Increment `%esp` by 4



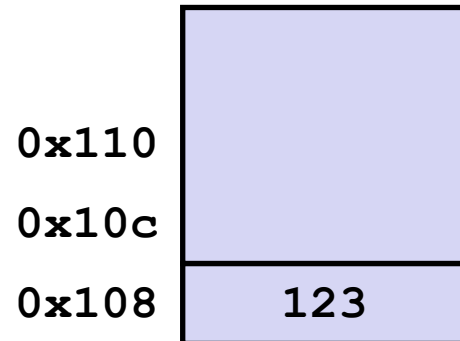
# Procedure Control Flow

- Use stack to support procedure (function) calls and return
- **Procedure call:** `call label`
  - Push *return address* on stack
  - Jump to *label*
- Return address: address of the *next* instruction right after call
- **Procedure return:** `ret`
  - Pop return address from stack
  - Jump to return address

# Procedure Call Example

|          |                |       |               |
|----------|----------------|-------|---------------|
| 804854e: | e8 3d 06 00 00 | call  | 80483c4 <sum> |
| 8048553: | 50             | pushl | %eax          |

before the call



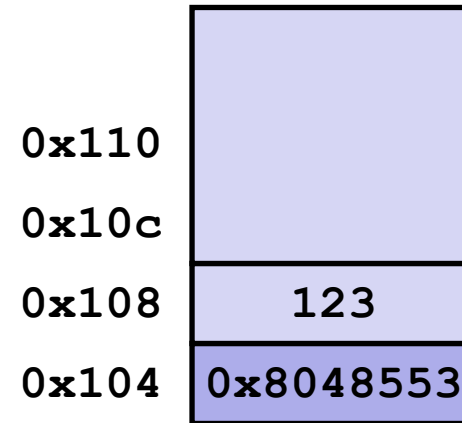
%esp

0x108

%eip

0x804854e

after the call



%esp

0x104

%eip

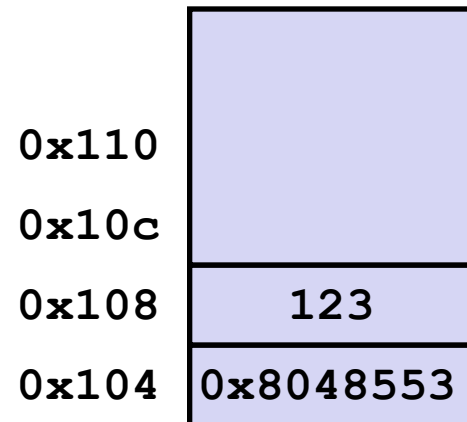
0x80483c4

*%eip: program counter*

# Procedure Return Example

|          |    |     |
|----------|----|-----|
| 80483ce: | c3 | ret |
|----------|----|-----|

before ret



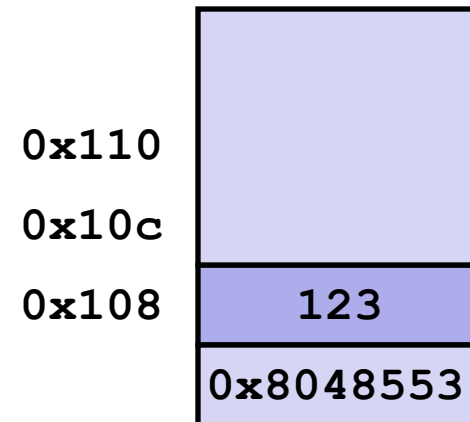
%esp

0x104

%eip

0x80483ce

after ret



%esp

0x108

%eip

0x8048553

*%eip: program counter*

