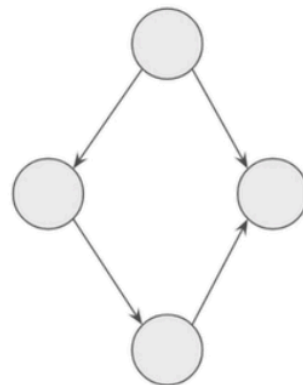
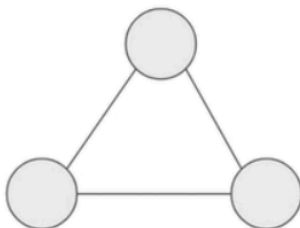


# Graph

类似LinkedList的概念，内存中不一定连续的数据，由各个节点的Reference串起来组成

- 可能有环
- 分为无向图和有向图
- 没有固定入口
- 可能有多个入口

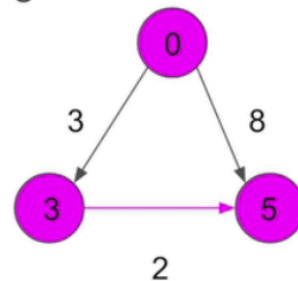
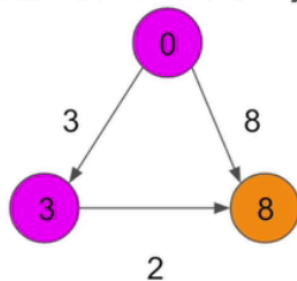
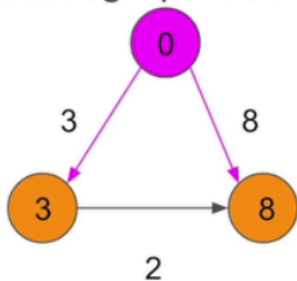


## BFS (Best-First Search)

针对Non-uniform cost graph的一种算法，核心思想是优先展开最“优”的节点

如何每次快速计算最“优” → Heap

最有名的graph中找最短路径的算法又称Dijkstra's Algorithm



# BFS模板

General Steps:

1. Initialize a Heap with all starting points marked with some initial costs, a HashSet to record visited nodes
2. While heap is not empty
  - a. Poll out one node
  - b. If it has already been expanded (visited), skip it
  - c. Otherwise mark the node as visited, update its cost
  - d. If this is the destination node, return
  - e. For all of its neighbors, offer them in to the heap with current node's cost + edge cost

Time:  $O((E + V) \log V)$       Space:  $O(V)$

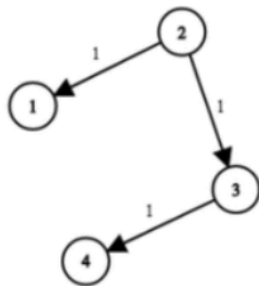
## 743. Network Delay Time

There are  $N$  network nodes, labelled  $1$  to  $N$ .

Given `times`, a list of travel times as **directed** edges `times[i] = (u, v, w)`, where  $u$  is the source node,  $v$  is the target node, and  $w$  is the time it takes for a signal to travel from source to target.

Now, we send a signal from a certain node  $K$ . How long will it take for all nodes to receive the signal? If it is impossible, return  $-1$ .

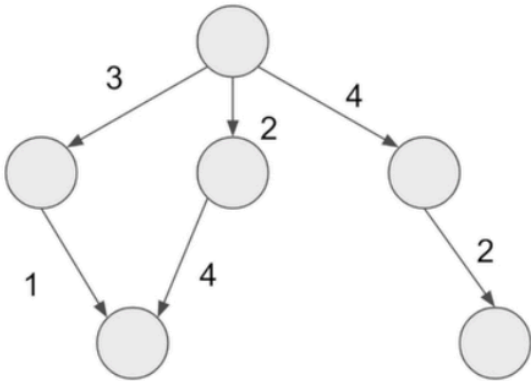
**Example 1:**



Input: `times = [[2,1,1],[2,3,1],[3,4,1]]`,  $N = 4$ ,  $K = 2$   
Output: 2

# 743. Network Delay Time

计算从初始节点到离初始节点最“远”节点的最短路径



## 743. Network Delay Time

建图, Adjacency List Map

1. For each edge (src, dst, cost)
  - a.  $\text{map}[\text{src}] = \{\text{dst}, \text{cost}\}$

```
Map<Integer, List<Cell>> map = new HashMap<>();
for (int[] time : times) {
    List<Cell> edges = map.getDefault(time[0], new ArrayList<>());
    edges.add(new Cell(time[1], time[2]));
    map.put(time[0], edges);
}
```

```
class Cell implements Comparable<Cell> {
    int node, time;
    Cell(int node, int time) {
        this.node = node;
        this.time = time;
    }

    public int compareTo(Cell c2) {
        return time - c2.time;
    }
}
```

## 743. Network Delay Time

1. Initialize Heap with all starting points marked with cost=0, a **HashMap** to record **visited nodes and their costs**
2. While heap is not empty
  - a. Poll out one node
  - b. If it has already been expanded (visited), skip it
  - c. Otherwise mark the node as visited, update its cost
  - d. If this is the destination node, return
  - e. For all of its neighbors, offer them in to the heap with current node's cost + edge cost (**delay time**)

```
public int networkDelayTime(int[][] times, int N, int k) {
    Map<Integer, List<Cell>> map = new HashMap<>();
    for (int[] time : times) {
        List<Cell> edges = map.getOrDefault(time[0], new ArrayList<>());
        edges.add(new Cell(time[1], time[2]));
        map.put(time[0], edges);
    }

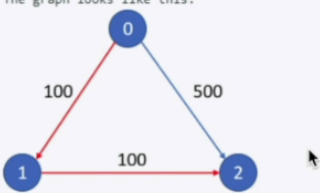
    Map<Integer, Integer> costs = new HashMap<>();
    PriorityQueue<Cell> heap = new PriorityQueue<>();
    heap.offer(new Cell(k, 0));
    while (!heap.isEmpty()) {
        Cell cur = heap.poll();
        if (costs.containsKey(cur.node))
            continue;
        costs.put(cur.node, cur.time);
        if (map.containsKey(cur.node)) {
            for (Cell nei : map.get(cur.node)) {
                if (!costs.containsKey(nei.node)) {
                    heap.offer(new Cell(nei.node, cur.time + nei.time));
                }
            }
        }
    }
    if (costs.size() != N)
        return -1;
    int res = 0;
    for (int x : costs.values())
        res = Math.max(res, x);
    return res;
}
```

# 787. Cheapest Flights Within K Stops

There are  $n$  cities connected by  $m$  flights. Each flight starts from city  $u$  and arrives at  $v$  with a price  $w$ .

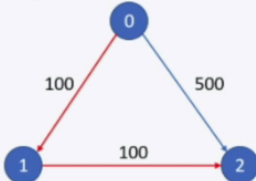
Now given all the cities and flights, together with starting city  $src$  and the destination  $dst$ , your task is to find the cheapest price from  $src$  to  $dst$  with up to  $k$  stops. If there is no such route, output  $-1$ .

**Example 1:**  
**Input:**  
`n = 3, edges = [[0,1,100],[1,2,100],[0,2,500]]`  
`src = 0, dst = 2, k = 1`  
**Output:** 200  
**Explanation:**  
The graph looks like this:



The cheapest price from city 0 to city 2 with at most 1 stop costs 200, as marked red in the picture.

**Example 2:**  
**Input:**  
`n = 3, edges = [[0,1,100],[1,2,100],[0,2,500]]`  
`src = 0, dst = 2, k = 0`  
**Output:** 500  
**Explanation:**  
The graph looks like this:



The cheapest price from city 0 to city 2 with at most 0 stop costs 500, as marked blue in the picture.

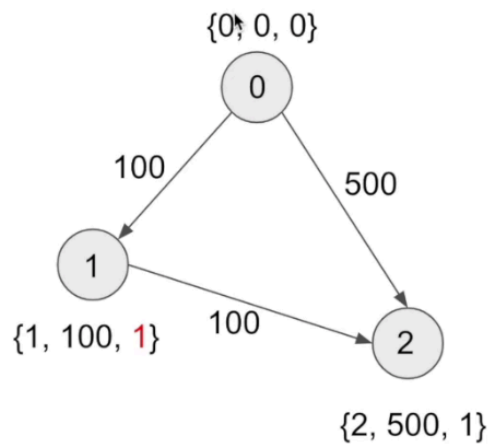
- Constraints:**
- The number of nodes  $n$  will be in range  $[1, 100]$ , with nodes labeled from  $0$  to  $n - 1$ .
  - The size of flights will be in range  $[0, n * (n - 1) / 2]$ .
  - The format of each flight will be  $(src, dst, price)$ .
  - The price of each flight will be in the range  $[1, 10000]$ .
  - $k$  is in the range of  $[0, n - 1]$ .
  - There will not be any duplicated flights or self cycles.



## 787. Cheapest Flights Within K Stops

以往我们Heap中存的节点信息有  $\{\text{node}, \text{cost}\}$ ，这次要多存一个stop的信息记录我们来到当前节点已经经过了多少次转机  $\rightarrow \{\text{node}, \text{cost}, \text{stop}\}$ ，如果一个node的stop数超出K，那就不能展开此node (termination state)

$k = 0$



## 787. Cheapest Flights Within K Stops

建图，Adjacency List Map

1. For each edge (src, dst, cost)
  - a.  $\text{map}[\text{src}] = \{\text{dst}, \text{cost}\}$

```
Map<Integer, List<int[]>> map = new HashMap<>();
for (int[] flight : flights) {
    List<int[]> to = map.getOrDefault(flight[0], new ArrayList<>());
    to.add(new int[] {flight[1], flight[2]});
    map.put(flight[0], to);
}
```

附件.rar

## 727. Cheapest Flights Within K Stops

1. Initialize Heap with starting city marked with cost=0, a HashSet to record visited cities
2. While heap is not empty
  - a. Poll out one node
  - b. If node.stop > k, skip it
  - c. Otherwise mark the node as visited, update its cost
  - d. If this is the destination node, return node.cost
  - e. For all of its neighbors, offer them in to the heap with current node's cost + edge cost (flight cost)

```
public int findCheapestPrice(int n, int[][] flights, int src, int dst, int K) {
    Map<Integer, List<int[]>> map = new HashMap<>();
    for (int[] flight : flights) {
        List<int[]> to = map.getOrDefault(flight[0], new ArrayList<>());
        to.add(new int[] {flight[1], flight[2]});
        map.put(flight[0], to);
    }

    PriorityQueue<Cell> heap = new PriorityQueue<>();
    heap.offer(new Cell(src, K, 0));

    while (!heap.isEmpty()) {
        Cell cur = heap.poll();
        if (cur.dst == dst)
            return cur.price;
        if (cur.stop >= 0 && map.containsKey(cur.dst)) {
            for (int[] next : map.get(cur.dst)) {
                heap.offer(new Cell(next[0], cur.stop - 1, cur.price + next[1]));
            }
        }
    }
    return -1;
}
```

```
class Cell implements Comparable<Cell> {
    int dst, stop, price;
    Cell(int dst, int stop, int price) {
        this.dst = dst;
        this.stop = stop;
        this.price = price;
    }

    public int compareTo(Cell other) {
        return price - other.price;
    }
}
```

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Ugly Number II (264)

Find K Pairs with Smallest Sums (373)

Swim in Rising Water (778)

Kth Smallest Element in a Sorted Matrix (378)