Math 3310 Module on Induction

Dr. Heavilin

Copyright © 2020 Justin Heavilin All Rights Reserved



Contents

2	Ind	uction	11
	2.1	Mathematical Induction	11
	2.2	The Least Number Principle (or <i>The Well Ordering Principle</i>)	15
		2.2.1 A Paper on The Least Number Principle (or <i>The Well Ordering Principle</i>)	17
	2.3	Using the Principle of Mathematical Induction (PMI)	18
		2.3.1 A Paper on Induction	25
		2.3.2 Another Paper on Induction	26

List of Theorems

2.1.2	Theorem (Principle of Mathematical Induction.)	12
2.1.4	Theorem (Principle of Complete $\{Strong\}$ Mathematical Induction.)	12
2.2.1	Theorem (The Least Number Principle)	15
2.3.1	Theorem (Prove 4 divides $(5^n - 1)$)	18
2.3.2	Theorem (Prove 4 divides $(5^n - 1)$)	19
2.3.3	Theorem (For $n \in \mathbb{N}$, $\sum_{i=0}^{n} i = \frac{n(n+1)}{2}$)	19
2.3.5	Theorem (All dogs are the same breed)	19

List of Definitions

2.2.2	Definition (Axiom)																																	-	15
-------	--------------------	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	---	----

List of Problems

2.1.1	Problem (Consider an island populated with only two types of people)	11
2.2.1	Problem (Suppose Q is a property of sets of natural numbers.)	16
2.2.2	Problem (Prove that no finite set can be put into 1-1 correspondence with any of its proper subsets.)	16
2.2.3	Problem (Suppose Vladamir is playing a game with the Countess.)	16
2.3.1	Problem (Any Finite Non-Empty Set of Billiard Balls Must Be The Same Color)	20
2.3.2	Problem	22
2.3.3	Problem	22
2.3.4	Problem	23
2.3.5	Problem	23
2.3.6	Problem	24
2.3.7	Problem	24

Chapter 2

Induction

2.1 Mathematical Induction

No discrete math course would be complete without a discussion on mathematical induction. It is a powerful tool used in proofs throughout mathematics. One of the best presentations of the Principle of Mathematical Induction was presented by Dr. Cangelosi in Math 4200. In this class we will take a less formal approach that I hope will still prove valuable. Let us begin with a reasoning exercise.



Problem 2.1.1 (Consider an island populated with only two types of people)

CONSIDER AN ISLAND POPULATED WITH ONLY TWO TYPES OF PEOPLE, ONE TYPE CAN ONLY TELL THE TRUTH, AND THE	ΗE
OTHER TYPE CAN ONLY LIE. WE WILL CALL THE TRUTH-TELLING PEOPLE Knights and the lying people Knaves. Of	٧E
OF THE INHABITANTS OF THIS ISLAND SAYS, "THIS IS NOT THE FIRST TIME I HAVE SAID WHAT I AM NOW SAYING."	
The question is, it that person a knight or a knave?	
We need to examine the line of reasoning that led you to the answer. It is an example of <i>inductive</i> reasoning. Here is anoth example of inductive reasoning.	.er
Example 2.1.1:	
Suppose the following: if it is raining in Logan on a given day, then it will rain on the following day. Moreover, it is raining today. What can you conclude?	ng

2.1. Mathematical Induction Math 3310

The idea of induction is stated more formally by mathematicians in the following theorem. Let me just note that the symbology $P: \mathbb{N} \to \{T, F\}$ reads something like, "...The property P maps the natural numbers onto the set true and false." Once you learn to read these mathematical hieroglyphics you can amaze people at cocktail parties and that special someone you are interested in impressing will soon be eating out of your hand. Trust me,...happens all the time.

Theorem 2.1.2 (Principle of Mathematical Induction.). If a property holds for the number zero, and <u>whenever</u> it holds for any number n it also holds for n + 1, then it holds for all natural numbers

$$P(0) = T \land (P(k) = T \rightarrow P(k+1) = T) \rightarrow P(n) = T \forall n \in \mathbb{N}.$$

Now we have presented the theorem, we don't expect to immediately understand it, know how to use it, or appreciate its implications. That will come. Instead, to begin I want to spent our time constructively discussing the philosophy so that we can cultivate a relationship to this reasoning. Once we do that, I suspect that all of the other parts will fall into place. Induction is sort of slippery, but at the same time natural - like a bar of oatmeal soap. So let's work through some more examples.

Example 2.1.3:

Suppose you live in the 1950's and the milkman delivers milk on a daily milk run. You decide to leave a note on the doorstep that reads

On any day that you leave milk, be sure to leave milk the following day as well! Thanks.

-Englebert Humperdink.

But after a month you still have no milk. You wait up for the milkman and ask why you don't have any milk. The reply is, "I did exactly what you instructed, I never left any milk and nor did I leave any milk on a day following that day." The milkman is logically correct and you are out of milk. If instead, your note had read,

On any day that you leave milk, be sure to leave milk the following day as well! Oh,...and please leave me milk today. Thanks.

-Englebert Humperdink.

you would have been receiving milk all along.

Strong Induction

There is a variant on induction called *Strong Induction* and the only difference on the surface is that instead of assuming P(0) (that is to say the property holds for n=0), we assume that the property holds for all numbers from zero up to n. Stated similarly to Theorem 2.1.2:

Theorem 2.1.4 (Principle of Complete $\{Strong\}$ Mathematical Induction.). Suppose if a property P of the natural numbers is true for all natural numbers less than n means it is true for P(n), then it also holds for all n.

$$P: \mathbb{N} \to \{T, F\} \quad \land \quad P(0) = T \quad \land \quad (P(k) = T \to P(n) = T \ \forall n < k) \quad \to \quad P(n) = T \ \forall \ n \in \mathbb{N}.$$

Interestingly enough, one can prove Theorem 2.1.2 using 2.1.4 and alternatively could prove Theorem 2.1.4 using 2.1.2. These two theorems are very similar. The preferred use is often context specific; some problems are simply presented in a manner lend to one form or another. I don't know of any *rule* governing the application of these theorems, one over another. However, if the property we are trying to prove involves more that once antecedent in the chain of logic, then usually we resort to strong induction. But you wont build a career from that skill set alone. Let's just have a look at an example.

2.1. Mathematical Induction Math 3310

Example 2.1.5 ($a_1 = 1, a_2 = 8 \land a_n = a_{n-1} + 2a_{n-2}, \ \forall \ n > 3$): Prove that for the sequence

$$a_n = a_{n-1} + 2a_{n-2}, (2.1)$$

with intial conditions $a_1 = 1$, $a_2 = 8$ has a closed form expression

$$a_n = 3 \cdot 2^{n-1} + 2(-1)^n. (2.2)$$

Proof. There are basically two steps to an induction proof, the base case(s) and the induction step. I am going to name Equation 2.2 P(n), i.e.

$$P(n): a_n = 3 \cdot 2^{n-1} + 2(-1)^n.$$

Steps of the proof

Running dialog with the monkeys in my head.

• Base Case.

P(1) and P(2)

Observe through brute force calculations, for both Equations 2.1 and 2.2 indeed $a_1 = 1$ and $a_2 = 8$. These base cases simply get the ball rolling, first to verify that we aren't chasing something completely crazy, but also to satisfy the hypothesis of Theorem 2.1.4.

• Induction Step.

$$a_k = 3 \cdot 2^{k-1} + 2(-1)^k.$$

$$a_{k+1} = a_k + 2a_{k-1}$$

$$a_{k+1} = \left(3 \cdot 2^{k-1} + 2(-1)^k\right) + 2\left(3 \cdot 2^{k-2} + 2(-1)^{k-1}\right)$$

$$a_{k+1} = 3 \cdot \left(2^{k-1} + 2^{k-1}\right) + 2\left((-1)^k + 2(-1)^{k-1}\right)$$

$$a_{k+1} = 3 \cdot 2^k + 2(-1)^{k+1}$$

Suppose this is true for some kThis is the definition of the recurrence.

Using P(n) for n = k and n = k - 1.

Some simplifications.

Algebraic manipulation.

Conclusion.

And we find P(n) for n = k + 1, so by the strong principle of induction, $P(n) \, \forall \, n \in \mathbb{N}$.

A lot could be said about this preceding example. Let's walk through it together. I will attempt to connect the proof to the pieces and parts of Theorem 2.1.2

 $\bullet \ \ P: \mathbb{N} \to \{T, F\} \quad \land \ P(0) = T \quad \land \quad (P(k) = T \to P(k+1) = T) \quad \to \quad P(n) = T \ \forall \ n \in \mathbb{N}.$

This red refers to Eq. 2.1. This statement is either true or false. P(n) is the mapping that takes each number n and decides if the equation holds for it. If so, then T, else F.

• $P: \mathbb{N} \to \{T, F\}$ \wedge P(0) = T \wedge $(P(k) = T \to P(k+1) = T) \to P(n) = T \forall n \in \mathbb{N}.$

Here we are talking about the base case (or cases). For our problem we considered n = 1 and n = 2. If P(n) isn't true for either of these values, then it is game over, because obviously the statetment isn't true for all n. This is the thread hanging off the sweater that, once we tug on it, the whole sweater falls apart.

 $\bullet \ P: \mathbb{N} \to \{T,F\} \quad \wedge \quad P(0) = T \quad \wedge \quad (P(k) = T \to P(k+1) = T) \quad \to \quad P(n) = T \ \forall \ n \in \mathbb{N}.$

Notice the parenthesis surrounding this entire statement. It means that the whole shooting match has to fly. In other words, when there is a k for which P(n) is true, then it must also be true that the NEXT case, P(k+1) must also be true. Think of this as the continuity of the sweater thread we are pulling on; as we pull, the entire thread is connected by the link of $k \to k+1$. In our example it is the induction step where we show algebraically that $a_{k+1} = 3 \cdot 2^k + 2(-1)^{k+1}$ FOLLOWS directly from $a_k = 3 \cdot 2^{k-1} + 2(-1)^k$

• $P: \mathbb{N} \to \{T, F\}$ \land P(0) = T \land $(P(k) = T \to P(k+1) = T) <math>\to$ $P(n) = T \forall n \in \mathbb{N}.$

Finally, having satisfied the conditions (premises) of theorem 2.1.2, we can conclude the the whole sweater is going to end up a ball of thread on the floor.

Fall 2019 Dr. Heavilin Online Version

2.1. Mathematical Induction Math 3310

Here are some examples to show how things fail.

Example 2.1.6 $(a_1 = 1, a_2 = 8 \land a_n = a_{n-1} + 2a_{n-2}, \forall n > 3)$:

Prove that for the sequence

$$a_n = a_{n-1} + 2a_{n-2}, (2.3)$$

with intial conditions $a_1 = 1$, $a_2 = 8$ that a_n has a closed form expression

$$a_n = 3 \cdot 2^{n-1}. (2.4)$$

Let's start this proof the same way we start ALL induction proofs, namely, state the premise, tell the reader you are proving by induction, and then establish some base cases.

Proof. We will attempt to prove by induction that $a_n = a_{n-1} + 2a_{n-2}$, with intial conditions $a_1 = 1, a_2 = 8$ has a closed form solution $P(n): a_n = 3 \cdot 2^{n-1}$. We begin by observing that for n = 3, $P(3) = a_3 = 3 \cdot 2^2 = 12$. However, we know that the recursion relation, $a_n = a_{n-1} + 2a_{n-2}$ yields $a_3 = 8 + 2 \cdot 1 = 9$. Since $P(3) \neq 9$ we have proved that $P(n): a_n = 3 \cdot 2^{n-1}$ is not the closed form solution for the recursion relation $a_n = a_{n-1} + 2a_{n-2}$

Interresting! This is actually a *proof by contradiction*, not really a complete proof by induction. With hindsight, now that we realize the statement is false, we can re-tool our proof in the form of a proof by contradition, and abandon the idea of proof by induction. Essentially, we don't use proof by induction to prove something is false.

Here is another more interesting example of failure.

Example 2.1.7 $(a_1 = 1 \text{ and } a_n = a_{n-1} + n, \forall n)$:

Prove that for the sequence

$$a_{n+1} = a_n + n + 1, (2.5)$$

with intial conditions $a_1 = 1$, a_n has a closed form expression

$$P(n): a_n = \frac{-2n^3 + 13n^2 - 21n + 12}{2}. (2.6)$$

Let us begin, as before, by stating what we want to prove and establishing some base cases.

Proof. We will attempt to prove by induction that the recursion relation, $a_n = a_{n-1} + n$, with intial condition $a_1 = 1$ has a closed form solution $P(n): a_n = \frac{-2n^3 + 13n^2 - 21n + 12}{2}$. We begin by observing that for n = 1, $P(1) = 1 = a_1$, $P(2) = 3 = a_2$, and $P(3) = 6 = a_3$. These are the base cases for our induction proof. Now that we have established that the relationship holds for at least "some" $n \in \mathbb{N}$, then next part of Theorem 2.1.2 begs us to suppose for some k, that if $P(k) = \frac{-2n^3 + 13n^2 - 21n + 12}{2}$ is true, then it is also true for the next k, ie, k + 1. Let's follow through the same algebraic form as in a previous example.

$$a_k = \frac{-2\,k^3 + 13\,k^2 - 21k + 12}{2}$$
 Suppose this is true for some k
$$a_{k+1} = a_k + k + 1$$
 Consdier $k+1$
$$a_{k+1} = \frac{-2\,k^3 + 13\,k^2 - 21k + 12}{2} + k + 1$$
 Employ our assumption from two lines above.
$$a_{k+1} = \frac{-2\,k^3 + 13\,k^2 - 19k + 14}{2}$$
 Do some algebra.

In order for this proof to be complete, we would expect that this result would be equivalent to Eq 2.6 when n = k + 1,

$$P(k+1): a_{k+1} = \frac{-2(k+1)^3 + 13(k+1)^2 - 21(k+1) + 12}{2}$$

$$= \frac{-2k^3 + 7k^2 - k + 2}{2}$$
Simplify
$$= \frac{-2k^3 + 13k^2 - 19k + 14}{2} + \frac{-5k^2 + 18k - 12}{2}$$
A despirate attempt to make things look that same.

However, there is just no way that we can legitamately arrange the terms to be what we need them to be. I have seen student go though obscene contortions of algebra, violating all the algebra they learned in high school just to get the answer to look like the back of the book. But the rational conclusion is that we simply can not satisfy the premise to Theorem 2.1.2, and as a result must conclude that the closed form solution is not true.

2.2 The Least Number Principle (or The Well Ordering Principle)

We continue in a chain of ideas that are strongly related to one another. I was never presented with both of these ideas hand-in-hand. But they are related and it is helpful to see them side by side, because one implies the other. In fact, each is used to prove that the other is true. It sounds like a house of card, but by assuming 2.1.2 as an axiom, we can prove 2.2.1. Alternatively, by assuming 2.2.1 we can prove 2.1.2.

Theorem 2.2.1 (The Least Number Principle). Every non-empty set of natural numbers contains a least element.

Proof. I realize it is perhaps considered *bad form* to present a proof in a COLOUR BY NUMBERS style. I do not mean to insult your intelligence, but for a person like me, I find is can be helpful to break something down into little bite-sized pieces so that I make sure I understand what is happening at each step along the way. So here goes!

The proof is by contradiction.

Steps of the proof

• Assume $S \neq \emptyset$, $S \subseteq \mathbb{N}$, and S has no least element.

• Let $P(n): n \in \mathbb{N} - S$.

• $\forall m, n \in \mathbb{N}, m < n : P(m).$

• $\forall n \in \mathbb{N}, P(n)$

• $S = \emptyset$

ullet S must contain a least element.

Running dialog with the monkeys in my head.

Suppose S is an arbitrary **non-empty** subset of N, such that S has no least element. Between you and me, we know this doesn't exist, but let's prove it doesn't exist.

The conventional venacular for this type of proof uses a notation requiring a little getting used to. The symbols P(n) are read to mean the property is true for n. So were we write something like, P(5), we mean "P(5) is true". For our purposes, the proposition is that 5 is not in S, or rather, $S \in \mathbb{N} - S$.

This is saying that if we grab another number, say m, then it can't possibly be in S, because if it were in the set S and if m < n, then the set S would have a lesser number than n. By assumption, S has no such element. So m can not be in S, or to use our notation, P(m) is true.

We have determined in the previous bullet that nothing smaller than n could be in S, so P(m) must be true for all m < n. Then by strong induction, our property P of the natural numbers is true for all natural numbers less than n, then it also holds for n, otherwise, n would be the least element of the set S.

Our original set must be empty, since everything is in $\mathbb{N} - S$.

To begin we supposed that S was non-empty, butt now come to find that only the empty set can have no least element. Therefore by contradiction, every non-empty set must have a least element.

Theorem 2.2.1 is also referred to as the Well Ordering Principle, but I think the words "least number" do a better job of describing the idea, so I will use those words. (Why do mathematicians go out of their way to make stuff cryptic?) If you stop and think about this, Theorem 2.2.1 makes senes. But here is the rub. Often there is a catch to proving something that makes total sense. For example, proving that 1+1=2 would probably take several class lectures....but I digress. Here is a notion that is as important as proving any of Theorems 2.1.2, 2.1.4, or 2.2.1. To prove any of them, you have to assume one of them as an axiom.

Definition 2.2.2 (Axiom):

Wikipedia states: An axiom (or postulate) is a statement that is taken to be true, to serve as a premise or starting point for further reasoning and arguments.

Remark 2.2.3:

Frankly, this is about as good as any other definition. Practically speaking, before you begin building a mathematics, you need to create a framework (using axioms and definitions). More to the point,...what is a *point*? Think about it for a minute. Has anyone proven to you that points exist? Mathematicians come up with a set of axioms so that they can start describing things, discussing things, proving things. Once a conjecture is proved, it becomes a theorem. One theorem leads to another and pretty soon you have a mathematical structure.

Problem 2.2.1 (Suppose Q is a property of sets of natural numbers.)

Suppose Q is a property of sets of natural numbers. And suppose Q satisfies the following

- $\bullet \ Q$ holds for the empty set.
- For any finite set A and any number $n \notin A$, then if Q holds for A, then Q holds for $A \cup \{n\}$.

Then Q holds for all finite sets of natural numbers.	
<i>Proof.</i> Prove that Q holds for all finite sets of natural numbers.	
Problem 2.2.2 (Prove that no finite set can be put into 1-1 correspondence with any of its proper subsetence of the proper subsets. (HINT: PROPER SUBSETS.) BY INDUCTING ON THE NUMBER OF ELEMENTS IN THE SET)	
Problem 2.2.3 (Suppose Vladamir is playing a game with the Countess.) Suppose Vladamir is playing a game with the Countess (these lovebirds are always playing games gether). He gives her a bag containing an infinite supply of ping pong balls used for BINGO (each has a positive number printed on it). For each number n there are infinitely many of those balls (bered n). Then Vladamir places a box with infinite capacity on the table. This box contains similar be as in the bag, however there are only finitely many, instead of infinitely many as in the bag. Each the Countess must remove one ball from the box, and replace it with a <i>finite</i> number of balls from bag - but those balls must be of a lower number. For example, she can pull out a 51 from the box replace it with a thousand 48's. If the box on the table gets emptied, the Countess has to drink it water and stand in the sun. Otherwise, Vladamir has to eat raw garlic. Is there a strategy by which the Countess can win? If so, how.	BALL NUM- BALLS DAY, I THE AND

2.2.1 A Paper on The Least Number Principle (or The Well Ordering Principle)

This is a fun little paper with some example problems (Theorems and Proofs). There is very little theory contained in this paper; its strength lies in the presentation of the Well Ordering Principle as an alternative approach to problems most students are told <u>must</u> be solved by induction (I am a big fan of subverting the dominant paradigm). I feel the more versitile your constructs are, the stronger thinker you will become. I love the idea of solving problems multiple ways, and in the case of *The Well Ordering Principle* and *The Principle of Mathematical Induction* we discover that they are distinct but related approaches. The former is a more set-theoretic perspective, while the latter is more number-theoretic in nature. At least that is just my interpretation, take it or leave it.[2]



92.30 The well-ordering principle for **N**Author(s): DES MacHALE
Source: The Mathematical Gazette, Vol. 92, No. 524 (July 2008), pp. 257-259
Published by: The Mathematical Association
Stable URL: https://www.jstor.org/stable/27821774
Accessed: 26-01-2020 23:19 UTC

JSTOR is a not-for-profit service that helps scholars, researchers, and students discover, use, and build upon a wide range of content in a trusted digital archive. We use information technology and tools to increase productivity and facilitate new forms of scholarship. For more information about JSTOR, please contact support@jstor.org.

Your use of the JSTOR archive indicates your acceptance of the Terms & Conditions of Use, available at https://about.jstor.org/terms



This content downloaded from 129.123.124.117 on Sun, 26 Jan 2020 23:19:02 UTC All use subject to https://about.jstor.org/terms

2.3 Using the Principle of Mathematical Induction (PMI)

Theroem 2.1.2 is used to prove statements of the form $(\forall n \in \mathbb{N} : P(n))$ where P(n) is some statement. Our goal is to prove that the set of all n where P(n) is true (Call this \mathcal{T}) is \mathbb{N} . To verify the hypothesis of the PMI, we must do two things:

- 1. Prove that $1 \in \mathcal{T}$ (i.e. show that P(1) is true)
- 2. Prove that if $k \in \mathcal{T}$, then $k+1 \in \mathcal{T}$ (i.e show that if P(k) is true, then P(k+1) is true.)

As review, step 1 is called the *base case* and step 2 is called the *inductive step*, which includes an induction hypothesis. The recursive nature of a proof by induction, i.e., that the truth of a proposition P(k+1) must be demonstrated from the truth of P(k), may seem confusing. However, many familiar functions on the natural numbers such as f(x) = x! or f(x) = 2x can be defined recursively. There is a nice little two-page article at the end of these notes that might be helpful.

When writing a proof by mathematical induction, we should remember to always keep the reader informed. This means that at the beginning of the proof, we should state that a "proof by induction" will be used, and early on in the proof we should define P(n). Remember,

"Don't do unto others what you would not want done unto you."

- Some long haired, bearded, sandal wearing, desert nomadic socialist.

Consdier the following examples. To prove the theorem that 4 divides $(5^n - 1)^n$ we need a base case, and the inductive step. The base case can come from exploring the statement for a few values of n. The inductive step (also called the inductive hypothesis) is assuming the k-th case, i.e. 4 divides $(5^k - 1)^n$ and exploring the relationship for k + 1. In this proof, the statement, P(n), is "4 divides $(5^n - 1)$ ".

Before we go any further, I advise you to grab a pencil and follow along with this problem on a separate piece of paper. Doing mathematics is like juggling machetes, in that it looks simple when I watch someone else do it, but when I try it on your own, I am glad I have health insurance.

Theorem 2.3.1 (Prove 4 divides $(5^n - 1)$). $\forall n \in \mathbb{N}$, let P(n) be "4 divides $(5^n - 1)$ ".

Proof by Mathematical Induction. To begin, observe that when n = 1, $(5^1 - 1) = 4$ and 4 divides 4. For the inductive step, we prove that $\forall k \in \mathbb{N}$, if P(k) is true, then P(k+1) is also true. So let $k \in \mathbb{N}$ and assume that P(k) is true, i.e assume that 4 divides $(5^k - 1)$. From the definition of divisibility, this means there exists an integer, m, such that $4m = 5^k - 1$, or $5^k = 4m + 1$. In order to show that P(k+1) is true, we must show that 4 divides $5^{k+1} - 1$.

$$5^{k+1} - 1 = 5 \cdot 5^k - 1$$
 Since $5^{k+1} = 5^k \cdot 5$.

$$= 5(4m+1) - 1$$
 Above we determined $5^k = 4m+1$.

$$= 20m+5-1$$
 Distribute.

$$= 20m+4$$
 Simplify.

$$= 4(5m+1)$$
 Factor 4.

We see that $(5m+1) \in \mathbb{Z}$. Therefore 4 divides the quantity $5^{k+1} - 1 = 4(5m+1)$. We have thus proved that P(k+1) is true, and by The Principle of Mathematical Induction, we have proved the proposition true for all $n \in \mathbb{N}$.

Notice how the structure of this proof differs slightly from our first example.							

Did you ever notice that all of your Calculus homework problems were, in and of themselves, **Theorems** - and solving the problem constituted a *proof*. I wonder if students understood that, they might have less anxiety about proofs in general.

Theorem 2.3.2 (Prove 4 divides $(5^n - 1)$). Proof by Mathematical Induction. $\forall n \in \mathbb{N}$, let P(n) be "4 divides $(5^n - 1)$ ". To begin, observe that when n = 1, $(5^1 - 1) = 4$ and 4 divides 4.

For the inductive step, we prove that $\forall k \in \mathbb{N}$, if P(k) is true, then P(k+1) is also true. So let $k \in \mathbb{N}$ and assume that P(k) is true, i.e assume that 4 divides $(5^k - 1)$. From the definition of divisibility, this means there exists an integer, m, such that $4m = 5^k - 1$, or $5^k = 4m + 1$.

In order to show that P(k+1) is true, we must show that 4 divides $5^{k+1}-1$.

$$5^{k+1}-1=5\cdot 5^k-1$$
 Since $5^{k+1}=5^k\cdot 5$.

$$=5\left(4m+1\right)-1$$
 Above we determined $5^k=4m+1$.

$$=20m+5-1$$
 Distribute.

$$=20m+4$$
 Simplify.

$$=4\left(5m+1\right)$$
 Factor 4.

We see that $(5m+1) \in \mathbb{Z}$. Therefore 4 divides the quantity $5^{k+1} - 1 = 4(5m+1)$. We have thus proved that P(k+1) is true, and by The Principle of Mathematical Induction, we have proved the proposition true for all $n \in \mathbb{N}$.

Theorem 2.3.3 (For $n \in \mathbb{N}$, $\sum_{i=0}^{n} i = \frac{n(n+1)}{2}$). Proof by Mathematical Induction. We will argue by way of induction that

if n = 1 and 2, then the relationship, P(n) is $\sum_{i=0}^{n} i = \frac{n(n+1)}{2}$ is true for all $n \in \mathbb{N}$. Consider the base case(s).

$$P(1) = 0 + 1 = \frac{1(1+1)}{2} = 1; \quad P(2) = 0 + 1 + 2 = \frac{2(2+1)}{2} = 3.$$

Now assume P(k) holds for some $k \in \mathbb{N}$. We will show this relationship implies $\sum_{i=0}^{k+1} i = \frac{(k+1)(k+2)}{2}$.

$$\sum_{i=0}^{k+1} i = \sum_{i=0}^{k} i + (k+1)$$
 Split off the last term in the summation.
$$= \frac{k(k+1)}{2} + (k+1)$$
 Induction hypothesis.
$$= \frac{(k+1)(k+2) + 2(k+1)}{2}$$
 Algebraic simplification.
$$= \frac{(k+1)(k+2)}{2}$$

The first equality follows from the induction assumption, and the others from algebraic manipulations. The relationship is true for n = 1, and if P(k) is true, then P(k+1) is true. So the relationship is true for all nonnegative integers by the Principle of Mathematical Induction.

Remark 2.3.4:

It is considered *bad form*, and usually marked incorrect in math classes, if a mathematician only provides a string of mathematical expressions under the impression that they speak for themselves and are sufficient to qualify as a proof. Do not just write a bunch of mathematical symbols down without sufficient expository narration to support, explain, and guide the reader through your thought process. It might even be helpful to assume I have the IQ of Cool Whip, and you need to hold my hand through the work you do so that I don't get lost.

Something to remember. If a proposition is false, the proposed proof is of course incorrect. For example:

Theorem 2.3.5 (All dogs are the same breed). *Proof.* We will prove this proposition using mathematical induction. For each natural number n, we let P(n) be "Any set of n dogs consists entirely of dogs of the same breed."

T $P(n)$ be the property that for every set of n Billiard balls, they are all the same color. Consider e set of one ball. This set is non-empty, and obviously the ball(s) are of one color, since there is ly one ball. So $P(1)$ is true. E now show that if $P(n)$ holds, then also $P(n+1)$ holds. So suppose n is a number for which P holds. We consider any set of $n+1$ balls, numbering them from 1 to $n+1$. The balls one through n all have e same color (by our assumption that said sets with n balls all have the same color). Likewise the lls gathered from 2 to $n+1$ are also a set of n balls, consequently this set has all the same colored lls. Thus all the balls are of the same color, proving that $P(n+1)$ holds if $P(n)$ holds. So by Theorem 1.2, the property holds for all billiard balls.	
what is wrong with this Theorem:	
THE SET OF ONE BALL. THIS SET IS NON-EMPTY, AND OBVIOUSLY THE BALL(S) ARE OF ONE COLOR, SINCE THERE ONLY ONE BALL. SO $P(1)$ IS TRUE. WE NOW SHOW THAT IF $P(n)$ HOLDS, THEN ALSO $P(n+1)$ HOLDS. SO SUPPOSE n IS A NUMBER FOR WHICH P HOLDS NOW CONSIDER ANY SET OF $n+1$ BALLS, NUMBERING THEM FROM 1 TO $n+1$. THE BALLS ONE THROUGH n ALL HATTHE SAME COLOR (BY OUR ASSUMPTION THAT SAID SETS WITH n BALLS ALL HAVE THE SAME COLOR). LIKEWISE THE BALLS GATHERED FROM 2 TO $n+1$ ARE ALSO A SET OF n BALLS, CONSEQUENTLY THIS SET HAS ALL THE SAME COLOR	E IS LDS. AVE THE RED

The following is an example of how you should construct your induction proofs. It is not the only way, but it is better than most.

Example 2.3.6:

Prove by induction that, for all $n \in \mathbb{Z}_+$, $\sum_{i=1}^n (-1)^i i^2 = \frac{(-1)^n n(n+1)}{2}$.

Proof. We will prove by induction that, for all $n \in \mathbb{Z}_+$,

$$\sum_{i=1}^{n} (-1)^{i} i^{2} = \frac{(-1)^{n} n(n+1)}{2}.$$
(2.7)

When n = 1, the left side of Eq (2.7) is (-1)12 = -1, and the right side is $\frac{(-1)1(1+1)}{2} = -1$, so both sides are equal and Eq (2.7) is true for n = 1. Also, when n = 2, the left side of Eq (2.7) is $(-1)1^2 + (-1)^2 +$

$$\sum_{i=1}^{k+1} (-1)^{i} i^{2} = \sum_{i=1}^{k} (-1)^{i} i^{2} + (-1)^{k+1} (k+1) 2$$
 Write the last term separately.
$$= \frac{(-1)^{k} k (k+1)}{2} (-1)^{k+1} (k+1) 2$$
 By the induction hypothesis.
$$= \frac{(-1)^{k} (k+1)}{2} (k-2(k+1))$$
 Algebraic manipulation.
$$= \frac{(-1)^{k} (k+1)}{2} (-k-2)$$

$$= \frac{(-1)^{k+1} (k+2)}{2}$$

Thus, if Eq (2.7) holds for n = k, then it also holds for n = k + 1, and the proof of the induction step is complete. By the principle of induction, Eq (2.7) is true for all $n \in \mathbb{Z}_+$.

Tove the following by induction (I lease type one solution per page to help with grading.).	Prove the following by induction	(Please type one solution pe	er page to help with grading.):
---	----------------------------------	------------------------------	---------------------------------

Problem 2.3.2

$\sum_{i=1}^{n} \frac{1}{i(i+1)} = \frac{n}{n+1}, \forall n \ge 1.$		

Problem 2.3.3

$\sum_{i=1}^{n} (2i - 1) = n^2, \forall n \ge 1.$		

Duoblom	•	9	/

Problem 2.3.4 $\prod_{i=2}^{n} \left(1 - \frac{1}{i^2}\right) = \frac{n+1}{2n}, \forall n \ge 2.$	
Problem 2.3.5 $n! > 2^n, \forall n \ge 4.$	

Problem 2.3.6	
A SET OF n ELEMENTS HAS 2^n SUBSETS.	
Problem 2.3.7	
$\overline{A_1 \cup \cdots \cup A_n} = \overline{A_1} \cap \cdots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \cdots \cup A_n} = \overline{A_1} \cap \cdots \cap \overline{A_n}$ HINT: USE DEMORGAN'S LAW, I.E. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \cdots \cup A_n} = \overline{A_1} \cap \cdots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
Problem 2.3.7 $\overline{A_1 \cup \cdots \cup A_n} = \overline{A_1} \cap \cdots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	
$\overline{A_1 \cup \dots \cup A_n} = \overline{A_1} \cap \dots \cap \overline{A_n}$ Hint: Use DeMorgan's Law, i.e. $\overline{(A \cup B)} = \overline{A} \cap \overline{B}$	

2.3.1 A Paper on Induction

There are so many papers that introduce induction, and I don't know if this is the best - but it is one of the shortest. So let's start here, and see how people "talk" about induction in mathematics. There is almost a *formularic* approach to some problems, but at the same time, induction can be tricky. There have been cases where a "proof" by induction turned out later to actually be false. Don't look for a formula, look for a way of thinking![1]





Mathematical Induction: If Student k Understands It, Will Student k + 1? Author(s): Judith L. Gersting Source: The Two-Year College Mathematics Journal, Vol. 6, No. 2 (May, 1975), pp. 18-19 Published by: Taylor & Francis, Ltd. on behalf of the Mathematical Association of American

Source: The Two-Year College Mathematics Journal, Vol. 6, No. 2 (May, 1975), pp. 18-19
Published by: Taylor & Francis, Ltd. on behalf of the Mathematical Association of America
Stable URL: https://www.jstor.org/stable/3026460
Accessed: 01-01-2019 18:33 UTC

JSTOR is a not-for-profit service that helps scholars, researchers, and students discover, use, and build upon a wide range of content in a trusted digital archive. We use information technology and tools to increase productivity and facilitate new forms of scholarship. For more information about JSTOR, please contact support@jstor.org.

Your use of the JSTOR archive indicates your acceptance of the Terms & Conditions of Use, available at https://about.jstor.org/terms



Mathematical Association of America, Taylor & Francis, Ltd. are collaborating with JSTOR to digitize, preserve and extend access to The Two-Year College Mathematics Journal

This content downloaded from 129.123.124.117 on Tue, 01 Jan 2019 18:33:16 UTC All use subject to https://about.jstor.org/terms

2.3.2 Another Paper on Induction

Ok,...one more.[3]





Mathematical Induction for Freshmen Author(s): Richard Morris
Source: National Mathematics Magazine, Vol. 12, No. 4 (Jan., 1938), pp. 183-187
Published by: Taylor & Francis, Ltd. on behalf of the Mathematical Association of America Stable URL: https://www.jstor.org/stable/3028083
Accessed: 03-09-2019 15:16 UTC

JSTOR is a not-for-profit service that helps scholars, researchers, and students discover, use, and build upon a wide range of content in a trusted digital archive. We use information technology and tools to increase productivity and facilitate new forms of scholarship. For more information about JSTOR, please contact support@jstor.org.

Your use of the JSTOR archive indicates your acceptance of the Terms & Conditions of Use, available at https://about.jstor.org/terms



Mathematical Association of America, Taylor & Francis, Ltd. are collaborating with JSTOR to digitize, preserve and extend access to National Mathematics Magazine

This content downloaded from 129.123.124.117 on Tue, 03 Sep 2019 15:16:07 UTC All use subject to https://about.jstor.org/terms

BIBLIOGRAPHY Math 3310

Bibliography

[1] Gersting Judith L. Mathematical induction: If student k understands it, will student k + 1?. The Two-Year College Mathematics Journal, 6(2):18, 1975.

- [2] Des MacHale. The well-ordering principle for n. The Mathematical Gazette, 92(524):257–259, 2008.
- [3] Richard Morris. Mathematical induction for freshmen. National Mathematics Magazine, 12(4):183–187, 1938.

Fall 2020 Dr. Heavilin Online Version