```
recursive procedure DFS(v \in V):
     Close(v)
     if v is uncovered then
          if M_v(v) = l_f then
               REFINE(v);
               for all w \sqsubseteq v : Close(w)
          EXPAND(v);
          for all children w of v: DFS(w)
procedure UNWIND:
     set V \leftarrow \{\epsilon\}, E \leftarrow \emptyset, \psi(\epsilon) \leftarrow \text{True}, \triangleright \leftarrow \emptyset
     while there exists an uncovered leaf v \in V:
          for all w \in V s.t. w \sqsubset v: Close(w);
          DFS(v)
```

Fig. 5. DFS unwinding strategy

for all $w \in V$ s.t. $w \prec v$ and $M_v(w) = M_v(v)$:

procedure $Close(v \in V)$:

COVER(v, w)