```
recursive procedure DFS(v \in V):
     Close(v)
     if v is uncovered then
         if M_v(v) = l_f then
              Refine(v):
              for all w \sqsubseteq v : Close(w)
          Expand(v);
         for all children w of v: DFS(w)
procedure Unwind:
     set V \leftarrow \{\epsilon\}, E \leftarrow \emptyset, \psi(\epsilon) \leftarrow \text{True}, \triangleright \leftarrow \emptyset
     while there exists an uncovered leaf v \in V:
         for all w \in V s.t. w \sqsubset v: CLOSE(w);
         DFS(v)
                                        Fig. 5. DFS unwinding strategy
```

for all $w \in V$ s.t. $w \prec v$ and $M_v(w) = M_v(v)$:

procedure $Close(v \in V)$:

Cover(v, w)