

Question 5:

$$(a) \varphi_X = (v_1 \vee v_2) \wedge (v_1 \vee v_3) \wedge (v_2 \vee v_3) \wedge (v_4 \vee v_5) \wedge (v_4 \vee v_7) \wedge (v_5 \vee v_6) \wedge (v_6 \vee v_7) \wedge (\neg v_1 \vee \neg v_2 \vee \neg v_3)$$

(b) Base Case:

cycle of length 3 (forms triangle)

vertex cover clauses: 3 clauses (1 for each edge)

triangle free clauses: 1 clause (to prevent triangle)

Total clauses = 3 + 1 = 4 clauses.

So for  $n = 3$ :

$$\frac{4(3)}{3} = 4$$

$\therefore$  base case holds

Inductive Step:

Assume for any 2-regular graph with  $k$  vertices, the boolean expression  $\varphi_G$  has at most  $\frac{4k}{3}$  clauses.  $(k \leq n)$

Need to prove that for a 2-regular graph with  $k+1$  vertices the boolean expression  $\varphi_G$  has at most  $\frac{4(k+1)}{3}$  clauses

2-regular graph  $G$  with  $n$  vertices. To form a 2-regular graph  $G'$  with  $n+1$  vertices, adding new vertex 'v' and connect to existing vertices  $u$  and  $w$  (ensuring remains 2-R).

- replaces one existing edge  $(u \rightarrow w)$  with 2 new edges  $(u \rightarrow v)$  and  $(v \rightarrow w)$
- addition of 'v' and transformation of edge can add at most 1 new clause to  $\varphi_G$ .
- bc new edges are connected to only 2 existing vertices.

which means:

$$\frac{4k}{3} + 1 \leq \frac{4(k+1)}{3}$$

$$\frac{4k+3}{3} \leq \frac{4k+4}{3}$$

since  $\frac{4k+3}{3}$  is less, this holds  $\therefore$  by induction, it is seen that the boolean expression  $\varphi_G$  for any 2-regular graph  $G$  with  $n$  vertices has at most  $\frac{4n}{3}$  clauses.

