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**Q1**

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```
/* Algorithm to recursively compute a+b, where a and b are non-neg integers */
add( int a, int b )
    if( b == 0 )
        return a;
    return add( a+1, b-1 );

/* Test case */
a = 3, b = 2
add( 3, 2 )
    return add( 4, 1 )
        return add( 5, 0 )
            return 5
= 5
```

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**Q2**

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We need to find the sum of the array by adding each item recursively, and then divide by the number of elements in the array ( $N = \text{array.length}$ ). There will be three cases, as follows:

1.  $i == 0$
2.  $i == N-1$  (base case)
3. all other values of  $i$

Thus, the Java/pseudocode will be as follows:

```
/* Java/pseudocode for average of array */
averageArray( int[] A, int i, int N )
    if( i == 0 ) // Case 1 - this case returns the average
        return ( A[i] + averageArray( A, i+1, N ) ) / N;
    else if( i == N-1 ) // Case 2 (base case)
        return A[i];
    else // Case 3
        return A[i] + averageArray( A, i+1, N );
```

Then one could simply call `averageArray(A, 0, A.length)` with the initial call having  $i = 0$  as a requirement for proper calculation of average.

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**Q3**

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There are three cases here as follows:

1. Case 1 (base case 1)  $\rightarrow n == 0$ , return `f0`
2. Case 2 (base case 2)  $\rightarrow n == 1$ , return `f1`

3. Case 3  $\rightarrow n > 1$ , return  $\text{gfib}(f0, f1, n-1) + \text{gfib}(f0, f1, n-2)$

The Java/pseudocode is as follows:

```
/* Java/pseudocode for generalized fibonacci */
gfib( int f0, int f1, int n )
    if( n == 0 )
        return f0;
    else if( n == 1 )
        return f1;
    else
        return gfib(f0, f1, n-1) + gfib(f0, f1, n-2);
```

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## Q4

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For this question I wrote and executed a Java program to compute Ackerman(2,2). The code is as follows:

```
public class Ackerman
{
    public static void main(String[] args)
    {
        System.out.println("Ackerman(2,2) = " + ack(2,2));
    }

    public static int ack(int m, int n)
    {
        if( m == 0 )
        {
            return n + 1;
        }
        else if( n == 0 )
        {
            return ack(m-1, 1);
        }
        else
        {
            return ack(m-1, ack(m, n-1));
        }
    }
}
```

The result of this is:

```
C:\Users\Sean\Documents\jhu-cs\Data Structures\Module 3>java Ackerman
Ackerman(2,2) = 7
```

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## Q5

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For this problem, it is simple to convert the recursive function to an iterative function using a

while loop. The Java/pseudocode for this is as follows:

```
int rec(int n)
    bool x = f(n);
    int a = n;
    while( x == FALSE )
        /* any group of statements that do not change the value of n */
        a = g(a);
        x = f(a);
    return 0;
```

In this way,  $f(n)$  is initially calculated using  $n$ . By setting a variable 'a' equal to  $n$ ,  $n$  remains unchanged and  $a$  can be updated to equal  $g(n)$  and so on recursively. Finally, when  $x = f(a)$  returns TRUE, the method returns 0.

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## Q6

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There are four obvious ways to implement the queue. They are...

- Sorted array
- Sorted linked list
- Unsorted array
- Unsorted linked list

For my ADT, I chose to use an unsorted array, with methods modified to account for the additional requirements (namely searching for highest priority item in array and shifting items for deletion).

### ADT PriorityQueue

#### Data

An empty unsorted array of values with a reference to the first (front) item, which is highest priority. The queue stores information on two parts of an item - the data and the priority

#### Methods

##### isEmpty

Input None

Precondition None

Process Check **if** the queue contains any data items

Postcondition None

Output Return **true** if queue is empty, and **false** otherwise

##### Insert

Input A data item to be stored in the queue

Precondition Item has an assigned priority

Process Store an item at the rear of a queue

Postcondition The queue contains one additional data item

Output None

##### Delete

Input None

Precondition Queue contains meaningful data values

```
    Process    Remove the highest priority element, accomplished by searching the
               queue
    Postcondition Shift all elements greater than index of removed element left by
               one element
    Output    Return deleted value
Peek
    Input     None
    Precondition Queue contains meaningful data values
    Process   Search queue for highest priority item
    Postcondition None
    Output    Return the value of the data item at the front of the queue (highest
               priority)
end ADT PriorityQueue
```