

Throughout this article we denote by $H = \langle -1, 2 \rangle$ the subgroup generated by -1 and 2 in \mathbb{F}_p^\times and set

$$\ell_0 = [\mathbb{F}_p^\times : H]$$

for the index of H in \mathbb{F}_p^\times . Notice that if $4 \nmid o_p(2)$ then by definition we have that $\mathcal{O}(p) = \ell_0$. Furthermore, Conjecture [A](#) is a non-empty statement if and only if $\mathcal{O}(p) = \ell_0 \geq 3$.

The idea behind Conjecture [A](#) is the following. Suppose that Conjecture [A](#) holds, then each triple (a_i, b_i, c_i) in the conjecture corresponds to a nonequi-difference codeword $\mathbf{x}_i = \{0, a_i, -c_i\}$ with difference set $\Delta(\mathbf{x}_i) = \{\pm a_i, \pm b_i, \pm c_i\}$. Hence, we have $\left\lfloor \frac{\mathcal{O}(p)}{3} \right\rfloor$ nonequi-difference codewords whose difference sets are disjoint. From the complement of $\cup_{i=1} \Delta(\mathbf{x}_i)$ in \mathbb{F}_p^\times , their algorithm then produces $\frac{p-1-2\mathcal{O}(p)}{4}$ equi-difference codewords and hence gives a CAC of size matching the upper bound given in [\(1\)](#).

As an illustration, we briefly discuss the case treated in [\[FLS14, Example 3\]](#) where the length $p = 31$. Note that $o_{31}(2) = 5$ and hence $\mathcal{O}(31) = 3$. Then Conjecture [A](#) predicts that there are $3 \left\lfloor \frac{\mathcal{O}(31)}{3} \right\rfloor = 3$ cosets and one element in each coset such that their sum is zero. One finds that the triple $(2, 3, -5)$ gives a solution and the corresponding codeword is $\{0, 2, 5\}$ whose difference set is just $\{\pm 2, \pm 3, \pm 5\}$ while $2, 3$ and -5 lie exactly in three distinct cosets of H in \mathbb{F}_p^\times . Moreover, there are six equi-difference codewords $\{0, 4, 8\}$, $\{0, 6, 12\}$, $\{0, 7, 14\}$, $\{0, 9, 18\}$, $\{0, 10, 20\}$ and $\{0, 15, 30\}$ produced by their algorithm. In total, one concludes that the size of an optimal CAC of length 31 is $M(31) = 7$.

Independently, in [\[MZS14\]](#) the authors proposed a conjecture which provides solutions to the existence of the triples (A_i, B_i, C_i) in Conjecture [A](#) in terms of the group structure of \mathbb{F}_p^\times/H .

Conjecture B ([\[MZS14, Conjecture\]](#)). Let p be an odd prime. If $\ell_0 \geq 3$, then there exist $b \in gH$ and $c \in g^2H$ such that

$$1 + b + c = 0 \quad \text{in } \mathbb{F}_p$$

for some generator g of \mathbb{F}_p^\times .

Remark 1.1. We see that Conjecture [B](#) implies Conjecture [A](#) by setting $A_1 = H$, $B_1 = gH$, $C_1 = g^2H$, $A_2 = g^3H$, $B_2 = g^4H$, $C_2 = g^5H, \dots, A_e = g^{3e-3}H$, $B_e = g^{3e-2}H$ and $C_e = g^{3e-1}H$ where $e = \left\lfloor \frac{\ell_0}{3} \right\rfloor$. Moreover, Conjecture [B](#) does not assume that $4 \nmid o_p(2)$.

Note that the subgroup $H = \langle -1, 2 \rangle$ consists of all the ℓ_0 -th power of elements of \mathbb{F}_p^\times . It follows that the elements b and c in Conjecture [B](#) are of the forms gy^{ℓ_0} and $g^2x^{\ell_0}$ respectively for some $x, y \in \mathbb{F}_p^\times$. Observe that if $\ell_0 \geq 3$ then any \mathbb{F}_p -rational solutions (x, y) to the the *diagonal equation* $g^2X^{\ell_0} + gY^{\ell_0} + 1 = 0$ must satisfy $xy \neq 0$ since $-1 \in H$ and g is a generator of \mathbb{F}_p^\times . Thus, any \mathbb{F}_p -rational solution gives a pair of elements b and c in Conjecture [B](#). So Conjecture [B](#) is equivalent to the following statement.