Algorithm 2 (Discretized MMGA) takes not only its learning rate η but a small value γ in measuring an approximate gradient as inputs. In each time step, each player measures the gradients of its payoff for each variable of its strategy (lines 2-6). Then, the player updates its strategy by the gradients (lines 7-10). Here, note that the strategy update is weighted by the probability $x^{a|s}$ (line 8) in order to correspond to Algorithm 1. Here, each of lines 3-5 and line 8 can be updated in parallel with respect to a and s.

4 Theoretical Analysis

4.1 Continuous-Time Equivalence of Algorithms

The following theorems provide a unified understanding of different algorithms. Theorem 1 and 2 are concerned with continualization of the two discrete algorithms. Surprisingly, Theorem 3 proves the correspondence between these different continualized algorithms by Theorem 1 and 2.

Theorem 1 (Coutinualized MMRD). Let $p^{a|s}$ be the expected distribution when X chooses a under state s;

$$p_{i'}^{a|s} := \begin{cases} y^{b|s} & (s_{i'} = abs^{-}) \\ 0 & (\text{otherwise}) \end{cases} . \tag{7}$$

In the limit of $\eta \to 0$, Algorithm 1 is continualized as dynamics

$$\dot{x}^{a|s_i}(\mathbf{x}, \mathbf{y}) = p_i^{\text{st}} x^{a|s_i} \left(\pi(\boldsymbol{p}^{a|s_i}, \mathbf{x}, \mathbf{y}) - \bar{\pi}^{s_i}(\mathbf{x}, \mathbf{y}) \right),$$
(8)

$$\bar{\pi}^{s_i}(\mathbf{x}, \mathbf{y}) = \sum_a x^{a|s_i} \pi(\mathbf{p}^{a|s_i}, \mathbf{x}, \mathbf{y}), \tag{9}$$

for all $a \in A$ and $s \in S$. Here, $\bar{\pi}^{s_i}$ is the expected payoff under state s_i .

Theorem 2 (Continualized MMGA). In the limit of $\gamma \to 0$ and $\eta \to 0$, Algorithm 2 is continualized as dynamics

$$\dot{x}^{a|s}(\mathbf{x}, \mathbf{y}) = x^{a|s} \frac{\partial}{\partial x^{a|s}} u^{\text{st}}(\text{Norm}(\mathbf{x}), \mathbf{y}), \tag{10}$$

for all $a \in \mathcal{A}$ and $s \in \mathcal{S}$.

See Appendix A.1 and A.2 for the proof of Theorems 1 and 2.

Theorem 3 (Equivalence between the algorithms). The dynamics Eqs. (8) and (10) are equivalent.

Proof Sketch. Let \mathbf{x}' be the strategy given by $x^{a|s} \leftarrow x^{a|s} + \gamma$ in \mathbf{x} for $a \in \mathcal{A}$ and $s \in \mathcal{S}$. Then, we consider the changes of the Markov transition matrix $d\mathbf{M} := \mathbf{M}(\operatorname{Norm}(\mathbf{x}'), \mathbf{y}) - \mathbf{M}(\mathbf{x}, \mathbf{y})$ and the stationary distribution $d\mathbf{p}^{\operatorname{st}} := \mathbf{p}^{\operatorname{st}}(\operatorname{Norm}(\mathbf{x}'), \mathbf{y}) - \mathbf{p}^{\operatorname{st}}(\mathbf{x}, \mathbf{y})$. By considering this changes in the stationary condition $\mathbf{p}^{\operatorname{st}} = \mathbf{M}\mathbf{p}^{\operatorname{st}}$, we get $d\mathbf{p}^{\operatorname{st}} = (\mathbf{E} - \mathbf{M})^{-1} d\mathbf{M}\mathbf{p}^{\operatorname{st}}$ in $O(\gamma)$. The right-hand (resp. left-hand) side of this equation corresponds to the continualized MMRD (resp. MMGA). See Appendix A.3 for the full proof.

For games with a general number of actions, the study [7] has proposed a gradient ascent algorithm in relation to replicator dynamics. In light of this study, Theorem 3 extends the relation to the multi-memory games. This extension is neither simple nor trivial. The relation between replicator dynamics and gradient ascent has been proved by directly calculating $u^{\text{st}} = p^{\text{st}} \cdot u$ [17]. In multi-memory games, however, $u^{\text{st}} = p^{\text{st}} \cdot u$ is too hard to calculate. Thus, as seen in the proof sketch, we proved the relation by considering a slight change in the stationary condition $p^{\text{st}} = Mp^{\text{st}}$, technically avoiding such a hard direct calculation.

4.2 Learning Dynamics Near Nash Equilibrium

Below, let us discuss the learning dynamics in multi-memory games, especially divergence from Nash equilibrium in zero-sum payoff matrices. In order to obtain a phenomenological insight into the learning dynamics simply, we assume one-memory two-action zero-sum games in Assumption 1.