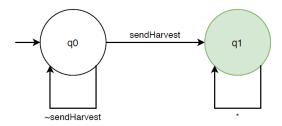
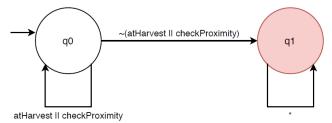
## Monitor

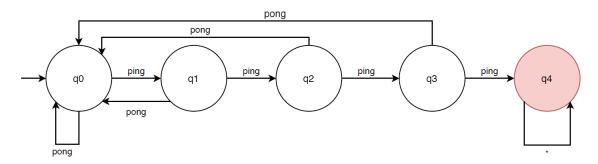
- 1. The Scout should at some point in time send a harvesting position to the Collector.
  - LTL: F(sendHarvest)
  - Not Safety: There is no bad prexif because for every infinite trace that is not in L we can always append {sendHarvest}. This would then be a good prefix.  $\emptyset^{\omega}$  does not have a bad prefix
  - Co-Safety: Every trace in L has sendHarvest at some point which is a good prefix.
  - Monitorable: Every finite trace can be extended to a good prefix by appending {sendHarvest} so we do not have an ugly prefix.
  - Automaton:



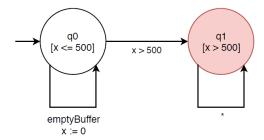
- 2. The Collector should always check its proximity sensors unless it is at a decent harvesting position.
  - LTL: G(atHarvest ∨ checkingProximity)
  - Safety: Every trace not in L has at least one prefix {¬atHarvest ⇒ ¬checkingProximity} which is a bad prefix.
  - Not Co-Safety: {atHarvest, checkProximity}<sup>ω</sup> does not have a good prefix.
    Every trace that is in L does not have a good prefix because we can always do {¬atHarvest ∧ ¬checkingProximity} in the next step.
  - Monitorable: Every finite trace can be extended to a bad prefix so it does not have an ugly prefix.
  - Automaton:



- 3. A robot should never ignore three consecutive PING messages.
  - LTL:  $G((ping \land X(\neg pong \land ping) \land XX(\neg pong \land ping)) \rightarrow XXX pong)$
  - Safety: Every trace not in L has at least one prefix x which is a bad prefix.
  - Not Co-Safety:  $\{pong\}^{\omega}$  does not have a good prefix. We can extend every good trace to a bad trace by appending  $\{ping, ping, ping, ping\}$ .
  - Monitorable: Every finite trace can be extended with the finite trace {ping, ping, ping, ping} to a bad prefix, therefore we do not have an ugly prefix.
  - Automaton:



- 4. When receiving a message, the robot copies the data over into a buffer. This buffer should be emptied at least every 500ms.
  - LTL: This specification cannot be expressed in LTL but only in TLTL.
  - TLTL:  $G(\triangleright_{\text{emptyBuffer}} \in [0, 500])$ . The time until emptyBuffer occurs lies within the interval [0, 500].
  - Safety: Every trace not in L has at least one prefix where the buffer is not emptied for more than 500ms, which is a bad prefix.
  - Not Co-Safety:  $\{(\text{emptyBuffer}, t)\}^{\omega}$  does not have a good prefix for all  $t \in \mathbb{N}$  with  $0 \le t \le 500$ . We can extend every such good trace to a bad trace by appending  $\{(\text{emptyBuffer}, t + 501)\}$ .
  - Monitorable: Every finite trace be extended to a bad prefix, therefore we do not have an ugly prefix.
  - Automaton:



## **SPI State Machine**

