

For $K1$

LDF only for sync. Tasks! Explain Modifications.

Alternative: In each timestep exponential calculations to determine schedule

$C2$

For RM, check if utilization $< \ln(2) \approx 0,69$ If yes, we know that schedulable, otherwise not sure.

If RM is deterministic, what it is for \wedge it is sufficient to show how RM fails as proof.

Alternative proof for this:

$$\left. \begin{array}{l} \text{Time when} \\ \text{to start task} \\ \text{Task finished} \end{array} \right\} \begin{array}{l} R_i^0 = C_i \\ R_i^{i+1} = C_i + \sum_{h=1}^{i-1} \left\lceil \frac{\pi_i^h}{T_h} \right\rceil C_h \end{array}$$

$$R_1^0 = 1 \quad R_2^0 = 1 \quad R_3^0 = 2 \quad R_4^0 = 1$$

$$R_1^1 = 1 \quad \checkmark$$

Because this task will never be delayed by any other task (sum = 0)

$$R_2^1 = 1 + \left\lceil \frac{1}{2} \right\rceil \cdot 1 = 2$$

$$R_3^1 = 2 + \left\lceil \frac{2}{2} \right\rceil \cdot 1 + \left\lceil \frac{3}{5} \right\rceil \cdot 1 = 4$$

$$R_4^1 = 1 + \left\lceil \frac{1}{2} \right\rceil \cdot 1 + \left\lceil \frac{1}{5} \right\rceil \cdot 1 + \left\lceil \frac{1}{8} \right\rceil \cdot 2 = 5$$

$$R_2^2 = 1 + \left\lceil \frac{2}{2} \right\rceil \cdot 1 = 2 \quad \checkmark$$

$$R_3^2 = 2 + \left\lceil \frac{4}{2} \right\rceil \cdot 1 + \left\lceil \frac{4}{5} \right\rceil \cdot 1 = 5$$

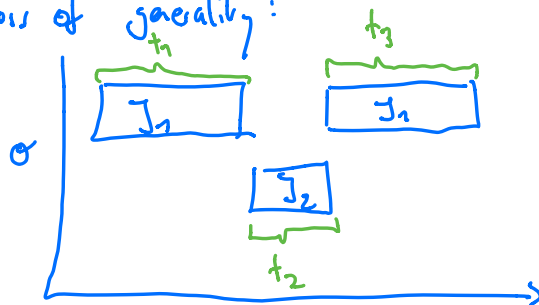
$$R_4^2 = 1 + \left\lceil \frac{5}{2} \right\rceil \cdot 1 + \left\lceil \frac{5}{5} \right\rceil \cdot 1 + \left\lceil \frac{5}{8} \right\rceil \cdot 2 = 7$$

$$R_3^3 = 2 + \left\lceil \frac{5}{2} \right\rceil \cdot 1 + \left\lceil \frac{5}{5} \right\rceil \cdot 1 = 6$$

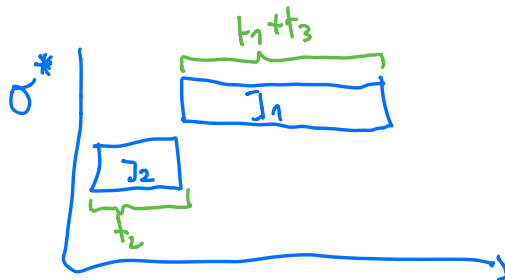
$R_4^4 = 21 \rightarrow$ miss deadline \rightarrow can't find RM schedule!

[C3] a) Assumption J_1, J_2 1 interrupt.

without loss of generality:



$$T_w(J_1, J_2, \sigma) = w_1(t_1 + t_2 + t_3) + w_2(t_1 + t_2)$$



$$T_w(J_1, J_2, \sigma^*) = w_1(t_1 + t_2 + t_3) + w_2(t_2)$$

$$\Rightarrow T_w(J_1, J_2, \sigma^*) \leq T_w(J_1, J_2, \sigma)$$

□

Assumption of 1 Interrupt: okay b.c. can apply iteratively

Assumption of 2 Tasks: we can combine tasks if there are more than 2.

[C3] b) Schedule tasks with highest Density first ($\frac{w}{c}$ or $w \cdot c$)

w	c
9	1
9	1
10	10

$$T_w = 10 \cdot 10 + 11 \cdot 9 + 12 \cdot 9 > 200$$
$$T_w^* = 1 \cdot 9 + 2 \cdot 9 + 12 \cdot 10 = 147$$

[C4] see our submission!