

1 Outline

- Extant CCG treatments of dynamic anaphora: [de Groote 2006](#), [Barker & Shan 2008](#). [Motivation for pursuing another approach. BS clearly no good ([Charlow 2010](#)). dG works well enough, and can be combined with BS regime, but...?]
- Continuized CCGs offer a grammar-wide generalization of scope-taking (“ubiquitous scopal pied piping”) using three operations: **lift**, **triv**, and **scope**. Any constituent can be a scope-taker.
- Proposal: replace **lift** and **triv** with options that countenance side effects ([Shan 2005](#)). Any side effects regime can be grafted onto a continuized CCG, by replacing **lift** and **triv** with monadic functors ([Moggi 1989](#), [Wadler 1992, 1994, 1995](#), [Shan 2002](#)).
- We provide a general technique for integrating a monadic approach to side effects with continuations-based approaches to scope in CCG. We relate our approach to the ContT monad transformer ([Liang et al. 1995](#)). Offers a type-theoretic way to track effects, integrate them into a well-developed CCG framework for scope-taking.
- Dynamic semantics is ([Shan 2001](#)):¹
 - State: ability to manipulate the discourse context, i.e. create discourse referents.
 - Nondeterminism: analogizes indefinites to referential expressions. Treats indefinites as referring expressions, though ones which refer indeterminately.
- Corollary: there is no need to settle on a single (“the”) grammar. Different and quite varied side effects regimes can be modularly grafted onto a simple applicative (“pure”) core. Lexical entries that would seem incongruous in a flat-footed standard perspective integrate seamlessly in a single grammar.
- Monads as a natural way to extend a continuations-based grammar with tools for dynamic binding and exceptional scope. In the end: you have functional application, plus the functors from whichever monads are implicated in a given language.
- The standard continuations-based perspective of [Barker 2002](#), [Shan & Barker 2006](#), [Barker & Shan 2014](#) is an instantiation of a more general perspective.

¹ NB: does not characterize all varieties of dynamic semantics. Dynamic treatments following [Groenendijk & Stokhof 1990](#) (e.g. [Zimmermann 1991](#), [Dekker 1993](#), [Szabolcsi 2003](#), [de Groote 2006](#)) provide a way for indefinites to extend their binding domain but do not treat indefinites as nondeterministic analogs of proper names.

- Standard dynamic techniques (DPL, DMG) not reducible to monads.
- Broader question: how this relates to the idea that continuations can simulate any monad ([Filinski 1994](#)). I don’t understand this result well enough to say anything.

2 Adding side effects to k

- Standard continuized grammar:
 - **lift**: $\lambda k. k\ x$
 - **triv**: $\lambda x. x$
 - **scope**: $\lambda k. m(\lambda f. n(\lambda x. k(f\ x)))$
- To do: insert figure with inference rules.
- Type-theoretic details here
- Adding side effects ([Wadler 1994, 1995](#), [Shan 2002](#)): monads
- Monad laws / punting
- Relating monads to continuized grammars:
 - Replace **lift** with \star
 - Replace **triv** with η
 - **scope** stays the same
- Two type constructors:
 - Bipartite Cont: $K\ a\ b ::= (a \rightarrow b) \rightarrow b$
 - Unary Monadic:

3 Finding the dynamic monad

- The meat of PLA ([Dekker 1994](#)): sentences are relations on stacks. Non-empty relations correspond to truth. Non-functional pairs in the relation correspond to nondeterminism introduced by indefinites (and perhaps disjunction).

$$\llbracket a \text{ linguist} \rrbracket = \lambda k s. \bigcup_{x \in \text{ling}} k\ x\ \widehat{s}x$$

- A different perspective on this: treating nondeterminism and state modification as side effects, within a functional programming setting for side effects.
- Monad for nondeterminism:

Definition 1 (The Set monad).

$$\begin{aligned} M\ a &::= a \rightarrow t \\ \eta\ x &::= \{x\} \\ m\ \star\ k &::= \bigcup_{x \in m} k\ x \end{aligned}$$

$$\frac{\Gamma \vdash f : b/a \quad \Delta \vdash e : a}{\Gamma \cdot \Delta \vdash fe : b} / \quad \frac{\Delta \vdash e : a \quad \Gamma \vdash f : a \setminus b}{\Delta \cdot \Gamma \vdash fe : b} \setminus \quad \frac{\Delta \vdash m : K(b/a)r \quad \Gamma \vdash n : Kar}{\Delta \cdot \Gamma \vdash S_{mn} : Kbr} // \quad \frac{\Delta \vdash m : Kar \quad \Gamma \vdash n : K(a \setminus b)r}{\Delta \cdot \Gamma \vdash S_{mn} : Kbr} //$$

$$\frac{\Gamma \vdash e : a}{\Gamma \vdash \lambda k. ke : Kar} \uparrow \quad \frac{\Gamma \vdash m : Krr}{\Gamma \vdash m(\lambda x. x) : r} \downarrow$$

Figure 1: Partial multimodal continuized grammar, no side effects.

$$\frac{\Gamma \vdash f : b/a \quad \Delta \vdash e : a}{\Gamma \cdot \Delta \vdash fe : b} / \quad \frac{\Delta \vdash e : a \quad \Gamma \vdash f : a \setminus b}{\Delta \cdot \Gamma \vdash fe : b} \setminus \quad \frac{\Delta \vdash m : K(b/a)r \quad \Gamma \vdash n : Kar}{\Delta \cdot \Gamma \vdash S_{mn} : Kbr} // \quad \frac{\Delta \vdash m : Kar \quad \Gamma \vdash n : K(a \setminus b)r}{\Delta \cdot \Gamma \vdash S_{mn} : Kbr} //$$

$$\frac{\Gamma \vdash e : Ma}{\Gamma \vdash \lambda k. e \star k : KaMr} \uparrow \quad \frac{\Gamma \vdash m : KrMr}{\Gamma \vdash m\eta : Mr} \downarrow$$

Figure 2: Partial multimodal continuized grammar, with side effects.

- Monad for state (generalization of monad for environment-sensitivity):

Definition 2 (The State monad).

$$\begin{aligned} Ma &::= s \rightarrow a \times s \\ \eta x &::= \lambda s. \langle x, s \rangle \\ m \star k &::= \lambda s. k (ms)_0 (ms)_1 \end{aligned}$$

- Use StateT to stitch the two together²

Definition 3 (The StateT monad transformer).

$$\begin{aligned} Ma &::= s \rightarrow L(a \times s) \\ \eta x &::= \lambda s. \eta_L \langle x, s \rangle \\ m \star k &::= \lambda s. m s \star_L \lambda \pi. k \pi_0 \pi_1 \end{aligned}$$

Definition 4 (The State_Set monad).

$$\begin{aligned} Ma &::= s \rightarrow (a \times s) \rightarrow t \\ \eta x &::= \lambda s. \{ \langle x, s \rangle \} \\ m \star k &::= \lambda s. \bigcup_{\pi \in ms} k \pi_0 \pi_1 \end{aligned}$$

- Static lexicon, dynamic lexicon
- Modular treatment of binding.

$$\begin{aligned} \text{Previous : } \mathbf{bind} m &:= \lambda k. m(\lambda x. k x x) \\ \text{Proposal : } \mathbf{bind} m &:= \lambda k. m(\lambda x s. k x \widehat{s} x) \end{aligned}$$

- Summing up: three combinators for “order-insensitive” (i.e. continuized combination). **unit**, **run**, **bind**

	lift m	M triv	bind M
Previous	$\lambda k. k m$	$M(\lambda x. x)$	$\lambda k. m(\lambda x. k x x)$
Proposal	$\lambda k. m \star k$	$M \eta$	$\lambda k. m(\lambda x s. k x \widehat{s} x)$

4 Examples

- Some upshots: no dynamic conjunction, completely standard model theory (cf. de Groote 2006). “Contexts of evaluation” are constructed on the fly.

de Groote 2001 Charlow 2014 Bumford to appear

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² Fn. about SetT

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