

## EVAN'S CONJECTURE

Let  $v$  be a  $b$ -bit word. Treat  $v$  as a binary number to get its value. For example,  $\text{val}(1, 1, 0, 0) = 1100_{\text{TWO}} = 12$ . We have  $v \prec w$  exactly when  $\text{val } v < \text{val } w$  and  $\text{val } \rho v < \text{val } \rho w$ .

- A *type I* transformation replaces a 0 in  $v$  with a 1.

Example:  $0010101 \rightarrow 0011101$ .

- A *type H* transformation replaces a substring in  $v$  of the form<sup>1</sup>  $011 \dots 10$  with its complement, i.e.,  $100 \dots 01$ .

Example:  $0011001 \rightarrow 0100101$ .

**Proposition 1.** *Suppose  $v, w \in W_b$  and  $w$  is derived from  $v$  by an I transformation, then  $v \prec w$ .*

*Proof.* Let  $v = (x, 0, y)$  and  $w = (x, 1, y)$ . Then  $\text{val } w = \text{val } v + 2^k > \text{val } v$  for some  $k$ . Likewise,  $\text{val } \rho w = \text{val } \rho v + 2^j > \text{val } \rho v$  for some  $j$ . Therefore  $v \prec w$ .  $\square$

**Proposition 2.** *Suppose  $v, w \in W_b$  and  $w$  is derived from  $v$  by an H transformation, then  $v \prec w$ .*

*Proof.* Let  $v = (x, 01 \dots 10, y)$  and  $w = (x, 10 \dots 01, y)$ . Note that  $10 \dots 01_{\text{TWO}} > 01 \dots 10_{\text{TWO}}$  and therefore  $\text{val } w > \text{val } v$ . Likewise,  $\text{val } \rho w > \text{val } \rho v$ . Therefore  $v \prec w$ .  $\square$

**Corollary 3.** *Let  $v, w \in W_b$ . If  $w$  is obtained from  $v$  by a series of I and H transformations, then  $v \preceq w$ .*  $\square$

The interesting part is the converse.

**Theorem 4.** *Let  $v, w \in W_b$  with  $v \prec w$ . Then  $w$  is obtained from  $v$  by a sequence of I and H transformations.*

*Proof.* The proof is by strong induction on  $b$ . It is easy to verify for  $W_1$ ,  $W_2$ , and  $W_3$ . Suppose the result has been shown for all  $W_i$  with  $i < b$ .

Let  $v, w \in W_b$ . We consider the first and last bits of both words, and break into cases. All told there are 16 possibilities: Four possibilities for  $v$ :  $(0, \bar{v}, 0)$ ,  $(0, \bar{v}, 1)$ ,  $(1, \bar{v}, 0)$ , and  $(1, \bar{v}, 1)$ . And likewise for  $w$ .

Cases as follows:

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<sup>1</sup>The substring begins and ends with a 0 and has one or more consecutive 1s in between.

1.  $v = (0, \bar{v}, 0)$  and  $w = (0, \bar{w}, 0)$ : Induction on  $\bar{v} \prec \bar{w}$ .
  2.  $v = (0, \bar{v}, 0)$  and  $w = (0, \bar{w}, 1)$ : Cannot happen because  $\rho v = (0, \rho \bar{v}, 0) \not\prec (1, \rho \bar{w}, 0) = \rho w$ .
  3.  $v = (0, \bar{v}, 0)$  and  $w = (1, \bar{w}, 0)$ :
  4.  $v = (0, \bar{v}, 0)$  and  $w = (1, \bar{w}, 1)$ :
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5.  $v = (0, \bar{v}, 1)$  and  $w = (0, \bar{w}, 0)$ :
  6.  $v = (0, \bar{v}, 1)$  and  $w = (0, \bar{w}, 1)$ :
  7.  $v = (0, \bar{v}, 1)$  and  $w = (1, \bar{w}, 0)$ :
  8.  $v = (0, \bar{v}, 1)$  and  $w = (1, \bar{w}, 1)$ :
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9.  $v = (1, \bar{v}, 0)$  and  $w = (0, \bar{w}, 0)$ :
  10.  $v = (1, \bar{v}, 0)$  and  $w = (0, \bar{w}, 1)$ :
  11.  $v = (1, \bar{v}, 0)$  and  $w = (1, \bar{w}, 0)$ :
  12.  $v = (1, \bar{v}, 0)$  and  $w = (1, \bar{w}, 1)$ :
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13.  $v = (1, \bar{v}, 1)$  and  $w = (0, \bar{w}, 0)$ :
14.  $v = (1, \bar{v}, 1)$  and  $w = (0, \bar{w}, 1)$ :
15.  $v = (1, \bar{v}, 1)$  and  $w = (1, \bar{w}, 0)$ :
16.  $v = (1, \bar{v}, 1)$  and  $w = (1, \bar{w}, 1)$ :

□

Recall that, in a poset,  $v$  is *covered* by  $w$  provided  $v \prec w$  and there is no element  $x$  with  $v \prec x \prec w$ .

**Corrolary 5.** For  $v, w \in W_b$ ,  $v$  is covered by  $w$  if and only if  $w$  is obtained from  $v$  by a type I or type H transformation. □

In particular, if  $w_1 \prec w_2 \prec \cdots \prec w_t$  is a maximum chain in  $W_b$ , then each  $w_{i+1}$  is obtained from  $w_i$  by a type I or a type H transformation.