### Game Theory

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Lecture 2

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### Lecture 1 recap

- Defined games in normal form
- Defined dominance notion
  - Iterative deletion
  - Does not always give a solution
- Defined best response and Nash equilibrium
  - Computed Nash equilibrium in some examples
- → Are some Nash equilibria better than others?
- → Can we always find a Nash equilibrium?

### Outline

- 1. Coordination games and Pareto optimality
- 2. Games with continuous action sets
  - Equilibrium computation and existence theorem
  - Example: Cournot duopoly

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### The Investment Game

- The players: you
- The strategies: each of you chose between investing nothing in a class project (\$0) or investing (\$10)
- Payoffs:
  - If you don't invest your payoff is \$0
  - If you invest you make a **net profit** of \$5 (gross profit = \$15; investment \$10) if more than 90% of the class chooses to invest. Otherwise, you lose \$10
- Choose your action (no communication!)

## Nash equilibrium

What are the Nash equilibria?

- Remark: to find Nash equilibria, we used a "guess and check method"
  - Checking is easy, guessing can be hard

## The Investment Game again

- Recall that:
  - Players: you
  - Strategies: invest \$0 or invest \$10
  - Payoffs:
    - If no invest  $\rightarrow$  \$0 \$5 net profit if  $\geq$  90% invest • If invest \$10  $\rightarrow$  -\$10 net profit if < 90% invest
- Let's play again! (no communication)
- We are heading toward an equilibrium
- There are certain cases in which playing converges in a natural sense to an equilibrium

#### Pareto domination

Is one equilibrium better than the other?

#### Definition: Pareto domination

A strategy profile s Pareto dominates strategy profile s' iif for all i,  $u_i(s) \ge u_i(s')$  and there exists j such that  $u_i(s) > u_i(s')$ ;

i.e., all players have at least as high payoffs and at least one player has strictly higher payoff.

In the investment game?

## Convergence to equilibrium in the Investment Game

- Why did we converge to the wrong NE?
- Remember when we started playing
  - We were more or less 50 % investing
- The starting point was already bad for the people who invested for them to lose confidence
- Then we just tumbled down
- What would have happened if we started with 95% of the class investing?

## Coordination game

- This is a *coordination game* 
  - We'd like everyone to coordinate their actions and <u>invest</u>
- Many other examples of coordination games
  - Party in a Villa
  - On-line Web Sites
  - Establishment of technological monopolies (Microsoft, HDTV)
  - Bank runs
- Unlike in prisoner's dilemma, <u>communication helps</u> in coordination games → <u>scope for leadership</u>
  - In prisoner's dilemma, a trusted third party (TTP) would need to impose players to adopt a strictly dominated strategy
  - In coordination games, a TTP just leads the crowd towards a better NE point (there is no dominated strategy)

### Battle of the sexes



Find the NE

Is there a NE better than the other(s)?

#### **Coordination Games**

- Pure coordination games: there is no conflict whether one NE is better than the other
  - E.g.: in the investment game, we all agreed that the NE with everyone investing was a "better" NE
- General coordination games: there is a source of conflict as players would agree to coordinate, but one NE is "better" for a player and not for the other
  - E.g.: Battle of the Sexes
- → Communication might fail in this case

### Pareto optimality

#### Definition: Pareto optimality

A strategy profile s is Pareto optimal if there does not exist a strategy profile s' that Pareto dominates s.

Battle of the sexes?

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## The partnership game (see exercise sheet 2)

- Two partners choose effort s<sub>i</sub> in S<sub>i</sub>=[0, 4]
- Share revenue and have quadratic costs

$$u_1(s_1, s_2) = \frac{1}{2} [4 (s_1 + s_2 + b s_1 s_2)] - s_1^2$$
  
 $u_2(s_1, s_2) = \frac{1}{2} [4 (s_1 + s_2 + b s_1 s_2)] - s_2^2$ 

• Best responses:

$$\hat{s}_1 = 1 + b s_2 = BR_1(s_2)$$
  
 $\hat{s}_2 = 1 + b s_1 = BR_2(s_1)$ 

## Finding the best response (with twice continuously differentiable utilities)

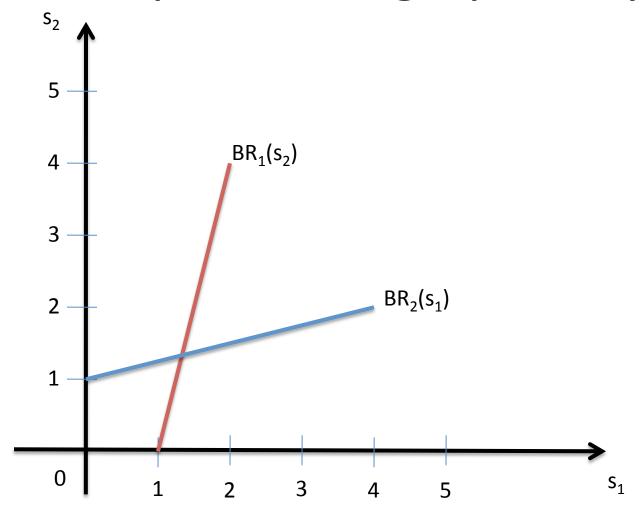
$$\frac{\partial u_1(s_1, s_2)}{\partial s_1} = 0$$

First order condition (FOC)

$$\frac{\partial^2 u_1(s_1, s_2)}{\partial^2 s_1} \le 0$$

- Second order condition (SOC)
- Remark: the SOC is automatically satisfied if  $u_i(s_i, s_{-i})$  is concave in  $s_i$  for all  $s_{-i}$  (very standard assumption)
- Remark 2: be careful with the borders!
  - Example  $u_1(s_1, s_2) = 10-(s_1+s_2)^2$
  - $S_1 = [0, 4]$ , what is the BR to  $S_2 = 2$ ?
  - Solving the FOC, what do we get?
  - When the FOC solution is outside S<sub>i</sub>, the BR is at the border

## Nash equilibrium graphically



• NE is fixed point of  $(s_1, s_2) \rightarrow (BR(s_2), BR(s_1))$ 

### Best response correspondence

- Definition: ŝ<sub>i</sub> is a BR to s<sub>-i</sub> if ŝ<sub>i</sub> solves max u<sub>i</sub>(s<sub>i</sub>, s<sub>-i</sub>)
- The BR to s<sub>-i</sub> may not be unique!
- BR(s<sub>-i</sub>): set of s<sub>i</sub> that solve max u<sub>i</sub>(s<sub>i</sub>, s<sub>-i</sub>)
- The definition can be written:  $\hat{s}_i$  is a BR to  $s_{-i}$  if  $\hat{s}_i \in BR_i(s_{-i}) = \underset{s_i}{\operatorname{argmax}} u_i(s_i, s_{-i})$
- Best response correspondence of i:  $s_{-i} \rightarrow BR_i(s_{-i})$
- (Correspondence = set-valued function)

## Nash equilibrium as a fixed point

- Game  $(N,(S_i)_{i\in N},(u_i)_{i\in N})$
- Let's define  $S = \times_{i \in N} S_i$  (set of strategy profiles) and the correspondence

$$B: S \to S$$
$$s \mapsto B(s) = \times_{i \in N} BR_i(s_{-i})$$

- For a given s, B(s) is the set of strategy profiles s' such that s<sub>i</sub>' is a BR to s<sub>-i</sub> for all i.
- A strategy profile  $s^*$  is a Nash eq. iif  $s^* \in B(s^*)$  (just a re-writing of the definition)

## Kakutani's fixed point theorem

#### Theorem: Kakutani's fixed point theorem

Let X be a compact convex subset of  $\mathbb{R}^n$  and let  $f: X \to X$  be a set-valued function for which:

- for all  $x \in X$ , the set f(x) is nonempty convex;
- the graph of f is closed.

Then there exists  $x^* \in X$  such that  $x^* \in f(x^*)$ 

### Closed graph (upper hemicontinuity)

- Definition: f has closed graph if for all sequences  $(x_n)$  and  $(y_n)$  such that  $y_n$  is in  $f(x_n)$  for all n,  $x_n \rightarrow x$  and  $y_n \rightarrow y$ , y is in f(x)
- Alternative definition: f has closed graph if for all x we have the following property: for any open neighbourhood V of f(x), there exists a neighbourhood U of x such that for all x in U, f(x) is a subset of V.
- Examples:

# Existence of (pure strategy) Nash equilibrium

#### Theorem: Existence of pure strategy NE

Suppose that the game  $(N,(S_i)_{i\in N},(u_i)_{i\in N})$  satisfies:

- The action set  $S_i$  of each player is a nonempty compact convex subset of  $\mathbb{R}^n$
- The utility  $u_i$  of each player is continuous in s (on s) and concave in  $s_i$  (on  $s_i$ )

Then, there exists a (pure strategy) Nash equilibrium.

Remark: the concave assumption can be relaxed

### **Proof**

- Define B as before. B satisfies the assumptions of Kakutani's fixed point theorem
- Therefore B has a fixed point which by definition is a Nash equilibrium!
- Now, we need to actually verify that B satisfies the assumptions of Kakutani's fixed point theorem!

## Example: the partnership game

- $N = \{1, 2\}$
- S = [0,4]x[0,4] compact convex
- Utilities are continuous and concave

$$u_1(s_1, s_2) = \frac{1}{2} [4 (s_1 + s_2 + b s_1 s_2)] - s_1^2$$
  
 $u_2(s_1, s_2) = \frac{1}{2} [4 (s_1 + s_2 + b s_1 s_2)] - s_2^2$ 

- Conclusion: there exists a NE!
- Ok, for this game, we already knew it!
- But the thm is much more general and applies to games where finding the equilibrium is much more difficult

## One more word on the partnership game before we move on

- We have found (see exercises) that
  - At Nash equilibrium:

$$s_1^* = s_2^* = 1/(1-b)$$

- To maximize the sum of utilities:

$$s_{1}^{W} = s_{2}^{W} = 1/(1/2-b) > s_{1}^{*}$$

- Sum of utilities called social welfare
- Both partners would be better off if they worked s<sup>W</sup><sub>1</sub> (with social planner, contract)
- Why do they work less than efficient?

## Externality

- At the margin, I bear the cost for the extra unit of effort I contribute, but I'm only reaping half of the induced profits, because of **profit sharing**
- This is known as an "externality"
- → When I'm figuring out the effort I have to put I don't take into account that other half of profit that goes to my partner
- →In other words, my effort benefits my partner, not just me
- Externalities are omnipresent: public good problems, free riding, etc. (see more in the netecon course)

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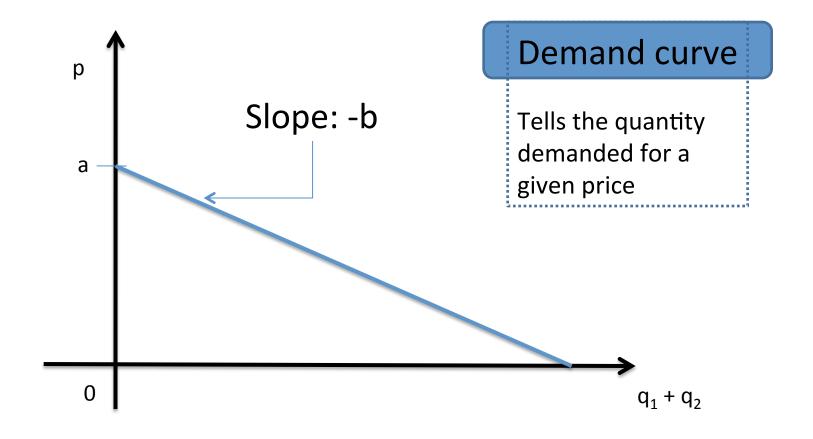
## **Cournot Duopoly**

- Example of application of games with continuous action set
- This game lies between two extreme cases in economics, in situations where firms (e.g. two companies) are competing on the same market
  - Perfect competition
  - Monopoly
- We're interested in understanding what happens in the middle
  - The game analysis will give us interesting economic insights on the duopoly market

## Cournot Duopoly: the game

- The players: 2 Firms, e.g., Coke and Pepsi
- Strategies: quantities players produce of <u>identical</u> products: q<sub>i</sub>, q<sub>-i</sub>
  - Products are perfect substitutes
- Cost of production: c \* q
  - Simple model of <u>constant marginal cost</u>
- Prices:  $p = a b (q_1 + q_2) = a bQ$ 
  - Market-clearing price

## Price in the Cournot duopoly



## Cournot Duopoly: payoffs

The payoffs: firms aim to <u>maximize profit</u>

$$u_1(q_1,q_2) = p * q_1 - c * q_1$$
  
 $p = a - b (q_1 + q_2)$ 

$$> u_1(q_1,q_2) = a * q_1 - b * q_1^2 - b * q_1 q_2 - c * q_1$$

The game is symmetric

$$> u_2(q_1,q_2) = a * q_2 - b * q_2^2 - b * q_1 q_2 - c * q_2$$

## Cournot Duopoly: best responses

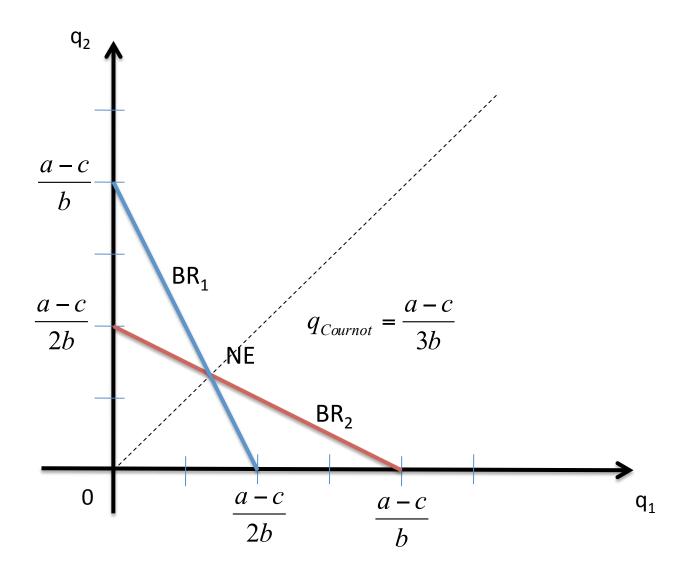
First order condition

$$a - 2bq_1 - bq_2 - c = 0$$

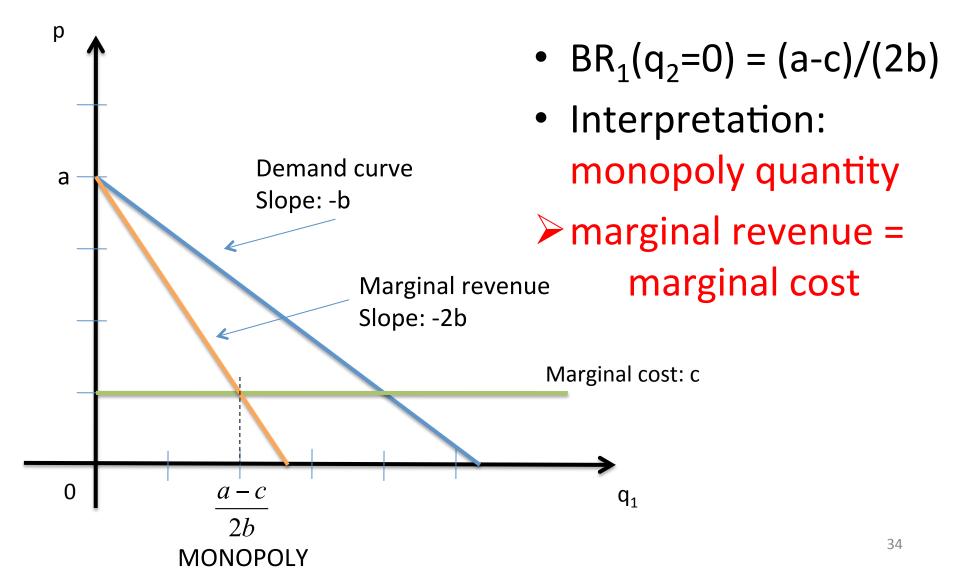
• Second order condition -2b < 0 [make sure it's a max]

$$\begin{cases} \hat{q}_1 = BR_1(q_2) = \frac{a-c}{2b} - \frac{q_2}{2} \\ \hat{q}_2 = BR_2(q_1) = \frac{a-c}{2b} - \frac{q_1}{2} \end{cases}$$

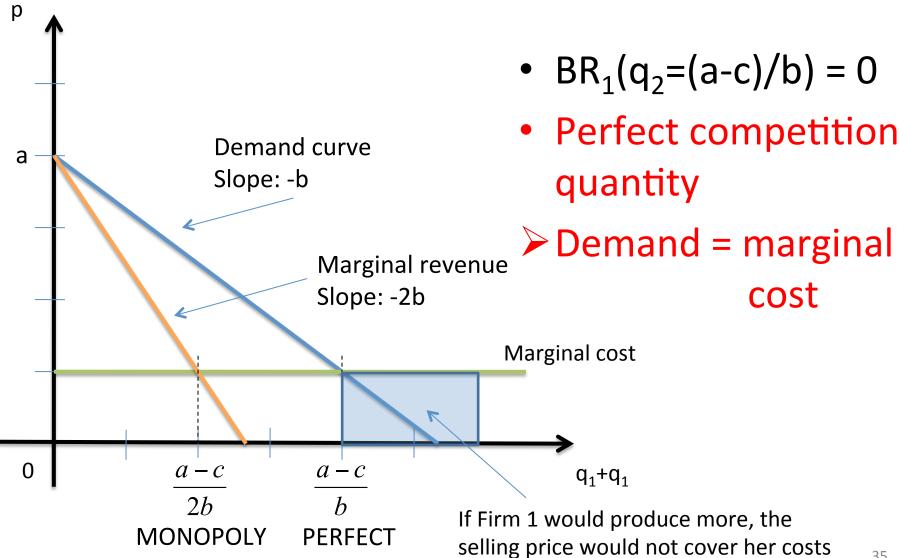
## Cournot Duopoly: best response diagram and Nash equilibrium



## Best response at $q_2=0$

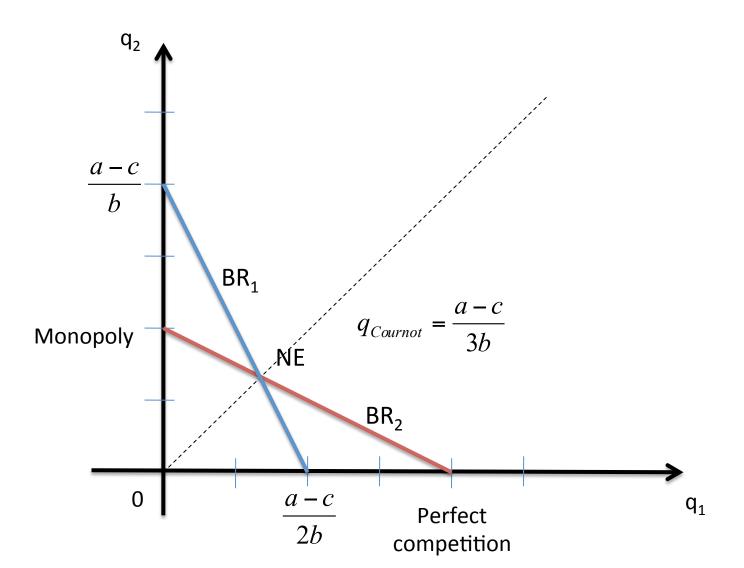


## When is $BR_1(q_2) = 0$ ?



**COMPETITION** 

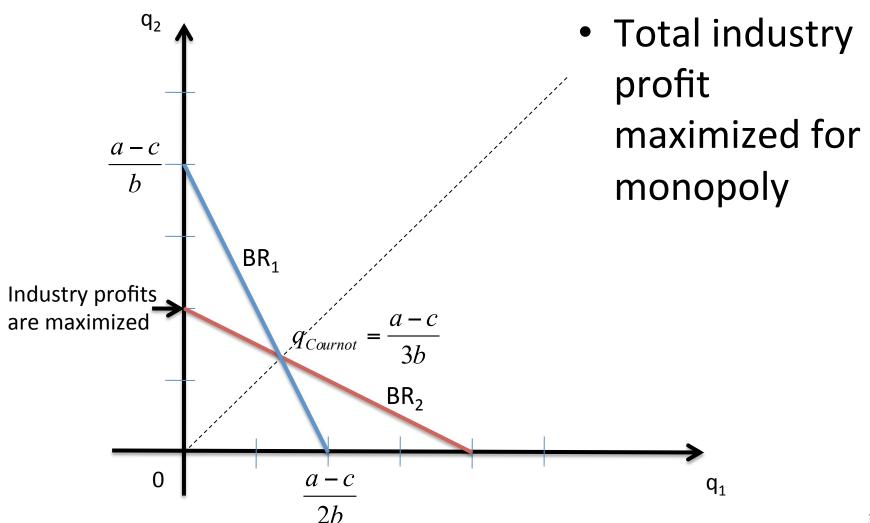
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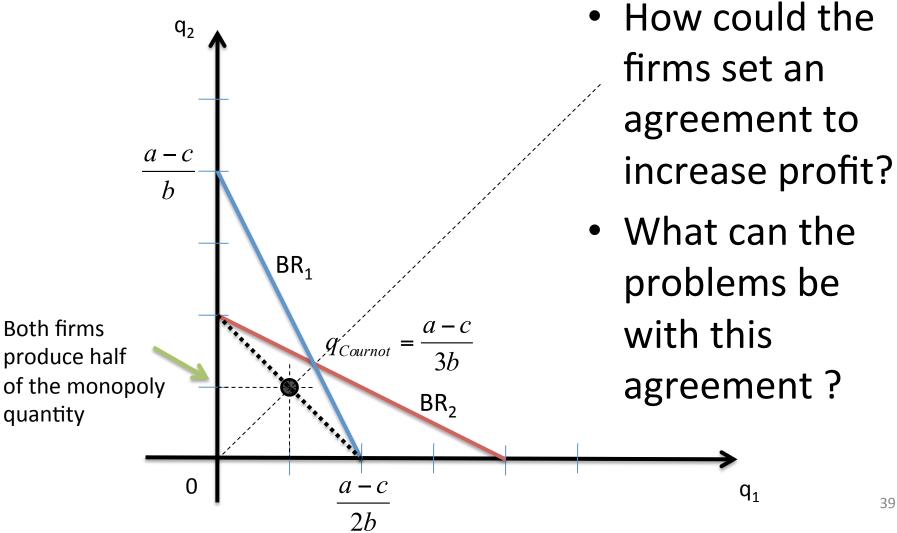
## Strategic substitutes/complements

- In Cournot duopoly: the more the other player does, the less I would do
- This is a game of **strategic substitutes** 
  - Note: of course the goods were substitutes
  - We're talking about strategies here
- In the partnership game, it was the opposite: the more the other player would the more I would do
- This is a game of **strategic complements**

### Cournot duopoly: Market perspective



## Cartel, agreement



## Cournot Duopoly: last observations

 How do quantities and prices we've encountered so far compare?

**QUANTITIES** 

$$\frac{a-c}{b}$$

$$\frac{2(a-c)}{3b}$$

$$\frac{a-c}{2b}$$

**PRICES** 

## Summary

- Coordination games
  - Pareto optimal NE sometimes exist
  - Scope for communication / leadership
- Games with continuous action sets (pure strategies)
  - Compute equilibrium with FOC, SOC
  - Equilibrium exists under concavity and continuity conditions
  - Cournot duopoly