### Game Theory

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Lecture 5

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### Lecture 3-4 recap

- Defined mixed strategy Nash equilibrium
- Proved existence of mixed strategy Nash equilibrium in finite games
- Discussed computation and interpretation of mixed strategies Nash equilibrium
- Defined another concept of equilibrium from evolutionary game theory
- → Today: introduce other solution concepts for simultaneous moves games
- → Introduce solutions for sequential moves games

### Outline

- Other solution concepts for simultaneous moves
  - Stability of equilibrium
    - Trembling-hand perfect equilibrium
  - Correlated equilibrium
  - Minimax theorem and zero-sum games
  - ε-Nash equilibrium
- The lender and borrower game: introduction and concepts from sequential moves

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### The Location Model

- Assume we have 2N players in this game (e.g., N=70)
  - Players have two types: tall and short
  - There are N tall players and N short players
- Players are people who need to decide in which town to live
- There are two towns: East town and West town
  - Each town can host no more than N players

#### Assume:

 If the number of people choosing a particular town is larger than the town capacity, the surplus will be redistributed randomly

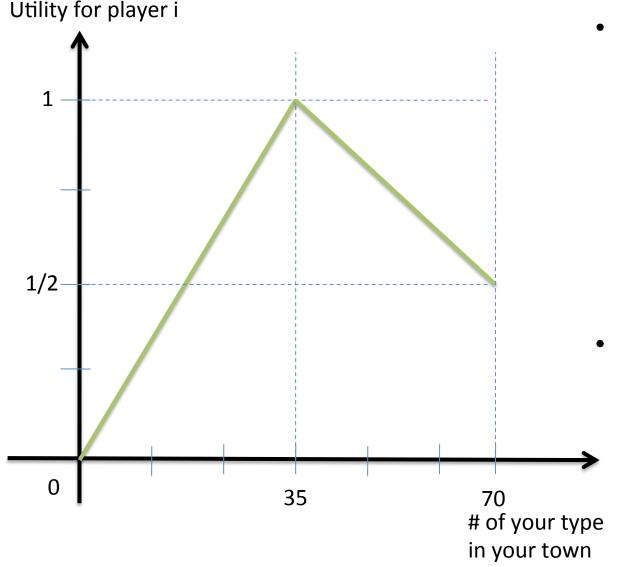
#### Game:

Players: 2N people

Strategies: East or West town

Payoffs

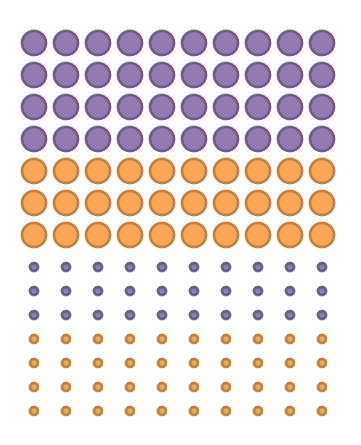
# The Location Model: payoffs



#### • The idea is:

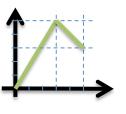
- If you are a small <u>minority</u> in your town you get a payoff of zero
- If you are in large
   <u>majority</u> in your town
   you get a payoff of ½
- If you are well
   <u>integrated</u> you get a payoff of 1
- People would like to live in mixed towns, but if they cannot, then they prefer to live in the majority town

#### Initial state



Assume the initial picture is this one

What will players do?

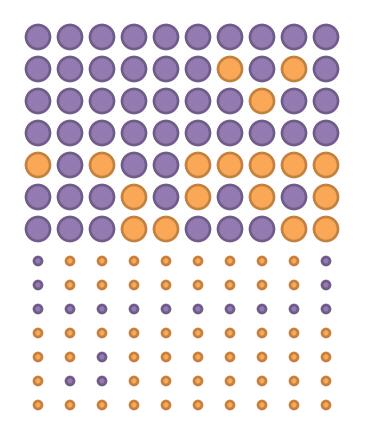




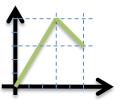
Tall player Short player



#### First iteration



- For tall players
  - There's a minority of east town "giants" to begin with
  - > switch to West town
- For short players
  - There's a minority of west town "dwarfs" to begin with
  - → switch to East town

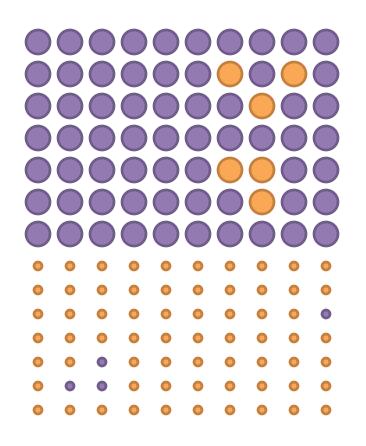




Tall player Short player West Town

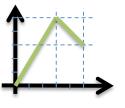
East Town

### Second iteration



Same trend

- Still a few players who did not understand
  - What is their payoff?

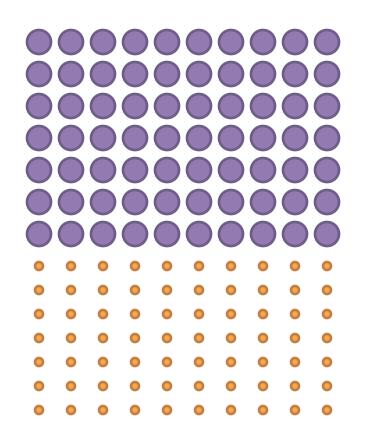




Tall player Short player



#### Last iteration



People got segregated

- But they would have preferred integrated towns!
  - Why? What happened?
  - People that started in a minority (even though not a "bad" minority)
     had incentives to deviate





Tall player



Short player



### The Location Model: Nash equilibria

- Two segregated NE:
  - Short, E; Tall, W
  - Short, W; Tall, E

Is there any other NE?

# Stability of equilibria

- The integrated equilibrium is not stable
  - If we move away from the 50% ratio, even a little bit, players have an incentive to deviate even more
  - We end up in one of the segregated equilibrium
- The segregated equilibria are stable
  - Introduce a small perturbation: players come back to segregation quickly
- Notion of stability in Physics: if you introduce a small perturbation, you come back to the initial state
- Tipping point:
  - Introduced by Grodzins (White flights in America)
  - Extended by Shelling (Nobel prize in 2005)

### Trembling-hand perfect equilibrium

#### Definition: Trembling-hand perfect equilibrium

A (mixed) strategy profile s is a trembling-hand perfect equilibrium if there exists a sequence  $s^{(0)}$ ,  $s^{(1)}$ , ... of fully mixed strategy profiles that converges towards s and such that for all k and all player i,  $s_i$  is a best response to  $s^{(k)}_{-i}$ .

- Fully-mixed strategy: positive probability on each action
- Informally: a player's action  $s_i$  must be BR not only to opponents equilibrium strategies  $s_{-i}$  but also to small perturbations of those  $s^{(k)}_{-i}$ .

### The Location Model

- The segregated equilibria are trembling-hand perfect
- The integrated equilibrium is not tremblinghand perfect

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# Example: battle of the sexes



- NE: (O, O), (S, S) and ((1/3, 2/3), (2/3, 1/3))
  - The mixed equilibrium has payoff 2/3 each
- Suppose the players can observe the outcome of a fair toss coin and condition their strategies on this outcome
  - New strategies possible: O if head, S if tails
  - Payoff 1.5 each
- The fair coin acts as a correlating device

### Correlated equilibrium: general case

- In the previous example: both players observe the exact same signal (outcome of the coin toss random variable)
- General case: each player receives a signal which can be correlated to the random variable (coin toss) and to the other players signal
- Model:
  - n random variables (one per player)
  - A joint distribution over the n RVs
  - Nature chooses according to the joint distribution and reveals to each player only his RV
  - → Agent can condition his action to his RV (his signal)

# Correlated equilibrium: definition

### Definition: Correlated equilibrium

A correlated equilibrium of the game (N,  $(A_i)$ ,  $(u_i)$ ) is a tuple  $(v, \pi, \sigma)$  where

- $v=(v_1, ..., v_n)$  is a tuple of random variables with domains  $(D_1, ..., D_n)$
- $\pi$  is a joint distribution over v
- $\sigma=(\sigma_1, ..., \sigma_n)$  is a vector of mappings  $\sigma_i: D_i \rightarrow A_i$  such that for all i and any mapping  $\sigma_i': D_i \rightarrow A_i$ ,

$$\sum_{d \in D_1 \times \cdots \times D_n} \pi(d) u(\sigma_1(d_1), \cdots, \sigma_i(d_i), \cdots, \sigma_n(d_n)) \geq \sum_{d \in D_1 \times \cdots \times D_n} \pi(d) u(\sigma_1(d_1), \cdots, \sigma_i'(d_i), \cdots, \sigma_n(d_n))$$

## Correlated vs Nash equilibrium

 The set of correlated equilibria contains the set of Nash equilibria

#### Theorem:

For every Nash equilibrium  $\sigma^*$ , there exists a correlated equilibrium (v,  $\pi$ ,  $\sigma$ ) such that for each player i, the distribution induced on  $A_i$  is  $\sigma_i$ .

• Proof: construct it with  $D_i=A_i$ , independent signals  $(\pi(d)=\sigma_1^*(d_1)x...x\sigma_n^*(d_n))$  and identity mappings  $\sigma_i$ 

# Correlated vs Nash equilibrium (2)

- Not all correlated equilibria correspond to a Nash equilibrium
- Example, the correlated equilibrium in the battle-of-sex game

→ Correlated equilibrium is a strictly weaker notion than NE

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### Maxmin strategy

Maximize "worst-case payoff"

#### Definition: Maxmin strategy

The maxmin strategy for player i is  $\underset{s_i}{\operatorname{arg max}} \min_{s_{-i}} u_i(s_i, s_{-i})$ 

Example

Attacker: Not attack

– Defender: Defend

Attack Not att 0,-1 0,0

defender

This is not a Nash equilibrium!

## Maxmin strategy: intuition

- Player i commits to strategy s<sub>i</sub> (possibly mixed)
- Player –i observe s<sub>i</sub> and choose s<sub>-i</sub> to minimize
   i's payoff

• Player i guarantees payoff at least equal to the maxmin value  $\max_{s_i} \min_{s_i} u_i(s_i, s_{-i})$ 

# Two players zero-sum games

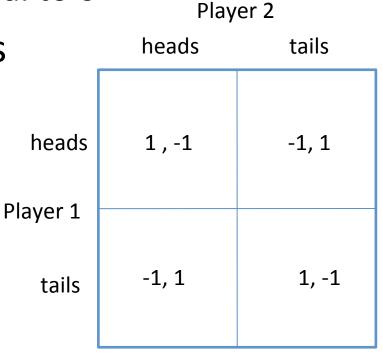
- Definition: a 2-players zero-sum game is a game where  $u_1(s)=-u_2(s)$  for all strategy profile s
  - Sum of payoffs constant equal to 0

Example: Matching pennies

• Define  $u(s)=u_1(s)$ 

Player 1: maximizer

– Player 2: minimizer



### Minimax theorem

#### Theorem: Minimax theorem (Von Neumann 1928)

For any two-player zero-sum game with finite action space:  $\max_{s_1} u(s_1, s_2) = \min_{s_2} \max_{s_1} u(s_1, s_2)$ 

- This quantity is called the value of the game
  - corresponds to the payoff of player 1 at NE
- Maxmin strategies ⇔ NE strategies
- Can be computed in polynomial time (through linear programming)

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## ε-Nash equilibrium

#### Definition: ε-Nash equilibrium

For  $\varepsilon>0$ , a strategy profile  $(s_1^*, s_2^*, ..., s_N^*)$  is an  $\varepsilon$ Nash equilibrium if, for each player i,  $u_i(s_i^*, s_{-i}^*) \ge u_i(s_i, s_{-i}^*) - \varepsilon \text{ for all } s_i \ne s_i^*$ 

- It is an approximate Nash equilibrium
  - Agents indifferent to small gains (could not gain more than  $\epsilon$  by unilateral deviation)
- A Nash equilibrium is an ε-Nash equilibrium for all ε!

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# "Cash in a Hat" game (1)

- Two players, 1 and 2
- Player 1 strategies: put \$0, \$1 or \$3 in a hat

Then, the hat is passed to player 2

 Player 2 strategies: either "match" (i.e., add the same amount of money in the hat) or take the cash

# "Cash in a Hat" game (2)

### Payoffs:

```
• Player 1: \begin{cases} \$0 \rightarrow \$0 \\ \$1 \rightarrow \text{ if match net profit } \$1, -\$1 \text{ if not } \$3 \rightarrow \text{ if match net profit } \$3, -\$3 \text{ if not } \end{cases}
```

Player 2: 

 Match \$1 → Net profit \$1.5

 Match \$3 → Net profit \$2
 Take the cash → \$ in the hat

### Lender & Borrower game

- The "cash in a hat" game is a toy version of the more general "lender and borrower" game:
  - Lenders: Banks, VC Firms, ...
  - Borrowers: entrepreneurs with project ideas
- The lender has to decide <u>how much money to</u> <u>invest</u> in the project
- After the money has been invested, the borrower could
  - Go forward with the project and work hard
  - Shirk, and run to Mexico with the money

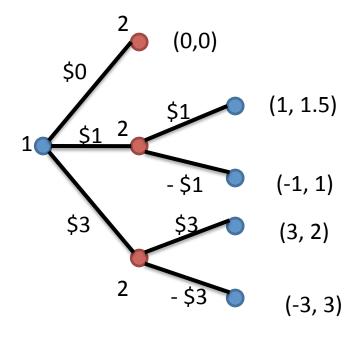
# Simultaneous vs. Sequential Moves

- What is different about this game wrt games studied until now?
- It is a sequential move game
  - Player chooses first, then player 2
- Timing is not the key
  - The key is that P2 observes P1's choice before choosing
  - And P1 knows that this is going to be the case

# Extensive form games

- A useful representation of such games is game trees also known as the extensive form
  - Each internal node of the tree will represent the ability of a player to make choices at a certain stage, and they are called <u>decision nodes</u>
  - Leafs of the tree are called <u>end nodes</u> and represent payoffs to both players
- Normal form games → matrices
- Extensive form games → trees

# "Cash in a hat" representation

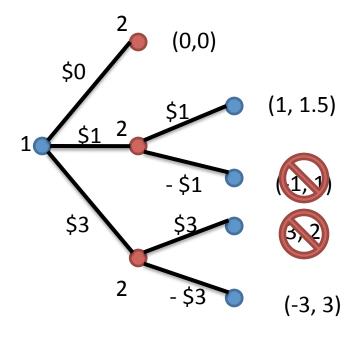


How to analyze such game?

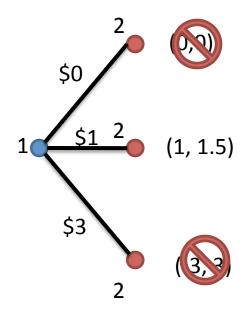
### **Backward Induction**

- Fundamental concept in game theory
- Idea: players that move early on in the game should <u>put</u> themselves in the shoes of other players playing later
  - → <u>anticipation</u>
- Look at the end of the tree and work back towards the root
  - Start with the last player and chose the strategies yielding higher payoff
  - This simplifies the tree
  - Continue with the before-last player and do the same thing
  - Repeat until you get to the root

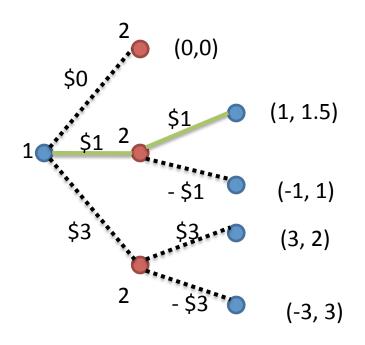
# Backward Induction in practice (1)



## Backward Induction in practice (2)



### Backward Induction in practice (3)



#### Outcome:

Player 1 chooses to invest \$1, Player 2 matches,

# The problem with the "lenders and borrowers" game

- It is not a disaster:
  - The lender doubled her money
  - The borrower was able to go ahead with a small scale project and make some money
- But, we would have liked to end up in another branch:
  - Larger project funded with \$3 and an outcome better for both the lender and the borrower
- Very similar to prisoner's dilemna
- What prevents us from getting to this latter good outcome?

#### Moral Hazard

- One player (the borrower) has incentives to do things that are not in the interests of the other player (the lender)
  - By giving a too big loan, the incentives for the borrower will be such that they will not be aligned with the incentives on the lender
  - Notice that moral hazard has also disadvantages for the borrower
- Example: Insurance companies offers "full-risk" policies
  - People subscribing for this policies may have no incentives to take care!
  - In practice, insurance companies force me to bear some deductible costs ("franchise")
- One party has incentive to take a risk because the cost is felt by another party
- How can we solve the Moral Hazard problem?

## Solution (1): Introduce laws

Today we have such laws: <u>bankruptcy laws</u>

- But, there are limits to the degree to which borrowers can be punished
  - The borrower can say: I can't repay, I'm bankrupt
  - And he/she's more or less allowed to have a fresh start

# Solution (2): Limits/restrictions on money

- Ask the borrowers a concrete plan (<u>business</u>)
   plan) on how he/she will spend the money
- This boils down to changing the order of play!

- Also faces some issues:
  - Lack of flexibility, which is the motivation to be an entrepreneur in the first place!
  - Problem of timing: it is sometimes hard to predict up-front all the expenses of a project

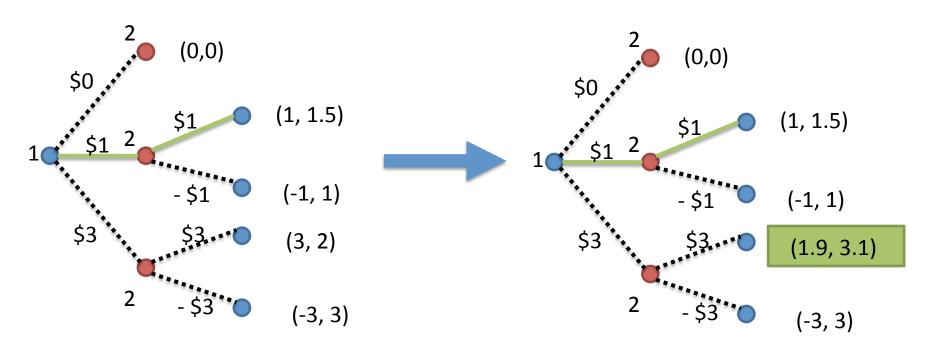
### Solution (3): Break the loan up

- Let the loan come in small installments
- If a borrower does well on the first installment, the lender will give a bigger installment next time

 It is similar to taking this one-shot game and turn it into a repeated game

## Solution (4): Change contract to avoid shirk -- **Incentives**

 The borrower could re-design the payoffs of the game in case the project is successful



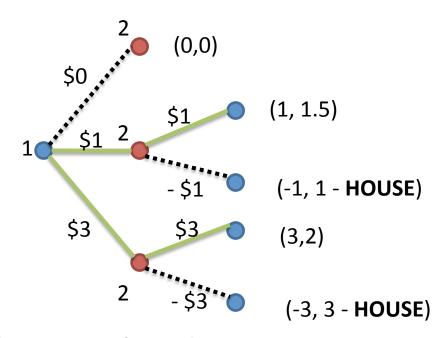
- Profit doesn't match investment but the outcome is better
  - Sometimes a smaller share of a larger pie can be bigger than a larger share of a smaller pie

### Absolute payoff vs ROI

- Previous example: larger absolute payoff in the new game on the right, but smaller return on investment (ROI)
- Which metric (absolute payoff or ROI) should an investment bank look at?

# Solution (5): Beyond incentives, collaterals

- The borrower could re-design the payoffs of the game in case the project is successful
  - Example: subtract house from run away payoffs



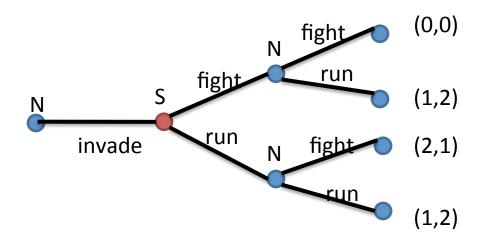
 Lowers the payoffs to borrower at some tree points, yet makes the borrower better off!

#### Collaterals

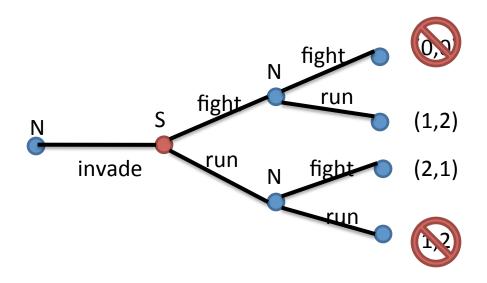
- They do hurt a player enough to change his/ her behavior
- → Lowering the payoffs at certain points of the game, does not mean that a player will be worse off!!

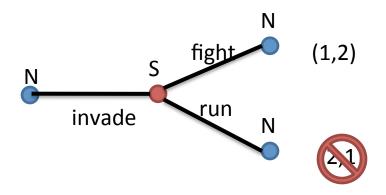
- Collaterals are part of a larger branch called commitment strategies
  - Next, an example of commitment strategies

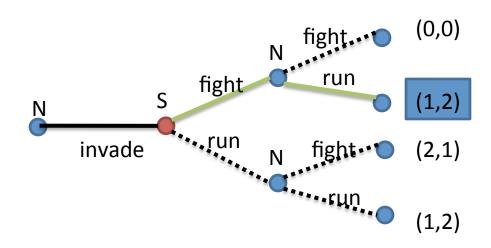
- Collaterals are part of a larger branch called commitment strategies
- Back in 1066, William the Conqueror lead an invasion from Normandy on the Sussex beaches
- We're talking about <u>military strategy</u>
- So basically we have two players (the armies) and the strategies available to the players are whether to "fight" or "run"



Let's analyze the game with Backward Induction





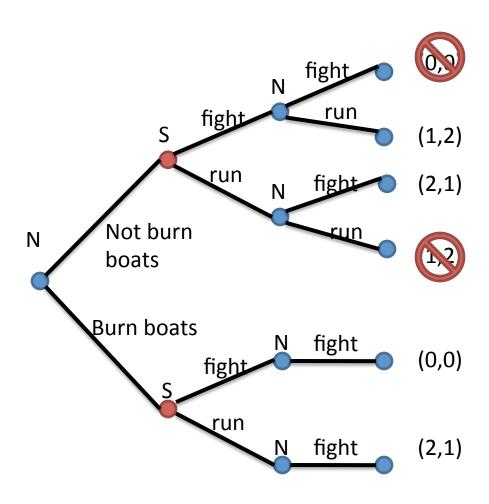


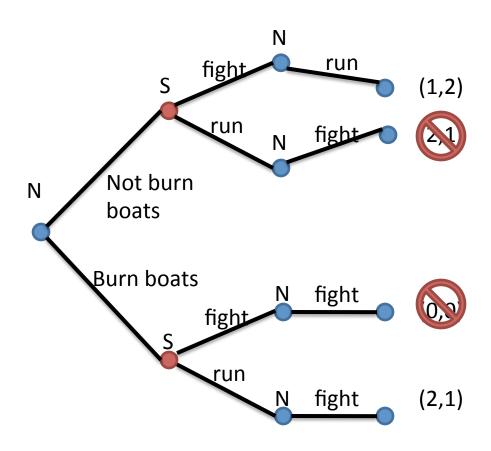
#### **Backward Induction tells us:**

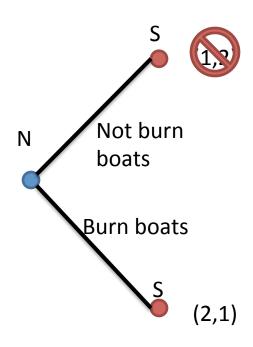
- Saxons will fight
- Normans will run away

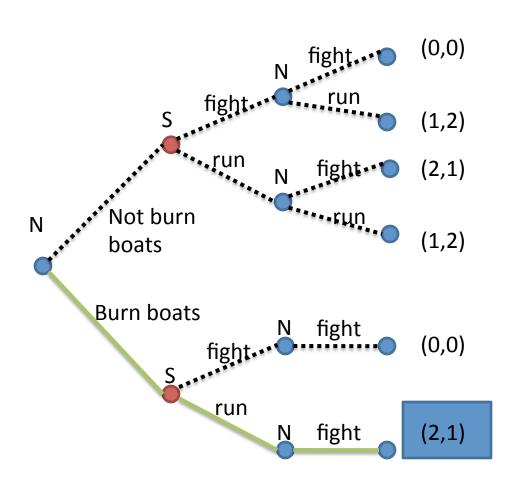


What did William the Conqueror do?









#### Commitment

 Sometimes, getting rid of choices can make me better off!

#### • **Commitment**:

- Fewer options change the behavior of others
- The other players <u>must know</u> about your commitments
  - Example: Dr. Strangelove movie