

# Game Theory

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## Lecture 4

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# Lecture 2-3 recap

- Proved existence of pure strategy Nash equilibrium in games with compact convex action sets and continuous concave utilities
  - Defined mixed strategy Nash equilibrium
  - Proved existence of mixed strategy Nash equilibrium in finite games
  - Discussed computation and interpretation of mixed strategies Nash equilibrium
- Nash equilibrium is not the only solution concept
- Today: Another solution concept: evolutionary stable strategies

# Outline

- Evolutionary stable strategies

# Evolutionary game theory

- Game theory  $\leftrightarrow$  evolutionary biology
- Idea:
  - *Relate strategies to phenotypes of genes*
  - *Relate payoffs to genetic fitness*
  - Strategies that do well “grow”, those that obtain lower payoffs “die out”
- Important note:
  - Strategies are **hardwired**, they are not chosen by players
- Assumptions:
  - Within species competition: no mixture of population

# Examples

- Using game theory to understand population dynamics
  - Evolution of species
  - Groups of lions deciding whether to attack in group an antelope
  - Ants deciding to respond to an attack of a spider
  - TCP variants, P2P applications
- Using evolution to interpret economic actions
  - Firms in a competitive market
  - Firms are bounded, they can't compute the best response, but have rules of thumbs and adopt hardwired (consistent) strategies
  - Survival of the fittest == rise of firms with low costs and high profits

# A simple model

- Assume simple game: two-player symmetric
- Assume **random tournaments**
  - Large population of individuals with hardwired strategies, pick two individuals at random and make them play the symmetric game
  - The player adopting the strategy yielding higher payoff will survive (and eventually gain new elements) whereas the player who “lost” the game will “die out”
- Start with entire population playing strategy  $s$
- Then introduce a **mutation**: a ***small*** group of individuals start playing strategy  $s'$
- Question: will the mutants survive and grow or die out?

# A simple example (1)

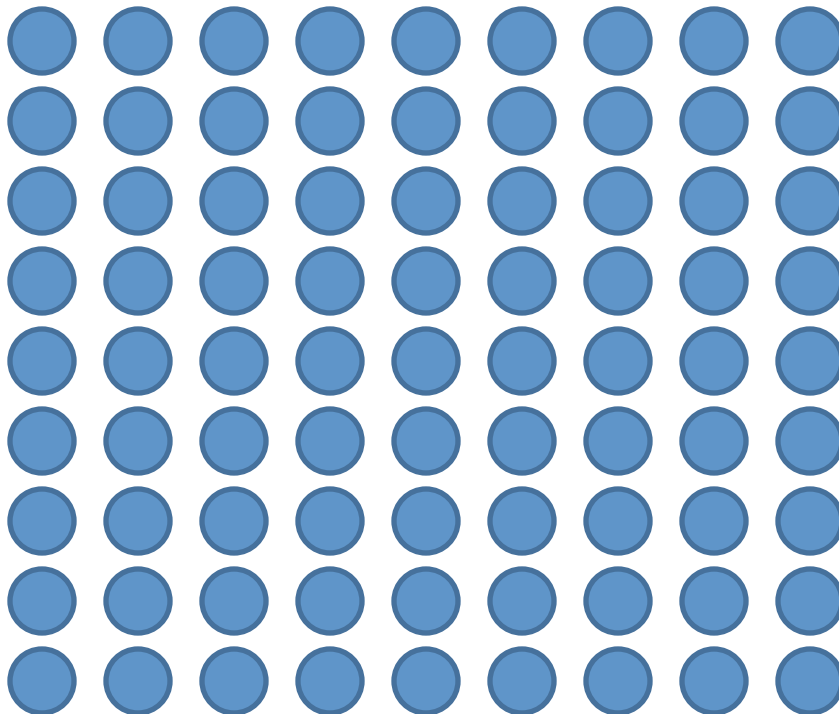
		Player 2	
		Cooperate	Defect
Player 1	C	2,2	0,3
	D	3,0	1,1
		$\varepsilon$	$1 - \varepsilon$

- Have you already seen this game?
- Examples:
  - Lions hunting in a cooperative group
  - Ants defending the nest in a cooperative group
- Question: ***is cooperation evolutionary stable?***

# A simple example (2)



Player strategy  
hardwired → C



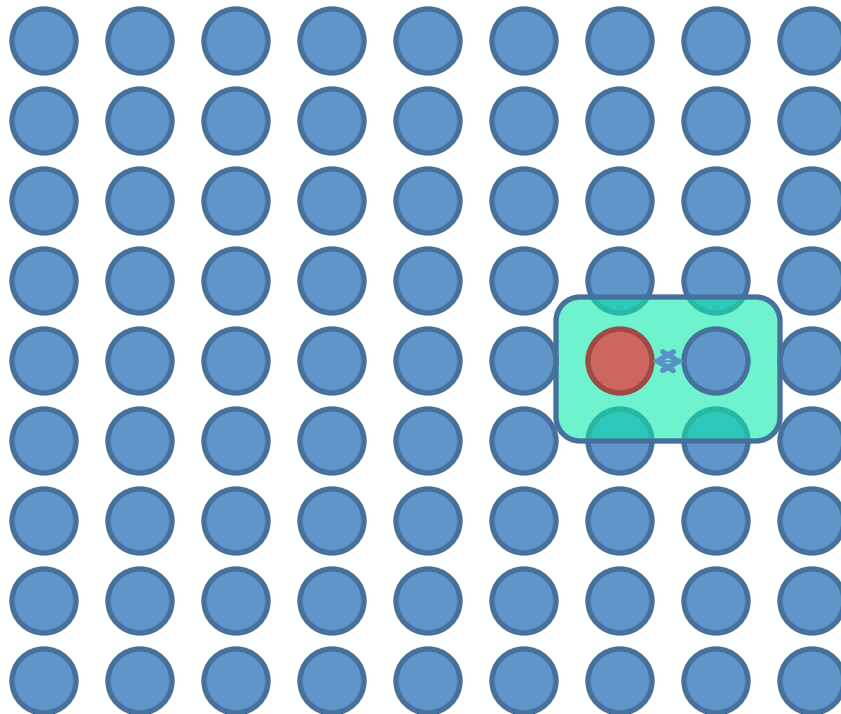
## ***“Spatial Game”***

All players are cooperative  
and get a payoff of 2

What happens with a  
mutation?



# A simple example (3)



 Player strategy  
hardwired  $\rightarrow$  C

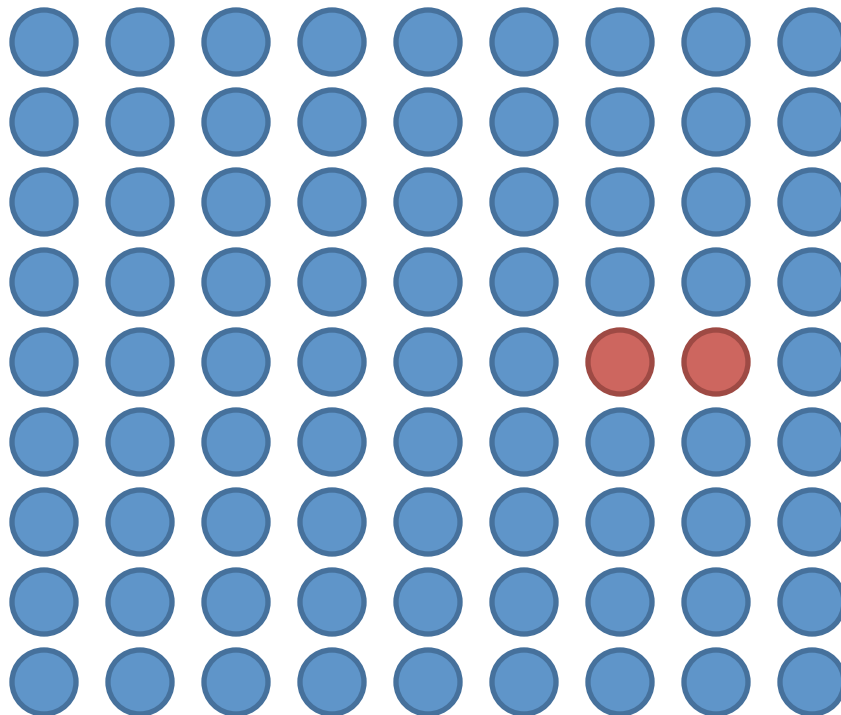
 Player strategy  
hardwired  $\rightarrow$  D

Focus your attention on this  
random “tournament”:

- Cooperating player will obtain a payoff of 0
- Defecting player will obtain a payoff of 3

Survival of the fittest:  
D wins over C

# A simple example (4)

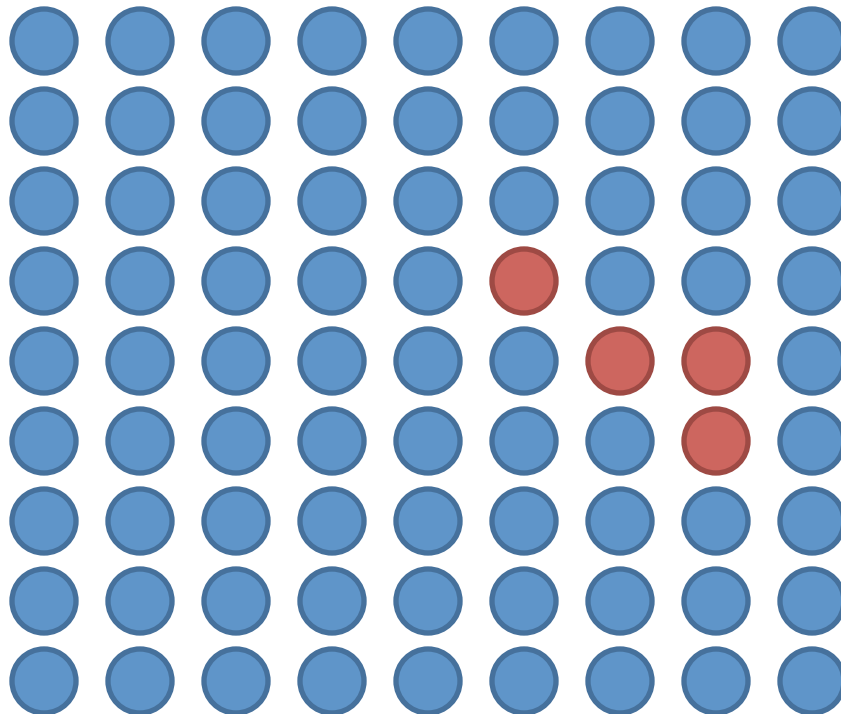


Player strategy  
hardwired → C



Player strategy  
hardwired → D

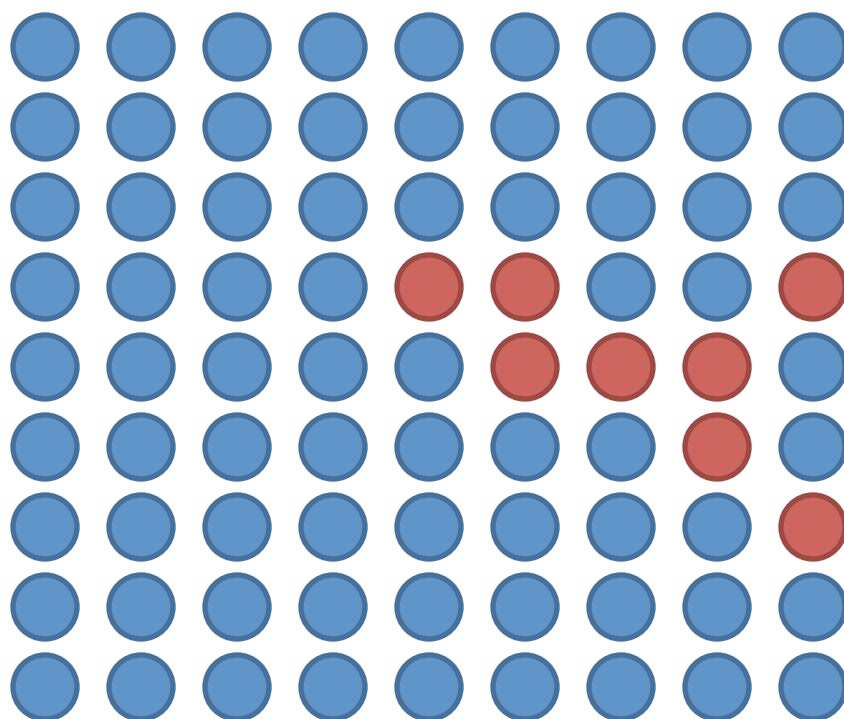
# A simple example (5)



● Player strategy  
hardwired → C

● Player strategy  
hardwired → D

# A simple example (6)



- Player strategy hardwired → C
- Player strategy hardwired → D

A small initial mutation is rapidly expanding instead of dying out

Eventually, C will die out

→ Conclusion: **C is not ES**

Remark: we have assumed asexual reproduction and no gene redistribution

# ESS Definition 1 [Maynard Smith 1972]

## Definition 1: Evolutionary stable strategy

In a symmetric 2-player game, the pure strategy  $\hat{s}$  is ES (in pure strategies) if there exists  $\varepsilon_0 > 0$  such that:

$$\underbrace{(1 - \varepsilon)[u(\hat{s}, \hat{s})] + \varepsilon[u(\hat{s}, s')]}_{\text{Payoff to ES } \hat{s}} > \underbrace{(1 - \varepsilon)[u(s', \hat{s})] + \varepsilon[u(s', s')]}_{\text{Payoff to mutant } s'}$$

for all possible deviations  $s'$  and for all mutation sizes  $\varepsilon < \varepsilon_0$ .

# ES strategies in the simple example

		Player 2	
		Cooperate	Defect
Player 1	C	2,2	0,3
	D	3,0	1,1
		1- $\epsilon$	$\epsilon$
		$\epsilon$	1- $\epsilon$

For C being a majority  
For D being a majority

- Is cooperation ES?

$$C \text{ vs. } [(1-\epsilon)C + \epsilon D] \rightarrow (1-\epsilon)2 + \epsilon 0 = 2(1-\epsilon)$$

$$D \text{ vs. } [(1-\epsilon)C + \epsilon D] \rightarrow (1-\epsilon)3 + \epsilon 1 = 3(1-\epsilon) + \epsilon$$

$$3(1-\epsilon) + \epsilon > 2(1-\epsilon)$$

→ **C is not ES** because the average payoff to C is lower than the average payoff to D

→ **A strictly dominated is never Evolutionarily Stable**  
– The strictly dominant strategy will be a successful mutation

# ES strategies in the simple example

		Player 2	
		Cooperate	Defect
Player 1	C	2,2	0,3
	D	3,0	1,1
		1- $\epsilon$	$\epsilon$
		$\epsilon$	1- $\epsilon$
		For C being a majority	
		For D being a majority	

- Is defection ES?

$$D \text{ vs. } [\epsilon C + (1-\epsilon)D] \rightarrow (1-\epsilon)1 + \epsilon 3 = (1-\epsilon) + 3\epsilon$$

$$C \text{ vs. } [\epsilon C + (1-\epsilon)D] \rightarrow (1-\epsilon)0 + \epsilon 2 = 2\epsilon$$

$$(1-\epsilon) + 3 > 2\epsilon$$

➔ **D is ES**: any mutation from D gets wiped out!

# Another example (1)

	a	b	c
a	2,2	0,0	0,0
b	0,0	0,0	1,1
c	0,0	1,1	0,0

- 2-players symmetric game with 3 strategies
- Is “c” ES?      c vs.  $[(1-\epsilon)c + \epsilon b] \rightarrow (1-\epsilon) 0 + \epsilon 1 = \epsilon$   
                          b vs.  $[(1-\epsilon)c + \epsilon b] \rightarrow (1-\epsilon) 1 + \epsilon 0 = 1 - \epsilon > \epsilon$

➔ “c” is not evolutionary stable, as “b” can invade it

- Note: “b”, the invader, is itself not ES!
  - It is not necessarily true that an invading strategy must itself be ES
  - But it still avoids dying out completely (grows to 50% here)



# Another example (3)

	a	b	c
a	2,2	0,0	0,0
b	0,0	0,0	1,1
c	0,0	1,1	0,0

- Is  $(c,c)$  a NE?

# Observation

- If  $s$  is **not Nash** (that is  $(s,s)$  is not a NE), then  $s$  is **not evolutionary stable** (ES)

Equivalently:

- If  $s$  is **ES**, then  $(s,s)$  is a **NE**

- Question: is the opposite true? That is:
  - If  $(s,s)$  is a **NE**, then  $s$  is **ES**

# Yet another example (1)

		Player 2	
		a	b
Player 1	a	1,1	0,0
	b	0,0	0,0
		$\varepsilon$	$1 - \varepsilon$

- NE of this game:  $(a,a)$  and  $(b,b)$
- Is  $b$  ES?
  - $b \rightarrow 0$
  - $a \rightarrow (1-\varepsilon) 0 + \varepsilon 1 = \varepsilon > 0$

→  $(b,b)$  is a NE, but it is not ES!

- This relates to the idea of a weak NE

→ If  $(s,s)$  is a **strict NE** then  $s$  is ES

# Strict Nash equilibrium

## Definition: Strict Nash equilibrium

A strategy profile  $(s_1^*, s_2^*, \dots, s_N^*)$  is a strict Nash Equilibrium if, for each player  $i$ ,

$$u_i(s_i^*, s_{-i}^*) > u_i(s_i, s_{-i}^*) \text{ for all } s_i \neq s_i^*$$

- Weak NE: the inequality is an equality for at least one alternative strategy
- Strict NE is sufficient but not necessary for ES

# ESS Definition 2

## Definition 2: Evolutionary stable strategy

In a symmetric 2-player game, the pure strategy  $\hat{s}$  is ES (in pure strategies) if:

A)  $(\hat{s}, \hat{s})$  is a symmetric Nash Equilibrium  
$$u(\hat{s}, \hat{s}) \geq u(s', \hat{s}) \quad \forall s'$$

**AND**

B) if  $u(\hat{s}, \hat{s}) = u(s', \hat{s})$  then  
$$u(\hat{s}, s') > u(s', s')$$

# Link between definitions 1 and 2

## Theorem

Definition 1  $\Leftrightarrow$  Definition 2

- Proof sketch:

# Recap: checking for ES strategies

- We have seen a definition that connects Evolutionary Stability to Nash Equilibrium
- By def 2, to check that  $\hat{s}$  is ES, we need to do:
  - First check if  $(\hat{s}, \hat{s})$  is a **symmetric** Nash Equilibrium
  - If it is a **strict** NE, we're done
  - Otherwise, we need to compare how  $\hat{s}$  performs against a mutation, and how a mutation performs against a mutation
  - If  $\hat{s}$  performs better, then we're done

# Example: Is “a” evolutionary stable?

		Player 2	
		a	b
Player 1	a	1,1	1,1
	b	1,1	0,0
		$\epsilon$	$1 - \epsilon$

- Is (a, a) a NE? Is it strict?
- Is “a” evolutionary stable?



# Evolution of social convention

- Evolution is often applied to social sciences
- Let's have a look at how driving to the left or right hand side of the road might evolve

	L	R
L	2,2	0,0
R	0,0	1,1

- What are the NE? are they strict? What are the ESS?
- Conclusion: we can have several ESS
  - They need not be equally good

# The game of Chicken

	a	b
a	0,0	2,1
b	1,2	0,0

- This is a **symmetric coordination game**
- Biology interpretation:
  - “a” : individuals that are aggressive
  - “b” : individuals that are non-aggressive
- What are the pure strategy NE?
  - They are not symmetric → no candidate for ESS

# The game of Chicken: mixed strategy

		NE	
		a	b
a		0,0	2,1
b		1,2	0,0

- What's the mixed strategy NE of this game?
  - Mixed strategy NE =  $[(2/3, 1/3), (2/3, 1/3)]$
  - This is a **symmetric** Nash Equilibrium
- Interpretation: there is an equilibrium in which 2/3 of the genes are aggressive and 1/3 are non-aggressive
- Is it a strict Nash equilibrium?
- Is it an ESS?

# Remark

- A mixed-strategy Nash equilibrium (with a support of at least 2 actions for one of the players) can never be a strict Nash equilibrium
- The definition of ESS is the same!

# ESS Definition 2bis

## Definition 2: Evolutionary stable strategy

In a symmetric 2-player game, the mixed strategy  $\hat{s}$  is ES (in mixed strategies) if:

A)  $(\hat{s}, \hat{s})$  is a symmetric Nash Equilibrium  
$$u(\hat{s}, \hat{s}) \geq u(s', \hat{s}) \quad \forall s'$$

**AND**

B) if  $u(\hat{s}, \hat{s}) = u(s', \hat{s})$  then  
$$u(\hat{s}, s') > u(s', s')$$

# The game of Chicken: ESS

	a	b
a	0,0	2,1
b	1,2	0,0

- Mixed strategy NE =  $[(2/3, 1/3), (2/3, 1/3)]$ .
- Is it an ESS? we need to check for all possible mixed mutations  $s'$ :  $u(\hat{s}, s') > u(s', s') \quad \forall s' \neq \hat{s}$
- Yes, it is (do it at home!)
- In many cases that arise in nature, the only equilibrium is a mixed equilibrium
  - It could mean that the gene itself is randomizing, which is plausible
  - It could be that there are actually two types surviving in the population (cf. our interpretation of mixed strategies)

# Hawks and doves

Dove



Hawk



# The Hawks and Dove game (1)

	H	D
H	$(v-c)/2, (v-c)/2$	$v, 0$
D	$0, v$	$v/2, v/2$

- More general game of aggression vs. non-aggression
  - The prize is food, and its value is  $v > 0$
  - There's a cost for fighting, which is  $c > 0$
- Note: we're still in the context of **within species competition**
  - So it's not a battle against two different animals, hawks and doves, we talk about strategies
    - "Act dovish vs. act hawkish"
- What are the ESS? How do they change with  $c, v$ ?



# The Hawks and Dove game (2)

	H	D
H	$(v-c)/2, (v-c)/2$	$v, 0$
D	$0, v$	$v/2, v/2$

- Can we have a ES population of doves?
- Is  $(D,D)$  a NE?
  - No, hence “D” is not ESS
  - Indeed, a mutation of hawks against doves would be profitable in that it would obtain a payoff of  $v$

# The Hawks and Dove game (3)

	H	D
H	$(v-c)/2, (v-c)/2$	$v, 0$
D	$0, v$	$v/2, v/2$

- Can we have a ES population of Hawks?
  - Is  $(H,H)$  a NE? It depends: it is a symmetric NE if  $(v-c)/2 \geq 0$
  - Case 1:  $v > c \rightarrow (H,H)$  is a **strict** NE  $\rightarrow$  “H” is ESS
  - Case 2:  $v = c \rightarrow (v-c)/2 = 0 \rightarrow u(H,H) = u(D,H)$  --  $(H, H)$  is a weak NE
    - Is  $u(H,D) = v$  larger than  $u(D,D) = v/2$ ? Yes  $\rightarrow$  “H” is ESS
- $\rightarrow$  H is ESS if  $v \geq c$
- If the prize is high and the cost for fighting is low, then you’ll see fights arising in nature

# The Hawks and Dove game (4)

	H	D
H	$(v-c)/2, (v-c)/2$	$v, 0$
D	$0, v$	$v/2, v/2$
	$\hat{s}$	$1 - \hat{s}$

- What if  $c > v$ ?
  - “H” is not ESS and “D” is not ESS (they are not NE)
- Step 1: find a mixed NE
- Step 2: verify the ESS condition

# The Hawks and Dove game: results

- In case  $v < c$  we have an evolutionarily stable state in which we have  $v/c$  hawks
  1. As  $v \nearrow$  we will have more hawks in ESS
  2. As  $c \nearrow$  we will have more doves in ESS
- By measuring the proportion of H and D, we can get the value of  $v/c$
- Payoff: 
$$E[u(D, \hat{s})] = E[u(H, \hat{s})] = 0 \frac{v}{c} + \left(1 - \frac{v}{c}\right) \frac{v}{2}$$

# One last example (1)

	R	P	S
R	1,1	$v,0$	$0,v$
P	$0,v$	1,1	$v,0$
S	$v,0$	$0,v$	1,1

- Assume  $1 < v < 2$ 
  - ~ Rock, paper, scissors
- Only NE:  $\hat{s} = (1/3, 1/3, 1/3)$  – mixed, not strict
- Is it an ESS?
  - Suppose  $s' = R$
  - $u(\hat{s}, R) = (1+v)/3 < 1$
  - $u(R, R) = 1$
- Conclusion: Not all games have an ESS!