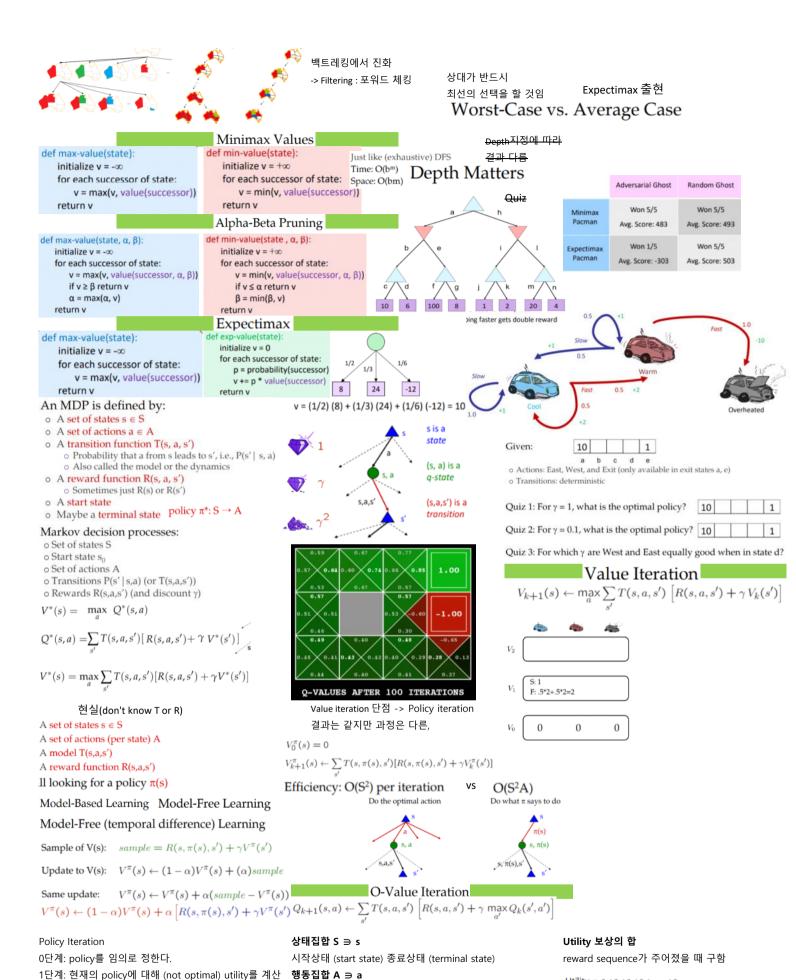
Problem이 무엇인지에 따라 2018년 4월 17일 화요일 오후 8:19 달라지는 것을 볼 수 있다. Optimal vs. complete planning Depth-First Search A search state keeps only the details needed for planning (abstraction) Breadth-First Search Planning vs. replanning o Problem: Pathing o Problem: Eat-All-Dots Uniform-Cost Search What are the states? o States: (x,y) location o States: {(x,y), dot booleans} Search Problems How many states? What are the actions? o Actions: NSEW o Actions: NSEW 용어미 o Successor: update location and possibly a dot boolean o Successor: update location How many successors from the start state? only What should the costs be? o A search problem consists of: o Goal test: is (x,y)=END o Goal test: dots all false o A state space **III** State space: o Citie o A successor function Successor function: Roads: Go to adjacent city with cost = distance (with actions, costs) example Start state: o A start state and a goal test o Arad Goal test: o A solution is a sequence of actions (a plan) o Is state == Bucharest? which transforms the start state to a goal state Solution? Search Trees and State Space Graphs Depth-First Search Breadth-First Search Uniform Cost Search never expand a state twice Tree A* Combining UCS and Greedy f(n) = g(n) + h(n)Greedy A* Graph Search 노드 확장 순서, 결과, 남아있는 노드 Consistency of Heuristics Uniform-cost $h(A) - h(C) \le cost(A \text{ to } C)$ Admissible Heuristics $0 \le h(n) \le h^*(n)$ function GRAPH-SEARCH(problem, fringe) return a solution, or failure $closed \leftarrow$ an empty set $fringe \leftarrow Insert(Make-Node(Initial-State[problem]), fringe)$ optimal No, optimal Only if costs are all 1 loop do if fringe is empty then return failure O(bs) O(bm) $O(b^{C*/\epsilon})$ $ode \leftarrow REMOVE\text{-}FRONT(fringe)$, takes time if GOAL-TEST(problem, STATE[node]) then return node space does the fringe take if STATE[node] is not in closed then
add STATE[node] to closed
for child-node in EXPAND(STATE[node], problem) do
fringe ← INSERT(child-node, fringe) $O(b^{C^*/\varepsilon})$ $O(b^s)$ LIFO stack FIFO queue priority queue cheapest node first: function TREE-SEARCH(problem, fringe) return a solution, or failure. No information about goal location ← Insert(make-node(initial-state[problem]), fringe) ->단점, 방향성X loop do greedy Object if fringe is empty then return failure $node \leftarrow REMOVE-FRONT(fringe)$ if GOAL-TEST(problem, STATE[node]) then return node for child-node in EXPAND(STATE[node], problem) do $fringe \leftarrow INSERT(child-node, fringe)$ end UCS and Constraint Satisfaction Problems Variables: WA, NT, Q, NSW, V, SA, T Variables: Q_k Domains: D = {red, green, blue} $\underline{\Psi}$ Domains: $\{1, 2, 3, ... N\}$ QUIZ > Constraints: adjacent regions must have Starting from X, where do you end up? different colors Constraints: Starting from Y, where do you end up? Backtracking Search 백트레킹:더 이상 전진불가 Hill Climbing Diagram Starting from Z, where do you end up? Idea 1: One variable at a time Idea 2: Check constraints as you go 풀기 Arc Consistency in a CSP Forward checking -> efficient Ordering: Minimum Remaining Values Least Constraining Value Structure - turns out trees are easy? BFS 아주 비효율 DFS vs. Backtracking 백트레킹에서 진화

-> Filtering : 포워드 체킹

상대가 반드시

최선의 선택을 할 것임

Expectimax 출현



전이함수(transition function) T(s, a, s') = P(s' | s, a)

보상함수(reward function) R(s, a, s')

현재 상태 s에서 행동 a를 했을 때, 다음상태가 s'이 될 확률

Utility = $r_0+r_1+r_2+r_3+ \cdots + r_{end}$

Utility = $\gamma^{0}r_{0} + \gamma^{1}r_{1} + \gamma^{2}r_{2} + \gamma^{3}r_{3} + \cdots + \gamma^{end}r_{end}$

인공지능 페이지 2

2단계: 계산한 utility를 이용하여 policy를 개선

개선한 policy를 현재의 policy로 한다.

3단계: 1~2단계를 policy가 수렴할 때까지 반복한다.