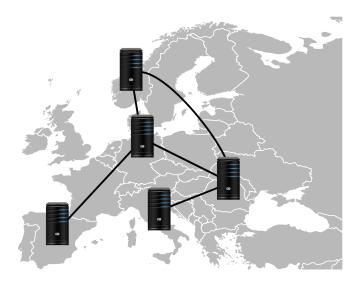
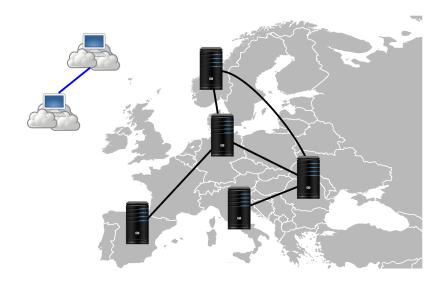
Specificity vs. Flexibility: On the Embedding Cost of a Virtual Network

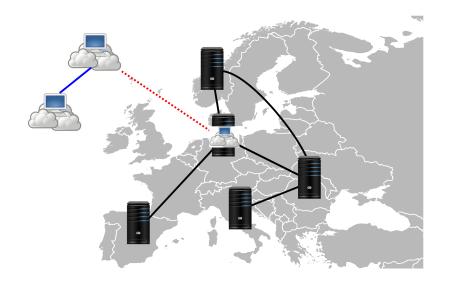
Arne Ludwig, Stefan Schmid, Anja Feldmann

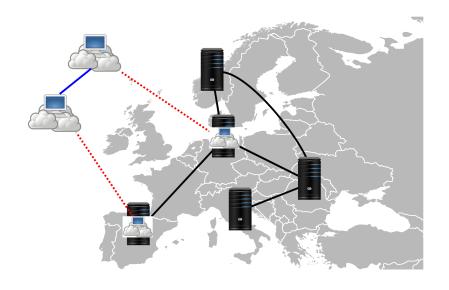
Telekom Innovation Laboratories & TU Berlin, Germany {arne,stefan,anja}@net.t-labs.tu-berlin.de

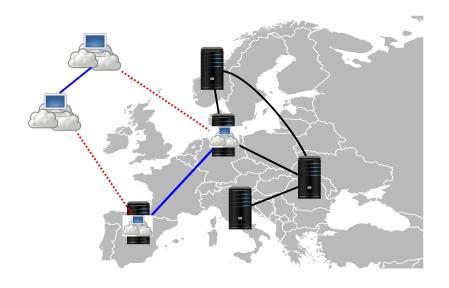
September 14, 2013

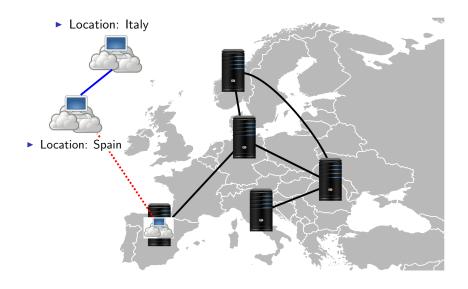


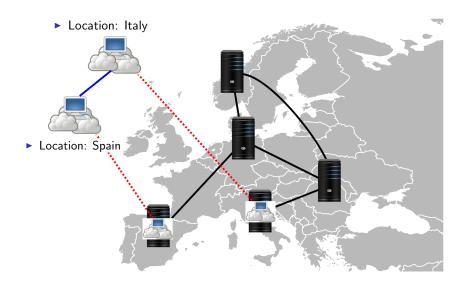


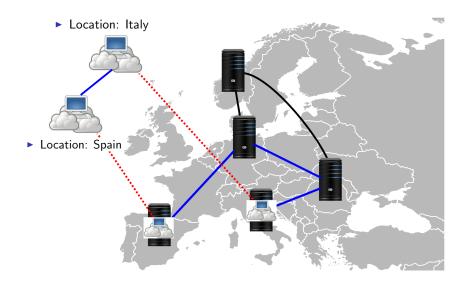




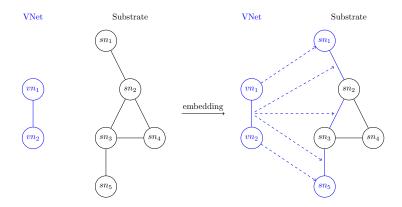








Standard Problem



How can flexibility be quantified?

Dimensions of Specifications - Properties

Focus on node properties

- Different properties:
 - Location
 - Virtualization technology
 - Operating system
 - etc.
- Substrate nodes have exactly one type per property
 - VNet node requests can specify multiple types
- All properties combined form a configuration
 - Example: {Italy, Xen, RedHat 7.3}

Specificity - Definition

Specificity $\sigma =$ percentage of lost alternatives

$$(\Rightarrow \sigma = \frac{\text{forbidden configurations}}{\text{all configurations -1}})$$



 $ightharpoonup \sigma = 0$: free choice of nodes



 $\sigma = 1$: only nodes with exactly defined types

VNet specificity: average specificity of its nodes

Price of Specificity (PoS) - Definition

- ▶ Cost_{σ}: cost under a given specificity $\sigma(VNet)$
- ► Cost₀: cost without specification constraints

Price of Specificity definition:

$$PoS = Cost_{\sigma}/Cost_{0}$$

PoS - Example



• Spec.: Spain + Italy ($\sigma = 1$)



▶ No specification ($\sigma = 0$)

PoS - Example

Cost metric: Number of hops



- ▶ Spec.: Spain + Italy $(\sigma = 1)$
- ▶ 3 hops



- ▶ No specification $(\sigma = 0)$
- ▶ 1 hop

PoS - Example

Cost metric: Number of hops



- ▶ Spec.: Spain + Italy ($\sigma = 1$)
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- ▶ 1 hop

$$PoS = \frac{3}{1} = 3$$

Overview

Introduction

- ► Embedding problem
- ► Specification, *PoS*

Overview

Introduction

- ► Embedding problem
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Upcoming

- ► Embedding algorithm
- ▶ Impact of different factors on the PoS

Optimal Algorithm

Constants:

```
\begin{array}{lll} \text{Substrate Vertices}: V_s & \text{Virtual Vertices}: V_r(r), r \in R \\ \text{Substrate Edges}: E_s : V_s \times V_s & \text{Virtual Edges}: E_s(r) : \rightarrow V_r(r) \times V_r(r), r \in R \\ \text{Unique}: uni\_checks, : V_s(s_1, s_2) \in E_s : (s_2, s_1) \not\in E_s \\ \text{SNode Capacity}: snc(s) \rightarrow \mathbb{R}^+, s \in V_s & \text{Unique}: uni\_checks, : V_r \in R, (v_1, v_2) \in E_r(r) : (v_2, v_1) \not\in E_r(r) \\ \text{Node Demand}: vnd(r, v) \rightarrow \mathbb{R}^+, r \in R, v \in V_r(r) \\ \text{VEdge Demand}: vld(r, e_r) \rightarrow \mathbb{R}^+, r \in R, e_r \in E_r(r) \\ \text{Edges-Reverse}: ER_s : \forall (s_1, s_2) \in E_s \exists (s_2, s_1) \in ER_s \wedge |E_s| = |ER_s| \\ \text{Migration Cost}: mig\_cost(r, v, s) \rightarrow \mathbb{R}^+ |V_r(r)| \times |V_r|, r \in R, v \in V_r(r), s \in V_s \\ \text{Possible Placements}: place(r, v, s) \rightarrow \{0, 1\}^{|V_r(r)| \times |V_r|}, r \in R, v \in V_r(r), s \in V_s \\ \end{array}
```

Requests : R

Variables:

```
Node Mapping : n\_map(r,v,s) \in \{0,1\}, r \in R, v \in V_v(r), s \in V_s
Flow Allocation : f\_alloc(r,e,eb) \geq 0, r \in R, e \in E_v(r), eb \in EB_s
```

Constraints:

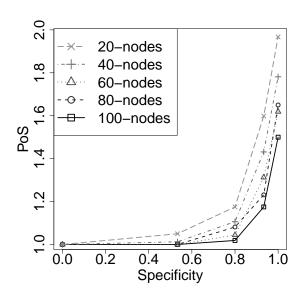
```
 \begin{array}{l} \text{Each Node Mapped}: \forall r \in R, v \in V_{v}(r): \sum_{s \in V_{v}} n\_map(r, v, s) \cdot place(r, v, s) = 1 \\ \text{Feasible}: \forall s \in V_{v}: \sum_{r \in R_{v} \in V_{v}(r)} n\_map(r, v, s) \cdot vnd(r, v) \leq snc(s) \\ \text{Guarantee Link Realization}: \forall r \in R, (v_{v}, v_{2}) \in E_{v}(r), s \in V_{s} \sum_{(s, s_{2}) \in V_{v} \times V_{v} \cap EB_{s}} f\_alloc(r, v_{1}, v_{2}, s, s_{2}) - \sum_{(s_{1}, s) \in V_{s} \times V_{v} \cap EB_{s}} f\_alloc(r, v_{1}, v_{2}, s, s_{2}) = vld(r, v_{1}, v_{2}) \cdot (n\_map(r, v_{1}, s) - n\_map(r, v_{2}, s)) \\ \text{Realize Flows}: \forall (s_{1}, s_{2}) \in E_{s} \sum_{r \in R_{v}(v_{1}, v_{2})} f\_alloc(r, v_{1}, v_{2}, s_{1}, s_{2}) + f\_alloc(r, v_{1}, v_{2}, s_{2}, s_{1}) \leq slc(s_{1}, s_{2}) \\ \end{array}
```

Objective function:

```
\text{Minimize Embedding Cost}: min: \sum\nolimits_{r \in R, (v_1, v_2) \in E_r(r), (s_1, s_2) \in E_s} f\_alloc(r, v_1, v_2, s_1, s_2) + f\_alloc(r, v_1, v_2, s_2, s_1)
```

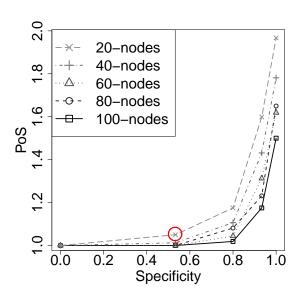
Optimal Algorithm

MIP (Mixed-integer program) Objective function: Minimize Link Cost Constraints to ensure feasibility Migration Different types of links Optimal embeddings



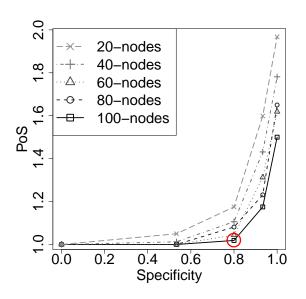
- 5-star VNet
- Node capacity of one
- ► Substrates created with Igen topology generator





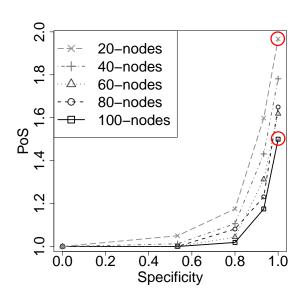
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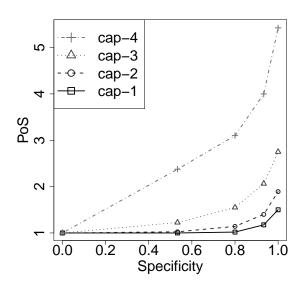




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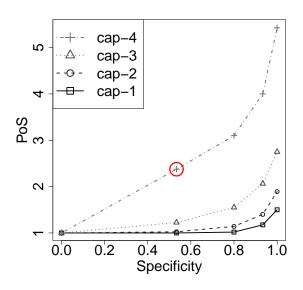
Impact of Node Capacity



- ▶ 5-star VNet
- Colocation allowed
- ➤ 100 nodes substrate created with *Igen*



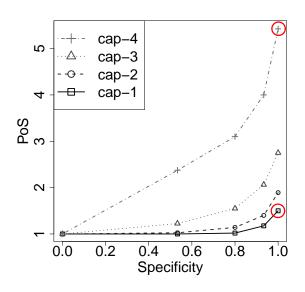
Impact of Node Capacity



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Impact of Node Capacity

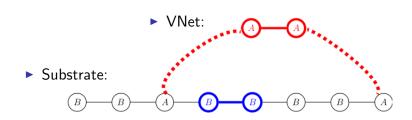


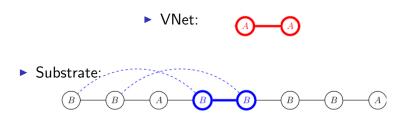
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Substrate:







► VNet:



Substrate:



► VNet:

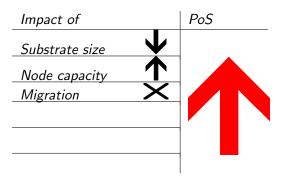


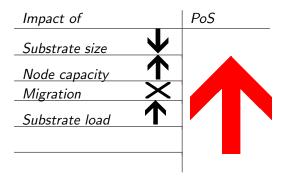
Substrate:

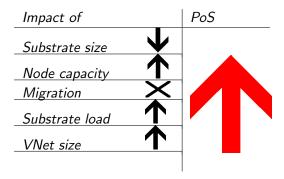


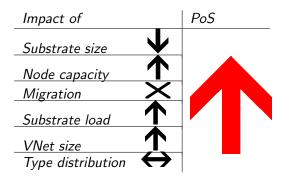
Impact of Migration

- Migration lowers average resource cost in general
- Depends on access policy
- Various impacts on PoS









Conclusion

- Impact of VNet specification on the embedding cost
- Optimal embeddings
- General embedding algorithm
- PoS, tool to adjust pricing and embedding (applied as a factor?)

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- Impact of VNet specification on the embedding cost
- Optimal embeddings
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- Prototype (open source)
- Specification language FLERD
- Project website*

^{*}www.net.t-labs.tu-berlin.de/~stefan/virtu.shtml