



## 2017 Winter midterm

Introduction to Theoretical Computer Science (Concordia University)



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DEPARTMENT OF COMPUTER SCIENCE & SOFTWARE ENGINEERING  
COMP335 INTRODUCTION TO THEORETICAL COMPUTER SCIENCE  
WINTER 2017

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Midterm Exam March 2

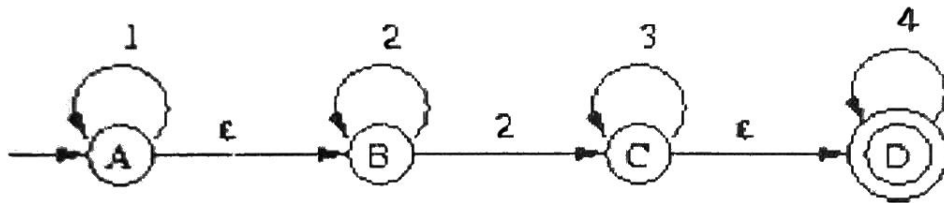
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- Your Name:
  - Your Student ID:
  - Your Signature:
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**Instructions:**

- There are 10 multiple-choice questions (Questions 1 – 10, each worth 3 points) and one proof question (Question 11, worth 9 points).
  - Answer questions 1 – 10 on the scan sheet.  
NOTE: Only the scan sheet will be graded. Anything written in the booklet will be ignored.
  - Answer Question 11 in the three boxes on page 6.  
NOTE: Only answers written in the three boxes will be graded. Anything else written in the booklet will be ignored.
  - Use provided scrap paper for your rough work. DO NOT hand in the scrap paper!
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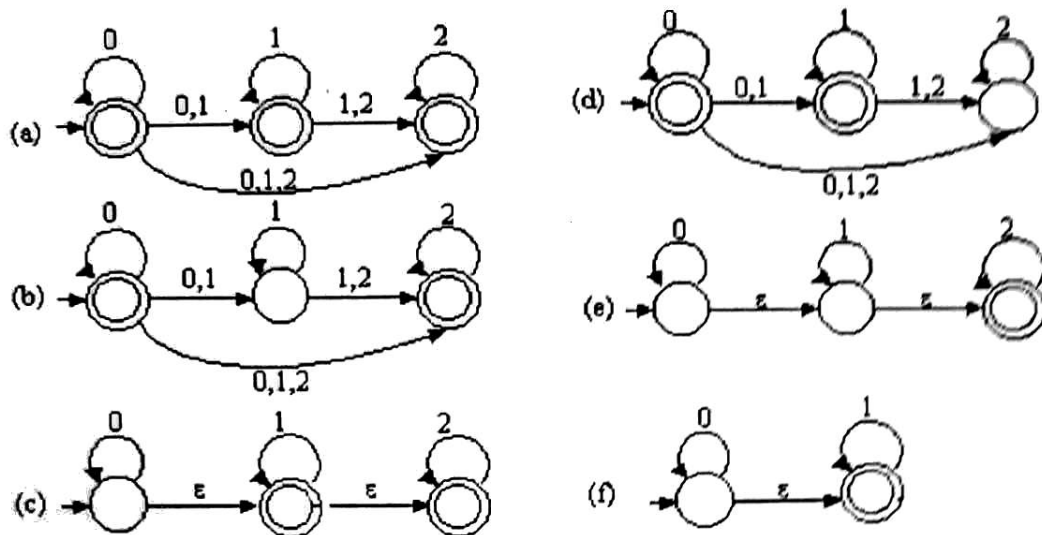
1. Let  $L_1 = \{a^i : i \geq 0\}$  and  $L_2 = \{a^i b^j : j \geq i \geq 0\}$ . Then the language  $L_1 L_2$  is
  - (a)  $\{a^{2i} b^j : i, j \geq 0\}$
  - (b)  $\{a^i b^j : i \geq j\}$
  - (c)  $\{a^i b^j : i, j \geq 0\}$
  - (d)  $\{a^i b^j : j \geq i\}$
2. Which of the following strings is NOT in the Kleene closure of the language  $\{011, 10, 110\}$ 
  - (a) 10111011
  - (b) 10110011
  - (c) 1001110
  - (d) 1010110
3. Consider the following  $\epsilon$ -NFA



When you convert this  $\epsilon$ -NFA to a DFA, which of the following would be a transition of the DFA?

- (a)  $\delta(\{A, B\}, 1) = \{A, B\}$
- (b)  $\delta(\{A, B\}, 1) = \{A\}$
- (c)  $\delta(\{A, B\}, 3) = \{C, D\}$
- (d)  $\delta(\{A, B\}, 2) = \{A, B, C\}$

4. Consider the following six  $\epsilon$ -NFA's



Which of these accept the same language?

- (a) (a) and (f)
- (b) (a) and (c)
- (c) (c) and (d)
- (d) (b) and (f)

5. Consider the two regular expressions

$$R = 0^* + 1^* \quad S = 01^* + 10^* + 1^*0 + (0^*1)^*$$

Then, consider the languages

$$L_1 = L(R) \setminus L(S), \quad L_2 = L(S) \setminus L(R), \quad L_3 = L(R) \cap L(S), \quad L_4 = \overline{L(R) \cup L(S)}.$$

and the strings

$$w_1 = 011, \quad w_2 = 111, \quad w_3 = 000, \quad w_4 = 1100$$

Which of the following is correct?

- (a)  $w_1 \in L_1, w_2 \in L_2, w_3 \in L_3, w_4 \in L_4$
- (b)  $w_3 \in L_1, w_1 \in L_2, w_2 \in L_3, w_4 \in L_4$
- (c)  $w_4 \in L_1, w_2 \in L_2, w_1 \in L_3, w_3 \in L_4$
- (d)  $w_2 \in L_1, w_1 \in L_2, w_4 \in L_3, w_3 \in L_4$

6. How many strings of length less than 4 is contained in  $L((a+b)^*b(c+cd)^*)$
- 10
  - 11
  - 12
  - 13
7. Which of the following regular expressions defines the complement of the language  $L((0+10)^*)$
- $(0+1)^*(11+1+\epsilon)^*$
  - $(0+1)^*(1+11)(0+1)^*$
  - $(0+10)^*1(\epsilon+11(0+1)^*)$
  - $(0+1)^*11(0+1)^*+(0+10)^*1$
8. Let  $h$  be a homomorphism from  $\{a, b, c\}$  to  $\{0, 1\}$ , where  $h(a) = 01$ ,  $h(b) = 0$ , and  $h(c) = 10$ . Which of the following strings is in  $h^{-1}(010010)$ .
- $bcab$
  - $abcb$
  - $bcba$
  - $abac$
9. Let  $A = (Q, \Sigma, \delta, q_0, \{q_f\})$  be an  $\epsilon$ -NFA that accepts language  $L(A)$ . Consider the following modifications of  $A$ .
- The automaton  $B$  constructed from  $A$  by adding  $\epsilon$ -transitions from  $q_0$  to every state, for which there is a path in  $A$  from  $q_0$  to that state.
  - The automaton  $C$  constructed from  $A$  by adding  $\epsilon$ -transitions to  $q_f$  from every state, for which there is a path in  $A$  to  $q_f$  from that state.
  - The automaton  $D$  constructed from  $A$  by doing both of the above modifications.

Here are three candidate languages:

- $L_1 = \{x : xy \in L(A) \text{ for some } y \in \Sigma^*\}$
- $L_2 = \{y : xyz \in L(A) \text{ for some } x, z \in \Sigma^*\}$
- $L_3 = \{y : xy \in L(A) \text{ for some } x \in \Sigma^*\}$

Which of the following is correct?

- $L(B) = L_1, L(C) = L_2, L(D) = L_3$
- $L(B) = L_3, L(C) = L_1, L(D) = L_2$
- $L(B) = L_3, L(C) = L_2, L(D) = L_3$
- $L(B) = L_2, L(C) = L_3, L(D) = L_1$

10. Let  $A$  be the following DFA.

	0	1
$\rightarrow A$	$E$	$B$
$* B$	$D$	$A$
$C$	$G$	$A$
$* D$	$G$	$E$
$E$	$A$	$D$
$F$	$B$	$E$
$* G$	$B$	$A$

When you minimize  $A$  using the table-filling algorithm, the following are the sets of indistinguishable (equivalent) states.

- (a)  $\{A, C\}, \{B, D, G\}, \{E, F\}$
- (b)  $\{A, C, E\}, \{B, D, G\}, \{F\}$
- (c)  $\{A\}, \{B, D, G\}, \{C, E, F\}$
- (d)  $\{A, E\}, \{B, D, G\}, \{C, F\}$

11. Let  $L$  be the language of those strings over  $\{0,1\}$  where the number of 0's differ from the number of 1's by at most 5. Complete the proof below, showing  $L$  is not regular.

**Proof:**

- Suppose to the contrary that  $L$  is regular  
 $\Rightarrow \exists$  DFA  $A$ , s.t.  $L(A) = L$ .
  - Let  $n$  be the number of states in  $A$ .  
 $\Rightarrow \forall w \in L(A)$ , if  $|w| \geq n$ , then  $w = xyz$ ,  
where  $x$ ,  $y$ , and  $z$  as in Pumping Lemma.
- (a) Choose a suitable  $w \in L$ , where  $|w| \geq n$ .

My solution:

$w =$

- (b) Find an  $i$ , such that  $xy^iz \in L(A)$  and  $xy^iz \notin L$ .

My solution:

$i =$

$xy^iz =$

- (c) Reason that  $xy^iz \notin L$

My reason:

Since  $xy^iz \in L(A) \Rightarrow L(A) \neq L$ .

- (e) contradicts (a):  
 $\Rightarrow$  (a) cannot be true  
 $\Rightarrow L$  cannot be regular