

# Bayes Nash equilibrium

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# Introduction

- so far
  - need to look at games of incomplete information (preference aggregation when preferences are private, auctions)
  - under certain assumptions decision makers can be modeled as expected utility maximizers
- still missing
  - how to react to information?
  - strategic interaction under uncertainty

## Bayes' rule: A simple example

- I know someone who lives in Munich. What is the probability that this person is male?
- I know someone who lives in Munich and who is 1.90 m tall. What is the probability that this person is male?
- I know someone who lives in Munich and has green eyes. What is the probability that this person is male?

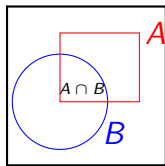
# Bayes' rule

## Bayes' rule

For two events  $A$  and  $B$

$$P(A|B) = \frac{P(B|A)P(A)}{P(B)}.$$

- easier to remember as  $P(A|B)P(B) = P(B|A)P(A)$   
which is also equal to  $P(A \cap B)$
- hence,  $P(A|B) = P(A \cap B)/P(B)$



## Bayes' rule: example

- an antigen test for a certain virus is 70% reliable at detecting an illness and 99.5% reliable at correctly reporting that somebody is healthy
- suppose about 80.000 people are currently infected with the virus
- suppose 80 million people live in Germany
- if a random person is takes a test and the test is positive, what is the probability that this person is infected?

# Bayes' rule: comments

- calculations are reasonably simple
- intuition often goes wrong when the prior is extremely skewed
- make sure to understand it as it will often loom in the background

# Independence

## Independence

Two random variables  $X$  and  $Y$  are independent if

$$P(X = x, Y = y) = P(X = x)P(Y = y).$$

- by Bayes' rule,  $P(X = x|Y = y) = P(X = x)$  if  $X$  and  $Y$  are independent
  - knowing  $Y$  does not affect my belief about  $X$
- independence will often be assumed to keep the models simple

## Games of incomplete information I: an example

- an incumbent decides whether to build a new plant (I for invest) at cost  $c$
- entrant simultaneously decides whether to enter (E)
- entrant does not know whether  $c$  is "low" ( $l$ ) or "high" ( $h$ )

Table: Payoffs with  $c = h$

	E	NE
I	0,-1	2,0
NI	2,1	3,0

Table: Payoffs with  $c = l$

	E	NE
I	1.5,-1	3.5,0
NI	2,1	3,0

- how to solve this game?



# Games of incomplete information II: an example

- entrant has to think about
  - how likely is it that incumbent has low cost or high cost?
  - what will incumbent do if he has high cost? what if he has low cost?
  - what should I do?
- incumbent with low cost has to think about
  - what will entrant do?
    - partly depends on what he thinks I would do if I had high costs. . .
- we will return to this example later on!

# Games of incomplete information III: general thoughts

- say two firms do not know the cost of the respective other firm
- the main trick:
  - add beliefs about costs of other firm (i.e. a probability distribution over possible costs)
  - maximize expected utility
- we might want to allow this belief to depend on own costs
  - e.g. a high cost firm may think it is more likely that the other firm has also high costs

# Games of incomplete information IV: formal description

- finite set of players:  $i = 1, \dots, N$
- each player has a set of pure strategies  $S_i$
- to capture uncertainty of other players:
  - player  $i$  has a **type**  $t_i$  from a set  $T_i$
  - player  $i$  knows his own type  $t_i$  but other players do not
- player  $i$  maximizes expected utility with Bernoulli utility function  $u_i : S \times T \rightarrow \mathfrak{R}$ 
  - $T = \times_{i=1}^N T_i$  is set of all type profiles
  - $S = \times_{i=1}^N S_i$  is set of all strategy profiles
  - actions and types of all players can affect  $i$ 's payoff
- each type of each player has a belief  $p_i(t_{-i}|t_i)$  about other players' types
  - $p_i(t_{-i}|t_i) \in [0, 1]$
  - $\sum_{t_{-i} \in T_{-i}} p_i(t_{-i}|t_i) = 1$  where  $T_{-i}$  is the set of type profiles of all players but  $i$

# Games of incomplete information V: formal description (short)

A  $N$ -player game of incomplete information can be denoted as  $G = (S_i, T_i, p_i, u_i)_{i=1}^N$  where

- $S_i$  is the strategy set of player  $i$
- $T_i$  is the type set of player  $i$
- $p_i$  assigns to each  $t_i \in T_i$  a belief over  $T_{-i}$
- $u_i : S \times T \rightarrow \mathfrak{R}$  is player  $i$ 's utility function.

If all  $S_i$  and  $T_i$  are finite, the  $G$  is called a *finite game of incomplete information*.

# Games of incomplete information VI: assumptions on beliefs

- usually, it is assumed that types have a joint distribution  $p$  (over  $T$ ) and beliefs are derived using Bayes' rule:

$$p_i(t_{-i}|t_i) = \frac{p(t_i, t_{-i})}{\sum_{t'_{-i} \in T_{-i}} p(t_i, t'_{-i})}$$

then  $p$  is called the *common prior*

- often we assume independence of types, i.e. the belief  $p_i(t_{-i}|t_i)$  is the same for all  $t_i$

# Bayesian Nash equilibrium I

each player maximizes expected utility given his type and others strategies

→ trick:

- think of each type of every type as an own player maximizing expected utility (with utility function  $u_i$  and beliefs  $p_i(t_{-i}|t_i)$ )
- a Bayesian Nash equilibrium consists of one strategy for each type of each player such that
  - the strategy of type  $t_i$  maximizes expected utility of player  $i$  given the strategies of the others and the belief  $p_i(\cdot|t_i)$

## Bayesian Nash equilibrium II (formal)

For Bayesian game  $G = (S_i, T_i, p_i, u_i)_{i=1}^N$ , define the auxiliary game of complete information  $G'$ :

- set of players is  $T_1 \cup T_2 \cup \dots \cup T_N$

- strategy set of player  $t_i$  is  $S_i$

- von Neumann-Morgenstern utility

$$v_{t_i}(s) = \mathbb{E}_{t_{-i} \in T_{-i}} [u_i(s(t_1), \dots, s(t_N), t, \dots, t_N)]$$

- where  $s(t_i)$  is the strategy of player  $t_i$  and

- $s = (s(t_1), \dots, s(t_N))$

- where expectation is taken using the belief  $p_i(\cdot | t_i)$

### Definition: Bayesian Nash equilibrium (BNE)

A (mixed) Bayesian Nash equilibrium of game  $G$  is a (mixed) Nash equilibrium of the corresponding auxiliary game  $G'$ .

## Bayesian Nash equilibrium III: back to example

- assume the belief  $p_E(l) = p_E(h) = 1/2$

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Table: Payoffs with  $c = l$

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- what is the optimal strategy for type  $h$ ?
- if type  $l$  invests with probability  $s(l)$ , what is the entrant's best response?
- if the entrant enters with probability  $s(e)$ , what is type  $l$ 's best response?



# public good example I

- $N$  guests at a garden party
- each guest has to decide whether to bring a speaker to play music,  $S_i = \{0, 1\}$
- payoff of player  $i$ :
  - zero if no one brings a speaker
  - $t_i$  if someone else brought a speaker
  - $t_i - 1/2$  if person  $i$  brought a speaker
- $t_i$  are independently distributed and 1 (high) with probability  $2/3$  and 0 (low) with probability  $1/3$
- we want to find a *symmetric BNE*, i.e. one where all high types use one strategy and all low types use one other strategy

## public good example II

- what is the optimal strategy of a low type?
- suppose all high types bring a speaker with probability  $\alpha$ 
  - for player  $i$ : what is the probability that no one else brings a speaker?
  - what is the expected payoff for a high type of player  $i$  when bringing the speaker?
  - what is the expected payoff for a high type of player  $i$  when not bringing the speaker?
- which value of  $\alpha$  gives a BNE?

