```
report.txt
+----+
| PROJECT : PintOS I |
DESIGN DOCUMENT
+----+
Sanjoy Chowdhury <Roll no 20172123>
Ranajit Saha < Roll no 20172119>
---- PRELIMINARIES ----
>> Sources referred -
>> [1] https://web.stanford.edu/class/cs140/projects/pintos/pintos.html
>> [2] https://pintosiiith.wordpress.com/2012/09/13/install-pintos-with-
gemu/
>> [3] http://arpith.xyz/2016/01/getting-started-with-pintos/
SETUP IDIOSYNCRASIES
>> [3] Ideally, QEMU should exit once PintOS powers off. It does not do
>> gets stuck at "Powering Off...". To rectify the issue the following
changes
>> were made to src/devices/shutdown.c
@@ -99,6 +99,8 @@
printf ("Powering off...\n");
serial flush ();
/* For ensuring proper shutdown on gemu */
outw(0x604, 0x0 | 0x2000);
/* This is a special power-off sequence supported by Bochs and
QEMU, but not by physical hardware. */
for (p = s; *p != ' \setminus 0'; p++)
HELLO WORLD
=========
>> "hello.c" was added to the "src/tests/threads" directory
>> Contents of file "src/tests/threads/hello.c"
/* Tests hello, which should print "Hello Pintos" to console and exit */
#include
#include
#include
#include
#include
#include
<stdio.h>
"tests/threads/tests.h"
"threads/malloc.h"
"threads/synch.h"
"threads/thread.h"
"devices/timer.h"
void
test hello (void)
printf("Hello Pintos\n");
pass ();
}
```

```
>> To ensure that "hello.c" was built with the kernel, we made changes to
>> "Make.tests", "tests.c" and "tests.h" in the "src/tests/threads"
directory
>> Changes to "Make.tests"
@@ -9,6 +9,7 @@
priority-fifo priority-preempt priority-sema priority-condvar
priority-donate-chain
\report.txt
mlfqs-load-1 mlfqs-load-60 mlfqs-load-avq mlfqs-recent-1 mlfqs-fair-2
+hello \
mlfqs-fair-20 mlfqs-nice-2 mlfqs-nice-10 mlfqs-block)
# Sources for tests.
@@ -36,6 +37,7 @@
tests/threads SRC +=
tests/threads SRC +=
tests/threads SRC +=
+tests/threads SRC +=
tests/threads/mlfqs-recent-1.c
tests/threads/mlfqs-fair.c
tests/threads/mlfqs-block.c
tests/threads/hello.c
MLFQS OUTPUTS =
tests/threads/mlfqs-load-1.output
>> Changes to "tests.c"
@@ -38,6 +38,7 @@
{"mlfqs-nice-2", test mlfqs nice 2},
{"mlfqs-nice-10", test mlfqs nice 10},
{"mlfqs-block", test mlfqs block},
{"hello", test hello},
static const char *test_name;
>> Changes to "tests.h"
@@ -32,6 +32,8 @@
extern test func test mlfqs nice 2;
extern test func test mlfqs nice 10;
extern test func test mlfqs block;
+/* Prints Hello world */
+extern test_func test_hello;
void msg (const char *, ...);
void fail (const char *, ...);
ALARM CLOCK
_____
---- ADDITIONAL CHANGES ----
>> The following changes were made to "src/Makefile.build" in order to
>> that the source file "src/threads/priority.c" is build with the
kernel.
@@ -15,6 +15,7 @@
threads SRC = threads/start.S
# Startup code.
threads SRC += threads/init.c
# Main program.
threads SRC += threads/thread.c
```

```
# Thread management core.
+threads SRC += threads/pqueue.c
# Priority queue management
threads SRC += threads/switch.S
# Thread switch routine.
threads SRC += threads/interrupt.c # Interrupt core.
threads SRC += threads/intr-stubs.S # Interrupt stubs.
---- DATA STRUCTURES ----
>> A generic priority-queue was implemented in "src/threads/pqueue.c"
>> (declarations in "src/threads/pqueue.h") to deal with the ordering of
>> thread wake up alarms.
>> Contents of "pqueue.h"
#ifndef PQUEUE H INCLUDED
#define PQUEUE H INCLUDED
#include <stdio.h>report.txt
#include <stdbool.h>
#include <inttypes.h>
#include "threads/malloc.h"
typedef struct heap64 elem t
void *data;
int64 t key;
} heap64 elem t;
typedef struct heap64 t
{
heap64 elem t *heap;
int n elem;
int max nelem;
} heap64 t;
extern heap64 t *
heap64_init (heap64_t *h, int max_nelem);
extern bool
heap64 is empty (heap64 t *h);
extern bool
heap64 is full (heap64 t *h);
#define HEAP_PARENT(i) ((i)>>1)
#define HEAP_LCHILD(i) ((i) <<1)</pre>
#define HEAP RCHILD(i) (((i) << 1) +1)
extern void
heap64 minpq heapify (heap64 t *h, int ind);
extern void
heap64 maxpq heapify (heap64 t *h, int ind);
extern void
heap64_peek (heap64_t *h);
extern void
heap64 minpq insert (heap64 t *h, void* elem);
extern void
heap64 maxpq insert (heap64 t *h, void *elem);
extern void
heap64 minpq pop (heap64 t *h);
extern void
heap64_maxpq_pop (heap64_t *h);
extern void
pr heap (heap64 t *h, void (*pr data) (void *));
#endif
>> Contents of "pqueue.c"
#include "threads/pqueue.h"
/* Allocates space for heap types's data and initializes it
```

```
* Caller must ensure/check that:
* h is not NULL
* h->heap is a valid memory address (i.e. NOT NULL) after the call
heap64 t *
heap64 init (heap64 t *h, int max nelem)
{report.txt
ASSERT ( h != NULL);
h->heap = malloc (max nelem * sizeof(heap64 elem t) );
h->max nelem = max nelem;
h->n elem = 0;
return h;
/* Returns true if heap has no elements
* Caller must ensure that:
* h is not NULL
* /
bool
heap64 is empty (heap64 t *h)
ASSERT (h != NULL);
return (h->n elem <= 0);
/* Returns true if heap is full
* Caller must ensure that:
* h is not NULL
*/
bool
heap64 is full (heap64 t *h)
ASSERT (h != NULL);
return (h->n elem >= h->max nelem);
/* Min heapifies the heap in h->heap from index ind (1 based)
* Caller must ensure that
* h is not NULL
*/
void
heap64 minpq heapify (heap64 t *h, int ind)
ASSERT (h != NULL);
int smallest;
while (1) {
int l = HEAP_LCHILD(ind);
int r = HEAP RCHILD(ind);
if (1 \le h-)n \in \mathbb{A} + h-)heap[1].key < h-)heap[ind].key)
smallest = 1;
else
smallest = ind;
if (r \le h-n elem \&\& h-heap[r].key < h-heap[smallest].key)
smallest = r;
if (smallest == ind)
break:
else {
heap64 elem t temp = h->heap[ind];
h->heap[ind] = h->heap[smallest];
h->heap[smallest] = temp;
ind = smallest;
```

```
}
}
}
void
heap64 maxpq heapify (heap64 t *h, int ind)
{report.txt
ASSERT (h != NULL);
int largest;
while (1) {
int l = HEAP LCHILD(ind);
int r = HEAP RCHILD(ind);
if (1 \le h- n elem \&\& h- heap[l].key > h- heap[ind].key)
largest = 1;
else
largest = ind;
if (r \le h->n elem \&\& h->heap[r].key > h->heap[largest].key)
largest = r;
if (largest == ind)
break;
else {
heap64 elem t temp = h->heap[ind];
h->heap[ind] = h->heap[largest];
h->heap[largest] = temp;
ind = largest;
/* Returns pointer to top element in heap
* Caller must ensure/check that:
* h is not NULL
* heap has at least one element
* Caller must cast the returned void * to their expected type
*/
void *
heap64 peek (heap64 t *h)
ASSERT (h != NULL);
return h->heap + 1;
/* Inserts into min heap
* Caller must ensure that
* h is not NULL
* heap has at least one element
* Caller must pass pointer to the element they are inserting
*/
void
heap64 minpq insert (heap64 t *h, void* elem)
ASSERT (h != NULL);
ASSERT (elem != NULL);
int i, p;
i = ++ (h->n \text{ elem});
h\rightarrow heap[i] = *((heap64 elem t *)elem);
while (i > 1 \&\& h-) heap[p = HEAP PARENT(i)].key > h-) heap[i].key {
heap64 elem t temp = h->heap[p];
h\rightarrow heap[p] = h\rightarrow heap[i];
h->heap[i] = temp;
i = p;
```

```
}
voidreport.txt
heap64 maxpq insert (heap64 t *h, void *elem)
ASSERT (h != NULL);
int i, p;
i = ++ (h->n \text{ elem});
h\rightarrow heap[i] = *((heap64 elem t *)elem);
while (i > 1 && h->heap[p = HEAP PARENT(i)].key < h->heap[i].key) {
heap64 elem t temp = h->heap[p];
h \rightarrow heap[p] = h \rightarrow heap[i];
h \rightarrow heap[i] = temp;
i = p;
}
}
void
heap64 minpq pop (heap64 t *h)
ASSERT (h != NULL);
h\rightarrow heap[1] = h\rightarrow heap[h\rightarrow n elem--];
heap64 minpq heapify(h, 1);
}
void
heap64 maxpq pop (heap64 t *h)
ASSERT (h != NULL);
h\rightarrow heap[1] = h\rightarrow heap[h\rightarrow n elem--];
heap64 maxpq heapify(h, 1);
void
pr heap (heap64 t *h, void (*pr data)(void *))
ASSERT (h != NULL);
int i;
for (i = 1; i <= h->n_elem; i++) {
printf ("(%d)[ key = %"PRId64"; { ", i, h->heap[i].key);
pr data (h->heap[i].data);
printf (" } ]\n");
}
>> The declarations for our priority queue of timer alarms was added to
>> "threads.h".
>> Changes to "thread.h"
@@ -4,6 +4,7 @@
#include <debug.h>
#include <list.h>
#include <stdint.h>
+#include "threads/pqueue.h"
/* States in a thread's life cycle. */
enum thread status
@@ -138,4 +139,25 @@
int thread_get_recent_cpu (void);
int thread get load avg (void);
+/* Timer-alarm structure */
+typedef struct alarm t {
+ struct thread *t;
/* Thread setting the alarm */
```

```
+ int64 t wk tm;
/* Wake-Up time (units: ticks since boot) */report.txt
+} alarm t;
+/* Priority queue of Alarms: Earliest alarm first */
+extern heap64 t timer pg;
+/* Ticks since boot (used for waking up threads) */
+extern int64 t glob tm;
+/* Return true if first thread has a higher priority, false otherwise */
+bool priority check(const struct list_elem *, const struct list_elem *,
void *);
+/* Maximum size of Timer-alarm Priority queue */
+#define TIMERPQ MAXSZ 50
+/* Maximum no. of processes that can be released per tick */
+#define MAX UNBLOCKS PER TICK 10
#endif /* threads/thread.h */
---- ALGORITHMS ----
>> Changes to "timer.c"
@@ -90,10 +90,23 @@
timer_sleep (int64 t ticks)
int64 t start = timer ticks ();
+ struct thread *t = thread current();
ASSERT (intr get level () == INTR ON);
- while (timer elapsed (start) < ticks)
thread yield ();
+ enum intr level old level = intr disable ();
+ /* Set alarm for the required wakeup time */
+ alarm t alm;
+ alm.t = t;
+ alm.wk tm = start + ticks;
+ /* Push alarm to the timer priority queue and block */
+ heap64_minpq_insert(&timer_pq, &alm);
+ thread block();
+ /* Sleep is over. Wakeup */
+ intr set level (old level);
/* Sleeps for approximately MS milliseconds. Interrupts must be
@@ -171,6 +184,7 @@
timer interrupt (struct intr frame *args UNUSED)
ticks++;
+ glob tm = ticks;
thread tick ();
>> Changes to "thread.c"
```

```
@@ -50,6 +50,10 @@
static long long kernel ticks;
static long long user_ticks;
+/* Priority Queue of alarms */
+heap64 t timer pq;
+int64_t glob_tm;
/* # of timer ticks in kernel threads. */
/* # of timer ticks in user programs. */report.txt
8
/* Scheduling. */
#define TIME_SLICE 4
static unsigned thread ticks;
/* # of timer ticks to give each thread. */
/* # of timer ticks since last yield. */
@@ -134,8 +149,31 @@
else
kernel ticks++;
+
+
/* Enforce preemption. */
if (++thread ticks >= TIME SLICE)
int n unblocked = 0; /* No. of threads unblocked in this run */
/* While there're threads to ublock and
* we haven't unblocked too many */
while (heap64 is empty (&timer pq) == false &&
n_unblocked < MAX_UNBLOCKS_PER_TICK) {</pre>
/* Get the thread with the earliest wake time */
alarm t alm = *((alarm t *)heap64 peek (&timer pq));
/* If it's before wake-up time break,
^{\star} otherwise wake the thread up ^{\star}/
if (alm.wk tm > glob tm)
break;
else {
```

```
thread unblock (alm.t);
/* Wake the thread up */
n unblocked++;
/* Increment the count of unblocks */
/* Enforce preemption upon time-slice expiry and
* preempt idle thread if any other threads were unblocked*/
if (++thread ticks >= TIME SLICE ||
(n unblocked > 0 && t == idle thread))
intr yield on return ();
}
>>
>>
>>
>>
>>
>>
>>
>>
>>
>>
>>
>>
>>
When timer sleep() is called we create an alarm object which stores a
pointer to the
calling thread and the time that it wants to wake up at (units - ticks
since OS boot)
We push this object in a min priority queue and block the calling thread.
When the timer interrupt() is called we update the glob tm variable with
the current
time (units ticks since OS boot). timer interrupt then calls
thread tick(). Inside
thread tick() we peek for the thread with the earliest wake up time (O(1)
operation
in the min heap). If the current time is past or equal to its wake up
time, we
iteratively wake up as many threads as we can upto the limit specified by
MAX UNBLOCKS PER TICK
These threads can be scheduled only when the scheduler gets the chance to
run again
which might be after the current thread expires its time slice. However
if the idle
thread is currently running and we have unblocked atleast one thread we
preempt the
idle thread.
---- SYNCHRONIZATION ----
>>
>>
>>
>>
In order to ensure that no race conditions occur during simultaneous
access to
thread sleep() we disable interrupts before any action statements occur.
us the chance to insert the alarm to the min heap and block the calling
thread
```

heap64 minpq pop (&timer pq);

```
atomically. Interrupts get enabled again by the next thread to be
scheduled.
---- RATIONALE ----
>>
>>
>>
>>
>>
Race conditions were avoided by disabling interrupts instead of using
another
mechanism like semaphores, locks, etc. because cooperation between
interacting threads could not be guaranteed in this case.
We had also considered maintaining a sleep interval variable inside each
thread structure which would get decremented at each timer tick allowing
thereport.txt
>> thread to wake up when that variable becomes zero. However, this
approach required
>> us to perform operations linear in the number of threads at each timer
tick (O(n)).
>> Our current approach ONLY takes amortized constant time per tick.
PRIORITY SCHEDULING A
---- DATA STRUCTURES ----
>>
>>
>>
>>
Ready list is now maintained as an ordered list on the basis of thread
priority. The highest priority thread will reside at the front of the
list.
This was done by utilising the functions list insert ordered() provided
"src/libs/kernel/list.c" by providing a comparator function
priority check()
---- ALGORITHMS ----
>> A comparator function priority check() was defined to ensure sorted
ordering
>> within the ready list.
>> Changes to "thread.c"
@@ -71,6 +75,17 @@
void thread schedule tail (struct thread *prev);
static tid t allocate tid (void);
+/* Return true if first thread has a higher priority, false otherwise */
+bool priority_check(const struct list elem * first,
const struct list elem * second, void *aux)
+ struct thread * first thread = list entry (first, struct thread, elem);
+ struct thread * second thread = list entry (second, struct thread,
+ if (first thread->priority > second thread->priority)
return true;
+ else
return false;
+ }
```

```
/* Initializes the threading system by transforming the code
that's currently running into a thread. This can't work in
general and it is possible in this case only because loader.S
>>
>>
>>
>>
Changes were made to ready list manipulation statements inside thread.c
to ensure that the list remains sorted and threads are scheduled
according to
their priorities by the scheduler.
Changes to "thread.c"
@@ -209,6 +247,20 @@
/* Add to run queue. */
thread unblock (t);
/* If the newly created process has higher priority it may execute before
this function finishes */
old level=intr disable();
if (!list empty(&ready list)) {
struct thread * current thread = thread current();
struct thread * list front thread = list entry (list front (&ready list),
struct thread, elem);
if (current thread->priority < list front thread->priority)
thread yield();
}
+
intr set level (old level);
return tid;
}report.txt
@@ -245,7 +297,7 @@
old level = intr disable ();
ASSERT (t->status == THREAD BLOCKED);
list push back (&ready list, &t->elem);
list insert ordered (&ready list, &t->elem, priority check, NULL);
t->status = THREAD READY;
intr set level (old level);
@@ -316,7 +368,7 @@
old level = intr disable ();
if (cur != idle thread)
list push back (&ready list, &cur->elem);
```

```
list insert ordered (&ready list, &cur->elem, priority check, NULL);
cur->status = THREAD READY;
schedule ();
intr set level (old level);
@@ -339,11 +391,25 @@
-/* Sets the current thread's priority to NEW PRIORITY. */
+/* Sets the current thread's priority to NEW PRIORITY.
+ * Yield if the current thread no longer has the
+ * highest priority among all the ready threads */
thread set priority (int new priority)
- thread current ()->priority = new priority;
+ enum intr_level old_level = intr_disable();
+ struct thread * current thread = thread current();
+ int current priority = current thread->priority;
+ current thread->priority = new priority;
+ if (new priority < current priority && !list empty ( &ready list)) {
struct thread * list_front_thread = list_entry (list_front (&ready list),
struct thread, elem);
if (current thread->priority < list front thread->priority)
thread yield();
+ }
+ intr set level (old level);
/* Returns the current thread's priority. */
---- SYNCHRONIZATION ----
>> No changes were made to the synchronization arrangements within the
functions
>> modified because we were already operating with interrupts disabled
---- RATIONALE ----
>> The approach taken by us allows us to implement the desired operations
>> very simple yet efficient manner, allowing us to reuse the
functionality already
>> available with the list implementation.
PRIORITY SCHEDULING B
---- ALGORITHM ----
>> Changes to "synch.c"report.txt
11
@@ -68,7 +68,7 @@
old level = intr disable ();
while (sema->value == 0)
{
list push back (&sema->waiters, &thread current ()->elem);
list insert ordered (&sema->waiters, &thread current ()->elem,
priority check, NULL);
```

```
thread_block ();
sema->value--;
@@ -113,10 +113,17 @@
ASSERT (sema != NULL);
+
old_level = intr_disable ();
if (!list empty (&sema->waiters))
if (!list_empty (&sema->waiters)) {
struct thread * current_thread = thread_current();
struct thread * list front thread = list entry (list front (&sema-
>waiters),
struct thread, elem);
thread unblock (list entry (list pop front (&sema->waiters),
struct thread, elem));
sema->value++;
sema->value++;
if (current thread->priority < list front thread->priority)
thread yield();
}
else
sema->value++;
intr set level (old_level);
---- RATIONALE ----
>> Similar to the approach used in thread priority scheduling we
converted
>> the semaphore waiting list to an ordered list based on the priority of
the threads
>> waiting on the semaphore.
----END-----END-----
_____
```