1 Join-Size Estimation

(1 P.)

1. Estimate the size of the join $R(a,b) \bowtie S(b,c)$ using histograms for R.b and S.b. Assume V(R,b) = V(S,b) = 30 and the histograms for both attributes give the frequencies of the four most common values, as below, and further assume that every value appearing in the relation with the smaller set of values (R in this case) will also appear in the set of values of the other relation.

	0	1	2	3	others
R.b	10	8	7	11	39

	0	1	2	4	others
S.b	12	8	10	8	52

How does this estimate compare with the simpler estimate, assuming that all 30 values are equally likely to occur, with T(R) = 75 and T(S) = 90?

2. Estimate the size of the natural join $R(a,b) \bowtie S(b,c)$ if we have the following histogram information. Give a lower and upper bound for the join size and explain under which circumstances they appear.

	b < 0	b = 0	b > 0
R	300	100	400
S	300	200	600

2 Join-Ordering: Dynamic Programming

(1 P.)

1. Manually create the DP-table for the relations A,B,C with cardinalities |A| = 100, |B| = 25, |C| = 80 and selectivities $f_{A,C} = 0.05$, $f_{B,C} = 0.3$ with C_{out} as cost function. Cross products are allowed this time. Please keep the replaced entries in the table and highlight the final ones.

2. Given the following DP-table with intermediate results and the query graph (with selectivities):

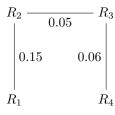
Relations	T	T	$C_{out}(T)$
$\{R_1\}$	R_1	40	0
$\{R_2\}$	R_2	10	0
$\{R_3\}$	R_3	20	0
$\{R_4\}$	R_4	30	0
$\{R_1, R_2\}$	$(R_1 \bowtie R_2)$	60	60
$\{R_1, R_3\}$	$(R_1 \times R_3)$	800	800
$\{R_1, R_4\}$	$(R_1 \times R_4)$	1200	1200
$\{R_2, R_3\}$	$(R_2 \bowtie R_3)$	10	10
$\{R_2, R_4\}$	$(R_2 \times R_4)$	300	300
$\{R_3, R_4\}$	$(R_3 \bowtie R_4)$	36	36
$\{R_1, R_2, R_3\}$	$((R_2 \bowtie R_3) \bowtie R_1)$	60	70
$\{R_1, R_2, R_4\}$	$((R_1 \bowtie R_2) \times R_4)$	1800	1860
$\{R_1, R_3, R_4\}$	$((R_3 \bowtie R_4) \times R_1)$	1440	1476
$\{R_2, R_3, R_4\}$	$((R_2 \bowtie R_3) \bowtie R_4)$	18	28

•
$$|R_1| = 40$$

•
$$|R_2| = 10$$

•
$$|R_3| = 20$$

•
$$|R_4| = 30$$



Calculate the optimal bushy join tree for the relations $\{R_1, R_2, R_3, R_4\}$ with the DP-algorithm shown in the lecture.

3 Simplifying queries

(1 P.)

Given the following queries from the uni_db schema from:

```
SELECT DISTINCT p.name,
  (SELECT MAX(position) FROM professors p2
    WHERE p2.PID=p.PID)
4 FROM professors p, lectures 1
5 WHERE
   EXISTS
     (SELECT *
      FROM
         (SELECT
9
            e.matrnr,
10
            (SELECT s.name FROM students s WHERE e.matrnr=s.matrnr),
            LID
12
         FROM exam e
13
         ORDER BY e.grade LIMIT 2
14
         ) t,
        lectures 12
      WHERE t.LID = 12.LID
17
        AND 12.LID = 1.LID
18
19
   AND p.PID = 1.heldby;
2.
1 SELECT
   s2.name,
_{\scriptscriptstyle 3} (WITH RECURSIVE t(LID) AS (
     SELECT exam.LID FROM students, exam
     WHERE students.matrnr = exam.matrnr AND students.matrnr = s2.
     matrnr
       AND exam.grade = (SELECT min(grade) FROM exam)
   UNION ALL
     SELECT prerequisites.required FROM t, prerequisites WHERE t.LID =
      prerequisites.lecture
9 )
10 SELECT count(*) FROM t) x
11 FROM
12 students s2;
```

Simplify (and optimize) these queries, as shown in the lecture. Let Postgres EXPLAIN the original and simplified query plans and interpret them.

4 Correctness of Unnesting

(1 P.)

Provide unnested queries for the given nested queries. Show through an example, by specifying contents of tables and corresponding results, why the type of join (e.g., INNER, LEFT OUTER) is important when unnesting a certain query. Considering the given queries, will the change in the type of the join impact the correctness of the query?

```
1.
    SELECT DISTINCT P.playerId
    FROM Player P
    WHERE (
        SELECT COUNT(G.id)
        FROM Game G
        WHERE G.playerId = P.playerId
      ) >10

2.
    SELECT DISTINCT P.name, (SELECT
        COUNT(*)
    FROM Game G
    WHERE P.playerId = G.playerId)
    FROM Player P
```