

Internet Research Task Force (IRTF)
Request for Comments: 9380
Category: Informational
ISSN: 2070-1721

A. Faz-Hernandez
Cloudflare, Inc.
S. Scott
Oso Security, Inc.
N. Sullivan
Cloudflare, Inc.
R. S. Wahby
Stanford University
C. A. Wood
Cloudflare, Inc.
August 2023

Hashing to Elliptic Curves

Abstract

This document specifies a number of algorithms for encoding or hashing an arbitrary string to a point on an elliptic curve. This document is a product of the Crypto Forum Research Group (CFRG) in the IRTF.

Status of This Memo

This document is not an Internet Standards Track specification; it is published for informational purposes.

This document is a product of the Internet Research Task Force (IRTF). The IRTF publishes the results of Internet-related research and development activities. These results might not be suitable for deployment. This RFC represents the consensus of the Crypto Forum Research Group of the Internet Research Task Force (IRTF). Documents approved for publication by the IRSG are not candidates for any level of Internet Standard; see Section 2 of RFC 7841.

Information about the current status of this document, any errata, and how to provide feedback on it may be obtained at <https://www.rfc-editor.org/info/rfc9380>.

Copyright Notice

Copyright (c) 2023 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust's Legal Provisions Relating to IETF Documents (<https://trustee.ietf.org/license-info>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document.

Table of Contents

1. Introduction
 - 1.1. Requirements Notation

- 2.1. Elliptic Curves
- 2.2. Terminology
 - 2.2.1. Mappings
 - 2.2.2. Encodings
 - 2.2.3. Random Oracle Encodings
 - 2.2.4. Serialization
 - 2.2.5. Domain Separation
- 3. Encoding Byte Strings to Elliptic Curves
 - 3.1. Domain Separation Requirements
- 4. Utility Functions
 - 4.1. The `sgn0` Function
- 5. Hashing to a Finite Field
 - 5.1. Efficiency Considerations in Extension Fields
 - 5.2. `hash_to_field` Implementation
 - 5.3. `expand_message`
 - 5.3.1. `expand_message_xmd`
 - 5.3.2. `expand_message_xof`
 - 5.3.3. Using DSTs Longer than 255 Bytes
 - 5.3.4. Defining Other `expand_message` Variants
- 6. Deterministic Mappings
 - 6.1. Choosing a Mapping Function
 - 6.2. Interface
 - 6.3. Notation
 - 6.4. Sign of the Resulting Point
 - 6.5. Exceptional Cases
 - 6.6. Mappings for Weierstrass Curves
 - 6.6.1. Shallue-van de Woestijne Method
 - 6.6.2. Simplified Shallue-van de Woestijne-Ulas Method
 - 6.6.3. Simplified SWU for $AB == 0$
 - 6.7. Mappings for Montgomery Curves
 - 6.7.1. Elligator 2 Method
 - 6.8. Mappings for Twisted Edwards Curves
 - 6.8.1. Rational Maps from Montgomery to Twisted Edwards Curves
 - 6.8.2. Elligator 2 Method
- 7. Clearing the Cofactor
- 8. Suites for Hashing
 - 8.1. Implementing a Hash-to-Curve Suite
 - 8.2. Suites for NIST P-256
 - 8.3. Suites for NIST P-384
 - 8.4. Suites for NIST P-521
 - 8.5. Suites for `curve25519` and `edwards25519`
 - 8.6. Suites for `curve448` and `edwards448`
 - 8.7. Suites for `secp256k1`
 - 8.8. Suites for BLS12-381
 - 8.8.1. BLS12-381 G1
 - 8.8.2. BLS12-381 G2
 - 8.9. Defining a New Hash-to-Curve Suite
 - 8.10. Suite ID Naming Conventions
- 9. IANA Considerations
- 10. Security Considerations
 - 10.1. Properties of Encodings
 - 10.2. Hashing Passwords
 - 10.3. Constant-Time Requirements
 - 10.4. `encode_to_curve`: Output Distribution and Indifferentiability
 - 10.5. `hash_to_field` Security

- 10.6. expand_message_xmd Security
- 10.7. Domain Separation for expand_message Variants
- 10.8. Target Security Levels
- 11. References
 - 11.1. Normative References
 - 11.2. Informative References
- Appendix A. Related Work
- Appendix B. Hashing to ristretto255
- Appendix C. Hashing to decaf448
- Appendix D. Rational Maps
 - D.1. Generic Mapping from Montgomery to Twisted Edwards
 - D.2. Mapping from Weierstrass to Montgomery
- Appendix E. Isogeny Maps for Suites
 - E.1. 3-Isogeny Map for secp256k1
 - E.2. 11-Isogeny Map for BLS12-381 G1
 - E.3. 3-Isogeny Map for BLS12-381 G2
- Appendix F. Straight-Line Implementations of Deterministic Mappings
 - F.1. Shallue-van de Woestijne Method
 - F.2. Simplified SWU Method
 - F.2.1. sqrt_ratio Subroutine
 - F.3. Elligator 2 Method
- Appendix G. Curve-Specific Optimized Sample Code
 - G.1. Interface and Projective Coordinate Systems
 - G.2. Elligator 2
 - G.2.1. curve25519 ($q = 5 \pmod{8}$, $K = 1$)
 - G.2.2. edwards25519
 - G.2.3. curve448 ($q = 3 \pmod{4}$, $K = 1$)
 - G.2.4. edwards448
 - G.2.5. Montgomery Curves with $q = 3 \pmod{4}$
 - G.2.6. Montgomery Curves with $q = 5 \pmod{8}$
 - G.3. Cofactor Clearing for BLS12-381 G2
- Appendix H. Scripts for Parameter Generation
 - H.1. Finding Z for the Shallue-van de Woestijne Map
 - H.2. Finding Z for Simplified SWU
 - H.3. Finding Z for Elligator 2
- Appendix I. sqrt and is_square Functions
 - I.1. sqrt for $q = 3 \pmod{4}$
 - I.2. sqrt for $q = 5 \pmod{8}$
 - I.3. sqrt for $q = 9 \pmod{16}$
 - I.4. Constant-Time Tonelli-Shanks Algorithm
 - I.5. is_square for $F = GF(p^2)$
- Appendix J. Suite Test Vectors
 - J.1. NIST P-256
 - J.1.1. P256_XMD:SHA-256_SSWU_RO_
 - J.1.2. P256_XMD:SHA-256_SSWU_NU_
 - J.2. NIST P-384
 - J.2.1. P384_XMD:SHA-384_SSWU_RO_
 - J.2.2. P384_XMD:SHA-384_SSWU_NU_
 - J.3. NIST P-521
 - J.3.1. P521_XMD:SHA-512_SSWU_RO_
 - J.3.2. P521_XMD:SHA-512_SSWU_NU_
 - J.4. curve25519
 - J.4.1. curve25519_XMD:SHA-512_ELL2_RO_
 - J.4.2. curve25519_XMD:SHA-512_ELL2_NU_
 - J.5. edwards25519

- J.5.1. edwards25519_XMD:SHA-512_ELL2_RO
- J.5.2. edwards25519_XMD:SHA-512_ELL2_NU
- J.6. curve448
 - J.6.1. curve448_XOF:SHAKE256_ELL2_RO
 - J.6.2. curve448_XOF:SHAKE256_ELL2_NU
- J.7. edwards448
 - J.7.1. edwards448_XOF:SHAKE256_ELL2_RO
 - J.7.2. edwards448_XOF:SHAKE256_ELL2_NU
- J.8. secp256k1
 - J.8.1. secp256k1_XMD:SHA-256_SSWU_RO
 - J.8.2. secp256k1_XMD:SHA-256_SSWU_NU
- J.9. BLS12-381 G1
 - J.9.1. BLS12381G1_XMD:SHA-256_SSWU_RO
 - J.9.2. BLS12381G1_XMD:SHA-256_SSWU_NU
- J.10. BLS12-381 G2
 - J.10.1. BLS12381G2_XMD:SHA-256_SSWU_RO
 - J.10.2. BLS12381G2_XMD:SHA-256_SSWU_NU

Appendix K. Expand Test Vectors

- K.1. expand_message_xmd(SHA-256)
- K.2. expand_message_xmd(SHA-256) (Long DST)
- K.3. expand_message_xmd(SHA-512)
- K.4. expand_message_xof(SHAKE128)
- K.5. expand_message_xof(SHAKE128) (Long DST)
- K.6. expand_message_xof(SHAKE256)

Acknowledgements

Contributors

Authors' Addresses

1. Introduction

Many cryptographic protocols require a procedure that encodes an arbitrary input, e.g., a password, to a point on an elliptic curve. This procedure is known as hashing to an elliptic curve, where the hashing procedure provides collision resistance and does not reveal the discrete logarithm of the output point. Prominent examples of cryptosystems that hash to elliptic curves include password-authenticated key exchanges [BM92] [J96] [BMP00] [p1363.2], Identity-Based Encryption [BF01], Boneh-Lynn-Shacham signatures [BLS01] [BLS-SIG], Verifiable Random Functions [MRV99] [VRF], and Oblivious Pseudorandom Functions [NR97] [OPRFs].

Unfortunately for implementors, the precise hash function that is suitable for a given protocol implemented using a given elliptic curve is often unclear from the protocol's description. Meanwhile, an incorrect choice of hash function can have disastrous consequences for security.

This document aims to bridge this gap by providing a comprehensive set of recommended algorithms for a range of curve types. Each algorithm conforms to a common interface: it takes as input an arbitrary-length byte string and produces as output a point on an elliptic curve. We provide implementation details for each algorithm, describe the security rationale behind each recommendation, and give guidance for elliptic curves that are not explicitly covered. We also present optimized implementations for internal functions used by these algorithms.

Readers wishing to quickly specify or implement a conforming hash function should consult Section 8, which lists recommended hash-to-curve suites and describes both how to implement an existing suite and how to specify a new one.

This document does not specify probabilistic rejection sampling methods, sometimes referred to as "try-and-increment" or "hunt-and-peck," because the goal is to specify algorithms that can plausibly be computed in constant time. Use of these probabilistic rejection methods is NOT RECOMMENDED, because they have been a perennial cause of side-channel vulnerabilities. See Dragonblood [VR20] as one example of this problem in practice, and see Appendix A for an informal description of rejection sampling methods and the timing side-channels they introduce.

This document represents the consensus of the Crypto Forum Research Group (CFRG).

1.1. Requirements Notation

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

2. Background

2.1. Elliptic Curves

The following is a brief definition of elliptic curves, with an emphasis on important parameters and their relation to hashing to curves. For further reference on elliptic curves, consult [CFADLNV05] or [W08].

Let F be the finite field $GF(q)$ of prime characteristic $p > 3$. (This document does not consider elliptic curves over fields of characteristic 2 or 3.) In most cases, F is a prime field, so $q = p$. Otherwise, F is an extension field, so $q = p^m$ for an integer $m > 1$. This document writes elements of extension fields in a primitive element or polynomial basis, i.e., as a vector of m elements of $GF(p)$ written in ascending order by degree. The entries of this vector are indexed in ascending order starting from 1, i.e., $x = (x_1, x_2, \dots, x_m)$. For example, if $q = p^2$ and the primitive element basis is $(1, I)$, then $x = (a, b)$ corresponds to the element $a + b * I$, where $x_1 = a$ and $x_2 = b$. (Note that all choices of basis are isomorphic, but certain choices may result in a more efficient implementation; this document does not make any particular assumptions about choice of basis.)

An elliptic curve E is specified by an equation in two variables and a finite field F . An elliptic curve equation takes one of several standard forms, including (but not limited to) Weierstrass, Montgomery, and Edwards.

The curve E induces an algebraic group of order n , meaning that the group has n distinct elements. (This document uses additive notation for the elliptic curve group operation.) Elements of an elliptic curve group are points with affine coordinates (x, y) satisfying the curve equation, where x and y are elements of F . In addition, all elliptic curve groups have a distinguished element, the identity point, which acts as the identity element for the group operation. On certain curves (including Weierstrass and Montgomery curves), the identity point cannot be represented as an (x, y) coordinate pair.

For security reasons, cryptographic applications of elliptic curves generally require using a (sub)group of prime order. Let G be such a subgroup of the curve of prime order r , where $n = h * r$. In this equation, h is an integer called the cofactor. An algorithm that takes as input an arbitrary point on the curve E and produces as output a point in the subgroup G of E is said to "clear the cofactor." Such algorithms are discussed in Section 7.

Certain hash-to-curve algorithms restrict the form of the curve equation, the characteristic of the field, or the parameters of the curve. For each algorithm presented, this document lists the relevant restrictions.

The table below summarizes quantities relevant to hashing to curves:

Symbol	Meaning	Relevance
F, q, p	A finite field F of characteristic p and $\#F = q = p^m$.	For prime fields, $q = p$; otherwise, $q = p^m$ and $m > 1$.
E	Elliptic curve.	E is specified by an equation and a field F .
n	Number of points on the elliptic curve E .	$n = h * r$, for h and r defined below.
G	A prime-order subgroup of the points on E .	G is a destination group to which byte strings are encoded.
r	Order of G .	r is a prime factor of n (usually, the largest such factor).
h	Cofactor, $h \geq 1$.	h is an integer satisfying $n = h * r$.

Table 1: Summary of Symbols and Their Definitions

2.2. Terminology

In this section, we define important terms used throughout the document.

2.2.1. Mappings

A mapping is a deterministic function from an element of the field F to a point on an elliptic curve E defined over F .

In general, the set of all points that a mapping can produce over all possible inputs may be only a subset of the points on an elliptic curve (i.e., the mapping may not be surjective). In addition, a mapping may output the same point for two or more distinct inputs (i.e., the mapping may not be injective). For example, consider a mapping from F to an elliptic curve having n points: if the number of elements of F is not equal to n , then this mapping cannot be bijective (i.e., both injective and surjective), since the mapping is defined to be deterministic.

Mappings may also be invertible, meaning that there is an efficient algorithm that, for any point P output by the mapping, outputs an x in F such that applying the mapping to x outputs P . Some of the mappings given in Section 6 are invertible, but this document does not discuss inversion algorithms.

2.2.2. Encodings

Encodings are closely related to mappings. Like a mapping, an encoding is a function that outputs a point on an elliptic curve. In contrast to a mapping, however, the input to an encoding is an arbitrary-length byte string.

This document constructs deterministic encodings by composing a hash function H_f with a deterministic mapping. In particular, H_f takes as input an arbitrary string and outputs an element of F . The deterministic mapping takes that element as input and outputs a point on an elliptic curve E defined over F . Since H_f takes arbitrary-length byte strings as inputs, it cannot be injective: the set of inputs is larger than the set of outputs, so there must be distinct inputs that give the same output (i.e., there must be collisions). Thus, any encoding built from H_f is also not injective.

Like mappings, encodings may be invertible, meaning that there is an efficient algorithm that, for any point P output by the encoding, outputs a string s such that applying the encoding to s outputs P . However, the instantiation of H_f used by all encodings specified in this document (Section 5) is not invertible; thus, those encodings are also not invertible.

In some applications of hashing to elliptic curves, it is important that encodings do not leak information through side channels. [VR20] is one example of this type of leakage leading to a security vulnerability. See Section 10.3 for further discussion.

2.2.3. Random Oracle Encodings

A random-oracle encoding satisfies a strong property: it can be

proved indifferentiable from a random oracle [MRH04] under a suitable assumption.

Both constructions described in Section 3 are indifferentiable from random oracles [MRH04] when instantiated following the guidelines in this document. The constructions differ in their output distributions: one gives a uniformly random point on the curve, the other gives a point sampled from a nonuniform distribution.

A random-oracle encoding with a uniform output distribution is suitable for use in many cryptographic protocols proven secure in the random-oracle model. See Section 10.1 for further discussion.

2.2.4. Serialization

A procedure related to encoding is the conversion of an elliptic curve point to a bit string. This is called serialization, and it is typically used for compactly storing or transmitting points. The inverse operation, deserialization, converts a bit string to an elliptic curve point. For example, [SEC1] and [p1363a] give standard methods for serialization and deserialization.

Deserialization is different from encoding in that only certain strings (namely, those output by the serialization procedure) can be deserialized. In contrast, this document is concerned with encodings from arbitrary strings to elliptic curve points. This document does not cover serialization or deserialization.

2.2.5. Domain Separation

Cryptographic protocols proven secure in the random-oracle model are often analyzed under the assumption that the random oracle only answers queries associated with that protocol (including queries made by adversaries) [BR93]. In practice, this assumption does not hold if two protocols use the same function to instantiate the random oracle. Concretely, consider protocols P1 and P2 that query a random-oracle R0: if P1 and P2 both query R0 on the same value x, the security analysis of one or both protocols may be invalidated.

A common way of addressing this issue is called domain separation, which allows a single random oracle to simulate multiple, independent oracles. This is effected by ensuring that each simulated oracle sees queries that are distinct from those seen by all other simulated oracles. For example, to simulate two oracles R01 and R02 given a single oracle R0, one might define

$$\begin{aligned} R01(x) &:= R0("R01" || x) \\ R02(x) &:= R0("R02" || x) \end{aligned}$$

where `||` is the concatenation operator. In this example, "R01" and "R02" are called domain separation tags (DSTs); they ensure that queries to R01 and R02 cannot result in identical queries to R0, meaning that it is safe to treat R01 and R02 as independent oracles.

In general, domain separation requires defining a distinct injective encoding for each oracle being simulated. In the above example,

"R01" and "R02" have the same length and thus satisfy this requirement when used as prefixes. The algorithms specified in this document take a different approach to ensuring injectivity; see Sections 5.3 and 10.7 for more details.

3. Encoding Byte Strings to Elliptic Curves

This section presents a general framework and interface for encoding byte strings to points on an elliptic curve. The constructions in this section rely on three basic functions:

- * The function `hash_to_field` hashes arbitrary-length byte strings to a list of one or more elements of a finite field F ; its implementation is defined in Section 5.

`hash_to_field(msg, count)`

Input:

- `msg`, a byte string containing the message to hash.
- `count`, the number of elements of F to output.

Output:

- `(u_0, ..., u_(count - 1))`, a list of field elements.

Steps: defined in Section 5.

- * The function `map_to_curve` calculates a point on the elliptic curve E from an element of the finite field F over which E is defined. Section 6 describes mappings for a range of curve families.

`map_to_curve(u)`

Input: `u`, an element of field F .

Output: `Q`, a point on the elliptic curve E .

Steps: defined in Section 6.

- * The function `clear_cofactor` sends any point on the curve E to the subgroup G of E . Section 7 describes methods to perform this operation.

`clear_cofactor(Q)`

Input: `Q`, a point on the elliptic curve E .

Output: `P`, a point in G .

Steps: defined in Section 7.

The two encodings (Section 2.2.2) defined in this section have the same interface and are both random-oracle encodings (Section 2.2.3). Both are implemented as a composition of the three basic functions above. The difference between the two is that their outputs are sampled from different distributions:

- * `encode_to_curve` is a nonuniform encoding from byte strings to points in G . That is, the distribution of its output is not uniformly random in G : the set of possible outputs of `encode_to_curve` is only a fraction of the points in G , and some

points in this set are more likely to be output than others. Section 10.4 gives a more precise definition of `encode_to_curve`'s output distribution.

`encode_to_curve(msg)`

Input: `msg`, an arbitrary-length byte string.
Output: `P`, a point in `G`.

Steps:

1. `u = hash_to_field(msg, 1)`
2. `Q = map_to_curve(u[0])`
3. `P = clear_cofactor(Q)`
4. return `P`

- * `hash_to_curve` is a uniform encoding from byte strings to points in `G`. That is, the distribution of its output is statistically close to uniform in `G`.

This function is suitable for most applications requiring a random oracle returning points in `G`, when instantiated with any of the `map_to_curve` functions described in Section 6. See Section 10.1 for further discussion.

`hash_to_curve(msg)`

Input: `msg`, an arbitrary-length byte string.
Output: `P`, a point in `G`.

Steps:

1. `u = hash_to_field(msg, 2)`
2. `Q0 = map_to_curve(u[0])`
3. `Q1 = map_to_curve(u[1])`
4. `R = Q0 + Q1` # Point addition
5. `P = clear_cofactor(R)`
6. return `P`

Each hash-to-curve suite in Section 8 instantiates one of these encoding functions for a specific elliptic curve.

3.1. Domain Separation Requirements

All uses of the encoding functions defined in this document **MUST** include domain separation (Section 2.2.5) to avoid interfering with other uses of similar functionality.

Applications that instantiate multiple, independent instances of either `hash_to_curve` or `encode_to_curve` **MUST** enforce domain separation between those instances. This requirement applies in both the case of multiple instances targeting the same curve and the case of multiple instances targeting different curves. (This is because the internal `hash_to_field` primitive (Section 5) requires domain separation to guarantee independent outputs.)

Domain separation is enforced with a domain separation tag (DST), which is a byte string constructed according to the following

requirements:

1. Tags **MUST** be supplied as the DST parameter to hash_to_field, as described in Section 5.
2. Tags **MUST** have nonzero length. A minimum length of 16 bytes is **RECOMMENDED** to reduce the chance of collisions with other applications.
3. Tags **SHOULD** begin with a fixed identification string that is unique to the application.
4. Tags **SHOULD** include a version number.
5. For applications that define multiple ciphersuites, each ciphersuite's tag **MUST** be different. For this purpose, it is **RECOMMENDED** to include a ciphersuite identifier in each tag.
6. For applications that use multiple encodings, to either the same curve or different curves, each encoding **MUST** use a different tag. For this purpose, it is **RECOMMENDED** to include the encoding's Suite ID (Section 8) in the domain separation tag. For independent encodings based on the same suite, each tag **SHOULD** also include a distinct identifier, e.g., "ENC1" and "ENC2".

As an example, consider a fictional application named Quux that defines several different ciphersuites, each for a different curve. A reasonable choice of tag is "QUUX-V<xx>-CS<yy>-<suiteID>", where <xx> and <yy> are two-digit numbers indicating the version and ciphersuite, respectively, and <suiteID> is the Suite ID of the encoding used in ciphersuite <yy>.

As another example, consider a fictional application named Baz that requires two independent random oracles to the same curve. Reasonable choices of tags for these oracles are "BAZ-V<xx>-CS<yy>-<suiteID>-ENC1" and "BAZ-V<xx>-CS<yy>-<suiteID>-ENC2", respectively, where <xx>, <yy>, and <suiteID> are as described above.

The example tags given above are assumed to be ASCII-encoded byte strings without null termination, which is the **RECOMMENDED** format. Other encodings can be used, but in all cases the encoding as a sequence of bytes **MUST** be specified unambiguously.

4. Utility Functions

Algorithms in this document use the utility functions described below, plus standard arithmetic operations (addition, multiplication, modular reduction, etc.) and elliptic curve point operations (point addition and scalar multiplication).

For security, implementations of these functions **SHOULD** be constant time: in brief, this means that execution time and memory access patterns **SHOULD NOT** depend on the values of secret inputs, intermediate values, or outputs. For such constant-time implementations, all arithmetic, comparisons, and assignments **MUST**

also be implemented in constant time. Section 10.3 briefly discusses constant-time security issues.

Guidance on implementing low-level operations (in constant time or otherwise) is beyond the scope of this document; readers should consult standard reference material [MOV96] [CFADLNV05].

- * **CMOV(a, b, c):** If *c* is False, CMOV returns *a*; otherwise, it returns *b*. For constant-time implementations, this operation must run in a time that is independent of the value of *c*.
- * **AND, OR, NOT, and XOR** are standard bitwise logical operators. For constant-time implementations, short-circuit operators **MUST** be avoided.
- * **is_square(x):** This function returns True whenever the value *x* is a square in the field *F*. By Euler's criterion, this function can be calculated in constant time as

$$\text{is_square}(x) := \begin{cases} \text{True,} & \text{if } x^{((q-1)/2)} \text{ is } 0 \text{ or } 1 \text{ in } F; \\ \text{False,} & \text{otherwise.} \end{cases}$$

In certain extension fields, *is_square* can be computed in constant time more quickly than by the above exponentiation. [AR13] and [S85] describe optimized methods for extension fields. Appendix I.5 gives an optimized straight-line method for $GF(p^2)$.

- * **sqrt(x):** The sqrt operation is a multi-valued function, i.e., there exist two roots of *x* in the field *F* whenever *x* is square (except when *x* = 0). To maintain compatibility across implementations while allowing implementors leeway for optimizations, this document does not require *sqrt()* to return a particular value. Instead, as explained in Section 6.4, any function that calls *sqrt* also specifies how to determine the correct root.

The preferred way of computing square roots is to fix a deterministic algorithm particular to *F*. We give several algorithms in Appendix I.

- * **sgn0(x):** This function returns either 0 or 1 indicating the "sign" of *x*, where *sgn0(x)* == 1 just when *x* is "negative". (In other words, this function always considers 0 to be positive.) Section 4.1 defines this function and discusses its implementation.
- * **inv0(x):** This function returns the multiplicative inverse of *x* in *F*, extended to all of *F* by fixing *inv0(0)* == 0. A straightforward way to implement *inv0* in constant time is to compute

$$\text{inv0}(x) := x^{(q-2)}.$$

Notice that on input 0, the output is 0 as required. Certain fields may allow faster inversion methods; detailed discussion of such methods is beyond the scope of this document.

- * **I2OSP and OS2IP:** These functions are used to convert a byte string to and from a non-negative integer as described in [RFC8017]. (Note that these functions operate on byte strings in big-endian byte order.)
- * **a || b:** denotes the concatenation of byte strings a and b. For example, "ABC" || "DEF" == "ABCDEF".
- * **substr(str, sbegin, slen):** For a byte string str, this function returns the slen-byte substring starting at position sbegin; positions are zero indexed. For example, substr("ABCDEFGH", 2, 3) == "CDE".
- * **len(str):** For a byte string str, this function returns the length of str in bytes. For example, len("ABC") == 3.
- * **strxor(str1, str2):** For byte strings str1 and str2, strxor(str1, str2) returns the bitwise XOR of the two strings. For example, strxor("abc", "XYZ") == "9;9" (the strings in this example are ASCII literals, but strxor is defined for arbitrary byte strings). In this document, strxor is only applied to inputs of equal length.

4.1. The sgn0 Function

This section defines a generic sgn0 implementation that applies to any field $F = GF(p^m)$. It also gives simplified implementations for the cases $F = GF(p)$ and $F = GF(p^2)$.

The definition of the sgn0 function for extension fields relies on the polynomial basis or vector representation of field elements, and iterates over the entire vector representation of the input element. As a result, sgn0 depends on the primitive polynomial used to define the polynomial basis; see Section 8 for more information about this basis, and see Section 2.1 for a discussion of representing elements of extension fields as vectors.

sgn0(x)

Parameters:

- F, a finite field of characteristic p and order $q = p^m$.
- p, the characteristic of F (see immediately above).
- m, the extension degree of F, $m \geq 1$ (see immediately above).

Input: x, an element of F.

Output: 0 or 1.

Steps:

1. sign = 0
2. zero = 1
3. for i in (1, 2, ..., m):
4. sign_i = x_i mod 2
5. zero_i = x_i == 0
6. sign = sign OR (zero AND sign_i) # Avoid short-circuit logic ops
7. zero = zero AND zero_i
8. return sign

When $m == 1$, `sgn0` can be significantly simplified:

`sgn0_m_eq_1(x)`

Input: x , an element of $\text{GF}(p)$.

Output: 0 or 1.

Steps:

1. return $x \bmod 2$

The case $m == 2$ is only slightly more complicated:

`sgn0_m_eq_2(x)`

Input: x , an element of $\text{GF}(p^2)$.

Output: 0 or 1.

Steps:

1. $\text{sign}_0 = x_0 \bmod 2$

2. $\text{zero}_0 = x_0 == 0$

3. $\text{sign}_1 = x_1 \bmod 2$

4. $s = \text{sign}_0 \text{ OR } (\text{zero}_0 \text{ AND } \text{sign}_1)$ # Avoid short-circuit logic ops

5. return s

5. Hashing to a Finite Field

The `hash_to_field` function hashes a byte string `msg` of arbitrary length into one or more elements of a field F . This function works in two steps: it first hashes the input byte string to produce a uniformly random byte string, and then interprets this byte string as one or more elements of F .

For the first step, `hash_to_field` calls an auxiliary function `expand_message`. This document defines two variants of `expand_message`: one appropriate for hash functions like SHA-2 [FIPS180-4] or SHA-3 [FIPS202], and another appropriate for extendable-output functions such as SHAKE128 [FIPS202]. Security considerations for each `expand_message` variant are discussed below (Sections 5.3.1 and 5.3.2).

Implementors **MUST NOT** use rejection sampling to generate a uniformly random element of F , to ensure that the `hash_to_field` function is amenable to constant-time implementation. The reason is that rejection sampling procedures are difficult to implement in constant time, and later well-meaning "optimizations" may silently render an implementation non-constant-time. This means that any `hash_to_field` function based on rejection sampling would be incompatible with constant-time implementation.

The `hash_to_field` function is also suitable for securely hashing to scalars. For example, when hashing to the scalar field for an elliptic curve (sub)group with prime order r , it suffices to instantiate `hash_to_field` with target field $\text{GF}(r)$.

The `hash_to_field` function is designed to be indifferentially private from a

random oracle [MRH04] when `expand_message` (Section 5.3) is modeled as a random oracle (see Section 10.5 for details about its indistinguishability). Ensuring indistinguishability requires care; to see why, consider a prime p that is close to $3/4 * 2^{256}$. Reducing a random 256-bit integer modulo this p yields a value that is in the range $[0, p / 3]$ with probability roughly $1/2$, meaning that this value is statistically far from uniform in $[0, p - 1]$.

To control bias, `hash_to_field` instead uses random integers whose length is at least $\text{ceil}(\log_2(p)) + k$ bits, where k is the target security level for the suite in bits. Reducing such integers mod p gives bias at most 2^{-k} for any p ; this bias is appropriate when targeting k -bit security. For each such integer, `hash_to_field` uses `expand_message` to obtain L uniform bytes, where

$$L = \text{ceil}((\text{ceil}(\log_2(p)) + k) / 8)$$

These uniform bytes are then interpreted as an integer via OS2IP. For example, for a 255-bit prime p , and $k = 128$ -bit security, $L = \text{ceil}((255 + 128) / 8) = 48$ bytes.

Note that k is an upper bound on the security level for the corresponding curve. See Section 10.8 for more details and Section 8.9 for guidelines on choosing k for a given curve.

5.1. Efficiency Considerations in Extension Fields

The `hash_to_field` function described in this section is inefficient for certain extension fields. Specifically, when hashing to an element of the extension field $\text{GF}(p^m)$, `hash_to_field` requires expanding `msg` into $m * L$ bytes (for L as defined above). For extension fields where $\log_2(p)$ is significantly smaller than the security level k , this approach is inefficient: it requires `expand_message` to output roughly $m * \log_2(p) + m * k$ bits, whereas $m * \log_2(p) + k$ bytes suffices to generate an element of $\text{GF}(p^m)$ with bias at most 2^{-k} . In such cases, applications MAY use an alternative `hash_to_field` function, provided it meets the following security requirements:

- * The function MUST output one or more field elements that are uniformly random except with bias at most 2^{-k} .
- * The function MUST NOT use rejection sampling.
- * The function SHOULD be amenable to straight-line implementations.

For example, Pornin [P20] describes a method for hashing to $\text{GF}(9767^{19})$ that meets these requirements while using fewer output bits from `expand_message` than `hash_to_field` would for that field.

5.2. `hash_to_field` Implementation

The following procedure implements `hash_to_field`.

The `expand_message` parameter to this function MUST conform to the requirements given in Section 5.3. Section 3.1 discusses the

REQUIRED method for constructing DST, the domain separation tag. Note that `hash_to_field` may fail (ABORT) if `expand_message` fails.

`hash_to_field(msg, count)`

Parameters:

- DST, a domain separation tag (see Section 3.1).
- F, a finite field of characteristic p and order $q = p^m$.
- p , the characteristic of F (see immediately above).
- m , the extension degree of F, $m \geq 1$ (see immediately above).
- $L = \lceil (\lceil \log_2(p) \rceil + k) / 8 \rceil$, where k is the security parameter of the suite (e.g., $k = 128$).
- `expand_message`, a function that expands a byte string and domain separation tag into a uniformly random byte string (see Section 5.3).

Input:

- `msg`, a byte string containing the message to hash.
- `count`, the number of elements of F to output.

Output:

- $(u_0, \dots, u_{(\text{count} - 1)})$, a list of field elements.

Steps:

1. `len_in_bytes = count * m * L`
2. `uniform_bytes = expand_message(msg, DST, len_in_bytes)`
3. `for i in (0, ..., count - 1):`
4. `for j in (0, ..., m - 1):`
5. `elm_offset = L * (j + i * m)`
6. `tv = substr(uniform_bytes, elm_offset, L)`
7. `e_j = OS2IP(tv) mod p`
8. `u_i = (e_0, ..., e_{(m - 1)})`
9. `return (u_0, ..., u_{(count - 1)})`

5.3. `expand_message`

`expand_message` is a function that generates a uniformly random byte string. It takes three arguments:

1. `msg`, a byte string containing the message to hash,
2. DST, a byte string that acts as a domain separation tag, and
3. `len_in_bytes`, the number of bytes to be generated.

This document defines the following two variants of `expand_message`:

- * `expand_message_xmd` (Section 5.3.1) is appropriate for use with a wide range of hash functions, including SHA-2 [FIPS180-4], SHA-3 [FIPS202], BLAKE2 [RFC7693], and others.
- * `expand_message_xof` (Section 5.3.2) is appropriate for use with extendable-output functions (XOFs), including functions in the SHAKE [FIPS202] or BLAKE2X [BLAKE2X] families.

These variants should suffice for the vast majority of use cases, but

other variants are possible; Section 5.3.4 discusses requirements.

5.3.1. `expand_message_xmd`

The `expand_message_xmd` function produces a uniformly random byte string using a cryptographic hash function H that outputs b bits. For security, H MUST meet the following requirements:

- * The number of bits output by H MUST be $b \geq 2 * k$, where k is the target security level in bits, and b MUST be divisible by 8. The first requirement ensures k -bit collision resistance; the second ensures uniformity of `expand_message_xmd`'s output.
- * H MAY be a Merkle-Damgaard hash function like SHA-2. In this case, security holds when the underlying compression function is modeled as a random oracle [CDMP05]. (See Section 10.6 for discussion.)
- * H MAY be a sponge-based hash function like SHA-3 or BLAKE2. In this case, security holds when the inner function is modeled as a random transformation or as a random permutation [BDPV08].
- * Otherwise, H MUST be a hash function that has been proved indistinguishable from a random oracle [MRH04] under a reasonable cryptographic assumption.

SHA-2 [FIPS180-4] and SHA-3 [FIPS202] are typical and RECOMMENDED choices. As an example, for the 128-bit security level, $b \geq 256$ bits and either SHA-256 or SHA3-256 would be an appropriate choice.

The hash function H is assumed to work by repeatedly ingesting fixed-length blocks of data. The length in bits of these blocks is called the input block size (s). As examples, $s = 1024$ for SHA-512 [FIPS180-4] and $s = 576$ for SHA3-512 [FIPS202]. For correctness, H requires $b \leq s$.

The following procedure implements `expand_message_xmd`.

`expand_message_xmd(msg, DST, len_in_bytes)`

Parameters:

- H , a hash function (see requirements above).
- b_in_bytes , $b / 8$ for b the output size of H in bits. For example, for $b = 256$, $b_in_bytes = 32$.
- s_in_bytes , the input block size of H , measured in bytes (see discussion above). For example, for SHA-256, $s_in_bytes = 64$.

Input:

- `msg`, a byte string.
- `DST`, a byte string of at most 255 bytes. See below for information on using longer DSTs.
- `len_in_bytes`, the length of the requested output in bytes, not greater than the lesser of $(255 * b_in_bytes)$ or $2^{16}-1$.

Output:

- `uniform_bytes`, a byte string.

Steps:

```
1. ell = ceil(len_in_bytes / b_in_bytes)
2. ABORT if ell > 255 or len_in_bytes > 65535 or len(DST) > 255
3. DST_prime = DST || I2OSP(len(DST), 1)
4. Z_pad = I2OSP(0, s_in_bytes)
5. l_i_b_str = I2OSP(len_in_bytes, 2)
6. msg_prime = Z_pad || msg || l_i_b_str || I2OSP(0, 1) || DST_prime
7. b_0 = H(msg_prime)
8. b_1 = H(b_0 || I2OSP(1, 1) || DST_prime)
9. for i in (2, ..., ell):
10.    b_i = H(strxor(b_0, b_(i - 1)) || I2OSP(i, 1) || DST_prime)
11. uniform_bytes = b_1 || ... || b_ell
12. return substr(uniform_bytes, 0, len_in_bytes)
```

Note that the string `Z_pad` (step 6) is prefixed to `msg` before computing `b_0` (step 7). This is necessary for security when `H` is a Merkle-Damgård hash, e.g., SHA-2 (see Section 10.6). Hashing this additional data means that the cost of computing `b_0` is higher than the cost of simply computing `H(msg)`. In most settings, this overhead is negligible, because the cost of evaluating `H` is much less than the other costs involved in hashing to a curve.

It is possible, however, to entirely avoid this overhead by taking advantage of the fact that `Z_pad` depends only on `H`, and not on the arguments to `expand_message_xmd`. To do so, first precompute and save the internal state of `H` after ingesting `Z_pad`. Then, when computing `b_0`, initialize `H` using the saved state. Further details are implementation dependent and are beyond the scope of this document.

5.3.2. `expand_message_xof`

The `expand_message_xof` function produces a uniformly random byte string using an extendable-output function (XOF) `H`. For security, `H` MUST meet the following criteria:

- * The collision resistance of `H` MUST be at least `k` bits.
- * `H` MUST be an XOF that has been proved indifferentiable from a random oracle under a reasonable cryptographic assumption.

The SHAKE XOF family [FIPS202] is a typical and RECOMMENDED choice. As an example, for 128-bit security, SHAKE128 would be an appropriate choice.

The following procedure implements `expand_message_xof`.

```
expand_message_xof(msg, DST, len_in_bytes)
```

Parameters:

- `H(m, d)`, an extendable-output function that processes input message `m` and returns `d` bytes.

Input:

- `msg`, a byte string.
- `DST`, a byte string of at most 255 bytes.

See below for information on using longer DSTs.

- `len_in_bytes`, the length of the requested output in bytes.

Output:

- `uniform_bytes`, a byte string.

Steps:

1. ABORT if `len_in_bytes > 65535` or `len(DST) > 255`
2. `DST_prime = DST || I2OSP(len(DST), 1)`
3. `msg_prime = msg || I2OSP(len_in_bytes, 2) || DST_prime`
4. `uniform_bytes = H(msg_prime, len_in_bytes)`
5. return `uniform_bytes`

5.3.3. Using DSTs Longer than 255 Bytes

The `expand_message` variants defined in this section accept domain separation tags of at most 255 bytes. If applications require a domain separation tag longer than 255 bytes, e.g., because of requirements imposed by an invoking protocol, implementors **MUST** compute a short domain separation tag by hashing, as follows:

- * For `expand_message_xmd` using hash function `H`, DST is computed as
`DST = H("H2C-OVERSIZE-DST-" || a_very_long_DST)`
- * For `expand_message_xof` using extendable-output function `H`, DST is computed as

`DST = H("H2C-OVERSIZE-DST-" || a_very_long_DST, ceil(2 * k / 8))`

Here, `a_very_long_DST` is the DST whose length is greater than 255 bytes, `"H2C-OVERSIZE-DST-"` is a 17-byte ASCII string literal, and `k` is the target security level in bits.

5.3.4. Defining Other `expand_message` Variants

When defining a new `expand_message` variant, the most important consideration is that hash-to-field models `expand_message` as a random oracle. Thus, implementors **SHOULD** prove indistinguishability from a random oracle under an appropriate assumption about the underlying cryptographic primitives; see Section 10.5 for more information.

In addition, `expand_message` variants:

- * **MUST** give collision resistance commensurate with the security level of the target elliptic curve.
- * **MUST** be built on primitives designed for use in applications requiring cryptographic randomness. As examples, a secure stream cipher is an appropriate primitive, whereas a Mersenne twister pseudorandom number generator [MT98] is not.
- * **MUST NOT** use rejection sampling.
- * **MUST** give independent values for distinct `(msg, DST, length)` inputs. Meeting this requirement is subtle. As a simplified

example, hashing `msg || DST` does not work, because in this case distinct `(msg, DST)` pairs whose concatenations are equal will return the same output (e.g., `("AB", "CDEF")` and `("ABC", "DEF")`). The variants defined in this document use a suffix-free encoding of DST to avoid this issue.

- * **MUST** use the domain separation tag DST to ensure that invocations of cryptographic primitives inside of `expand_message` are domain-separated from invocations outside of `expand_message`. For example, if the `expand_message` variant uses a hash function `H`, an encoding of DST **MUST** be added either as a prefix or a suffix of the input to each invocation of `H`. Adding DST as a suffix is the **RECOMMENDED** approach.
- * **SHOULD** read `msg` exactly once, for efficiency when `msg` is long.

In addition, each `expand_message` variant **MUST** specify a unique `EXP_TAG` that identifies that variant in a Suite ID. See Section 8.10 for more information.

6. Deterministic Mappings

The mappings in this section are suitable for implementing either nonuniform or uniform encodings using the constructions in Section 3. Certain mappings restrict the form of the curve or its parameters. For each mapping presented, this document lists the relevant restrictions.

Note that mappings in this section are not interchangeable: different mappings will almost certainly output different points when evaluated on the same input.

6.1. Choosing a Mapping Function

This section gives brief guidelines on choosing a mapping function for a given elliptic curve. Note that the suites given in Section 8 are recommended mappings for the respective curves.

If the target elliptic curve is a Montgomery curve (Section 6.7), the Elligator 2 method (Section 6.7.1) is recommended. Similarly, if the target elliptic curve is a twisted Edwards curve (Section 6.8), the twisted Edwards Elligator 2 method (Section 6.8.2) is recommended.

The remaining cases are Weierstrass curves. For curves supported by the Simplified Shallue-van de Woestijne-Ulas (SWU) method (Section 6.6.2), that mapping is the recommended one. Otherwise, the Simplified SWU method for `AB == 0` (Section 6.6.3) is recommended if the goal is best performance, while the Shallue-van de Woestijne method (Section 6.6.1) is recommended if the goal is simplicity of implementation. (The reason for this distinction is that the Simplified SWU method for `AB == 0` requires implementing an isogeny map in addition to the mapping function, while the Shallue-van de Woestijne method does not.)

The Shallue-van de Woestijne method (Section 6.6.1) works with any curve and may be used in cases where a generic mapping is required.

Note, however, that this mapping is almost always more computationally expensive than the curve-specific recommendations above.

6.2. Interface

The generic interface shared by all mappings in this section is as follows:

$$(x, y) = \text{map_to_curve}(u)$$

The input u and outputs x and y are elements of the field F . The affine coordinates (x, y) specify a point on an elliptic curve defined over F . Note, however, that the point (x, y) is not a uniformly random point.

6.3. Notation

As a rough guide, the following conventions are used in pseudocode:

- * All arithmetic operations are performed over a field F , unless explicitly stated otherwise.
- * u : the input to the mapping function. This is an element of F produced by the `hash_to_field` function.
- * (x, y) , (s, t) , (v, w) : the affine coordinates of the point output by the mapping. Indexed variables (e.g., x_1, y_2, \dots) are used for candidate values.
- * tv_1, tv_2, \dots : reusable temporary variables.
- * c_1, c_2, \dots : constant values, which can be computed in advance.

6.4. Sign of the Resulting Point

In general, elliptic curves have equations of the form $y^2 = g(x)$. The mappings in this section first identify an x such that $g(x)$ is square, then take a square root to find y . Since there are two square roots when $g(x) \neq 0$, this may result in an ambiguity regarding the sign of y .

When necessary, the mappings in this section resolve this ambiguity by specifying the sign of the y -coordinate in terms of the input to the mapping function. Two main reasons support this approach: first, this covers elliptic curves over any field in a uniform way, and second, it gives implementors leeway in optimizing square-root implementations.

6.5. Exceptional Cases

Mappings may have exceptional cases, i.e., inputs u on which the mapping is undefined. These cases must be handled carefully, especially for constant-time implementations.

For each mapping in this section, we discuss the exceptional cases

and show how to handle them in constant time. Note that all implementations SHOULD use `inv0` (Section 4) to compute multiplicative inverses, to avoid exceptional cases that result from attempting to compute the inverse of 0.

6.6. Mappings for Weierstrass Curves

The mappings in this section apply to a target curve E defined by the equation

$$y^2 = g(x) = x^3 + A * x + B$$

where $4 * A^3 + 27 * B^2 \neq 0$.

6.6.1. Shallue-van de Woestijne Method

Shallue and van de Woestijne [SW06] describe a mapping that applies to essentially any elliptic curve. (Note, however, that this mapping is more expensive to evaluate than the other mappings in this document.)

The parameterization given below is for Weierstrass curves; its derivation is detailed in [W19]. This parameterization also works for Montgomery curves (Section 6.7) and twisted Edwards curves (Section 6.8) via the rational maps given in Appendix D: first, evaluate the Shallue-van de Woestijne mapping to an equivalent Weierstrass curve, then map that point to the target Montgomery or twisted Edwards curve using the corresponding rational map.

Preconditions: A Weierstrass curve $y^2 = x^3 + A * x + B$.

Constants:

- * A and B , the parameter of the Weierstrass curve.
- * Z , a non-zero element of F meeting the below criteria. Appendix H.1 gives a Sage script [SAGE] that outputs the RECOMMENDED Z .
 1. $g(Z) \neq 0$ in F .
 2. $-(3 * Z^2 + 4 * A) / (4 * g(Z)) \neq 0$ in F .
 3. $-(3 * Z^2 + 4 * A) / (4 * g(Z))$ is square in F .
 4. At least one of $g(Z)$ and $g(-Z / 2)$ is square in F .

Sign of y : Inputs u and $-u$ give the same x -coordinate for many values of u . Thus, we set $\text{sgn0}(y) == \text{sgn0}(u)$.

Exceptions: The exceptional cases for u occur when $(1 + u^2 * g(Z)) * (1 - u^2 * g(Z)) == 0$. The restrictions on Z given above ensure that implementations that use `inv0` to invert this product are exception free.

Operations:

```

1. tv1 = u^2 * g(Z)
2. tv2 = 1 + tv1
3. tv1 = 1 - tv1
4. tv3 = inv0(tv1 * tv2)
5. tv4 = sqrt(-g(Z) * (3 * Z^2 + 4 * A))    # can be precomputed
6. If sgn0(tv4) == 1, set tv4 = -tv4        # sgn0(tv4) MUST equal 0
7. tv5 = u * tv1 * tv3 * tv4
8. tv6 = -4 * g(Z) / (3 * Z^2 + 4 * A)    # can be precomputed
9. x1 = -Z / 2 - tv5
10. x2 = -Z / 2 + tv5
11. x3 = Z + tv6 * (tv2^2 * tv3)^2
12. If is_square(g(x1)), set x = x1 and y = sqrt(g(x1))
13. Else If is_square(g(x2)), set x = x2 and y = sqrt(g(x2))
14. Else set x = x3 and y = sqrt(g(x3))
15. If sgn0(u) != sgn0(y), set y = -y
16. return (x, y)

```

Appendix F.1 gives an example straight-line implementation of this mapping.

6.6.2. Simplified Shallue-van de Woestijne-Ulas Method

The function `map_to_curve_simple_swu(u)` implements a simplification of the Shallue-van de Woestijne-Ulas mapping [U07] described by Brier et al. [BCIMRT10], which they call the "simplified SWU" map. Wahby and Boneh [WB19] generalize and optimize this mapping.

Preconditions: A Weierstrass curve $y^2 = x^3 + A * x + B$ where $A \neq 0$ and $B \neq 0$.

Constants:

- * A and B, the parameters of the Weierstrass curve.
- * Z, an element of F meeting the below criteria. Appendix H.2 gives a Sage script [SAGE] that outputs the RECOMMENDED Z. The criteria are as follows:
 1. Z is non-square in F,
 2. $Z \neq -1$ in F,
 3. the polynomial $g(x) - Z$ is irreducible over F, and
 4. $g(B / (Z * A))$ is square in F.

Sign of y: Inputs u and -u give the same x-coordinate. Thus, we set $\text{sgn0}(y) == \text{sgn0}(u)$.

Exceptions: The exceptional cases are values of u such that $Z^2 * u^4 + Z * u^2 == 0$. This includes $u == 0$ and may include other values that depend on Z. Implementations must detect this case and set $x1 = B / (Z * A)$, which guarantees that $g(x1)$ is square by the condition on Z given above.

Operations:

1. $tv1 = \text{inv0}(Z^2 * u^4 + Z * u^2)$
2. $x1 = (-B / A) * (1 + tv1)$
3. If $tv1 == 0$, set $x1 = B / (Z * A)$
4. $gx1 = x1^3 + A * x1 + B$
5. $x2 = Z * u^2 * x1$
6. $gx2 = x2^3 + A * x2 + B$
7. If $\text{is_square}(gx1)$, set $x = x1$ and $y = \text{sqrt}(gx1)$
8. Else set $x = x2$ and $y = \text{sqrt}(gx2)$
9. If $\text{sgn0}(u) \neq \text{sgn0}(y)$, set $y = -y$
10. return (x, y)

Appendix F.2 gives a general and optimized straight-line implementation of this mapping. For more information on optimizing this mapping, see Section 4 of [WB19] or the example code found at [hash2curve-repo].

6.6.3. Simplified SWU for $AB == 0$

Wahby and Boneh [WB19] show how to adapt the Simplified SWU mapping to Weierstrass curves having $A == 0$ or $B == 0$, which the mapping of Section 6.6.2 does not support. (The case $A == B == 0$ is excluded because $y^2 = x^3$ is not an elliptic curve.)

This method applies to curves like secp256k1 [SEC2] and to pairing-friendly curves in the Barreto-Lynn-Scott family [BLS03], Barreto-Naehrig family [BN05], and other families.

This method requires finding another elliptic curve E' given by the equation

$$y'^2 = g'(x') = x'^3 + A' * x' + B'$$

that is isogenous to E and has $A' \neq 0$ and $B' \neq 0$. (See [WB19], Appendix A, for one way of finding E' using [SAGE].) This isogeny defines a map $\text{iso_map}(x', y')$ given by a pair of rational functions. iso_map takes as input a point on E' and produces as output a point on E .

Once E' and iso_map are identified, this mapping works as follows: on input u , first apply the Simplified SWU mapping to get a point on E' , then apply the isogeny map to that point to get a point on E .

Note that iso_map is a group homomorphism, meaning that point addition commutes with iso_map . Thus, when using this mapping in the `hash_to_curve` construction discussed in Section 3, one can effect a small optimization by first mapping u_0 and u_1 to E' , adding the resulting points on E' , and then applying iso_map to the sum. This gives the same result while requiring only one evaluation of iso_map .

Preconditions: An elliptic curve E' with $A' \neq 0$ and $B' \neq 0$ that is isogenous to the target curve E with isogeny map iso_map from E' to E .

Helper functions:

* `map_to_curve_simple_swu` is the mapping of Section 6.6.2 to E'

* `iso_map` is the isogeny map from E' to E

Sign of y : For this map, the sign is determined by `map_to_curve_simple_swu`. No further sign adjustments are necessary.

Exceptions: `map_to_curve_simple_swu` handles its exceptional cases. Exceptional cases of `iso_map` are inputs that cause the denominator of either rational function to evaluate to zero; such cases MUST return the identity point on E .

Operations:

```
1. (x', y') = map_to_curve_simple_swu(u)    # (x', y') is on E'
2.  (x, y) = iso_map(x', y')               # (x, y) is on E
3. return (x, y)
```

See [hash2curve-repo] or Section 4.3 of [WB19] for details on implementing the isogeny map.

6.7. Mappings for Montgomery Curves

The mapping defined in this section applies to a target curve M defined by the equation

$$K * t^2 = s^3 + J * s^2 + s$$

6.7.1. Elligator 2 Method

Bernstein, Hamburg, Krasnova, and Lange give a mapping that applies to any curve with a point of order 2 [BHKL13], which they call Elligator 2.

Preconditions: A Montgomery curve $K * t^2 = s^3 + J * s^2 + s$ where $J \neq 0$, $K \neq 0$, and $(J^2 - 4) / K^2$ is non-zero and non-square in F .

Constants:

* J and K , the parameters of the elliptic curve.

* Z , a non-square element of F . Appendix H.3 gives a Sage script [SAGE] that outputs the RECOMMENDED Z .

Sign of t : This mapping fixes the sign of t as specified in [BHKL13]. No additional adjustment is required.

Exceptions: The exceptional case is $Z * u^2 == -1$, i.e., $1 + Z * u^2 == 0$. Implementations must detect this case and set $x_1 = -(J / K)$. Note that this can only happen when $q = 3 \pmod{4}$.

Operations:

```
1. x1 = -(J / K) * inv0(1 + Z * u^2)
2. If x1 == 0, set x1 = -(J / K)
```

```

3. gx1 = x1^3 + (J / K) * x1^2 + x1 / K^2
4. x2 = -x1 - (J / K)
5. gx2 = x2^3 + (J / K) * x2^2 + x2 / K^2
6. If is_square(gx1), set x = x1, y = sqrt(gx1) with sgn0(y) == 1.
7. Else set x = x2, y = sqrt(gx2) with sgn0(y) == 0.
8. s = x * K
9. t = y * K
10. return (s, t)

```

Appendix F.3 gives an example straight-line implementation of this mapping. Appendix G.2 gives optimized straight-line procedures that apply to specific classes of curves and base fields.

6.8. Mappings for Twisted Edwards Curves

Twisted Edwards curves (a class of curves that includes Edwards curves) are given by the equation

$$a * v^2 + w^2 = 1 + d * v^2 * w^2$$

with $a \neq 0$, $d \neq 0$, and $a \neq d$ [BBJLP08].

These curves are closely related to Montgomery curves (Section 6.7): every twisted Edwards curve is birationally equivalent to a Montgomery curve ([BBJLP08], Theorem 3.2). This equivalence yields an efficient way of hashing to a twisted Edwards curve: first, hash to an equivalent Montgomery curve, then transform the result into a point on the twisted Edwards curve via a rational map. This method of hashing to a twisted Edwards curve thus requires identifying a corresponding Montgomery curve and rational map. We describe how to identify such a curve and map immediately below.

6.8.1. Rational Maps from Montgomery to Twisted Edwards Curves

There are two ways to select a Montgomery curve and rational map for use when hashing to a given twisted Edwards curve. The selected Montgomery curve and rational map **MUST** be specified as part of the hash-to-curve suite for a given twisted Edwards curve; see Section 8.

1. When hashing to a standardized twisted Edwards curve for which a corresponding Montgomery form and rational map are also standardized, the standard Montgomery form and rational map **SHOULD** be used to ensure compatibility with existing software.

In certain cases, e.g., `edwards25519` [RFC7748], the sign of the rational map from the twisted Edwards curve to its corresponding Montgomery curve is not given explicitly. In this case, the sign **MUST** be fixed such that applying the rational map to the twisted Edwards curve's base point yields the Montgomery curve's base point with correct sign. (For `edwards25519`, see [RFC7748] and [Err4730].)

When defining new twisted Edwards curves, a Montgomery equivalent and rational map **SHOULD** also be specified, and the sign of the rational map **SHOULD** be stated explicitly.

2. When hashing to a twisted Edwards curve that does not have a standardized Montgomery form or rational map, the map given in Appendix D SHOULD be used.

6.8.2. Elligator 2 Method

Preconditions: A twisted Edwards curve E and an equivalent Montgomery curve M meeting the requirements in Section 6.8.1.

Helper functions:

- * `map_to_curve_elligator2` is the mapping of Section 6.7.1 to the curve M .
- * `rational_map` is a function that takes a point (s, t) on M and returns a point (v, w) on E . This rational map should be chosen as defined in Section 6.8.1.

Sign of t (and v): For this map, the sign is determined by `map_to_curve_elligator2`. No further sign adjustments are required.

Exceptions: The exceptions for the Elligator 2 mapping are as given in Section 6.7.1. The exceptions for the rational map are as given in Section 6.8.1. No other exceptions are possible.

The following procedure implements the Elligator 2 mapping for a twisted Edwards curve. (Note that the output point is denoted (v, w) because it is a point on the target twisted Edwards curve.)

`map_to_curve_elligator2_edwards(u)`

Input: u , an element of F .

Output: (v, w) , a point on E .

1. $(s, t) = \text{map_to_curve_elligator2}(u)$ # (s, t) is on M
2. $(v, w) = \text{rational_map}(s, t)$ # (v, w) is on E
3. return (v, w)

7. Clearing the Cofactor

The mappings of Section 6 always output a point on the elliptic curve, i.e., a point in a group of order $h * r$ (Section 2.1). Obtaining a point in G may require a final operation commonly called "clearing the cofactor," which takes as input any point on the curve and produces as output a point in the prime-order (sub)group G (Section 2.1).

The cofactor can always be cleared via scalar multiplication by h . For elliptic curves where $h = 1$, i.e., the curves with a prime number of points, no operation is required. This applies, for example, to the NIST curves P-256, P-384, and P-521 [FIPS186-4].

In some cases, it is possible to clear the cofactor via a faster method than scalar multiplication by h . These methods are equivalent to (but usually faster than) multiplication by some scalar h_{eff} whose value is determined by the method and the curve. Examples of

fast cofactor clearing methods include the following:

- * For certain pairing-friendly curves having subgroup G_2 over an extension field, Scott et al. [SBCDK09] describe a method for fast cofactor clearing that exploits an efficiently computable endomorphism. Fuentes-Castañeda et al. [FKR11] propose an alternative method that is sometimes more efficient. Budroni and Pintore [BP17] give concrete instantiations of these methods for Barreto-Lynn-Scott pairing-friendly curves [BLS03]. This method is described for the specific case of BLS12-381 in Appendix G.3.
- * Wahby and Boneh ([WB19], Section 5) describe a trick due to Scott for fast cofactor clearing on any elliptic curve for which the prime factorization of h and the structure of the elliptic curve group meet certain conditions.

The `clear_cofactor` function is parameterized by a scalar `h_eff`. Specifically,

$$\text{clear_cofactor}(P) := h_eff * P$$

where $*$ represents scalar multiplication. When a curve does not support a fast cofactor clearing method, `h_eff` = `h` and the cofactor **MUST** be cleared via scalar multiplication.

When a curve admits a fast cofactor clearing method, `clear_cofactor` **MAY** be evaluated either via that method or via scalar multiplication by the equivalent `h_eff`; these two methods give the same result. Note that in this case scalar multiplication by the cofactor `h` does not generally give the same result as the fast method and **MUST NOT** be used.

8. Suites for Hashing

This section lists recommended suites for hashing to standard elliptic curves.

A hash-to-curve suite fully specifies the procedure for hashing byte strings to points on a specific elliptic curve group. Section 8.1 describes how to implement a suite. Applications that require hashing to an elliptic curve should use either an existing suite or a new suite specified as described in Section 8.9.

All applications using a hash-to-curve suite **MUST** choose a domain separation tag (DST) in accordance with the guidelines in Section 3.1. In addition, applications whose security requires a random oracle that returns uniformly random points on the target curve **MUST** use a suite whose encoding type is `hash_to_curve`; see Section 3 and immediately below for more information.

A hash-to-curve suite comprises the following parameters:

- * Suite ID, a short name used to refer to a given suite. Section 8.10 discusses the naming conventions for Suite IDs.
- * encoding type, either `uniform` (`hash_to_curve`) or `nonuniform`

(`encode_to_curve`). See Section 3 for definitions of these encoding types.

- * `E`, the target elliptic curve over a field `F`.
- * `p`, the characteristic of the field `F`.
- * `m`, the extension degree of the field `F`. If $m > 1$, the suite **MUST** also specify the polynomial basis used to represent extension field elements.
- * `k`, the target security level of the suite in bits. (See Section 10.8 for discussion.)
- * `L`, the length parameter for `hash_to_field` (Section 5).
- * `expand_message`, one of the variants specified in Section 5.3 plus any parameters required for the specified variant (for example, `H`, the underlying hash function).
- * `f`, a mapping function from Section 6.
- * `h_eff`, the scalar parameter for `clear_cofactor` (Section 7).

In addition to the above parameters, the mapping `f` may require additional parameters `Z`, `M`, `rational_map`, `E'`, or `iso_map`. When applicable, these **MUST** be specified.

The table below lists suites **RECOMMENDED** for some elliptic curves. The corresponding parameters are given in the following subsections. Applications instantiating cryptographic protocols whose security analysis relies on a random oracle that outputs points with a uniform distribution **MUST NOT** use a nonuniform encoding. Moreover, applications that use a nonuniform encoding **SHOULD** carefully analyze the security implications of nonuniformity. When the required encoding is not clear, applications **SHOULD** use a uniform encoding for security.

E	Suites	Section
NIST P-256	P256_XMD:SHA-256_SSWU_RO P256_XMD:SHA-256_SSWU_NU	8.2
NIST P-384	P384_XMD:SHA-384_SSWU_RO P384_XMD:SHA-384_SSWU_NU	8.3
NIST P-521	P521_XMD:SHA-512_SSWU_RO P521_XMD:SHA-512_SSWU_NU	8.4
curve25519	curve25519_XMD:SHA-512_ELL2_RO curve25519_XMD:SHA-512_ELL2_NU	8.5
edwards25519	edwards25519_XMD:SHA-512_ELL2_RO edwards25519_XMD:SHA-512_ELL2_NU	8.5

curve448	curve448_X0F:SHAKE256_ELL2_R0_ curve448_X0F:SHAKE256_ELL2_NU_	8.6
edwards448	edwards448_X0F:SHAKE256_ELL2_R0_ edwards448_X0F:SHAKE256_ELL2_NU_	8.6
secp256k1	secp256k1_XMD:SHA-256_SSWU_R0_ secp256k1_XMD:SHA-256_SSWU_NU_	8.7
BLS12-381 G1	BLS12381G1_XMD:SHA-256_SSWU_R0_ BLS12381G1_XMD:SHA-256_SSWU_NU_	8.8
BLS12-381 G2	BLS12381G2_XMD:SHA-256_SSWU_R0_ BLS12381G2_XMD:SHA-256_SSWU_NU_	8.8

Table 2: Suites for hashing to elliptic curves.

8.1. Implementing a Hash-to-Curve Suite

A hash-to-curve suite requires the following functions. Note that some of these require utility functions from Section 4.

1. Base field arithmetic operations for the target elliptic curve, e.g., addition, multiplication, and square root.
2. Elliptic curve point operations for the target curve, e.g., point addition and scalar multiplication.
3. The `hash_to_field` function; see Section 5. This includes the `expand_message` variant (Section 5.3) and any constituent hash function or XOF.
4. The suite-specified mapping function; see the corresponding subsection of Section 6.
5. A cofactor clearing function; see Section 7. This may be implemented as scalar multiplication by `h_eff` or as a faster equivalent method.
6. The desired encoding function; see Section 3. This is either `hash_to_curve` or `encode_to_curve`.

8.2. Suites for NIST P-256

This section defines ciphersuites for the NIST P-256 elliptic curve [FIPS186-4].

`P256_XMD:SHA-256_SSWU_R0_` is defined as follows:

- * encoding type: `hash_to_curve` (Section 3)
- * E: $y^2 = x^3 + A * x + B$, where
 - $A = -3$

- B = 0x5ac635d8aa3a93e7b3ebbd55769886bc651d06b0cc53b0f63bce3c3e27d2604b
- * p: $2^{256} - 2^{224} + 2^{192} + 2^{96} - 1$
- * m: 1
- * k: 128
- * expand_message: expand_message_xmd (Section 5.3.1)
- * H: SHA-256
- * L: 48
- * f: Simplified SWU method (Section 6.6.2)
- * Z: -10
- * h_eff: 1

P256_XMD:SHA-256_SSWU_NU_ is identical to P256_XMD:SHA-256_SSWU_RO_, except that the encoding type is encode_to_curve (Section 3).

An optimized example implementation of the Simplified SWU mapping to P-256 is given in Appendix F.2.

8.3. Suites for NIST P-384

This section defines ciphersuites for the NIST P-384 elliptic curve [FIPS186-4].

P384_XMD:SHA-384_SSWU_RO_ is defined as follows:

- * encoding type: hash_to_curve (Section 3)
- * E: $y^2 = x^3 + A * x + B$, where
 - A = -3
 - B = 0xb3312fa7e23ee7e4988e056be3f82d19181d9c6efe8141120314088f5013875ac656398d8a2ed19d2a85c8edd3ec2aef
- * p: $2^{384} - 2^{128} - 2^{96} + 2^{32} - 1$
- * m: 1
- * k: 192
- * expand_message: expand_message_xmd (Section 5.3.1)
- * H: SHA-384
- * L: 72
- * f: Simplified SWU method (Section 6.6.2)

* Z: -12

* h_eff: 1

P384_XMD:SHA-384_SSWU_NU_ is identical to P384_XMD:SHA-384_SSWU_RO_, except that the encoding type is encode_to_curve (Section 3).

An optimized example implementation of the Simplified SWU mapping to P-384 is given in Appendix F.2.

8.4. Suites for NIST P-521

This section defines ciphersuites for the NIST P-521 elliptic curve [FIPS186-4].

P521_XMD:SHA-512_SSWU_RO_ is defined as follows:

* encoding type: hash_to_curve (Section 3)

* E: $y^2 = x^3 + A * x + B$, where

- $A = -3$

- $B = 0x51953eb9618e1c9a1f929a21a0b68540eea2da725b99b315f3b8b489918ef109e156193951ec7e937b1652c0bd3bb1bf073573df883d2c34f1ef451fd46b503f00$

* p: $2^{521} - 1$

* m: 1

* k: 256

* expand_message: expand_message_xmd (Section 5.3.1)

* H: SHA-512

* L: 98

* f: Simplified SWU method (Section 6.6.2)

* Z: -4

* h_eff: 1

P521_XMD:SHA-512_SSWU_NU_ is identical to P521_XMD:SHA-512_SSWU_RO_, except that the encoding type is encode_to_curve (Section 3).

An optimized example implementation of the Simplified SWU mapping to P-521 is given in Appendix F.2.

8.5. Suites for curve25519 and edwards25519

This section defines ciphersuites for curve25519 and edwards25519 [RFC7748]. Note that these ciphersuites MUST NOT be used when

hashing to ristretto255 [ristretto255-decaf448]. See Appendix B for information on how to hash to that group.

curve25519_XMD:SHA-512_ELL2_R0_ is defined as follows:

- * encoding type: hash_to_curve (Section 3)
- * E: $K * t^2 = s^3 + J * s^2 + s$, where
 - $J = 486662$
 - $K = 1$
- * p: $2^{255} - 19$
- * m: 1
- * k: 128
- * expand_message: expand_message_xmd (Section 5.3.1)
- * H: SHA-512
- * L: 48
- * f: Elligator 2 method (Section 6.7.1)
- * Z: 2
- * h_eff: 8

edwards25519_XMD:SHA-512_ELL2_R0_ is identical to curve25519_XMD:SHA-512_ELL2_R0_, except for the following parameters:

- * E: $a * v^2 + w^2 = 1 + d * v^2 * w^2$, where
 - $a = -1$
 - $d = 0x52036cee2b6ffe738cc740797779e89800700a4d4141d8ab75eb4dca135978a3$
- * f: Twisted Edwards Elligator 2 method (Section 6.8.2)
- * M: curve25519, defined in [RFC7748], Section 4.1
- * rational_map: the birational maps defined in [RFC7748], Section 4.1

curve25519_XMD:SHA-512_ELL2_NU_ is identical to curve25519_XMD:SHA-512_ELL2_R0_, except that the encoding type is encode_to_curve (Section 3).

edwards25519_XMD:SHA-512_ELL2_NU_ is identical to edwards25519_XMD:SHA-512_ELL2_R0_, except that the encoding type is encode_to_curve (Section 3).

Optimized example implementations of the above mappings are given in Appendix G.2.1 and Appendix G.2.2.

8.6. Suites for curve448 and edwards448

This section defines ciphersuites for curve448 and edwards448 [RFC7748]. Note that these ciphersuites **MUST NOT** be used when hashing to decaf448 [ristretto255-decaf448]. See Appendix C for information on how to hash to that group.

curve448_XOF:SHAKE256_ELL2_R0_ is defined as follows:

- * encoding type: hash_to_curve (Section 3)
- * E: $K * t^2 = s^3 + J * s^2 + s$, where
 - $J = 156326$
 - $K = 1$
- * p: $2^{448} - 2^{224} - 1$
- * m: 1
- * k: 224
- * expand_message: expand_message_xof (Section 5.3.2)
- * H: SHAKE256
- * L: 84
- * f: Elligator 2 method (Section 6.7.1)
- * Z: -1
- * h_eff: 4

edwards448_XOF:SHAKE256_ELL2_R0_ is identical to curve448_XOF:SHAKE256_ELL2_R0_, except for the following parameters:

- * E: $a * v^2 + w^2 = 1 + d * v^2 * w^2$, where
 - $a = 1$
 - $d = -39081$
- * f: Twisted Edwards Elligator 2 method (Section 6.8.2)
- * M: curve448, defined in [RFC7748], Section 4.2
- * rational_map: the 4-isogeny map defined in [RFC7748], Section 4.2

curve448_XOF:SHAKE256_ELL2_NU_ is identical to curve448_XOF:SHAKE256_ELL2_R0_, except that the encoding type is encode_to_curve (Section 3).

edwards448_X0F:SHAKE256_ELL2_NU_ is identical to edwards448_X0F:SHAKE256_ELL2_R0_, except that the encoding type is encode_to_curve (Section 3).

Optimized example implementations of the above mappings are given in Appendix G.2.3 and Appendix G.2.4.

8.7. Suites for secp256k1

This section defines ciphersuites for the secp256k1 elliptic curve [SEC2].

secp256k1_XMD:SHA-256_SSWU_R0_ is defined as follows:

- * encoding type: hash_to_curve (Section 3)
- * E: $y^2 = x^3 + 7$
- * p: $2^{256} - 2^{32} - 2^9 - 2^8 - 2^7 - 2^6 - 2^4 - 1$
- * m: 1
- * k: 128
- * expand_message: expand_message_xmd (Section 5.3.1)
- * H: SHA-256
- * L: 48
- * f: Simplified SWU for $AB == 0$ (Section 6.6.3)
- * Z: -11
- * E': $y'^2 = x'^3 + A' * x' + B'$, where
 - A': 0x3f8731abdd661adca08a5558f0f5d272e953d363cb6f0e5d405447c01a444533
 - B': 1771
- * iso_map: the 3-isogeny map from E' to E given in Appendix E.1
- * h_eff: 1

secp256k1_XMD:SHA-256_SSWU_NU_ is identical to secp256k1_XMD:SHA-256_SSWU_R0_, except that the encoding type is encode_to_curve (Section 3).

An optimized example implementation of the Simplified SWU mapping to the curve E' isogenous to secp256k1 is given in Appendix F.2.

8.8. Suites for BLS12-381

This section defines ciphersuites for groups G1 and G2 of the

BLS12-381 elliptic curve [BLS12-381].

8.8.1. BLS12-381 G1

BLS12381G1_XMD:SHA-256_SSWU_R0_ is defined as follows:

- * encoding type: hash_to_curve (Section 3)
- * $E: y^2 = x^3 + 4$
- * $p: 0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabfffeb153ffffb9fefffffffffaaab$
- * $m: 1$
- * $k: 128$
- * expand_message: expand_message_xmd (Section 5.3.1)
- * $H: \text{SHA-256}$
- * $L: 64$
- * $f: \text{Simplified SWU for } AB == 0 \text{ (Section 6.6.3)}$
- * $Z: 11$
- * $E': y'^2 = x'^3 + A' * x' + B'$, where
 - $A' = 0x144698a3b8e9433d693a02c96d4982b0ea985383ee66a8d8e8981aefd881ac98936f8da0e0f97f5cf428082d584c1d$
 - $B' = 0x12e2908d11688030018b12e8753eee3b2016c1f0f24f4070a0b9c14fcef35ef55a23215a316ceaa5d1cc48e98e172be0$
- * iso_map: the 11-isogeny map from E' to E given in Appendix E.2
- * $h_{\text{eff}}: 0xd201000000010001$

BLS12381G1_XMD:SHA-256_SSWU_NU_ is identical to BLS12381G1_XMD:SHA-256_SSWU_R0_, except that the encoding type is encode_to_curve (Section 3).

Note that the h_{eff} values for these suites are chosen for compatibility with the fast cofactor clearing method described by Scott ([WB19], Section 5).

An optimized example implementation of the Simplified SWU mapping to the curve E' isogenous to BLS12-381 G1 is given in Appendix F.2.

8.8.2. BLS12-381 G2

BLS12381G2_XMD:SHA-256_SSWU_R0_ is defined as follows:

- * encoding type: hash_to_curve (Section 3)

- * $E: y^2 = x^3 + 4 * (1 + I)$
- * base field F is $GF(p^m)$, where
 - p : 0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabfffeb153ffffb9fefffffffffaaab
 - m : 2
 - $(1, I)$ is the basis for F , where $I^2 + 1 == 0$ in F
- * k : 128
- * `expand_message`: `expand_message_xmd` (Section 5.3.1)
- * H : SHA-256
- * L : 64
- * f : Simplified SWU for $AB == 0$ (Section 6.6.3)
- * Z : $-(2 + I)$
- * E' : $y'^2 = x'^3 + A' * x' + B'$, where
 - $A' = 240 * I$
 - $B' = 1012 * (1 + I)$
- * `iso_map`: the isogeny map from E' to E given in Appendix E.3
- * `h_eff`: 0xbc69f08f2ee75b3584c6a0ea91b352888e2a8e9145ad7689986ff031508ffe1329c2f178731db956d82bf015d1212b02ec0ec69d7477c1ae954cbc06689f6a359894c0adebbf6b4e8020005aaa95551

`BLS12381G2_XMD:SHA-256_SSWU_NU` is identical to `BLS12381G2_XMD:SHA-256_SSWU_R0`, except that the encoding type is `encode_to_curve` (Section 3).

Note that the `h_eff` values for these suites are chosen for compatibility with the fast cofactor clearing method described by Budroni and Pintore ([BP17], Section 4.1) and are summarized in Appendix G.3.

An optimized example implementation of the Simplified SWU mapping to the curve E' isogenous to BLS12-381 G2 is given in Appendix F.2.

8.9. Defining a New Hash-to-Curve Suite

For elliptic curves not listed elsewhere in Section 8, a new hash-to-curve suite can be defined by the following:

1. E , F , p , and m are determined by the elliptic curve and its base field.
2. k is an upper bound on the target security level of the suite

(Section 10.8). A reasonable choice of k is $\text{ceil}(\log_2(r) / 2)$, where r is the order of the subgroup G of the curve E (Section 2.1).

3. Choose encoding type, either `hash_to_curve` or `encode_to_curve` (Section 3).
4. Compute L as described in Section 5.
5. Choose an `expand_message` variant from Section 5.3 plus any underlying cryptographic primitives (e.g., a hash function H).
6. Choose a mapping following the guidelines in Section 6.1, and select any required parameters for that mapping.
7. Choose h_{eff} to be either the cofactor of E or, if a fast cofactor clearing method is to be used, a value appropriate to that method as discussed in Section 7.
8. Construct a Suite ID following the guidelines in Section 8.10.

8.10. Suite ID Naming Conventions

Suite IDs MUST be constructed as follows:

`CURVE_ID || "_" || HASH_ID || "_" || MAP_ID || "_" || ENC_VAR || "_"`

The fields `CURVE_ID`, `HASH_ID`, `MAP_ID`, and `ENC_VAR` are ASCII-encoded strings of at most 64 characters each. Fields MUST contain only ASCII characters between 0x21 and 0x7E (inclusive), except that underscore (i.e., 0x5F) is not allowed.

As indicated above, each field (including the last) is followed by an underscore ("_", ASCII 0x5F). This helps to ensure that Suite IDs are prefix free. Suite IDs MUST include the final underscore and MUST NOT include any characters after the final underscore.

Suite ID fields MUST be chosen as follows:

- * `CURVE_ID`: a human-readable representation of the target elliptic curve.
- * `HASH_ID`: a human-readable representation of the `expand_message` function and any underlying hash primitives used in `hash_to_field` (Section 5). This field MUST be constructed as follows:

`EXP_TAG || ":" || HASH_NAME`

`EXP_TAG` indicates the `expand_message` variant:

- "XMD" for `expand_message_xmd` (Section 5.3.1).
- "XOF" for `expand_message_xof` (Section 5.3.2).

`HASH_NAME` is a human-readable name for the underlying hash primitive. As examples:

1. For `expand_message_xof` (Section 5.3.2) with SHAKE128, `HASH_ID` is "XOF:SHAKE128".
2. For `expand_message_xmd` (Section 5.3.1) with SHA3-256, `HASH_ID` is "XMD:SHA3-256".

Suites that use an alternative `hash_to_field` function that meets the requirements in Section 5.1 MUST indicate this by appending a tag identifying that function to the `HASH_ID` field, separated by a colon (":", ASCII 0x3A).

- * `MAP_ID`: a human-readable representation of the `map_to_curve` function as defined in Section 6. These are defined as follows:
 - "SVDW" for Shallue and van de Woestijne (Section 6.6.1).
 - "SSWU" for Simplified SWU (Sections 6.6.2 and 6.6.3).
 - "ELL2" for Elligator 2 (Sections 6.7.1 and 6.8.2).
- * `ENC_VAR`: a string indicating the encoding type and other information. The first two characters of this string indicate whether the suite represents a `hash_to_curve` or an `encode_to_curve` operation (Section 3), as follows:
 - If `ENC_VAR` begins with "R0", the suite uses `hash_to_curve`.
 - If `ENC_VAR` begins with "NU", the suite uses `encode_to_curve`.
 - `ENC_VAR` MUST NOT begin with any other string.

`ENC_VAR` MAY also be used to encode other information used to identify variants, for example, a version number. The RECOMMENDED way to do so is to add one or more subfields separated by colons. For example, "R0:V02" is an appropriate `ENC_VAR` value for the second version of a uniform encoding suite, while "R0:V02:F0001:BAR17" might be used to indicate a variant of that suite.

9. IANA Considerations

This document has no IANA actions.

10. Security Considerations

This section contains additional security considerations about the hash-to-curve mechanisms described in this document.

10.1. Properties of Encodings

Each encoding type (Section 3) accepts an arbitrary byte string and maps it to a point on the curve sampled from a distribution that depends on the encoding type. It is important to note that using a nonuniform encoding or directly evaluating one of the mappings of Section 6 produces an output that is easily distinguished from a

uniformly random point. Applications that use a nonuniform encoding **SHOULD** carefully analyze the security implications of nonuniformity. When the required encoding is not clear, applications **SHOULD** use a uniform encoding.

Both encodings given in Section 3 can output the identity element of the group G . The probability that either encoding function outputs the identity element is roughly $1/r$ for a random input, which is negligible for cryptographically useful elliptic curves. Further, it is computationally infeasible to find an input to either encoding function whose corresponding output is the identity element. (Both of these properties hold when the encoding functions are instantiated with a hash_to_field function that follows all guidelines in Section 5.) Protocols that use these encoding functions **SHOULD NOT** add a special case to detect and "fix" the identity element.

When the hash_to_curve function (Section 3) is instantiated with a hash_to_field function that is indifferentiable from a random oracle (Section 5), the resulting function is indifferentiable from a random oracle ([MRH04] [BCIMRT10] [FFSTV13] [LBB19] [H20]). In many cases, such a function can be safely used in cryptographic protocols whose security analysis assumes a random oracle that outputs uniformly random points on an elliptic curve. As Ristenpart et al. discuss in [RSS11], however, not all security proofs that rely on random oracles continue to hold when those oracles are replaced by indifferentiable functionalities. This limitation should be considered when analyzing the security of protocols relying on the hash_to_curve function.

10.2. Hashing Passwords

When hashing passwords using any function described in this document, an adversary who learns the output of the hash function (or potentially any intermediate value, e.g., the output of hash_to_field) may be able to carry out a dictionary attack. To mitigate such attacks, it is recommended to first execute a more costly key derivation function (e.g., PBKDF2 [RFC8018], scrypt [RFC7914], or Argon2 [RFC9106]) on the password, then hash the output of that function to the target elliptic curve. For collision resistance, the hash underlying the key derivation function should be chosen according to the guidelines listed in Section 5.3.1.

10.3. Constant-Time Requirements

Constant-time implementations of all functions in this document are **STRONGLY RECOMMENDED** for all uses, to avoid leaking information via side channels. It is especially important to use a constant-time implementation when inputs to an encoding are secret values; in such cases, constant-time implementations are **REQUIRED** for security against timing attacks (e.g., [VR20]). When constant-time implementations are required, all basic operations and utility functions must be implemented in constant time, as discussed in Section 4. In some applications (e.g., embedded systems), leakage through other side channels (e.g., power or electromagnetic side channels) may be pertinent. Defending against such leakage is outside the scope of this document, because the nature of the leakage and the appropriate defense depend on the application.

10.4. `encode_to_curve`: Output Distribution and Indifferentiability

The `encode_to_curve` function (Section 3) returns points sampled from a distribution that is statistically far from uniform. This distribution is bounded roughly as follows: first, it includes at least one eighth of the points in G , and second, the probability of points in the distribution varies by at most a factor of four. These bounds hold when `encode_to_curve` is instantiated with any of the `map_to_curve` functions in Section 6.

The bounds above are derived from several works in the literature. Specifically:

- * Shallue and van de Woestijne [SW06] and Fouque and Tibouchi [FT12] derive bounds on the Shallue-van de Woestijne mapping (Section 6.6.1).
- * Fouque and Tibouchi [FT10] and Tibouchi [T14] derive bounds for the Simplified SWU mapping (Sections 6.6.2 and 6.6.3).
- * Bernstein et al. [BHK13] derive bounds for the Elligator 2 mapping (Sections 6.7.1 and 6.8.2).

Indifferentiability of `encode_to_curve` follows from an argument similar to the one given by Brier et al. [BCIMRT10]; we briefly sketch this argument as follows. Consider an ideal random oracle $H_c()$ that samples from the distribution induced by the `map_to_curve` function called by `encode_to_curve`, and assume for simplicity that the target elliptic curve has cofactor 1 (a similar argument applies for non-unity cofactors). Indifferentiability holds just if it is possible to efficiently simulate the "inner" random oracle in `encode_to_curve`, namely, `hash_to_field`. The simulator works as follows: on a fresh query `msg`, the simulator queries $H_c(msg)$ and receives a point P in the image of `map_to_curve` (if `msg` is the same as a prior query, the simulator just returns the value it gave in response to that query). The simulator then computes the possible preimages of P under `map_to_curve`, i.e., elements u of F such that `map_to_curve(u) == P` (Tibouchi [T14] shows that this can be done efficiently for the Shallue-van de Woestijne and Simplified SWU maps, and Bernstein et al. show the same for Elligator 2). The simulator selects one such preimage at random and returns this value as the simulated output of the "inner" random oracle. By hypothesis, $H_c()$ samples from the distribution induced by `map_to_curve` on a uniformly random input element of F , so this value is uniformly random and induces the correct point P when passed through `map_to_curve`.

10.5. `hash_to_field` Security

The `hash_to_field` function, defined in Section 5, is indifferentiable from a random oracle [MRH04] when `expand_message` (Section 5.3) is modeled as a random oracle. Since indifferentiability proofs are composable, this also holds when `expand_message` is proved indifferentiable from a random oracle relative to an underlying primitive that is modeled as a random oracle. When following the guidelines in Section 5.3, both variants of `expand_message` defined in

that section meet this requirement (see also Section 10.6).

We very briefly sketch the indistinguishability argument for `hash_to_field`. Notice that each integer mod p that `hash_to_field` returns (i.e., each element of the vector representation of F) is a member of an equivalence class of roughly 2^k integers of length $\log_2(p) + k$ bits, all of which are equal modulo p . For each integer mod p that `hash_to_field` returns, the simulator samples one member of this equivalence class at random and outputs the byte string returned by `I2OSP`. (Notice that this is essentially the inverse of the `hash_to_field` procedure.)

10.6. `expand_message_xmd` Security

The `expand_message_xmd` function, defined in Section 5.3.1, is indistinguishable from a random oracle [MRH04] when one of the following holds:

1. H is indistinguishable from a random oracle,
2. H is a sponge-based hash function whose inner function is modeled as a random transformation or random permutation [BDPV08], or
3. H is a Merkle-Damgaard hash function whose compression function is modeled as a random oracle [CDMP05].

For cases (1) and (2), the indistinguishability of `expand_message_xmd` follows directly from the indistinguishability of H .

For case (3), i.e., where H is a Merkle-Damgaard hash function, indistinguishability follows from [CDMP05], Theorem 5. In particular, `expand_message_xmd` computes b_0 by prefixing the message with one block of zeros plus auxiliary information (length, counter, and DST). Then, each of the output blocks b_i , $i \geq 1$ in `expand_message_xmd` is the result of invoking H on a unique, prefix-free encoding of b_0 . This is true, first because the length of the input to all such invocations is equal and fixed by the choice of H and DST, and second because each such input has a unique suffix (because of the inclusion of the counter byte `I2OSP(i, 1)`).

The essential difference between the construction discussed in [CDMP05] and `expand_message_xmd` is that the latter hashes a counter appended to `strxor(b_0, b_{i-1})` (`{#hashtofield-expand-xmd}`, step 10) rather than to b_0 . This approach increases the Hamming distance between inputs to different invocations of H , which reduces the likelihood that nonidealities in H affect the distribution of the b_i values.

We note that `expand_message_xmd` can be used to instantiate a general-purpose indistinguishable functionality with variable-length output based on any hash function meeting one of the above criteria. Applications that use `expand_message_xmd` outside of `hash_to_field` should ensure domain separation by picking a distinct value for DST.

10.7. Domain Separation for `expand_message` Variants

As discussed in Section 2.2.5, the purpose of domain separation is to ensure that security analyses of cryptographic protocols that query multiple independent random oracles remain valid even if all of these random oracles are instantiated based on one underlying function H .

The `expand_message` variants in this document (Section 5.3) ensure domain separation by appending a suffix-free-encoded domain separation tag `DST_prime` to all strings hashed by H , an underlying hash or extendable-output function. (Other `expand_message` variants that follow the guidelines in Section 5.3.4 are expected to behave similarly, but these should be analyzed on a case-by-case basis.) For security, applications that use the same function H outside of `expand_message` should enforce domain separation between those uses of H and `expand_message`, and they should separate all of these from uses of H in other applications.

This section suggests four methods for enforcing domain separation from `expand_message` variants, explains how each method achieves domain separation, and lists the situations in which each is appropriate. These methods share a high-level structure: the application designer fixes a tag `DST_ext` distinct from `DST_prime` and augments calls to H with `DST_ext`. Each method augments calls to H differently, and each may impose additional requirements on `DST_ext`.

These methods can be used to instantiate multiple domain-separated functions (e.g., H_1 and H_2) by selecting distinct `DST_ext` values for each (e.g., `DST_ext1`, `DST_ext2`).

1. (Suffix-only domain separation.) This method is useful when domain-separating invocations of H from `expand_message_xmd` or `expand_message_xof`. It is not appropriate for domain-separating `expand_message` from HMAC- H [RFC2104]; for that purpose, see method 4.

To instantiate a suffix-only domain-separated function H_{so} , compute

$$H_{so}(msg) = H(msg \parallel DST_ext)$$

`DST_ext` should be suffix-free encoded (e.g., by appending one byte encoding the length of `DST_ext`) to make it infeasible to find distinct (msg, DST_ext) pairs that hash to the same value.

This method ensures domain separation because all distinct invocations of H have distinct suffixes, since `DST_ext` is distinct from `DST_prime`.

2. (Prefix-suffix domain separation.) This method can be used in the same cases as the suffix-only method.

To instantiate a prefix-suffix domain-separated function H_{ps} , compute

$$H_{ps}(msg) = H(DST_ext \parallel msg \parallel I2OSP(0, 1))$$

`DST_ext` should be prefix-free encoded (e.g., by adding a one-byte

prefix that encodes the length of DST_ext) to make it infeasible to find distinct (msg, DST_ext) pairs that hash to the same value.

This method ensures domain separation because appending the byte I2OSP(0, 1) ensures that inputs to H inside Hps are distinct from those inside expand_message. Specifically, the final byte of DST_prime encodes the length of DST, which is required to be nonzero (Section 3.1, requirement 2), and DST_prime is always appended to invocations of H inside expand_message.

3. (Prefix-only domain separation.) This method is only useful for domain-separating invocations of H from expand_message_xmd. It does not give domain separation for expand_message_xof or HMAC-H.

To instantiate a prefix-only domain-separated function Hpo, compute

```
Hpo(msg) = H(DST_ext || msg)
```

In order for this method to give domain separation, DST_ext should be at least b bits long, where b is the number of bits output by the hash function H. In addition, at least one of the first b bits must be nonzero. Finally, DST_ext should be prefix-free encoded (e.g., by adding a one-byte prefix that encodes the length of DST_ext) to make it infeasible to find distinct (msg, DST_ext) pairs that hash to the same value.

This method ensures domain separation as follows. First, since DST_ext contains at least one nonzero bit among its first b bits, it is guaranteed to be distinct from the value Z_pad (Section 5.3.1, step 4), which ensures that all inputs to H are distinct from the input used to generate b_0 in expand_message_xmd. Second, since DST_ext is at least b bits long, it is almost certainly distinct from the values b_0 and strxor(b_0, b_(i - 1)), and therefore all inputs to H are distinct from the inputs used to generate b_i, i >= 1, with high probability.

4. (XMD-HMAC domain separation.) This method is useful for domain-separating invocations of H inside HMAC-H (i.e., HMAC [RFC2104] instantiated with hash function H) from expand_message_xmd. It also applies to HKDF-H (i.e., HKDF [RFC5869] instantiated with hash function H), as discussed below.

Specifically, this method applies when HMAC-H is used with a non-secret key to instantiate a random oracle based on a hash function H (note that expand_message_xmd can also be used for this purpose; see Section 10.6). When using HMAC-H with a high-entropy secret key, domain separation is not necessary; see discussion below.

To choose a non-secret HMAC key DST_key that ensures domain separation from expand_message_xmd, compute

```
DST_key_preimage = "DERIVE-HMAC-KEY-" || DST_ext || I2OSP(0, 1)
```

`DST_key = H(DST_key_preimage)`

Then, to instantiate the random oracle `Hro` using HMAC-H, compute

`Hro(msg) = HMAC-H(DST_key, msg)`

The trailing zero byte in `DST_key_preimage` ensures that this value is distinct from inputs to `H` inside `expand_message_xmd` (because all such inputs have suffix `DST_prime`, which cannot end with a zero byte as discussed above). This ensures domain separation because, with overwhelming probability, all inputs to `H` inside of HMAC-H using key `DST_key` have prefixes that are distinct from the values `Z_pad`, `b_0`, and `strxor(b_0, b_(i - 1))` inside of `expand_message_xmd`.

For uses of HMAC-H that instantiate a private random oracle by fixing a high-entropy secret key, domain separation from `expand_message_xmd` is not necessary. This is because, similarly to the case above, all inputs to `H` inside HMAC-H using this secret key almost certainly have distinct prefixes from all inputs to `H` inside `expand_message_xmd`.

Finally, this method can be used with HKDF-H [RFC5869] by fixing the salt input to HKDF-Extract to `DST_key`, computed as above. This ensures domain separation for HKDF-Extract by the same argument as for HMAC-H using `DST_key`. Moreover, assuming that the input keying material (IKM) supplied to HKDF-Extract has sufficiently high entropy (say, commensurate with the security parameter), the HKDF-Expand step is domain-separated by the same argument as for HMAC-H with a high-entropy secret key (since a pseudorandom key is exactly that).

10.8. Target Security Levels

Each ciphersuite specifies a target security level (in bits) for the underlying curve. This parameter ensures the corresponding `hash_to_field` instantiation is conservative and correct. We stress that this parameter is only an upper bound on the security level of the curve and is neither a guarantee nor endorsement of its suitability for a given application. Mathematical and cryptographic advancements may reduce the effective security level for any curve.

11. References

11.1. Normative References

- [Err4730] RFC Errata, "Erratum ID 4730", RFC 7748, July 2016, <<https://www.rfc-editor.org/errata/eid4730>>.
- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, DOI 10.17487/RFC2119, March 1997, <<https://www.rfc-editor.org/info/rfc2119>>.
- [RFC7748] Langley, A., Hamburg, M., and S. Turner, "Elliptic Curves for Security", RFC 7748, DOI 10.17487/RFC7748, January

2016, <<https://www.rfc-editor.org/info/rfc7748>>.

[RFC8017] Moriarty, K., Ed., Kaliski, B., Jonsson, J., and A. Rusch, "PKCS #1: RSA Cryptography Specifications Version 2.2", RFC 8017, DOI 10.17487/RFC8017, November 2016, <<https://www.rfc-editor.org/info/rfc8017>>.

[RFC8174] Leiba, B., "Ambiguity of Uppercase vs Lowercase in RFC 2119 Key Words", BCP 14, RFC 8174, DOI 10.17487/RFC8174, May 2017, <<https://www.rfc-editor.org/info/rfc8174>>.

11.2. Informative References

[AFQTZ14] Aranha, D. F., Fouque, P.-A., Qian, C., Tibouchi, M., and J. C. Zapalowicz, "Binary Elligator Squared", In Selected Areas in Cryptography - SAC 2014, pages 20-37, DOI 10.1007/978-3-319-13051-4_2, November 2014, <https://doi.org/10.1007/978-3-319-13051-4_2>.

[AR13] Adj, G. and F. Rodríguez-Henríquez, "Square Root Computation over Even Extension Fields", In IEEE Transactions on Computers. vol 63 issue 11, pages 2829-2841, DOI 10.1109/TC.2013.145, November 2014, <<https://doi.org/10.1109/TC.2013.145>>.

[BBJLP08] Bernstein, D. J., Birkner, P., Joye, M., Lange, T., and C. Peters, "Twisted Edwards Curves", In AFRICACRYPT 2008, pages 389-405, DOI 10.1007/978-3-540-68164-9_26, June 2008, <https://doi.org/10.1007/978-3-540-68164-9_26>.

[BCIMRT10] Brier, E., Coron, J.-S., Icart, T., Madore, D., Randriam, H., and M. Tibouchi, "Efficient Indifferentiable Hashing into Ordinary Elliptic Curves", In Advances in Cryptology - CRYPTO 2010, pages 237-254, DOI 10.1007/978-3-642-14623-7_13, August 2010, <https://doi.org/10.1007/978-3-642-14623-7_13>.

[BDPV08] Bertoni, G., Daemen, J., Peeters, M., and G. Van Assche, "On the Indifferentiability of the Sponge Construction", In Advances in Cryptology - EUROCRYPT 2008, pages 181-197, DOI 10.1007/978-3-540-78967-3_11, April 2008, <https://doi.org/10.1007/978-3-540-78967-3_11>.

[BF01] Boneh, D. and M. Franklin, "Identity-Based Encryption from the Weil Pairing", In Advances in Cryptology - CRYPTO 2001, pages 213-229, DOI 10.1007/3-540-44647-8_13, August 2001, <https://doi.org/10.1007/3-540-44647-8_13>.

[BHKL13] Bernstein, D. J., Hamburg, M., Krasnova, A., and T. Lange, "Elligator: elliptic-curve points indistinguishable from uniform random strings", In Proceedings of the 2013 ACM SIGSAC Conference on Computer and Communications Security, pages 967-980, DOI 10.1145/2508859.2516734, November 2013, <<https://doi.org/10.1145/2508859.2516734>>.

[BLAKE2X] Aumasson, J.-P., Neves, S., Wilcox-O'Hearn, Z., and C.

Winnerlein, "BLAKE2X", December 2016,
<<https://blake2.net/blake2x.pdf>>.

- [BLMP19] Bernstein, D. J., Lange, T., Martindale, C., and L. Panny, "Quantum Circuits for the CSIDH: Optimizing Quantum Evaluation of Isogenies", In Advances in Cryptology - EUROCRYPT 2019, pages 409-441, DOI 10.1007/978-3-030-17656-3, May 2019, <https://doi.org/10.1007/978-3-030-17656-3_15>.
- [BLS-SIG] Boneh, D., Gorbunov, S., Wahby, R. S., Wee, H., Wood, C. A., and Z. Zhang, "BLS Signatures", Work in Progress, Internet-Draft, draft-irtf-cfrg-bls-signature-05, 16 June 2022, <<https://datatracker.ietf.org/doc/html/draft-irtf-cfrg-bls-signature-05>>.
- [BLS01] Boneh, D., Lynn, B., and H. Shacham, "Short Signatures from the Weil Pairing", In Journal of Cryptology, vol 17, pages 297-319, DOI 10.1007/s00145-004-0314-9, July 2004, <<https://doi.org/10.1007/s00145-004-0314-9>>.
- [BLS03] Barreto, P. S. L. M., Lynn, B., and M. Scott, "Constructing Elliptic Curves with Prescribed Embedding Degrees", In Security in Communication Networks, pages 257-267, DOI 10.1007/3-540-36413-7_19, September 2002, <https://doi.org/10.1007/3-540-36413-7_19>.
- [BLS12-381]

Bowe, S., "BLS12-381: New zk-SNARK Elliptic Curve Construction", March 2017,
<<https://electriccoin.co/blog/new-snark-curve/>>.
- [BM92] Bellare, S. M. and M. Merritt, "Encrypted key exchange: password-based protocols secure against dictionary attacks", In IEEE Symposium on Security and Privacy - Oakland 1992, pages 72-84, DOI 10.1109/RISP.1992.213269, May 1992, <<https://doi.org/10.1109/RISP.1992.213269>>.
- [BMP00] Boyko, V., MacKenzie, P., and S. Patel, "Provably Secure Password-Authenticated Key Exchange Using Diffie-Hellman", In Advances in Cryptology - EUROCRYPT 2000, pages 156-171, DOI 10.1007/3-540-45539-6_12, May 2000, <https://doi.org/10.1007/3-540-45539-6_12>.
- [BN05] Barreto, P. S. L. M. and M. Naehrig, "Pairing-Friendly Elliptic Curves of Prime Order", In Selected Areas in Cryptography 2005, pages 319-331, DOI 10.1007/11693383_22, 2006, <https://doi.org/10.1007/11693383_22>.
- [BP17] Budroni, A. and F. Pintore, "Efficient hash maps to \mathbb{G}_2 on BLS curves", Cryptology ePrint Archive, Paper 2017/419, May 2017, <<https://eprint.iacr.org/2017/419>>.
- [BR93] Bellare, M. and P. Rogaway, "Random oracles are practical: a paradigm for designing efficient protocols", In

Proceedings of the 1993 ACM Conference on Computer and Communications Security, pages 62-73,
DOI 10.1145/168588.168596, December 1993,
<<https://doi.org/10.1145/168588.168596>>.

- [C93] Cohen, H., "A Course in Computational Algebraic Number Theory", Springer-Verlag, ISBN 9783642081422,
DOI 10.1007/978-3-662-02945-9, 1993,
<<https://doi.org/10.1007/978-3-662-02945-9>>.
- [CDMP05] Coron, J.-S., Dodis, Y., Malinaud, C., and P. Puniya, "Merkle-Damgård Revisited: How to Construct a Hash Function", In Advances in Cryptology -- CRYPTO 2005, pages 430-448, DOI 10.1007/11535218_26, August 2005,
<https://doi.org/10.1007/11535218_26>.
- [CFADLNV05] Cohen, H., Frey, G., Avanzi, R., Doche, C., Lange, T., Nguyen, K., and F. Vercauteren, "Handbook of Elliptic and Hyperelliptic Curve Cryptography", Chapman and Hall / CRC, ISBN 9781584885184, 2005,
<<https://www.crcpress.com/9781584885184>>.
- [CK11] Couveignes, J.-M. and J.-G. Kammerer, "The geometry of flex tangents to a cubic curve and its parameterizations", In Journal of Symbolic Computation, vol 47 issue 3, pages 266-281, DOI 10.1016/j.jsc.2011.11.003, March 2012,
<<https://doi.org/10.1016/j.jsc.2011.11.003>>.
- [F11] Farashahi, R. R., "Hashing into Hessian Curves", In AFRICACRYPT 2011, pages 278-289,
DOI 10.1007/978-3-642-21969-6_17, July 2011,
<https://doi.org/10.1007/978-3-642-21969-6_17>.
- [FFSTV13] Farashahi, R. R., Fouque, P.-A., Shparlinski, I. E., Tibouchi, M., and J. F. Voloch, "Indifferentiable deterministic hashing to elliptic and hyperelliptic curves", In Mathematics of Computation. vol 82, pages 491-512, DOI 10.1090/S0025-5718-2012-02606-8, 2013,
<<https://doi.org/10.1090/S0025-5718-2012-02606-8>>.
- [FIPS180-4] National Institute of Standards and Technology (NIST), "Secure Hash Standard (SHS)", FIPS 180-4,
DOI 10.6028/NIST.FIPS.180-4, August 2015,
<<https://nvlpubs.nist.gov/nistpubs/FIPS/NIST.FIPS.180-4.pdf>>.
- [FIPS186-4] National Institute of Standards and Technology (NIST), "Digital Signature Standard (DSS)", FIPS 186-4,
DOI 10.6028/NIST.FIPS.186-4, July 2013,
<<https://nvlpubs.nist.gov/nistpubs/FIPS/NIST.FIPS.186-4.pdf>>.
- [FIPS202] National Institute of Standards and Technology (NIST),

"SHA-3 Standard: Permutation-Based Hash and Extendable-Output Functions", FIPS 202, DOI 10.6028/NIST.FIPS.202, August 2015, <<https://nvlpubs.nist.gov/nistpubs/FIPS/NIST.FIPS.202.pdf>>.

- [FJT13] Fouque, P.-A., Joux, A., and M. Tibouchi, "Injective Encodings to Elliptic Curves", In ACISP 2013, pages 203-218, DOI 10.1007/978-3-642-39059-3_14, 2013, <https://doi.org/10.1007/978-3-642-39059-3_14>.
- [FKR11] Fuentes-Castañeda, L., Knapp, E., and F. Rodriguez-Henriquez, "Faster Hashing to G²", In Selected Areas in Cryptography, pages 412-430, DOI 10.1007/978-3-642-28496-0_25, August 2011, <https://doi.org/10.1007/978-3-642-28496-0_25>.
- [FSV09] Farashahi, R. R., Shparlinski, I. E., and J. F. Voloch, "On hashing into elliptic curves", In Journal of Mathematical Cryptology, vol 3 no 4, pages 353-360, DOI 10.1515/JMC.2009.022, March 2009, <<https://doi.org/10.1515/JMC.2009.022>>.
- [FT10] Fouque, P.-A. and M. Tibouchi, "Estimating the Size of the Image of Deterministic Hash Functions to Elliptic Curves", In Progress in Cryptology - LATINCRYPT 2010, pages 81-91, DOI 10.1007/978-3-642-14712-8_5, August 2010, <https://doi.org/10.1007/978-3-642-14712-8_5>.
- [FT12] Fouque, P.-A. and M. Tibouchi, "Indifferentiable Hashing to Barreto--Naehrig Curves", In Progress in Cryptology - LATINCRYPT 2012, pages 1-17, DOI 10.1007/978-3-642-33481-8_1, 2012, <https://doi.org/10.1007/978-3-642-33481-8_1>.
- [H20] Hamburg, M., "Indifferentiable hashing from Elligator 2", Cryptology ePrint Archive, Paper 2020/1513, 2020, <<https://eprint.iacr.org/2020/1513>>.
- [hash2curve-repo] "Hashing to Elliptic Curves", commit 664b135, June 2022, <<https://github.com/cfrg/draft-irtf-cfrg-hash-to-curve>>.
- [Icart09] Icart, T., "How to Hash into Elliptic Curves", In Advances in Cryptology - CRYPTO 2009, pages 303-316, DOI 10.1007/978-3-642-03356-8_18, August 2009, <https://doi.org/10.1007/978-3-642-03356-8_18>.
- [J96] Jablon, D. P., "Strong password-only authenticated key exchange", In SIGCOMM Computer Communication Review, vol 26 issue 5, pages 5-26, DOI 10.1145/242896.242897, October 1996, <<https://doi.org/10.1145/242896.242897>>.
- [jubjub-fq] "zkcrypto/jubjub - fq.rs", 2019, <<https://github.com/zkcrypto/jubjub/pull/18>>.

- [KLR10] Kammerer, J.-G., Lercier, R., and G. Renault, "Encoding Points on Hyperelliptic Curves over Finite Fields in Deterministic Polynomial Time", In Pairing-Based Cryptography - Pairing 2010, pages 278-297, DOI 10.1007/978-3-642-17455-1_18, 2010, <https://doi.org/10.1007/978-3-642-17455-1_18>.
- [L13] Langley, A., "Implementing Elligator for Curve25519", December 2013, <<https://www.imperialviolet.org/2013/12/25/elligator.html>>.
- [LBB19] Lipp, B., Blanchet, B., and K. Bhargavan, "A Mechanised Cryptographic Proof of the WireGuard Virtual Private Network Protocol", In INRIA Research Report 9269, April 2019, <<https://hal.inria.fr/hal-02100345/>>.
- [MOV96] Menezes, A. J., van Oorschot, P. C., and S. A. Vanstone, "Handbook of Applied Cryptography", CRC Press, ISBN 9780849385230, October 1996, <<http://cacr.uwaterloo.ca/hac/>>.
- [MRH04] Maurer, U., Renner, R., and C. Holenstein, "Indifferentiability, Impossibility Results on Reductions, and Applications to the Random Oracle Methodology", In TCC 2004: Theory of Cryptography, pages 21-39, DOI 10.1007/978-3-540-24638-1_2, February 2004, <https://doi.org/10.1007/978-3-540-24638-1_2>.
- [MRV99] Micali, S., Rabin, M., and S. Vadhan, "Verifiable random functions", 40th Annual Symposium on Foundations of Computer Science (Cat. No.99CB37039), pages 120-130, DOI 10.1109/SFFCS.1999.814584, October 1999, <<https://doi.org/10.1109/SFFCS.1999.814584>>.
- [MT98] Matsumoto, M. and T. Nishimura, "Mersenne twister: A 623-dimensionally equidistributed uniform pseudo-random number generator", In ACM Transactions on Modeling and Computer Simulation (TOMACS), vol 8 issue 1, pages 3-30, DOI 10.1145/272991.272995, January 1998, <<https://doi.org/10.1145/272991.272995>>.
- [NR97] Naor, M. and O. Reingold, "Number-theoretic constructions of efficient pseudo-random functions", In Proceedings 38th Annual Symposium on Foundations of Computer Science, pages 458-467, DOI 10.1109/SFCS.1997.646134, October 1997, <<https://doi.org/10.1109/SFCS.1997.646134>>.
- [OPRFs] Davidson, A., Faz-Hernandez, A., Sullivan, N., and C. A. Wood, "Oblivious Pseudorandom Functions (OPRFs) using Prime-Order Groups", Work in Progress, Internet-Draft, draft-irtf-cfrg-voprf-21, 21 February 2023, <<https://datatracker.ietf.org/doc/html/draft-irtf-cfrg-voprf-21>>.
- [p1363.2] IEEE, "IEEE Standard Specification for Password-Based Public-Key Cryptography Techniques", IEEE 1363.2-2008,

September 2008,
<https://standards.ieee.org/standard/1363_2-2008.html>.

- [p1363a] IEEE, "IEEE Standard Specifications for Public-Key Cryptography - Amendment 1: Additional Techniques", IEEE 1363a-2004, March 2004,
<<https://standards.ieee.org/standard/1363a-2004.html>>.
- [P20] Pornin, T., "Efficient Elliptic Curve Operations On Microcontrollers With Finite Field Extensions", Cryptology ePrint Archive, Paper 2020/009, 2020,
<<https://eprint.iacr.org/2020/009>>.
- [RCB16] Renes, J., Costello, C., and L. Batina, "Complete Addition Formulas for Prime Order Elliptic Curves", In Advances in Cryptology - EUROCRYPT 2016, pages 403-428, DOI 10.1007/978-3-662-49890-3_16, April 2016,
<https://doi.org/10.1007/978-3-662-49890-3_16>.
- [RFC2104] Krawczyk, H., Bellare, M., and R. Canetti, "HMAC: Keyed-Hashing for Message Authentication", RFC 2104, DOI 10.17487/RFC2104, February 1997,
<<https://www.rfc-editor.org/info/rfc2104>>.
- [RFC5869] Krawczyk, H. and P. Eronen, "HMAC-based Extract-and-Expand Key Derivation Function (HKDF)", RFC 5869, DOI 10.17487/RFC5869, May 2010,
<<https://www.rfc-editor.org/info/rfc5869>>.
- [RFC7693] Saarinen, M., Ed. and J. Aumasson, "The BLAKE2 Cryptographic Hash and Message Authentication Code (MAC)", RFC 7693, DOI 10.17487/RFC7693, November 2015,
<<https://www.rfc-editor.org/info/rfc7693>>.
- [RFC7914] Percival, C. and S. Josefsson, "The scrypt Password-Based Key Derivation Function", RFC 7914, DOI 10.17487/RFC7914, August 2016, <<https://www.rfc-editor.org/info/rfc7914>>.
- [RFC8018] Moriarty, K., Ed., Kaliski, B., and A. Rusch, "PKCS #5: Password-Based Cryptography Specification Version 2.1", RFC 8018, DOI 10.17487/RFC8018, January 2017,
<<https://www.rfc-editor.org/info/rfc8018>>.
- [RFC9106] Biryukov, A., Dinu, D., Khovratovich, D., and S. Josefsson, "Argon2 Memory-Hard Function for Password Hashing and Proof-of-Work Applications", RFC 9106, DOI 10.17487/RFC9106, September 2021,
<<https://www.rfc-editor.org/info/rfc9106>>.
- [ristretto255-decaf448] de Valence, H., Grigg, J., Hamburg, M., Lovecruft, I., Tankersley, G., and F. Valsorda, "The ristretto255 and decaf448 Groups", Work in Progress, Internet-Draft, draft-irtf-cfrg-ristretto255-decaf448-07, 3 April 2023,
<<https://datatracker.ietf.org/doc/html/draft-irtf-cfrg-ristretto255-decaf448-07>>.

- [RSS11] Ristenpart, T., Shacham, H., and T. Shrimpton, "Careful with Composition: Limitations of the Indifferentiability Framework", In Advances in Cryptology - EUROCRYPT 2011, pages 487-506, DOI 10.1007/978-3-642-20465-4_27, May 2011, <https://doi.org/10.1007/978-3-642-20465-4_27>.
- [S05] Skałba, M., "Points on elliptic curves over finite fields", In Acta Arithmetica, vol 117 no 3, pages 293-301, DOI 10.4064/aa117-3-7, 2005, <<https://doi.org/10.4064/aa117-3-7>>.
- [S85] Schoof, R., "Elliptic curves over finite fields and the computation of square roots mod p", In Mathematics of Computation, vol 44 issue 170, pages 483-494, DOI 10.1090/S0025-5718-1985-0777280-6, April 1985, <<https://doi.org/10.1090/S0025-5718-1985-0777280-6>>.
- [SAGE] The Sage Developers, "SageMath, the Sage Mathematics Software System", <<https://www.sagemath.org>>.
- [SBCDK09] Scott, M., Benger, N., Charlemagne, M., Dominguez Perez, L. J., and E. J. Kachisa, "Fast Hashing to G2 on Pairing-Friendly Curves", In Pairing-Based Cryptography - Pairing 2009, pages 102-113, DOI 10.1007/978-3-642-03298-1_8, August 2009, <https://doi.org/10.1007/978-3-642-03298-1_8>.
- [SEC1] Standards for Efficient Cryptography Group (SECG), "SEC 1: Elliptic Curve Cryptography", May 2009, <<http://www.secg.org/sec1-v2.pdf>>.
- [SEC2] Standards for Efficient Cryptography Group (SECG), "SEC 2: Recommended Elliptic Curve Domain Parameters", January 2010, <<http://www.secg.org/sec2-v2.pdf>>.
- [SS04] Schinzel, A. and M. Skałba, "On equations $y^2 = x^n + k$ in a finite field", In Bulletin Polish Academy of Sciences. Mathematics, vol 52 no 3, pages 223-226, DOI 10.4064/ba52-3-1, 2004, <<https://doi.org/10.4064/ba52-3-1>>.
- [SW06] Shallue, A. and C. E. van de Woestijne, "Construction of Rational Points on Elliptic Curves over Finite Fields", In Algorithmic Number Theory - ANTS 2006, pages 510-524, DOI 10.1007/11792086_36, July 2006, <https://doi.org/10.1007/11792086_36>.
- [T14] Tibouchi, M., "Elligator Squared: Uniform Points on Elliptic Curves of Prime Order as Uniform Random Strings", In Financial Cryptography and Data Security - FC 2014, pages 139-156, DOI 10.1007/978-3-662-45472-5_10, November 2014, <https://doi.org/10.1007/978-3-662-45472-5_10>.
- [TK17] Tibouchi, M. and T. Kim, "Improved elliptic curve hashing and point representation", In Designs, Codes, and

Cryptography, vol 82, pages 161-177,
DOI 10.1007/s10623-016-0288-2, January 2017,
<<https://doi.org/10.1007/s10623-016-0288-2>>.

- [U07] Ulas, M., "Rational Points on Certain Hyperelliptic Curves over Finite Fields", In Bulletin Polish Academy of Science. Mathematics, vol 55 no 2, pages 97-104, DOI 10.4064/ba55-2-1, July 2007, <<https://doi.org/10.4064/ba55-2-1>>.
- [VR20] Vanhoef, M. and E. Ronen, "Dragonblood: Analyzing the Dragonfly Handshake of WPA3 and EAP-pwd", In IEEE Symposium on Security & Privacy (SP), May 2020, <<https://eprint.iacr.org/2019/383>>.
- [VRF] Goldberg, S., Reyzin, L., Papadopoulos, D., and J. Včelák, "Verifiable Random Functions (VRFs)", Work in Progress, Internet-Draft, draft-irtf-cfrg-vrf-15, 9 August 2022, <<https://datatracker.ietf.org/doc/html/draft-irtf-cfrg-vrf-15>>.
- [W08] Washington, L. C., "Elliptic Curves: Number Theory and Cryptography, Second Edition", Chapman and Hall / CRC, ISBN 9781420071467, April 2008, <<https://www.crcpress.com/9781420071467>>.
- [W19] Wahby, R. S., "An explicit, generic parameterization for the Shallue--van de Woestijne map", commit e2a625f, March 2020, <https://github.com/cfrg/draft-irtf-cfrg-hash-to-curve/blob/draft-irtf-cfrg-hash-to-curve-14/doc/svdw_params.pdf>.
- [WB19] Wahby, R. S. and D. Boneh, "Fast and simple constant-time hashing to the BLS12-381 elliptic curve", In IACR Transactions on Cryptographic Hardware and Embedded Systems, vol 2019 issue 4, Cryptology ePrint Archive, Paper 2019/403, DOI 10.13154/tches.v2019.i4.154-179, August 2019, <<https://eprint.iacr.org/2019/403>>.

Appendix A. Related Work

The problem of mapping arbitrary bit strings to elliptic curve points has been the subject of both practical and theoretical research. This section briefly describes the background and research results that underlie the recommendations in this document. This section is provided for informational purposes only.

A naive but generally insecure method of mapping a string msg to a point on an elliptic curve E having n points is to first fix a point P that generates the elliptic curve group, and a hash function H_n from bit strings to integers less than n ; then compute $H_n(msg) * P$, where the $*$ operator represents scalar multiplication. The reason this approach is insecure is that the resulting point has a known discrete log relationship to P . Thus, except in cases where this method is specified by the protocol, it must not be used; doing so risks catastrophic security failures.

Boneh et al. [BLS01] describe an encoding method they call MapToGroup, which works roughly as follows: first, use the input string to initialize a pseudorandom number generator, then use the generator to produce a value x in F . If x is the x -coordinate of a point on the elliptic curve, output that point. Otherwise, generate a new value x in F and try again. Since a random value x in F has probability about $1/2$ of corresponding to a point on the curve, the expected number of tries is just two. However, the running time of this method, which is generally referred to as a probabilistic try-and-increment algorithm, depends on the input string. As such, it is not safe to use in protocols sensitive to timing side channels, as was exemplified by the Dragonblood attack [VR20].

Schinzel and Skalba [SS04] introduce a method of constructing elliptic curve points deterministically, for a restricted class of curves and a very small number of points. Skalba [S05] generalizes this construction to more curves and more points on those curves. Shallue and van de Woestijne [SW06] further generalize and simplify Skalba's construction, yielding concretely efficient maps to a constant fraction of the points on almost any curve. Fouque and Tibouchi [FT12] give a parameterization of this mapping for Barreto-Naehrig pairing-friendly curves [BN05].

Ulas [U07] describes a simpler version of the Shallue-van de Woestijne map, and Brier et al. [BCIMRT10] give a further simplification, which the authors call the "Simplified SWU" map. That simplified map applies only to fields of characteristic $p = 3 \pmod{4}$; Wahby and Boneh [WB19] generalize to fields of any characteristic and give further optimizations.

Boneh and Franklin give a deterministic algorithm mapping to certain supersingular curves over fields of characteristic $p = 2 \pmod{3}$ [BF01]. Icart gives another deterministic algorithm that maps to any curve over a field of characteristic $p = 2 \pmod{3}$ [Icart09]. Several extensions and generalizations follow this work, including [FSV09], [FT10], [KLR10], [F11], and [CK11].

Following the work of Farashahi [F11], Fouque et al. [FJT13] describe a mapping to curves over fields of characteristic $p = 3 \pmod{4}$ having a number of points divisible by 4. Bernstein et al. [BHK13] optimize this mapping and describe a related mapping that they call "Elligator 2," which applies to any curve over a field of odd characteristic having a point of order 2. This includes Curve25519 and Curve448, both of which are CFRG-recommended curves [RFC7748]. Bernstein et al. [BLMP19] extend the Elligator 2 map to a class of supersingular curves over fields of characteristic $p = 3 \pmod{4}$.

An important caveat regarding all of the above deterministic mapping functions is that none of them map to the entire curve, but rather to some fraction of the points. This means that they cannot be used directly to construct a random oracle that outputs points on the curve.

Brier et al. [BCIMRT10] give two solutions to this problem. The first, which Brier et al. prove applies to Icart's method, computes

$f(H_0(\text{msg})) + f(H_1(\text{msg}))$ for two distinct hash functions H_0 and H_1 from bit strings to F and a mapping f from F to the elliptic curve E . The second, which applies to essentially all deterministic mappings but is more costly, computes $f(H_0(\text{msg})) + H_2(\text{msg}) * P$, where P is a generator of the elliptic curve group, H_2 is a hash from bit strings to integers modulo r , and r is the order of the elliptic curve group.

Farashahi et al. [FFSTV13] improve the analysis of the first method, showing that it applies to essentially all deterministic mappings. Tibouchi and Kim [TK17] further refine the analysis and describe additional optimizations.

Complementary to the problem of mapping from bit strings to elliptic curve points, Bernstein et al. [BHK13] study the problem of mapping from elliptic curve points to uniformly random bit strings, giving solutions for a class of curves that includes Montgomery and twisted Edwards curves. Tibouchi [T14] and Aranha et al. [AFQTZ14] generalize these results. This document does not deal with this complementary problem.

Appendix B. Hashing to ristretto255

ristretto255 [ristretto255-decaf448] provides a prime-order group based on curve25519 [RFC7748]. This section describes `hash_to_ristretto255`, which implements a random-oracle encoding to this group that has a uniform output distribution (Section 2.2.3) and the same security properties and interface as the `hash_to_curve` function (Section 3).

The ristretto255 API defines a one-way map ([ristretto255-decaf448], Section 4.3.4); this section refers to that map as `ristretto255_map`.

The `hash_to_ristretto255` function MUST be instantiated with an `expand_message` function that conforms to the requirements given in Section 5.3. In addition, it MUST use a domain separation tag constructed as described in Section 3.1, and all domain separation recommendations given in Section 10.7 apply when implementing protocols that use `hash_to_ristretto255`.

`hash_to_ristretto255(msg)`

Parameters:

- `DST`, a domain separation tag (see discussion above).
- `expand_message`, a function that expands a byte string and domain separation tag into a uniformly random byte string (see discussion above).
- `ristretto255_map`, the one-way map from the ristretto255 API.

Input: `msg`, an arbitrary-length byte string.

Output: `P`, an element of the ristretto255 group.

Steps:

1. `uniform_bytes = expand_message(msg, DST, 64)`
2. `P = ristretto255_map(uniform_bytes)`
3. return `P`

Since `hash_to_ristretto255` is not a hash-to-curve suite, it does not have a Suite ID. If a similar identifier is needed, it **MUST** be constructed following the guidelines in Section 8.10, with the following parameters:

- * `CURVE_ID`: "ristretto255"
- * `HASH_ID`: as described in Section 8.10
- * `MAP_ID`: "R255MAP"
- * `ENC_VAR`: "R0"

For example, if `expand_message` is `expand_message_xmd` using SHA-512, the **REQUIRED** identifier is:

`ristretto255_XMD:SHA-512_R255MAP_R0_`

Appendix C. Hashing to decaf448

Similar to `ristretto255`, `decaf448` [`ristretto255-decaf448`] provides a prime-order group based on `curve448` [RFC7748]. This section describes `hash_to_decaf448`, which implements a random-oracle encoding to this group that has a uniform output distribution (Section 2.2.3) and the same security properties and interface as the `hash_to_curve` function (Section 3).

The `decaf448` API defines a one-way map ([`ristretto255-decaf448`], Section 5.3.4); this section refers to that map as `decaf448_map`.

The `hash_to_decaf448` function **MUST** be instantiated with an `expand_message` function that conforms to the requirements given in Section 5.3. In addition, it **MUST** use a domain separation tag constructed as described in Section 3.1, and all domain separation recommendations given in Section 10.7 apply when implementing protocols that use `hash_to_decaf448`.

`hash_to_decaf448(msg)`

Parameters:

- `DST`, a domain separation tag (see discussion above).
- `expand_message`, a function that expands a byte string and domain separation tag into a uniformly random byte string (see discussion above).
- `decaf448_map`, the one-way map from the `decaf448` API.

Input: `msg`, an arbitrary-length byte string.

Output: `P`, an element of the `decaf448` group.

Steps:

1. `uniform_bytes = expand_message(msg, DST, 112)`
2. `P = decaf448_map(uniform_bytes)`
3. return `P`

Since `hash_to_decaf448` is not a hash-to-curve suite, it does not have a Suite ID. If a similar identifier is needed, it **MUST** be

constructed following the guidelines in Section 8.10, with the following parameters:

- * CURVE_ID: "decaf448"
- * HASH_ID: as described in Section 8.10
- * MAP_ID: "D448MAP"
- * ENC_VAR: "R0"

For example, if `expand_message` is `expand_message_xof` using SHAKE256, the REQUIRED identifier is:

`decaf448_XOF:SHAKE256_D448MAP_R0_`

Appendix D. Rational Maps

This section gives rational maps that can be used when hashing to twisted Edwards or Montgomery curves.

Given a twisted Edwards curve, Appendix D.1 shows how to derive a corresponding Montgomery curve and how to map from that curve to the twisted Edwards curve. This mapping may be used when hashing to twisted Edwards curves as described in Section 6.8.

Given a Montgomery curve, Appendix D.2 shows how to derive a corresponding Weierstrass curve and how to map from that curve to the Montgomery curve. This mapping can be used to hash to Montgomery or twisted Edwards curves via the Shallue-van de Woestijne method (Section 6.6.1) or Simplified SWU method (Section 6.6.2), as follows:

- * For Montgomery curves, first map to the Weierstrass curve, then convert to Montgomery coordinates via the mapping.
- * For twisted Edwards curves, compose the mapping from Weierstrass to Montgomery with the mapping from Montgomery to twisted Edwards (Appendix D.1) to obtain a Weierstrass curve and a mapping to the target twisted Edwards curve. Map to this Weierstrass curve, then convert to Edwards coordinates via the mapping.

D.1. Generic Mapping from Montgomery to Twisted Edwards

This section gives a generic birational map between twisted Edwards and Montgomery curves.

The map in this section is a simplified version of the map given in [BBJLP08], Theorem 3.2. Specifically, this section's map handles exceptional cases in a simplified way that is geared towards hashing to a twisted Edwards curve's prime-order subgroup.

The twisted Edwards curve

$$a * v^2 + w^2 = 1 + d * v^2 * w^2$$

is birationally equivalent to the Montgomery curve

$$K * t^2 = s^3 + J * s^2 + s$$

which has the form required by the Elligator 2 mapping of Section 6.7.1. The coefficients of the Montgomery curve are

$$* J = 2 * (a + d) / (a - d)$$

$$* K = 4 / (a - d)$$

The rational map from the point (s, t) on the above Montgomery curve to the point (v, w) on the twisted Edwards curve is given by

$$* v = s / t$$

$$* w = (s - 1) / (s + 1)$$

This mapping is undefined when $t == 0$ or $s == -1$, i.e., when the denominator of either of the above rational functions is zero. Implementations MUST detect exceptional cases and return the value $(v, w) = (0, 1)$, which is the identity point on all twisted Edwards curves.

The following straight-line implementation of the above rational map handles the exceptional cases.

`monty_to_edw_generic(s, t)`

Input: (s, t) , a point on the curve $K * t^2 = s^3 + J * s^2 + s$.

Output: (v, w) , a point on an equivalent twisted Edwards curve.

```

1. tv1 = s + 1
2. tv2 = tv1 * t           # (s + 1) * t
3. tv2 = inv0(tv2)         # 1 / ((s + 1) * t)
4. v = tv2 * tv1           # 1 / t
5. v = v * s               # s / t
6. w = tv2 * t             # 1 / (s + 1)
7. tv1 = s - 1
8. w = w * tv1             # (s - 1) / (s + 1)
9. e = tv2 == 0
10. w = CM0V(w, 1, e)     # handle exceptional case
11. return (v, w)
```

For completeness, we also give the inverse relations. (Note that this map is not required when hashing to twisted Edwards curves.) The coefficients of the twisted Edwards curve corresponding to the above Montgomery curve are

$$* a = (J + 2) / K$$

$$* d = (J - 2) / K$$

The rational map from the point (v, w) on the twisted Edwards curve to the point (s, t) on the Montgomery curve is given by

$$* s = (1 + w) / (1 - w)$$

$$* \quad t = (1 + w) / (v * (1 - w))$$

The mapping is undefined when $v == 0$ or $w == 1$. When the goal is to map into the prime-order subgroup of the Montgomery curve, it suffices to return the identity point on the Montgomery curve in the exceptional cases.

D.2. Mapping from Weierstrass to Montgomery

The rational map from the point (s, t) on the Montgomery curve

$$K * t^2 = s^3 + J * s^2 + s$$

to the point (x, y) on the equivalent Weierstrass curve

$$y^2 = x^3 + A * x + B$$

is given by

$$* \quad A = (3 - J^2) / (3 * K^2)$$

$$* \quad B = (2 * J^3 - 9 * J) / (27 * K^3)$$

$$* \quad x = (3 * s + J) / (3 * K)$$

$$* \quad y = t / K$$

The inverse map, from the point (x, y) to the point (s, t) , is given by

$$* \quad s = (3 * K * x - J) / 3$$

$$* \quad t = y * K$$

This mapping can be used to apply the Shallue-van de Woestijne method (Section 6.6.1) or Simplified SWU method (Section 6.6.2) to Montgomery curves.

Appendix E. Isogeny Maps for Suites

This section specifies the isogeny maps for the secp256k1 and BLS12-381 suites listed in Section 8.

These maps are given in terms of affine coordinates. Wahby and Boneh ([WB19], Section 4.3) show how to evaluate these maps in a projective coordinate system (Appendix G.1), which avoids modular inversions.

Refer to [hash2curve-repo] for a Sage [SAGE] script that constructs these isogenies.

E.1. 3-Isogeny Map for secp256k1

This section specifies the isogeny map for the secp256k1 suite listed in Section 8.7.

The 3-isogeny map from (x', y') on E' to (x, y) on E is given by the following rational functions:

* $x = x_num / x_den$, where

$$- x_num = k_(1,3) * x'^3 + k_(1,2) * x'^2 + k_(1,1) * x' + k_(1,0)$$

$$- x_den = x'^2 + k_(2,1) * x' + k_(2,0)$$

* $y = y' * y_num / y_den$, where

$$- y_num = k_(3,3) * x'^3 + k_(3,2) * x'^2 + k_(3,1) * x' + k_(3,0)$$

$$- y_den = x'^3 + k_(4,2) * x'^2 + k_(4,1) * x' + k_(4,0)$$

The constants used to compute x_num are as follows:

$$* k_(1,0) = 0x8e38e38e38e38e38e38e38e38e38e38e38e38e38e38e38daaaaa8c7$$

$$* k_(1,1) = 0x7d3d4c80bc321d5b9f315cea7fd44c5d595d2fc0bf63b92dffff1044f17c6581$$

$$* k_(1,2) = 0x534c328d23f234e6e2a413deca25caece4506144037c40314ecbd0b53d9dd262$$

$$* k_(1,3) = 0x8e38e38e38e38e38e38e38e38e38e38e38e38e38e38e38daaaaa88c$$

The constants used to compute x_den are as follows:

$$* k_(2,0) = 0xd35771193d94918a9ca34ccbb7b640dd86cd409542f8487d9fe6b745781eb49b$$

$$* k_(2,1) = 0xedadc6f64383dc1df7c4b2d51b54225406d36b641f5e41bbc52a56612a8c6d14$$

The constants used to compute y_num are as follows:

$$* k_(3,0) = 0x4bda12f684bda12f684bda12f684bda12f684bda12f684b8e38e23c$$

$$* k_(3,1) = 0xc75e0c32d5cb7c0fa9d0a54b12a0a6d5647ab046d686da6fdffc90fc201d71a3$$

$$* k_(3,2) = 0x29a6194691f91a73715209ef6512e576722830a201be2018a765e85a9ecee931$$

$$* k_(3,3) = 0x2f684bda12f684bda12f684bda12f684bda12f684bda12f38e38d84$$

The constants used to compute y_den are as follows:

$$* k_(4,0) =$$

0xfffeffff93b

- * $k_{(4,1)} = 0x7a06534bb8bdb49fd5e9e6632722c2989467c1bfc8e8d978dfb425d2685c2573$
- * $k_{(4,2)} = 0x6484aa716545ca2cf3a70c3fa8fe337e0a3d21162f0d6299a7bf8192bfd2a76f$

E.2. 11-Isogeny Map for BLS12-381 G1

The 11-isogeny map from (x', y') on E' to (x, y) on E is given by the following rational functions:

- * $x = x_num / x_den$, where
 - $x_num = k_{(1,11)} * x'^{11} + k_{(1,10)} * x'^{10} + k_{(1,9)} * x'^9 + \dots + k_{(1,0)}$
 - $x_den = x'^{10} + k_{(2,9)} * x'^9 + k_{(2,8)} * x'^8 + \dots + k_{(2,0)}$
- * $y = y' * y_num / y_den$, where
 - $y_num = k_{(3,15)} * x'^{15} + k_{(3,14)} * x'^{14} + k_{(3,13)} * x'^{13} + \dots + k_{(3,0)}$
 - $y_den = x'^{15} + k_{(4,14)} * x'^{14} + k_{(4,13)} * x'^{13} + \dots + k_{(4,0)}$

The constants used to compute x_num are as follows:

- * $k_{(1,0)} = 0x11a05f2b1e833340b809101dd99815856b303e88a2d7005ff2627b56cdb4e2c85610c2d5f2e62d6eaeac1662734649b7$
- * $k_{(1,1)} = 0x17294ed3e943ab2f0588bab22147a81c7c17e75b2f6a8417f565e33c70d1e86b4838f2a6f318c356e834eef1b3cb83bb$
- * $k_{(1,2)} = 0xd54005db97678ec1d1048c5d10a9a1bce032473295983e56878e501ec68e25c958c3e3d2a09729fe0179f9dac9edcb0$
- * $k_{(1,3)} = 0x1778e7166fcc6db74e0609d307e55412d7f5e4656a8dbf25f1b33289f1b330835336e25ce3107193c5b388641d9b6861$
- * $k_{(1,4)} = 0xe99726a3199f4436642b4b3e4118e5499db995a1257fb3f086eeb65982fac18985a286f301e77c451154ce9ac8895d9$
- * $k_{(1,5)} = 0x1630c3250d7313ff01d1201bf7a74ab5db3cb17dd952799b9ed3ab9097e68f90a0870d2dcae73d19cd13c1c66f652983$
- * $k_{(1,6)} = 0xd6ed6553fe44d296a3726c38ae652bfb11586264f0f8ce19008e218f9c86b2a8da25128c1052ecadd7f225a139ed84$
- * $k_{(1,7)} = 0x17b81e7701abdbe2e8743884d1117e53356de5ab275b4db1a682c62ef0f2753339b7c8f8c8f475af9ccb5618e3f0c88e$
- * $k_{(1,8)} = 0x80d3cf1f9a78fc47b90b33563be990dc43b756ce79f5574a2c596c928c5d1de4fa295f296b74e956d71986a8497e317$

- * $k_{(1,9)} = 0x169b1f8e1bcfa7c42e0c37515d138f22dd2ecb803a0c5c99676314baf4bb1b7fa3190b2edc0327797f241067be390c9e$
- * $k_{(1,10)} = 0x10321da079ce07e272d8ec09d2565b0dfa7dccdde6787f96d50af36003b14866f69b771f8c285decca67df3f1605fb7b$
- * $k_{(1,11)} = 0x6e08c248e260e70bd1e962381edee3d31d79d7e22c837bc23c0bf1bc24c6b68c24b1b80b64d391fa9c8ba2e8ba2d229$

The constants used to compute x_{den} are as follows:

- * $k_{(2,0)} = 0x8ca8d548cff19ae18b2e62f4bd3fa6f01d5ef4ba35b48ba9c9588617fc8ac62b558d681be343df8993cf9fa40d21b1c$
- * $k_{(2,1)} = 0x12561a5deb559c4348b4711298e536367041e8ca0cf0800c0126c2588c48bf5713daa8846cb026e9e5c8276ec82b3bff$
- * $k_{(2,2)} = 0xb2962fe57a3225e8137e629bfff2991f6f89416f5a718cd1fca64e00b11aceacd6a3d0967c94fedcfcc239ba5cb83e19$
- * $k_{(2,3)} = 0x3425581a58ae2fec83aafef7c40eb545b08243f16b1655154cca8abc28d6fd04976d5243eecf5c4130de8938dc62cd8$
- * $k_{(2,4)} = 0x13a8e162022914a80a6f1d5f43e7a07dffdfc759a12062bb8d6b44e833b306da9bd29ba81f35781d539d395b3532a21e$
- * $k_{(2,5)} = 0xe7355f8e4e667b955390f7f0506c6e9395735e9ce9cad4d0a43bcef24b8982f7400d24bc4228f11c02df9a29f6304a5$
- * $k_{(2,6)} = 0x772caacf16936190f3e0c63e0596721570f5799af53a1894e2e073062aede9cea73b3538f0de06cec2574496ee84a3a$
- * $k_{(2,7)} = 0x14a7ac2a9d64a8b230b3f5b074cf01996e7f63c21bca68a81996e1cdf9822c580fa5b9489d11e2d311f7d99bbdcc5a5e$
- * $k_{(2,8)} = 0xa10ecf6ada54f825e920b3dafc7a3cce07f8d1d7161366b74100da67f39883503826692abba43704776ec3a79a1d641$
- * $k_{(2,9)} = 0x95fc13ab9e92ad4476d6e3eb3a56680f682b4ee96f7d03776df533978f31c1593174e4b4b7865002d6384d168ecdd0a$

The constants used to compute y_{num} are as follows:

- * $k_{(3,0)} = 0x90d97c81ba24ee0259d1f094980dcfa11ad138e48a869522b52af6c956543d3cd0c7aee9b3ba3c2be9845719707bb33$
- * $k_{(3,1)} = 0x134996a104ee5811d51036d776fb46831223e96c254f383d0f906343eb67ad34d6c56711962fa8bfe097e75a2e41c696$
- * $k_{(3,2)} = 0xcc786baa966e66f4a384c86a3b49942552e2d658a31ce2c344be4b91400da7d26d521628b00523b8dfe240c72de1f6$
- * $k_{(3,3)} = 0x1f86376e8981c217898751ad8746757d42aa7b90eeb791c09e4a3ec03251cf9de405aba9ec61deca6355c77b0e5f4cb$

- * $k_{(3,4)} = 0x8cc03fdefe0ff135caf4fe2a21529c4195536fbe3ce50b879833fd221351adc2ee7f8dc099040a841b6daecf2e8fedb$
- * $k_{(3,5)} = 0x16603fca40634b6a2211e11db8f0a6a074a7d0d4afadb7bd76505c3d3ad5544e203f6326c95a807299b23ab13633a5f0$
- * $k_{(3,6)} = 0x4ab0b9bcfac1bbcb2c977d027796b3ce75bb8ca2be184cb5231413c4d634f3747a87ac2460f415ec961f8855fe9d6f2$
- * $k_{(3,7)} = 0x987c8d5333ab86fde9926bd2ca6c674170a05bfe3bdd81ffd038da6c26c842642f64550fedfe935a15e4ca31870fb29$
- * $k_{(3,8)} = 0x9fc4018bd96684be88c9e221e4da1bb8f3abd16679dc26c1e8b6e6a1f20cabe69d65201c78607a360370e577bdba587$
- * $k_{(3,9)} = 0xe1bba7a1186bdb5223abde7ada14a23c42a0ca7915af6fe06985e7ed1e4d43b9b3f7055dd4eba6f2bafaaebca731c30$
- * $k_{(3,10)} = 0x19713e47937cd1be0dfd0b8f1d43fb93cd2fcbcb6caf493fd1183e416389e61031bf3a5cce3fbafce813711ad011c132$
- * $k_{(3,11)} = 0x18b46a908f36f6deb918c143fed2edcc523559b8aaf0c2462e6bfe7f911f643249d9cdf41b44d606ce07c8a4d0074d8e$
- * $k_{(3,12)} = 0xb182cac101b9399d155096004f53f447aa7b12a3426b08ec02710e807b4633f06c851c1919211f20d4c04f00b971ef8$
- * $k_{(3,13)} = 0x245a394ad1eca9b72fc00ae7be315dc757b3b080d4c158013e6632d3c40659cc6cf90ad1c232a6442d9d3f5db980133$
- * $k_{(3,14)} = 0x5c129645e44cf1102a159f748c4a3fc5e673d81d7e86568d9ab0f5d396a7ce46ba1049b6579afb7866b1e715475224b$
- * $k_{(3,15)} = 0x15e6be4e990f03ce4ea50b3b42df2eb5cb181d8f84965a3957add4fa95af01b2b665027efec01c7704b456be69c8b604$

The constants used to compute y_{den} are as follows:

- * $k_{(4,0)} = 0x16112c4c3a9c98b252181140fad0eae9601a6de578980be6eec3232b5be72e7a07f3688ef60c206d01479253b03663c1$
- * $k_{(4,1)} = 0x1962d75c2381201e1a0cbd6c43c348b885c84ff731c4d59ca4a10356f453e01f78a4260763529e3532f6102c2e49a03d$
- * $k_{(4,2)} = 0x58df3306640da276faaae7d6e8eb15778c4855551ae7f310c35a5dd279cd2eca6757cd636f96f891e2538b53dbf67f2$
- * $k_{(4,3)} = 0x16b7d288798e5395f20d23bf89edb4d1d115c5dbddbcd30e123da489e726af41727364f2c28297ada8d26d98445f5416$
- * $k_{(4,4)} = 0xbe0e079545f43e4b00cc912f8228ddcc6d19c9f0f69bbb0542eda0fc9dec916a20b15dc0fd2ededda39142311a5001d$
- * $k_{(4,5)} = 0x8d9e5297186db2d9fb266eaac783182b70152c65550d881c5ecd87b6f0f5a6449f38db9dfa9cce202c6477faaf9b7ac$

- * $k_{(4,6)} = 0x166007c08a99db2fc3ba8734ace9824b5eecfdfa8d0cf8ef5dd365bc400a0051d5fa9c01a58b1fb93d1a1399126a775c$
- * $k_{(4,7)} = 0x16a3ef08be3ea7ea03bcddfabb66ff6ee5a4375efa1f4fd7feb34fd206357132b920f5b00801dee460ee415a15812ed9$
- * $k_{(4,8)} = 0x1866c8ed336c61231a1be54fd1d74cc4f9fb0ce4c6af5920abc5750c4bf39b4852cfe2f7bb9248836b233d9d55535d4a$
- * $k_{(4,9)} = 0x167a55cda70a6e1cea820597d94a84903216f763e13d87bb5308592e7ea7d4fbc7385ea3d529b35e346ef48bb8913f55$
- * $k_{(4,10)} = 0x4d2f259eea405bd48f010a01ad2911d9c6dd039bb61a6290e591b36e636a5c871a5c29f4f83060400f8b49cba8f6aa8$
- * $k_{(4,11)} = 0xaccbb67481d033ff5852c1e48c50c477f94ff8aefce42d28c0f9a88cea7913516f968986f7ebbea9684b529e2561092$
- * $k_{(4,12)} = 0xad6b9514c767fe3c3613144b45f1496543346d98adf02267d5cee9a00d9b8693000763e3b90ac11e99b138573345cc$
- * $k_{(4,13)} = 0x2660400eb2e4f3b628bdd0d53cd76f2bf565b94e72927c1cb748df27942480e420517bd8714cc80d1fadc1326ed06f7$
- * $k_{(4,14)} = 0xe0fa1d816ddc03e6b24255e0d7819c171c40f65e273b853324efcd6356caa205ca2f570f13497804415473a1d634b8f$

E.3. 3-Isogeny Map for BLS12-381 G2

The 3-isogeny map from (x', y') on E' to (x, y) on E is given by the following rational functions:

- * $x = x_{\text{num}} / x_{\text{den}}$, where
 - $x_{\text{num}} = k_{(1,3)} * x'^3 + k_{(1,2)} * x'^2 + k_{(1,1)} * x' + k_{(1,0)}$
 - $x_{\text{den}} = x'^2 + k_{(2,1)} * x' + k_{(2,0)}$
- * $y = y' * y_{\text{num}} / y_{\text{den}}$, where
 - $y_{\text{num}} = k_{(3,3)} * x'^3 + k_{(3,2)} * x'^2 + k_{(3,1)} * x' + k_{(3,0)}$
 - $y_{\text{den}} = x'^3 + k_{(4,2)} * x'^2 + k_{(4,1)} * x' + k_{(4,0)}$

The constants used to compute x_{num} are as follows:

- * $k_{(1,0)} = 0x5c759507e8e333ebb5b7a9a47d7ed8532c52d39fd3a042a88b58423c50ae15d5c2638e343d9c71c6238aaaaaaaa97d6 + 0x5c759507e8e333ebb5b7a9a47d7ed8532c52d39fd3a042a88b58423c50ae15d5c2638e343d9c71c6238aaaaaaaa97d6 * I$
- * $k_{(1,1)} = 0x11560bf17baa99bc32126fced787c88f984f87adf7ae0c7f9a208c6b4f20a4181472aaa9cb8d555526a9fffffffffc71a * I$

- * $k_{(1,2)} = 0x11560bf17baa99bc32126fced787c88f984f87adf7ae0c7f9a208c6b4f20a4181472aaa9cb8d555526a9fffffffffc71e + 0x8ab05f8bdd54cde190937e76bc3e447cc27c3d6fbd7063fcd104635a790520c0a395554e5c6aaaa9354ffffffffffe38d * I$
- * $k_{(1,3)} = 0x171d6541fa38ccfaed6dea691f5fb614cb14b4e7f4e810aa22d6108f142b85757098e38d0f671c7188e2aaaaaaaa5ed1$

The constants used to compute x_{den} are as follows:

- * $k_{(2,0)} = 0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabffffeb153ffffb9feffffffffffaa63 * I$
- * $k_{(2,1)} = 0xc + 0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabffffeb153ffffb9feffffffffffaa9f * I$

The constants used to compute y_{num} are as follows:

- * $k_{(3,0)} = 0x1530477c7ab4113b59a4c18b076d11930f7da5d4a07f649bf54439d87d27e500fc8c25ebf8c92f6812cfc71c71c6d706 + 0x1530477c7ab4113b59a4c18b076d11930f7da5d4a07f649bf54439d87d27e500fc8c25ebf8c92f6812cfc71c71c6d706 * I$
- * $k_{(3,1)} = 0x5c759507e8e33ebb5b7a9a47d7ed8532c52d39fd3a042a88b58423c50ae15d5c2638e343d9c71c6238aaaaaaaa97be * I$
- * $k_{(3,2)} = 0x11560bf17baa99bc32126fced787c88f984f87adf7ae0c7f9a208c6b4f20a4181472aaa9cb8d555526a9fffffffffc71c + 0x8ab05f8bdd54cde190937e76bc3e447cc27c3d6fbd7063fcd104635a790520c0a395554e5c6aaaa9354ffffffffffe38f * I$
- * $k_{(3,3)} = 0x124c9ad43b6cf79bfbf7043de3811ad0761b0f37a1e26286b0e977c69aa274524e79097a56dc4bd9e1b371c71c718b10$

The constants used to compute y_{den} are as follows:

- * $k_{(4,0)} = 0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabffffeb153ffffb9feffffffffffa8fb + 0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabffffeb153ffffb9feffffffffffa8fb * I$
- * $k_{(4,1)} = 0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabffffeb153ffffb9feffffffffffa9d3 * I$
- * $k_{(4,2)} = 0x12 + 0x1a0111ea397fe69a4b1ba7b6434bacd764774b84f38512bf6730d2a0f6b0f6241eabffffeb153ffffb9feffffffffffaa99 * I$

Appendix F. Straight-Line Implementations of Deterministic Mappings

This section gives straight-line implementations of the mappings of Section 6. These implementations are generic, i.e., they are defined for any curve and field. Appendix G gives further implementations that are optimized for specific classes of curves and fields.

F.1. Shallue-van de Woestijne Method

This section gives a straight-line implementation of the Shallue-van de Woestijne method for any Weierstrass curve of the form given in Section 6.6. See Section 6.6.1 for information on the constants used in this mapping.

Note that the constant c_3 below MUST be chosen such that $\text{sgn}_0(c_3) = 0$. In other words, if the square-root computation returns a value cx such that $\text{sgn}_0(cx) = 1$, set $c_3 = -cx$; otherwise, set $c_3 = cx$.

`map_to_curve_svdw(u)`

Input: u , an element of F .

Output: (x, y) , a point on E .

Constants:

1. $c_1 = g(Z)$
2. $c_2 = -Z / 2$
3. $c_3 = \text{sqrt}(-g(Z) * (3 * Z^2 + 4 * A))$ # $\text{sgn}_0(c_3)$ MUST equal 0
4. $c_4 = -4 * g(Z) / (3 * Z^2 + 4 * A)$

Steps:

1. $tv1 = u^2$
2. $tv1 = tv1 * c_1$
3. $tv2 = 1 + tv1$
4. $tv1 = 1 - tv1$
5. $tv3 = tv1 * tv2$
6. $tv3 = \text{inv}_0(tv3)$
7. $tv4 = u * tv1$
8. $tv4 = tv4 * tv3$
9. $tv4 = tv4 * c_3$
10. $x1 = c_2 - tv4$
11. $gx1 = x1^2$
12. $gx1 = gx1 + A$
13. $gx1 = gx1 * x1$
14. $gx1 = gx1 + B$
15. $e1 = \text{is_square}(gx1)$
16. $x2 = c_2 + tv4$
17. $gx2 = x2^2$
18. $gx2 = gx2 + A$
19. $gx2 = gx2 * x2$
20. $gx2 = gx2 + B$
21. $e2 = \text{is_square}(gx2) \text{ AND NOT } e1$ # Avoid short-circuit logic ops
22. $x3 = tv2^2$
23. $x3 = x3 * tv3$
24. $x3 = x3^2$
25. $x3 = x3 * c_4$
26. $x3 = x3 + Z$
27. $x = \text{CMOV}(x3, x1, e1)$ # $x = x1$ if $gx1$ is square, else $x = x3$
28. $x = \text{CMOV}(x, x2, e2)$ # $x = x2$ if $gx2$ is square and $gx1$ is not
29. $gx = x^2$
30. $gx = gx + A$
31. $gx = gx * x$
32. $gx = gx + B$
33. $y = \text{sqrt}(gx)$
34. $e3 = \text{sgn}_0(u) == \text{sgn}_0(y)$
35. $y = \text{CMOV}(-y, y, e3)$ # Select correct sign of y

36. return (x, y)

F.2. Simplified SWU Method

This section gives a straight-line implementation of the Simplified SWU method for any Weierstrass curve of the form given in Section 6.6. See Section 6.6.2 for information on the constants used in this mapping.

This optimized, straight-line procedure applies to any base field. The `sqrt_ratio` subroutine is defined in Appendix F.2.1.

`map_to_curve_simple_swu(u)`

Input: `u`, an element of F .

Output: (x, y) , a point on E .

Steps:

```
1.  tv1 = u^2
2.  tv1 = Z * tv1
3.  tv2 = tv1^2
4.  tv2 = tv2 + tv1
5.  tv3 = tv2 + 1
6.  tv3 = B * tv3
7.  tv4 = CMOV(Z, -tv2, tv2 != 0)
8.  tv4 = A * tv4
9.  tv2 = tv3^2
10. tv6 = tv4^2
11. tv5 = A * tv6
12. tv2 = tv2 + tv5
13. tv2 = tv2 * tv3
14. tv6 = tv6 * tv4
15. tv5 = B * tv6
16. tv2 = tv2 + tv5
17.  x = tv1 * tv3
18. (is_gx1_square, y1) = sqrt_ratio(tv2, tv6)
19.  y = tv1 * u
20.  y = y * y1
21.  x = CMOV(x, tv3, is_gx1_square)
22.  y = CMOV(y, y1, is_gx1_square)
23. e1 = sgn0(u) == sgn0(y)
24.  y = CMOV(-y, y, e1)
25.  x = x / tv4
26. return (x, y)
```

F.2.1. `sqrt_ratio` Subroutine

This section defines three variants of the `sqrt_ratio` subroutine used by the above procedure. The first variant can be used with any field; the others are optimized versions for specific fields.

The routines given in this section depend on the constant Z from the Simplified SWU map. For correctness, `sqrt_ratio` and `map_to_curve_simple_swu` MUST use the same value for Z .

F.2.1.1. `sqrt_ratio` for Any Field

`sqrt_ratio(u, v)`

Parameters:

- F , a finite field of characteristic p and order $q = p^m$.
- Z , the constant from the Simplified SWU map.

Input: u and v , elements of F , where $v \neq 0$.

Output: (b, y) , where

- $b = \text{True}$ and $y = \sqrt{u / v}$ if (u / v) is square in F , and
- $b = \text{False}$ and $y = \sqrt{Z * (u / v)}$ otherwise.

Constants:

1. c_1 , the largest integer such that 2^{c_1} divides $q - 1$.
2. $c_2 = (q - 1) / (2^{c_1})$ # Integer arithmetic
3. $c_3 = (c_2 - 1) / 2$ # Integer arithmetic
4. $c_4 = 2^{c_1} - 1$ # Integer arithmetic
5. $c_5 = 2^{(c_1 - 1)}$ # Integer arithmetic
6. $c_6 = Z^{c_2}$
7. $c_7 = Z^{((c_2 + 1) / 2)}$

Procedure:

1. $tv_1 = c_6$
2. $tv_2 = v^{c_4}$
3. $tv_3 = tv_2^2$
4. $tv_3 = tv_3 * v$
5. $tv_5 = u * tv_3$
6. $tv_5 = tv_5^{c_3}$
7. $tv_5 = tv_5 * tv_2$
8. $tv_2 = tv_5 * v$
9. $tv_3 = tv_5 * u$
10. $tv_4 = tv_3 * tv_2$
11. $tv_5 = tv_4^{c_5}$
12. $isQR = tv_5 == 1$
13. $tv_2 = tv_3 * c_7$
14. $tv_5 = tv_4 * tv_1$
15. $tv_3 = \text{CMOV}(tv_2, tv_3, isQR)$
16. $tv_4 = \text{CMOV}(tv_5, tv_4, isQR)$
17. for i in $(c_1, c_1 - 1, \dots, 2)$:
 18. $tv_5 = i - 2$
 19. $tv_5 = 2^{tv_5}$
 20. $tv_5 = tv_4^{tv_5}$
 21. $e_1 = tv_5 == 1$
 22. $tv_2 = tv_3 * tv_1$
 23. $tv_1 = tv_1 * tv_1$
 24. $tv_5 = tv_4 * tv_1$
 25. $tv_3 = \text{CMOV}(tv_2, tv_3, e_1)$
 26. $tv_4 = \text{CMOV}(tv_5, tv_4, e_1)$
27. return $(isQR, tv_3)$

F.2.1.2. Optimized `sqrt_ratio` for $q = 3 \bmod 4$

`sqrt_ratio_3mod4(u, v)`

Parameters:

- F , a finite field of characteristic p and order $q = p^m$,

where $q = 3 \bmod 4$.

- Z , the constant from the Simplified SWU map.

Input: u and v , elements of F , where $v \neq 0$.

Output: (b, y) , where

$b = \text{True}$ and $y = \sqrt{u / v}$ if (u / v) is square in F , and
 $b = \text{False}$ and $y = \sqrt{Z * (u / v)}$ otherwise.

Constants:

1. $c1 = (q - 3) / 4$ # Integer arithmetic
2. $c2 = \sqrt{-Z}$

Procedure:

1. $tv1 = v^2$
2. $tv2 = u * v$
3. $tv1 = tv1 * tv2$
4. $y1 = tv1^{c1}$
5. $y1 = y1 * tv2$
6. $y2 = y1 * c2$
7. $tv3 = y1^2$
8. $tv3 = tv3 * v$
9. $isQR = tv3 == u$
10. $y = \text{CMOV}(y2, y1, isQR)$
11. return $(isQR, y)$

F.2.1.3. Optimized $\sqrt{\text{ratio}}$ for $q = 5 \bmod 8$

$\sqrt{\text{ratio}}_{5\text{mod}8}(u, v)$

Parameters:

- F , a finite field of characteristic p and order $q = p^m$,
where $q = 5 \bmod 8$.
- Z , the constant from the Simplified SWU map.

Input: u and v , elements of F , where $v \neq 0$.

Output: (b, y) , where

$b = \text{True}$ and $y = \sqrt{u / v}$ if (u / v) is square in F , and
 $b = \text{False}$ and $y = \sqrt{Z * (u / v)}$ otherwise.

Constants:

1. $c1 = (q - 5) / 8$
2. $c2 = \sqrt{-1}$
3. $c3 = \sqrt{Z / c2}$

Steps:

1. $tv1 = v^2$
2. $tv2 = tv1 * v$
3. $tv1 = tv1^2$
4. $tv2 = tv2 * u$
5. $tv1 = tv1 * tv2$
6. $y1 = tv1^{c1}$
7. $y1 = y1 * tv2$
8. $tv1 = y1 * c2$
9. $tv2 = tv1^2$
10. $tv2 = tv2 * v$
11. $e1 = tv2 == u$

```

12. y1 = CM0V(y1, tv1, e1)
13. tv2 = y1^2
14. tv2 = tv2 * v
15. isQR = tv2 == u
16. y2 = y1 * c3
17. tv1 = y2 * c2
18. tv2 = tv1^2
19. tv2 = tv2 * v
20. tv3 = Z * u
21. e2 = tv2 == tv3
22. y2 = CM0V(y2, tv1, e2)
23. y = CM0V(y2, y1, isQR)
24. return (isQR, y)

```

F.3. Elligator 2 Method

This section gives a straight-line implementation of the Elligator 2 method for any Montgomery curve of the form given in Section 6.7. See Section 6.7.1 for information on the constants used in this mapping.

Appendix G.2 gives optimized straight-line procedures that apply to specific classes of curves and base fields, including curve25519 and curve448 [RFC7748].

map_to_curve_elligator2(u)

Input: u, an element of F.

Output: (s, t), a point on M.

Constants:

1. $c1 = J / K$
2. $c2 = 1 / K^2$

Steps:

```

1. tv1 = u^2
2. tv1 = Z * tv1                # Z * u^2
3. e1 = tv1 == -1                # exceptional case: Z * u^2 == -1
4. tv1 = CM0V(tv1, 0, e1)        # if tv1 == -1, set tv1 = 0
5. x1 = tv1 + 1
6. x1 = inv0(x1)
7. x1 = -c1 * x1                # x1 = -(J / K) / (1 + Z * u^2)
8. gx1 = x1 + c1
9. gx1 = gx1 * x1
10. gx1 = gx1 + c2
11. gx1 = gx1 * x1              # gx1 = x1^3 + (J / K) * x1^2 + x1 / K^2
12. x2 = -x1 - c1
13. gx2 = tv1 * gx1
14. e2 = is_square(gx1)         # If is_square(gx1)
15. x = CM0V(x2, x1, e2)        # then x = x1, else x = x2
16. y2 = CM0V(gx2, gx1, e2)     # then y2 = gx1, else y2 = gx2
17. y = sqrt(y2)
18. e3 = sgn0(y) == 1
19. y = CM0V(y, -y, e2 XOR e3)  # fix sign of y
20. s = x * K
21. t = y * K

```

22. return (s, t)

Appendix G. Curve-Specific Optimized Sample Code

This section gives sample implementations optimized for some of the elliptic curves listed in Section 8. Sample Sage code [SAGE] for each algorithm can also be found in [hash2curve-repo].

G.1. Interface and Projective Coordinate Systems

The sample code in this section uses a different interface than the mappings of Section 6. Specifically, each mapping function in this section has the following signature:

$(x_n, x_d, y_n, y_d) = \text{map_to_curve}(u)$

The resulting affine point (x, y) is given by $(x_n / x_d, y_n / y_d)$.

The reason for this modified interface is that it enables further optimizations when working with points in a projective coordinate system. This is desirable, for example, when the resulting point will be immediately multiplied by a scalar, since most scalar multiplication algorithms operate on projective points.

Projective coordinates are also useful when implementing random-oracle encodings (Section 3). One reason is that, in general, point addition is faster using projective coordinates. Another reason is that, for Weierstrass curves, projective coordinates allow using complete addition formulas [RCB16]. This is especially convenient when implementing a constant-time encoding, because it eliminates the need for a special case when $Q_0 == Q_1$, which incomplete addition formulas usually do not handle.

The following are two commonly used projective coordinate systems and the corresponding conversions:

- * A point (X, Y, Z) in homogeneous projective coordinates corresponds to the affine point $(x, y) = (X / Z, Y / Z)$; the inverse conversion is given by $(X, Y, Z) = (x, y, 1)$. To convert (x_n, x_d, y_n, y_d) to homogeneous projective coordinates, compute $(X, Y, Z) = (x_n * y_d, y_n * x_d, x_d * y_d)$.
- * A point (X', Y', Z') in Jacobian projective coordinates corresponds to the affine point $(x, y) = (X' / Z'^2, Y' / Z'^3)$; the inverse conversion is given by $(X', Y', Z') = (x, y, 1)$. To convert (x_n, x_d, y_n, y_d) to Jacobian projective coordinates, compute $(X', Y', Z') = (x_n * x_d * y_d^2, y_n * y_d^2 * x_d^3, x_d * y_d)$.

G.2. Elligator 2

G.2.1. curve25519 ($q = 5 \pmod{8}$), $K = 1$)

The following is a straight-line implementation of Elligator 2 for curve25519 [RFC7748] as specified in Section 8.5.

This implementation can also be used for any Montgomery curve with $K = 1$ over $\text{GF}(q)$ where $q = 5 \pmod{8}$.

map_to_curve_elligator2_curve25519(u)

Input: u, an element of F.

Output: (xn, xd, yn, yd) such that (xn / xd, yn / yd) is a point on curve25519.

Constants:

1. c1 = (q + 3) / 8 # Integer arithmetic
2. c2 = 2^c1
3. c3 = sqrt(-1)
4. c4 = (q - 5) / 8 # Integer arithmetic

Steps:

1. tv1 = u^2
2. tv1 = 2 * tv1
3. xd = tv1 + 1 # Nonzero: -1 is square (mod p), tv1 is not
4. x1n = -J # x1 = x1n / xd = -J / (1 + 2 * u^2)
5. tv2 = xd^2
6. gxd = tv2 * xd # gxd = xd^3
7. gx1 = J * tv1 # x1n + J * xd
8. gx1 = gx1 * x1n # x1n^2 + J * x1n * xd
9. gx1 = gx1 + tv2 # x1n^2 + J * x1n * xd + xd^2
10. gx1 = gx1 * x1n # x1n^3 + J * x1n^2 * xd + x1n * xd^2
11. tv3 = gxd^2
12. tv2 = tv3^2 # gxd^4
13. tv3 = tv3 * gxd # gxd^3
14. tv3 = tv3 * gx1 # gx1 * gxd^3
15. tv2 = tv2 * tv3 # gx1 * gxd^7
16. y11 = tv2^c4 # (gx1 * gxd^7)^((p - 5) / 8)
17. y11 = y11 * tv3 # gx1 * gxd^3 * (gx1 * gxd^7)^((p - 5) / 8)
18. y12 = y11 * c3
19. tv2 = y11^2
20. tv2 = tv2 * gxd
21. e1 = tv2 == gx1
22. y1 = CMOV(y12, y11, e1) # If g(x1) is square, this is its sqrt
23. x2n = x1n * tv1 # x2 = x2n / xd = 2 * u^2 * x1n / xd
24. y21 = y11 * u
25. y21 = y21 * c2
26. y22 = y21 * c3
27. gx2 = gx1 * tv1 # g(x2) = gx2 / gxd = 2 * u^2 * g(x1)
28. tv2 = y21^2
29. tv2 = tv2 * gxd
30. e2 = tv2 == gx2
31. y2 = CMOV(y22, y21, e2) # If g(x2) is square, this is its sqrt
32. tv2 = y1^2
33. tv2 = tv2 * gxd
34. e3 = tv2 == gx1
35. xn = CMOV(x2n, x1n, e3) # If e3, x = x1, else x = x2
36. y = CMOV(y2, y1, e3) # If e3, y = y1, else y = y2
37. e4 = sgn0(y) == 1 # Fix sign of y
38. y = CMOV(y, -y, e3 XOR e4)
39. return (xn, xd, y, 1)

G.2.2. edwards25519

The following is a straight-line implementation of Elligator 2 for edwards25519 [RFC7748] as specified in Section 8.5. The subroutine `map_to_curve_elligator2_curve25519` is defined in Appendix G.2.1.

Note that the sign of the constant `c1` below is chosen as specified in Section 6.8.1, i.e., applying the rational map to the edwards25519 base point yields the curve25519 base point (see erratum [Err4730]).

`map_to_curve_elligator2_edwards25519(u)`

Input: `u`, an element of F .

Output: (x_n, x_d, y_n, y_d) such that $(x_n / x_d, y_n / y_d)$ is a point on edwards25519.

Constants:

1. `c1 = sqrt(-486664)` # `sgn0(c1)` MUST equal 0

Steps:

```
1. (xMn, xMd, yMn, yMd) = map_to_curve_elligator2_curve25519(u)
2. xn = xMn * yMd
3. xn = xn * c1
4. xd = xMd * yMn      # xn / xd = c1 * xM / yM
5. yn = xMn - xMd
6. yd = xMn + xMd      # (n / d - 1) / (n / d + 1) = (n - d) / (n + d)
7. tv1 = xd * yd
8. e = tv1 == 0
9. xn = CMOV(xn, 0, e)
10. xd = CMOV(xd, 1, e)
11. yn = CMOV(yn, 1, e)
12. yd = CMOV(yd, 1, e)
13. return (xn, xd, yn, yd)
```

G.2.3. curve448 ($q = 3 \pmod{4}$, $K = 1$)

The following is a straight-line implementation of Elligator 2 for curve448 [RFC7748] as specified in Section 8.6.

This implementation can also be used for any Montgomery curve with $K = 1$ over $GF(q)$ where $q = 3 \pmod{4}$.

`map_to_curve_elligator2_curve448(u)`

Input: `u`, an element of F .

Output: (x_n, x_d, y_n, y_d) such that $(x_n / x_d, y_n / y_d)$ is a point on curve448.

Constants:

1. `c1 = (q - 3) / 4` # Integer arithmetic

Steps:

```
1. tv1 = u^2
2. e1 = tv1 == 1
3. tv1 = CMOV(tv1, 0, e1) # If  $Z * u^2 == -1$ , set  $tv1 = 0$ 
4. xd = 1 - tv1
```

```

5.  x1n = -J
6.  tv2 = xd^2
7.  gxd = tv2 * xd      # gxd = xd^3
8.  gx1 = -J * tv1      # x1n + J * xd
9.  gx1 = gx1 * x1n     # x1n^2 + J * x1n * xd
10. gx1 = gx1 + tv2     # x1n^2 + J * x1n * xd + xd^2
11. gx1 = gx1 * x1n     # x1n^3 + J * x1n^2 * xd + x1n * xd^2
12. tv3 = gxd^2
13. tv2 = gx1 * gxd     # gx1 * gxd
14. tv3 = tv3 * tv2     # gx1 * gxd^3
15. y1 = tv3^c1        # (gx1 * gxd^3)^((p - 3) / 4)
16. y1 = y1 * tv2       # gx1 * gxd * (gx1 * gxd^3)^((p - 3) / 4)
17. x2n = -tv1 * x1n    # x2 = x2n / xd = -1 * u^2 * x1n / xd
18. y2 = y1 * u
19. y2 = CM0V(y2, 0, e1)
20. tv2 = y1^2
21. tv2 = tv2 * gxd
22. e2 = tv2 == gx1
23. xn = CM0V(x2n, x1n, e2) # If e2, x = x1, else x = x2
24. y = CM0V(y2, y1, e2)   # If e2, y = y1, else y = y2
25. e3 = sgn0(y) == 1      # Fix sign of y
26. y = CM0V(y, -y, e2 XOR e3)
27. return (xn, xd, y, 1)

```

G.2.4. edwards448

The following is a straight-line implementation of Elligator 2 for edwards448 [RFC7748] as specified in Section 8.6. The subroutine `map_to_curve_elligator2_curve448` is defined in Appendix G.2.3.

`map_to_curve_elligator2_edwards448(u)`

Input: u , an element of F .

Output: (x_n, x_d, y_n, y_d) such that $(x_n / x_d, y_n / y_d)$ is a point on edwards448.

Steps:

```

1. (xn, xd, yn, yd) = map_to_curve_elligator2_curve448(u)
2. xn2 = xn^2
3. xd2 = xd^2
4. xd4 = xd2^2
5. yn2 = yn^2
6. yd2 = yd^2
7. xEn = xn2 - xd2
8. tv2 = xEn - xd2
9. xEn = xEn * xd2
10. xEn = xEn * yd
11. xEn = xEn * yn
12. xEn = xEn * 4
13. tv2 = tv2 * xn2
14. tv2 = tv2 * yd2
15. tv3 = 4 * yn2
16. tv1 = tv3 + yd2
17. tv1 = tv1 * xd4
18. xEd = tv1 + tv2
19. tv2 = tv2 * xn

```

```

20. tv4 = xn * xd4
21. yEn = tv3 - yd2
22. yEn = yEn * tv4
23. yEn = yEn - tv2
24. tv1 = xn2 + xd2
25. tv1 = tv1 * xd2
26. tv1 = tv1 * xd
27. tv1 = tv1 * yn2
28. tv1 = -2 * tv1
29. yEd = tv2 + tv1
30. tv4 = tv4 * yd2
31. yEd = yEd + tv4
32. tv1 = xEd * yEd
33. e = tv1 == 0
34. xEn = CMOV(xEn, 0, e)
35. xEd = CMOV(xEd, 1, e)
36. yEn = CMOV(yEn, 1, e)
37. yEd = CMOV(yEd, 1, e)
38. return (xEn, xEd, yEn, yEd)

```

G.2.5. Montgomery Curves with $q = 3 \pmod{4}$

The following is a straight-line implementation of Elligator 2 that applies to any Montgomery curve defined over $\text{GF}(q)$ where $q = 3 \pmod{4}$.

For curves where $K = 1$, the implementation given in Appendix G.2.3 gives identical results with slightly reduced cost.

map_to_curve_elligator2_3mod4(u)

Input: u , an element of F .

Output: (x_n, x_d, y_n, y_d) such that $(x_n / x_d, y_n / y_d)$ is a point on the target curve.

Constants:

```

1. c1 = (q - 3) / 4          # Integer arithmetic
2. c2 = K^2

```

Steps:

```

1. tv1 = u^2
2. e1 = tv1 == 1
3. tv1 = CMOV(tv1, 0, e1) # If  $Z * u^2 == -1$ , set  $tv1 = 0$ 
4. xd = 1 - tv1
5. xd = xd * K
6. x1n = -J                  #  $x1 = x1n / xd = -J / (K * (1 + 2 * u^2))$ 
7. tv2 = xd^2
8. gxd = tv2 * xd
9. gxd = gxd * c2           #  $gxd = xd^3 * K^2$ 
10. gx1 = x1n * K
11. tv3 = xd * J
12. tv3 = gx1 + tv3         #  $x1n * K + xd * J$ 
13. gx1 = gx1 * tv3         #  $K^2 * x1n^2 + J * K * x1n * xd$ 
14. gx1 = gx1 + tv2         #  $K^2 * x1n^2 + J * K * x1n * xd + xd^2$ 
15. gx1 = gx1 * x1n        #  $K^2 * x1n^3 + J * K * x1n^2 * xd + x1n * xd^2$ 
16. tv3 = gxd^2

```

```

17. tv2 = gx1 * gxd      # gx1 * gxd
18. tv3 = tv3 * tv2      # gx1 * gxd^3
19. y1 = tv3^c1          # (gx1 * gxd^3)^((q - 3) / 4)
20. y1 = y1 * tv2        # gx1 * gxd * (gx1 * gxd^3)^((q - 3) / 4)
21. x2n = -tv1 * x1n     # x2 = x2n / xd = -1 * u^2 * x1n / xd
22. y2 = y1 * u
23. y2 = CM0V(y2, 0, e1)
24. tv2 = y1^2
25. tv2 = tv2 * gxd
26. e2 = tv2 == gx1
27. xn = CM0V(x2n, x1n, e2) # If e2, x = x1, else x = x2
28. xn = xn * K
29. y = CM0V(y2, y1, e2)   # If e2, y = y1, else y = y2
30. e3 = sgn0(y) == 1      # Fix sign of y
31. y = CM0V(y, -y, e2 XOR e3)
32. y = y * K
33. return (xn, xd, y, 1)

```

G.2.6. Montgomery Curves with $q = 5 \pmod{8}$

The following is a straight-line implementation of Elligator 2 that applies to any Montgomery curve defined over $\text{GF}(q)$ where $q = 5 \pmod{8}$.

For curves where $K = 1$, the implementation given in Appendix G.2.1 gives identical results with slightly reduced cost.

map_to_curve_elligator2_5mod8(u)

Input: u , an element of F .

Output: (x_n, x_d, y_n, y_d) such that $(x_n / x_d, y_n / y_d)$ is a point on the target curve.

Constants:

```

1. c1 = (q + 3) / 8      # Integer arithmetic
2. c2 = 2^c1
3. c3 = sqrt(-1)
4. c4 = (q - 5) / 8      # Integer arithmetic
5. c5 = K^2

```

Steps:

```

1. tv1 = u^2
2. tv1 = 2 * tv1
3. xd = tv1 + 1          # Nonzero: -1 is square (mod p), tv1 is not
4. xd = xd * K
5. x1n = -J              # x1 = x1n / xd = -J / (K * (1 + 2 * u^2))
6. tv2 = xd^2
7. gxd = tv2 * xd
8. gxd = gxd * c5        # gxd = xd^3 * K^2
9. gx1 = x1n * K
10. tv3 = xd * J
11. tv3 = gx1 + tv3      # x1n * K + xd * J
12. gx1 = gx1 * tv3      # K^2 * x1n^2 + J * K * x1n * xd
13. gx1 = gx1 + tv2      # K^2 * x1n^2 + J * K * x1n * xd + xd^2
14. gx1 = gx1 * x1n      # K^2 * x1n^3 + J * K * x1n^2 * xd + x1n * xd^2
15. tv3 = gxd^2

```

```

16. tv2 = tv3^2          # gxd^4
17. tv3 = tv3 * gxd      # gxd^3
18. tv3 = tv3 * gx1      # gx1 * gxd^3
19. tv2 = tv2 * tv3      # gx1 * gxd^7
20. y11 = tv2^c4         # (gx1 * gxd^7)^((q - 5) / 8)
21. y11 = y11 * tv3      # gx1 * gxd^3 * (gx1 * gxd^7)^((q - 5) / 8)
22. y12 = y11 * c3
23. tv2 = y11^2
24. tv2 = tv2 * gxd
25. e1 = tv2 == gx1
26. y1 = CMOV(y12, y11, e1) # If g(x1) is square, this is its sqrt
27. x2n = x1n * tv1      # x2 = x2n / xd = 2 * u^2 * x1n / xd
28. y21 = y11 * u
29. y21 = y21 * c2
30. y22 = y21 * c3
31. gx2 = gx1 * tv1      # g(x2) = gx2 / gxd = 2 * u^2 * g(x1)
32. tv2 = y21^2
33. tv2 = tv2 * gxd
34. e2 = tv2 == gx2
35. y2 = CMOV(y22, y21, e2) # If g(x2) is square, this is its sqrt
36. tv2 = y1^2
37. tv2 = tv2 * gxd
38. e3 = tv2 == gx1
39. xn = CMOV(x2n, x1n, e3) # If e3, x = x1, else x = x2
40. xn = xn * K
41. y = CMOV(y2, y1, e3)   # If e3, y = y1, else y = y2
42. e4 = sgn0(y) == 1      # Fix sign of y
43. y = CMOV(y, -y, e3 XOR e4)
44. y = y * K
45. return (xn, xd, y, 1)

```

G.3. Cofactor Clearing for BLS12-381 G2

The curve BLS12-381, whose parameters are defined in Section 8.8.2, admits an efficiently computable endomorphism, ψ , that can be used to speed up cofactor clearing for G2 [SBCDK09] [FKR11] [BP17] (see also Section 7). This section implements the endomorphism ψ and a fast cofactor clearing method described by Budroni and Pintore [BP17].

The functions in this section operate on points whose coordinates are represented as ratios, i.e., (x_n, x_d, y_n, y_d) corresponds to the point $(x_n / x_d, y_n / y_d)$; see Appendix G.1 for further discussion of projective coordinates. When points are represented in affine coordinates, one can simply ignore the denominators ($x_d == 1$ and $y_d == 1$).

The following function computes the Frobenius endomorphism for an element of $F = GF(p^2)$ with basis $(1, I)$, where $I^2 + 1 == 0$ in F . (This is the base field of the elliptic curve E defined in Section 8.8.2.)

frobenius(x)

Input: x , an element of $GF(p^2)$.

Output: a , an element of $GF(p^2)$.

Notation: $x = x_0 + I * x_1$, where x_0 and x_1 are elements of $GF(p)$.

Steps:

1. $a = x_0 - I * x_1$
2. return a

The following function computes the endomorphism ψ for points on the elliptic curve E defined in Section 8.8.2.

$\psi(x_n, x_d, y_n, y_d)$

Input: P , a point $(x_n / x_d, y_n / y_d)$ on the curve E (see above).

Output: Q , a point on the same curve.

Constants:

1. $c_1 = 1 / (1 + I)^{((p - 1) / 3)}$ # in $GF(p^2)$
2. $c_2 = 1 / (1 + I)^{((p - 1) / 2)}$ # in $GF(p^2)$

Steps:

1. $q_{xn} = c_1 * \text{frobenius}(x_n)$
2. $q_{xd} = \text{frobenius}(x_d)$
3. $q_{yn} = c_2 * \text{frobenius}(y_n)$
4. $q_{yd} = \text{frobenius}(y_d)$
5. return $(q_{xn}, q_{xd}, q_{yn}, q_{yd})$

The following function efficiently computes $\psi(\psi(P))$.

$\psi_2(x_n, x_d, y_n, y_d)$

Input: P , a point $(x_n / x_d, y_n / y_d)$ on the curve E (see above).

Output: Q , a point on the same curve.

Constants:

1. $c_1 = 1 / 2^{((p - 1) / 3)}$ # in $GF(p^2)$

Steps:

1. $q_{xn} = c_1 * x_n$
2. $q_{yn} = -y_n$
3. return $(q_{xn}, x_d, q_{yn}, y_d)$

The following function maps any point on the elliptic curve E (Section 8.8.2) into the prime-order subgroup G_2 . This function returns a point equal to $h_{\text{eff}} * P$, where h_{eff} is the parameter given in Section 8.8.2.

$\text{clear_cofactor_bls12381_g2}(P)$

Input: P , a point $(x_n / x_d, y_n / y_d)$ on the curve E (see above).

Output: Q , a point in the subgroup G_2 of BLS12-381.

Constants:

1. $c_1 = -15132376222941642752$ # the BLS parameter for BLS12-381
i.e., $-0xd201000000010000$

Notation: in this procedure, $+$ and $-$ represent elliptic curve point

addition and subtraction, respectively, and $*$ represents scalar multiplication.

Steps:

1. $t_1 = c_1 * P$
2. $t_2 = \text{psi}(P)$
3. $t_3 = 2 * P$
4. $t_3 = \text{psi}_2(t_3)$
5. $t_3 = t_3 - t_2$
6. $t_2 = t_1 + t_2$
7. $t_2 = c_1 * t_2$
8. $t_3 = t_3 + t_2$
9. $t_3 = t_3 - t_1$
10. $Q = t_3 - P$
11. return Q

Appendix H. Scripts for Parameter Generation

This section gives Sage scripts [SAGE] used to generate parameters for the mappings of Section 6.

H.1. Finding Z for the Shallue-van de Woestijne Map

The below function outputs an appropriate Z for the Shallue-van de Woestijne map (Section 6.6.1).

```
# Arguments:
# - F, a field object, e.g., F = GF(2^521 - 1)
# - A and B, the coefficients of the curve  $y^2 = x^3 + A * x + B$ 
def find_z_svdw(F, A, B, init_ctr=1):
    g = lambda x: F(x)^3 + F(A) * F(x) + F(B)
    h = lambda Z: -(F(3) * Z^2 + F(4) * A) / (F(4) * g(Z))
    # NOTE: if init_ctr=1 fails to find Z, try setting it to F.gen()
    ctr = init_ctr
    while True:
        for Z_cand in (F(ctr), F(-ctr)):
            # Criterion 1:
            #    $g(Z) \neq 0$  in F.
            if g(Z_cand) == F(0):
                continue
            # Criterion 2:
            #    $-(3 * Z^2 + 4 * A) / (4 * g(Z)) \neq 0$  in F.
            if h(Z_cand) == F(0):
                continue
            # Criterion 3:
            #    $-(3 * Z^2 + 4 * A) / (4 * g(Z))$  is square in F.
            if not is_square(h(Z_cand)):
                continue
            # Criterion 4:
            #   At least one of  $g(Z)$  and  $g(-Z / 2)$  is square in F.
            if is_square(g(Z_cand)) or is_square(g(-Z_cand / F(2))):
                return Z_cand
        ctr += 1
```

H.2. Finding Z for Simplified SWU

The below function outputs an appropriate Z for the Simplified SWU map (Section 6.6.2).

```
# Arguments:
# - F, a field object, e.g., F = GF(2^521 - 1)
# - A and B, the coefficients of the curve  $y^2 = x^3 + A * x + B$ 
def find_z_sswu(F, A, B):
    R.<xx> = F[]                                # Polynomial ring over F
    g = xx^3 + F(A) * xx + F(B)                #  $y^2 = g(x) = x^3 + A * x + B$ 
    ctr = F.gen()
    while True:
        for Z_cand in (F(ctr), F(-ctr)):
            # Criterion 1: Z is non-square in F.
            if is_square(Z_cand):
                continue
            # Criterion 2:  $Z \neq -1$  in F.
            if Z_cand == F(-1):
                continue
            # Criterion 3:  $g(x) - Z$  is irreducible over F.
            if not (g - Z_cand).is_irreducible():
                continue
            # Criterion 4:  $g(B / (Z * A))$  is square in F.
            if is_square(g(B / (Z_cand * A))):
                return Z_cand
        ctr += 1
```

H.3. Finding Z for Elligator 2

The below function outputs an appropriate Z for the Elligator 2 map (Section 6.7.1).

```
# Argument:
# - F, a field object, e.g., F = GF(2^255 - 19)
def find_z_ell2(F):
    ctr = F.gen()
    while True:
        for Z_cand in (F(ctr), F(-ctr)):
            # Z must be a non-square in F.
            if is_square(Z_cand):
                continue
            return Z_cand
        ctr += 1
```

Appendix I. sqrt and is_square Functions

This section defines special-purpose sqrt functions for the three most common cases, $q = 3 \pmod{4}$, $q = 5 \pmod{8}$, and $q = 9 \pmod{16}$, plus a generic constant-time algorithm that works for any prime modulus.

In addition, it gives an optimized is_square method for $GF(p^2)$.

I.1. sqrt for $q = 3 \pmod{4}$

```
sqrt_3mod4(x)
```


Parameters:

- F , a finite field of characteristic p and order $q = p^m$.

Input: x , an element of F .

Output: z , an element of F such that $(z^2) == x$, if x is square in F .

Constants:

1. $c1 = (q + 1) / 4$ # Integer arithmetic

Procedure:

1. return x^{c1}

I.2. sqrt for $q = 5 \pmod{8}$

sqrt_5mod8(x)

Parameters:

- F , a finite field of characteristic p and order $q = p^m$.

Input: x , an element of F .

Output: z , an element of F such that $(z^2) == x$, if x is square in F .

Constants:

1. $c1 = \text{sqrt}(-1)$ in F , i.e., $(c1^2) == -1$ in F

2. $c2 = (q + 3) / 8$ # Integer arithmetic

Procedure:

1. $tv1 = x^{c2}$

2. $tv2 = tv1 * c1$

3. $e = (tv1^2) == x$

4. $z = \text{CM0V}(tv2, tv1, e)$

5. return z

I.3. sqrt for $q = 9 \pmod{16}$

sqrt_9mod16(x)

Parameters:

- F , a finite field of characteristic p and order $q = p^m$.

Input: x , an element of F .

Output: z , an element of F such that $(z^2) == x$, if x is square in F .

Constants:

1. $c1 = \text{sqrt}(-1)$ in F , i.e., $(c1^2) == -1$ in F

2. $c2 = \text{sqrt}(c1)$ in F , i.e., $(c2^2) == c1$ in F

3. $c3 = \text{sqrt}(-c1)$ in F , i.e., $(c3^2) == -c1$ in F

4. $c4 = (q + 7) / 16$ # Integer arithmetic

Procedure:

1. $tv1 = x^{c4}$

2. $tv2 = c1 * tv1$

3. $tv3 = c2 * tv1$

4. $tv4 = c3 * tv1$

5. $e1 = (tv2^2) == x$

6. $e2 = (tv3^2) == x$

```

7. tv1 = CM0V(tv1, tv2, e1) # Select tv2 if (tv2^2) == x
8. tv2 = CM0V(tv4, tv3, e2) # Select tv3 if (tv3^2) == x
9. e3 = (tv2^2) == x
10. z = CM0V(tv1, tv2, e3) # Select the sqrt from tv1 and tv2
11. return z

```

I.4. Constant-Time Tonelli-Shanks Algorithm

This algorithm is a constant-time version of the classic Tonelli-Shanks algorithm ([C93], Algorithm 1.5.1) due to Sean Bowe, Jack Grigg, and Eirik Ogilvie-Wigley [jubjub-fq], adapted and optimized by Michael Scott.

This algorithm applies to $GF(p)$ for any p . Note, however, that the special-purpose algorithms given in the prior sections are faster, when they apply.

`sqrt_ts_ct(x)`

Parameters:

- F , a finite field of characteristic p and order $q = p^m$.

Input x , an element of F .

Output: z , an element of F such that $z^2 == x$, if x is square in F .

Constants:

1. c_1 , the largest integer such that 2^{c_1} divides $q - 1$.
2. $c_2 = (q - 1) / (2^{c_1})$ # Integer arithmetic
3. $c_3 = (c_2 - 1) / 2$ # Integer arithmetic
4. c_4 , a non-square value in F
5. $c_5 = c_4^{c_2}$ in F

Procedure:

1. $z = x^{c_3}$
2. $t = z * z$
3. $t = t * x$
4. $z = z * x$
5. $b = t$
6. $c = c_5$
7. for i in $(c_1, c_1 - 1, \dots, 2)$:
8. for j in $(1, 2, \dots, i - 2)$:
9. $b = b * b$
10. $e = b == 1$
11. $zt = z * c$
12. $z = \text{CM0V}(zt, z, e)$
13. $c = c * c$
14. $tt = t * c$
15. $t = \text{CM0V}(tt, t, e)$
16. $b = t$
17. return z

I.5. `is_square` for $F = GF(p^2)$

The following `is_square` method applies to any field $F = GF(p^2)$ with basis $(1, I)$ represented as described in Section 2.1, i.e., an element $x = (x_1, x_2) = x_1 + x_2 * I$.

Other optimizations of this type are possible in other extension fields; see, for example, [AR13] for more information.

`is_square(x)`

Parameters:

- F , an extension field of characteristic p and order $q = p^2$ with basis $(1, I)$.

Input: x , an element of F .

Output: True if x is square in F , and False otherwise.

Constants:

1. $c1 = (p - 1) / 2$ # Integer arithmetic

Procedure:

```
1.  $tv1 = x_1^2$ 
2.  $tv2 = I * x_2$ 
3.  $tv2 = tv2^2$ 
4.  $tv1 = tv1 - tv2$ 
5.  $tv1 = tv1^{c1}$ 
6.  $e1 = tv1 \neq -1$  # Note:  $-1$  in  $F$ 
7. return  $e1$ 
```

Appendix J. Suite Test Vectors

This section gives test vectors for each suite defined in Section 8. The test vectors in this section were generated using code that is available from [hash2curve-repo].

Each test vector in this section lists values computed by the appropriate encoding function, with variable names defined as in Section 3. For example, for a suite whose encoding type is random oracle, the test vector gives the value for msg , u , $Q0$, $Q1$, and the output point P .

J.1. NIST P-256

J.1.1. P256_XMD:SHA-256_SSWU_RO_

```
suite    = P256_XMD:SHA-256_SSWU_RO_
dst      = QUUX-V01-CS02-with-P256_XMD:SHA-256_SSWU_RO_

msg      =
P.x      = 2c15230b26dbc6fc9a37051158c95b79656e17a1a920b11394ca91
          c44247d3e4
P.y      = 8a7a74985cc5c776cdfe4b1f19884970453912e9d31528c060be9a
          b5c43e8415
u[0]     = ad5342c66a6dd0ff080df1da0ea1c04b96e0330dd89406465eeba1
          1582515009
u[1]     = 8c0f1d43204bd6f6ea70ae8013070a1518b43873bcd850aafa0a9e
          220e2eea5a
Q0.x     = ab640a12220d3ff283510ff3f4b1953d09fad35795140b1c5d64f3
          13967934d5
Q0.y     = dccb558863804a881d4fff3455716c836cef230e5209594ddd33d8
```


627e4352b1
 Q1.x = 418ac3d85a5ccc4ea8dec14f750a3a9ec8b85176c95a7022f39182
 6794eb5a75
 Q1.y = fd6604f69e9d9d2b74b072d14ea13050db72c932815523305cb9e8
 07cc900aff

 msg = a512_aa
 aa
 aa
 aa
 aa
 aa
 aa
 aa
 aa
 aa
 aa
 P.x = 457ae2981f70ca85d8e24c308b14db22f3e3862c5ea0f652ca38b5
 e49cd64bc5
 P.y = ecb9f0eadc9aead232dabc53235368c1394c78de05dd96893eefa6
 2b0f4757dc
 u[0] = 4ebc95a6e839b1ae3c63b847798e85cb3c12d3817ec6ebc10af6ee
 51adb29fec
 u[1] = 4e21af88e22ea80156aff790750121035b3eefaa96b425a8716e0d
 20b4e269ee
 Q0.x = d88b989ee9d1295df413d4456c5c850b8b2fb0f5402cc5c4c7e815
 412e926db8
 Q0.y = bb4a1edef506cf16def96afff41b16fc74f6dbd55c2210e5b8f01
 1ba32f4f40
 Q1.x = a281e34e628f3a4d2a53fa87ff973537d68ad4fbc28d3be5e8d9f6
 a2571c5a4b
 Q1.y = f6ed88a7aab56a488100e6f1174fa9810b47db13e86be999644922
 961206e184

J.1.1.2. P256_XMD:SHA-256_SSWU_NU_

suite = P256_XMD:SHA-256_SSWU_NU_
 dst = QUUX-V01-CS02-with-P256_XMD:SHA-256_SSWU_NU_

 msg =
 P.x = f871caad25ea3b59c16cf87c1894902f7e7b2c822c3d3f73596c5a
 ce8ddd14d1
 P.y = 87b9ae23335bee057b99bac1e68588b18b5691af476234b8971bc4
 f011ddc99b
 u[0] = b22d487045f80e9edcb0ecc8d4bf77833e2bf1f3a54004d7df1d57
 f4802d311f
 Q.x = f871caad25ea3b59c16cf87c1894902f7e7b2c822c3d3f73596c5a
 ce8ddd14d1
 Q.y = 87b9ae23335bee057b99bac1e68588b18b5691af476234b8971bc4
 f011ddc99b

 msg = abc
 P.x = fc3f5d734e8dce41ddac49f47dd2b8a57257522a865c124ed02b92
 b5237befa4
 P.y = fe4d197ecf5a62645b9690599e1d80e82c500b22ac705a0b421fac
 7b47157866
 u[0] = c7f96eadac763e176629b09ed0c11992225b3a5ae99479760601cb

[illegible]

J.2. NIST P-384

J.2.1. P384_XMD:SHA-384_SSWU_RO_

```

suite = P384_XMD:SHA-384_SSWU_RO
dst   = QUUX-V01-CS02-with-P384_XMD:SHA-384_SSWU_RO_

msg =
P.x = eb9fe1b4f4e14e71140803c1d99d0a93cd823d2b024040f9c067a8e
    ca1f5a2eeac9ad604973527a356f3fa3aeff0e4d83
P.y = 0c21708cff382b7f4643c07b105c2eaec2cead93a917d825601e63
    c8f21f6abd9abc22c93c2bed6f235954b25048bb1a
u[0] = 25c8d7dc1acd4ee617766693f7f8829396065d1b447eedb155871f
    effd9c6653279ac7e5c46edb7010a0e4ff64c9f3b4
u[1] = 59428be4ed69131df59a0c6a8e188d2d4ece3f1b2a3a02602962b4
    7efa4d7905945b1e2cc80b36aa35c99451073521ac
Q0.x = e4717e29eef38d862bee4902a7d21b44efb58c464e3e1f0d03894d
    94de310f8fffc6de86786dd3e15a1541b18d4eb2846
Q0.y = 6b95a6e639822312298a47526bb77d9cd7bcf76244c991c8cd7007
    5e2ee6e8b9a135c4a37e3c0768c7ca871c0ceb53d4
Q1.x = 509527cfc0750eedc53147e6d5f78596c8a3b7360e0608e2fab056
    3a1670d58d8ae107c9f04bcf90e89489ace5650efd
Q1.y = 33337b13cb35e173fdea4cb9e8cce915d836ff57803dbbeb7998aa
    49d17df2ff09b67031773039d09fbd9305a1566bc4

msg = abc
P.x = e02fc1a5f44a7519419dd314e29863f30df55a514da2d655775a81
    d413003c4d4e7fd59af0826dfaad4200ac6f60abe1
P.y = 01f638d04d98677d65bef99aef1a12a70a4cbb9270ec55248c0453
    0d8bc1f8f90f8a6a859a7c1f1ddccedf8f96d675f6
u[0] = 53350214cb6bef0b51abb791b1c4209a2b4c16a0c67e1ab1401017
    fad774cd3b3f9a8bcdff7f6229dd8dd5a075cb149a0
u[1] = c0473083898f63e03f26f14877a2407bd60c75ad491e7d26cbc6cc
    5ce815654075ec6b6898c7a41d74ceaf720a10c02e
Q0.x = fc853b69437aee9a19d5acf96a4ee4c5e04cf7b53406dfaa2afbdd
    7ad2351b7f554e4bbc6f5db4177d4d44f933a8f6ee
Q0.y = 7e042547e01834c9043b10f3a8221c4a879cb156f04f72bfccab0c
    047a304e30f2aa8b2e260d34c4592c0c33dd0c6482
Q1.x = 57912293709b3556b43a2dfb137a315d256d573b82ded120ef8c78
    2d607c05d930d958e50cb6dc1cc480b9afc38c45f1
Q1.y = de9387dab0eef0bda219c6f168a92645a84665c4f2137c14270fb4
    24b7532ff84843c3da383ceea24c47fa343c227bb8

msg = abcdef0123456789
P.x = bdecc1c1d870624965f19505be50459d363c71a699a496ab672f9a
    5d6b78676400926fbceee6fcd1780fe86e62b2aa89
P.y = 57cf1f99b5ee00f3c201139b3bfe4dd30a653193778d89a0accc5e
    0f47e46e4e4b85a0595da29c9494c1814acafe183c
u[0] = aab7fb87238cf6b2ab56cdcca7e028959bb2ea599d34f68484139d
    de85ec6548a6e48771d17956421bdb7790598ea52e
u[1] = 26e8d833552d7844d167833ca5a87c35bcfaa5a0d86023479fb28e
    5cd6075c18b168bf1f5d2a0ea146d057971336d8d1
Q0.x = 0ceece45b73f89844671df962ad2932122e878ad2259e650626924
    e4e7f132589341dec1480ebcbbbe3509d11fb570b7
Q0.y = fafd71a3115298f6be4ae5c6dfc96c400cfb55760f185b7b03f3fa
    45f3f91eb65d27628b3c705cafd0466fafa54883ce
Q1.x = dea1be8d3f9be4cbf4fab9d71d549dde76875b5d9b876832313a08
    3ec81e528cbc2a0a1d0596b3bcb0ba77866b129776
Q1.y = eb15fe71662214fb03b65541f40d3eb0f4cf5c3b559f647da138c9

```



```

aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x    = af129727a4207a8cb9e9dce656d88f79fce25edbcea350499d65e9
        bf1204537bdde73c7cefb752a6ed5ebcd44e183302
P.y    = ce68a3d5e161b2e6a968e4ddaa9e51504ad1516ec170c7eef3ca6b
        5327943eca95d90b23b009ba45f58b72906f2a99e2
u[0]   = 7b01ce9b8c5a60d9fbc202d6dde92822e46915d8c17e03fcb92ece
        1ed6074d01e149fc9236def40d673de903c1d4c166
Q.x    = af129727a4207a8cb9e9dce656d88f79fce25edbcea350499d65e9
        bf1204537bdde73c7cefb752a6ed5ebcd44e183302
Q.y    = ce68a3d5e161b2e6a968e4ddaa9e51504ad1516ec170c7eef3ca6b
        5327943eca95d90b23b009ba45f58b72906f2a99e2

```

J.3. NIST P-521

J.3.1. P521_XMD:SHA-512_SSWU_RO_

```

suite  = P521_XMD:SHA-512_SSWU_RO_
dst    = QUUX-V01-CS02-with-P521_XMD:SHA-512_SSWU_RO_

msg    =
P.x    = 00fd767cebb2452030358d0e9cf907f525f50920c8f607889a6a35
        680727f64f4d66b161fafeb2654bea0d35086bec0a10b30b14adef
        3556ed9f7f1bc23cecc9c088
P.y    = 0169ba78d8d851e930680322596e39c78f4fe31b97e57629ef6460
        ddd68f8763fd7bd767a4e94a80d3d21a3c2ee98347e024fc73ee1c
        27166dc3fe5eeef782be411d
u[0]   = 01e5f09974e5724f25286763f00ce76238c7a6e03dc396600350ee
        2c4135fb17dc555be99a4a4bae0fd303d4f66d984ed7b6a3ba3860
        93752a855d26d559d69e7e9e
u[1]   = 00ae593b42ca2ef93ac488e9e09a5fe5a2f6fb330d18913734ff60
        2f2a761fcaaf5f596e790bcc572c9140ec03f6cccc38f767f1c197
        5a0b4d70b392d95a0c7278aa
Q0.x   = 00b70ae99b6339fffac19cb9bfde2098b84f75e50ac1e80d6acb95
        4e4534af5f0e9c4a5b8a9c10317b8e6421574bae2b133b4f2b8c6c
        e4b3063da1d91d34fa2b3a3c
Q0.y   = 007f368d98a4ddbfb381fb354de40e44b19e43bb11a1278759f4ea7
        b485e1b6db33e750507c071250e3e443c1aaed61f2c28541bb54b1
        b456843eda1eb15ec2a9b36e
Q1.x   = 01143d0e9cddcdacd6a9aaf1bcf8d218c0afc45d4451239e821f5
        d2a56df92be942660b532b2aa59a9c635ae6b30e803c45a6ac8714
        32452e685d661cd41cf67214
Q1.y   = 00ff75515df265e996d702a5380defffab1a6d2bc232234c7bcffa
        433cd8aa791fbc8dcf667f08818bffa739ae25773b32073213cae9
        a0f2a917a0b1301a242dda0c

msg    = abc
P.x    = 002f89a1677b28054b50d15e1f81ed6669b5a2158211118ebdef8a
        6efc77f8ccaa528f698214e4340155abc1fa08f8f613ef14a04371
        7503d57e267d57155cf784a4
P.y    = 010e0be5dc8e753da8ce51091908b72396d3deed14ae166f66d8eb
        f0a4e7059ead169ea4bead0232e9b700dd380b316e9361cfdba55a
        08c73545563a80966ecbb86d
u[0]   = 003d00c37e95f19f358adeeaa47288ec39998039c3256e13c2a4c0
        0a7cb61a34c8969472960150a27276f2390eb5e53e47ab193351c2
        d2d9f164a85c6a5696d94fe8

```



```

4ad7356a04e08729ad9411f5
Q0.x = 0005eac7b0b81e38727efcab1e375f6779aea949c3e409b53a1d37
aa2acbac87a7e6ad24aafb3c52f82f7f0e21b872e88c55e17b7fa
21ce08a94ea2121c42c2eb73
Q0.y = 00a173b6a53a7420dbd61d4a21a7c0a52de7a5c6ce05f31403bef7
47d16cc8604a039a73bdd6e114340e55dacd6bea8e217ffbadfb8c
292afa3e1b2afc839a6ce7bb
Q1.x = 01881e3c193a69e4d88d8180a6879b74782a0bc7e529233e9f84bf
7f17d2f319c36920ffba26f9e57a1e045cc7822c834c239593b6e1
42a694aa00c757b0db79e5e8
Q1.y = 01558b16d396d866e476e001f2dd0758927655450b84e12f154032
c7c2a6db837942cd9f44b814f79b4d729996ced61eec61d85c6751
39cbffe3fbf071d2c21cfecb

msg = a512_aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = 00c12bc3e28db07b6b4d2a2b1167ab9e26fc2fa85c7b0498a17b03
47edf52392856d7e28b8fa7a2dd004611159505835b687ecf1a764
857e27e9745848c436ef3925
P.y = 01cd287df9a50c22a9231beb452346720bb163344a41c5f5a24e83
35b6ccc595fd436aea89737b1281aecb411eb835f0b939073fdd1d
d4d5a2492e91ef4a3c55bcdb
u[0] = 0033d06d17bc3b9a3efc081a05d65805a14a3050a0dd4dfb488461
8eb5c73980a59c5a246b18f58ad022dd3630faa22889fbb8ba1593
466515e6ab4aeb7381c26334
u[1] = 0092290ab99c3fea1a5b8fb2ca49f859994a04faee3301cefab312
d34227f6a2d0c3322cf76861c6a3683bdaa2dd2a6daa5d6906c663
e065338b2344d20e313f1114
Q0.x = 00041f6eb92af8777260718e4c22328a7d74203350c6c8f5794d99
d5789766698f459b83d5068276716f01429934e40af3d1111a2278
0b1e07e72238d2207e5386be
Q0.y = 001c712f0182813942b87cab8e72337db017126f52ed797dd23458
4ac9ae7e80dfe7abea11db02cf1855312eae1447dbaecc9d7e8c88
0a5e76a39f6258074e1bc2e0
Q1.x = 0125c0b69bcf55eab49280b14f707883405028e05c927cd7625d4e
04115bd0e0e6323b12f5d43d0d6d2eff16dbc244542f84ec05891
1260dc3bb6512ab5db285fbd
Q1.y = 008bddfb803b3f4c761458eb5f8a0aee3e1f7f68e9d7424405fa69
172919899317fb6ac1d6903a432d967d14e0f80af63e7035aaae0c
123e56862ce969456f99f102

```

J.3.2. P521_XMD:SHA-512_SSWU_NU_

```

suite = P521_XMD:SHA-512_SSWU_NU_
dst   = QUUX-V01-CS02-with-P521_XMD:SHA-512_SSWU_NU_

msg   =
P.x   = 01ec604b4e1e3e4c7449b7a41e366e876655538acf51fd40d08b97

```

	be066f7d020634e906b1b6942f9174b417027c953d75fb6ec64b8c ee2a3672d4f1987d13974705
P.y	= 00944fc439b4aad2463e5c9cfa0b0707af3c9a42e37c5a57bb4ecd 12fef9fb21508568aedcdd8d2490472df4bbafd79081c81e99f4da 3286eddf19be47e9c4cf0e91
u[0]	= 01e4947fe62a4e47792cee2798912f672fff820b2556282d9843b4 b465940d7683a986f93ccb0e9a191fbc09a6e770a564490d2a4ae5 1b287ca39f69c3d910ba6a4f
Q.x	= 01ec604b4e1e3e4c7449b7a41e366e876655538acf51fd40d08b97 be066f7d020634e906b1b6942f9174b417027c953d75fb6ec64b8c ee2a3672d4f1987d13974705
Q.y	= 00944fc439b4aad2463e5c9cfa0b0707af3c9a42e37c5a57bb4ecd 12fef9fb21508568aedcdd8d2490472df4bbafd79081c81e99f4da 3286eddf19be47e9c4cf0e91
msg	= abc
P.x	= 00c720ab56aa5a7a4c07a7732a0a4e1b909e32d063ae1b58db5f0e b5e09f08a9884bff55a2bef4668f715788e692c18c1915cd034a6b 998311fcf46924ce66a2be9a
P.y	= 003570e87f91a4f3c7a56be2cb2a078ffc153862a53d5e03e5dad5 bccc6c529b8bab0b7dbb157499e1949e4edab21cf5d10b782bc1e9 45e13d7421ad8121dbc72b1d
u[0]	= 0019b85ef78596efc84783d42799e80d787591fe7432dee1d9fa2b 7651891321be732ddf653fa8fefa34d86fb728db569d36b5b6ed39 83945854b2fc2dc6a75aa25b
Q.x	= 00c720ab56aa5a7a4c07a7732a0a4e1b909e32d063ae1b58db5f0e b5e09f08a9884bff55a2bef4668f715788e692c18c1915cd034a6b 998311fcf46924ce66a2be9a
Q.y	= 003570e87f91a4f3c7a56be2cb2a078ffc153862a53d5e03e5dad5 bccc6c529b8bab0b7dbb157499e1949e4edab21cf5d10b782bc1e9 45e13d7421ad8121dbc72b1d
msg	= abcdef0123456789
P.x	= 00bcacf32a968ff7971b3bbd9ce8edfbee1309e2019d7ff373c3838 7a782b005dce6ceffccfed5c6511c8f7f312f343f3a891029c585 8f45ee0bf370aba25fc990cc
P.y	= 00923517e767532d82cb8a0b59705eec2b7779ce05f9181c7d5d5e 25694ef8ebd4696343f0bc27006834d2517215ecf79482a84111f5 0c1bae25044fe1dd77744bbd
u[0]	= 01dba0d7fa26a562ee8a9014ebc2cca4d66fd9de036176aca8fc11 ef254cd1bc208847ab7701dbca7af328b3f601b11a1737a899575a 5c14f4dca5aaca45e9935e07
Q.x	= 00bcacf32a968ff7971b3bbd9ce8edfbee1309e2019d7ff373c3838 7a782b005dce6ceffccfed5c6511c8f7f312f343f3a891029c585 8f45ee0bf370aba25fc990cc
Q.y	= 00923517e767532d82cb8a0b59705eec2b7779ce05f9181c7d5d5e 25694ef8ebd4696343f0bc27006834d2517215ecf79482a84111f5 0c1bae25044fe1dd77744bbd
msg	= q128_qqq qq qq
P.x	= 001ac69014869b6c4ad7aa8c443c255439d36b0e48a0f57b03d6fe 9c40a66b4e2eaed2a93390679a5cc44b3a91862b34b673f0e92c83 187da02bf3db967d867ce748
P.y	= 00d5603d530e4d62b30fccfa1d90c2206654d74291c1db1c25b86a

```

051ee3fffc294e5d56f2e776853406bd09206c63d40f37ad882952
4cf89ad70b5d6e0b4a3b7341
u[0] = 00844da980675e1244cb209dcf3ea0aabec23bd54b2cda69fff86e
b3acc318bf3d01bae96e9cd6f4c5ceb5539df9a7ad7fcc5e9d5469
6081ba9782f3a0f6d14987e3
Q.x = 001ac69014869b6c4ad7aa8c443c255439d36b0e48a0f57b03d6fe
9c40a66b4e2eae2a93390679a5cc44b3a91862b34b673f0e92c83
187da02bf3db967d867ce748
Q.y = 00d5603d530e4d62b30fccfa1d90c2206654d74291c1db1c25b86a
051ee3fffc294e5d56f2e776853406bd09206c63d40f37ad882952
4cf89ad70b5d6e0b4a3b7341

msg = a512_aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = 01801de044c517a80443d2bd4f503a9e6866750d2f94a22970f62d
721f96e4310e4a828206d9cdeaa8f2d476705cc3bbc490a6165c68
7668f15ec178a17e3d27349b
P.y = 0068889ea2e1442245fe42bfda9e58266828c0263119f35a61631a
3358330f3bb84443fcb54fcd53a1d097fccbe310489b74ee143fc2
938959a83a1f7dd4a6fd395b
u[0] = 01aab1fb7e5cd44ba4d9f32353a383cb1bb9eb763ed40b32bdd5f6
66988970205998c0e44af6e2b5f6f8e48e969b3f649cae3c6ab463
e1b274d968d91c02f00cce91
Q.x = 01801de044c517a80443d2bd4f503a9e6866750d2f94a22970f62d
721f96e4310e4a828206d9cdeaa8f2d476705cc3bbc490a6165c68
7668f15ec178a17e3d27349b
Q.y = 0068889ea2e1442245fe42bfda9e58266828c0263119f35a61631a
3358330f3bb84443fcb54fcd53a1d097fccbe310489b74ee143fc2
938959a83a1f7dd4a6fd395b

```

J.4. curve25519

J.4.1. curve25519_XMD:SHA-512_ELL2_R0_

```

suite = curve25519_XMD:SHA-512_ELL2_R0_
dst   = QUUX-V01-CS02-with-curve25519_XMD:SHA-512_ELL2_R0_

msg =
P.x = 2de3780abb67e861289f5749d16d3e217ffa722192d16bbd9d1bfb
9d112b98c0
P.y = 3b5dc2a498941a1033d176567d457845637554a2fe7a3507d21abd
1c1bd6e878
u[0] = 005fe8a7b8fef0a16c105e6cadf5a6740b3365e18692a9c05bfb4
d97f645a6a
u[1] = 1347edbec6a2b5d8c02e058819819bee177077c9d10a4ce165aab0
fd0252261a
Q0.x = 36b4df0c864c64707cbf6cf36e9ee2c09a6cb93b28313c169be295
61bb904f98

```



```

u[0] = 46f5b22494bfeaa7f232cc8d054be68561af50230234d7d1d63d1d
9abeca8da5
Q.x = 7d56d1e08cb0ccb92baf069c18c49bb5a0dcd927eff8dcf75ca921
ef7f3e6eeb
Q.y = 404d9a7dc25c9c05c44ab9a94590e7c3fe2dcec74533a0b24b188a
5d5dacf429

msg = abcdef0123456789
P.x = 31ad08a8b0deeb2a4d8b0206ca25f567ab4e042746f792f4b7973f
3ae2096c52
P.y = 405070c28e78b4fa269427c82827261991b9718bd6c6e95d627d70
1a53c30db1
u[0] = 235fe40c443766ce7e18111c33862d66c3b33267efa50d50f9e8e5
d252a40aaa
Q.x = 3fbe66b9c9883d79e8407150e7c2a1c8680bee496c62fabe4619a7
2b3cabe90f
Q.y = 08ec476147c9a0a3ff312d303dbbd076abb7551e5fce82b48ab14b
433f8d0a7b

msg = q128_qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq
P.x = 027877759d155b1997d0d84683a313eb78bdb493271d935b622900
459d52ceaa
P.y = 54d691731a53baa30707f4a87121d5169fb5d587d70fb0292b5830
dedbec4c18
u[0] = 001e92a544463bda9bd04ddbe3d6eed248f82de32f522669efc5dd
ce95f46f5b
Q.x = 227e0bb89de700385d19ec40e857db6e6a3e634b1c32962f370d26
f84ff19683
Q.y = 5f86ff3851d262727326a32c1bf7655a03665830fa7f1b8b1e5a09
d85bc66e4a

msg = a512_aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = 5fd892c0958d1a75f54c3182a18d286efab784e774d1e017ba2fb2
52998b5dc1
P.y = 750af3c66101737423a4519ac792fb93337bd74ee751f19da4cf1e
94f4d6d0b8
u[0] = 1a68a1af9f663592291af987203393f707305c7bac9c8d63d6a729
bdc553dc19
Q.x = 3bcd651ee54d5f7b6013898aab251ee8ecc0688166fce6e9548d38
472f6bd196
Q.y = 1bb36ad9197299f111b4ef21271c41f4b7ecf5543db8bb5931307e
bdb2eaa465

```

J.5. edwards25519

J.5.1. edwards25519_XMD:SHA-512_ELL2_RO_

```
suite = edwards25519_XMD:SHA-512_ELL2_RO_
dst   = QUUX-V01-CS02-with-edwards25519_XMD:SHA-512_ELL2_RO_

msg   =
P.x   = 3c3da6925a3c3c268448dcabb47ccde5439559d9599646a8260e47
      b1e4822fc6
P.y   = 09a6c8561a0b22bef63124c588ce4c62ea83a3c899763af26d7953
      02e115dc21
u[0]  = 03fef4813c8cb5f98c6eef88fae174e6e7d5380de2b007799ac7ee
      712d203f3a
u[1]  = 780bddd137290c8f589dc687795aafae35f6b674668d92bf92ae7
      93e6a60c75
Q0.x  = 6549118f65bb617b9e8b438decedc73c496eaed496806d3b2eb9ee
      60b88e09a7
Q0.y  = 7315bcc8cf47ed68048d22bad602c6680b3382a08c7c5d3f439a97
      3fb4cf9feb
Q1.x  = 31dcfc5c58aa1bee6e760bf78cbe71c2bead8cebb2e397ece0f37a
      3da19c9ed2
Q1.y  = 7876d81474828d8a5928b50c82420b2bd0898d819e9550c5c82c39
      fc9bafa196

msg   = abc
P.x   = 608040b42285cc0d72cbb3985c6b04c935370c7361f4b7fbdb1ae7
      f8c1a8ecad
P.y   = 1a8395b88338f22e435bbd301183e7f20a5f9de643f11882fb237f
      88268a5531
u[0]  = 5081955c4141e4e7d02ec0e36becffaa1934df4d7a270f70679c78
      f9bd57c227
u[1]  = 005bdc17a9b378b6272573a31b04361f21c371b256252ae5463119
      aa0b925b76
Q0.x  = 5c1525bd5d4b4e034512949d187c39d48e8cd84242aa4758956e4a
      dc7d445573
Q0.y  = 2bf426cf7122d1a90abc7f2d108befc2ef415ce8c2d09695a74072
      40faa01f29
Q1.x  = 37b03bba828860c6b459ddad476c83e0f9285787a269df2156219b
      7e5c86210c
Q1.y  = 285ebf5412f84d0ad7bb4e136729a9ffd2195d5b8e73c0dc85110c
      e06958f432

msg   = abcdef0123456789
P.x   = 6d7fabf47a2dc03fe7d47f7dddd21082c5fb8f86743cd020f3fb14
      7d57161472
P.y   = 53060a3d140e7fbcd a641ed3cf42c88a75411e648a1add71217f70
      ea8ec561a6
u[0]  = 285ebaa3be701b79871bcb6e225ecc9b0b32dff2d60424b4c50642
      636a78d5b3
u[1]  = 2e253e6a0ef658fedb8e4bd6a62d1544fd6547922acb3598ec6b36
      9760b81b31
Q0.x  = 3ac463dd7fddb773b069c5b2b01c0f6b340638f54ee3bd92d452fc
      ec3015b52d
Q0.y  = 7b03ba1e8db9ec0b390d5c90168a6a0b7107156c994c674b61fe69
      6cbeb46baf
Q1.x  = 0757e7e904f5e86d2d2f4acf7e01c63827fde2d363985aa7432106
      f1b3a444ec
```



```

aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x      = 6e5e1f37e99345887fc12111575fc1c3e36df4b289b8759d23af14
          d774b66bfff
P.y      = 2c90c3d39eb18ff291d33441b35f3262cdd307162cc97c31bfcc7a
          4245891a37
u[0]     = 3cb0178a8137cefa5b79a3a57c858d7eeaa787b2781be4a362a2f
          0750d24fa0
Q.x      = 3e6368cff6e88a58e250c54bd27d2c989ae9b3acb6067f2651ad28
          2ab8c21cd9
Q.y      = 38fb39f1566ca118ae6c7af42810c0bb9767ae5960abb5a8ca7925
          30bfb9447d

```

J.6. curve448

J.6.1. curve448_X0F:SHAKE256_ELL2_R0_

```

suite    = curve448_X0F:SHAKE256_ELL2_R0_
dst      = QUUX-V01-CS02-with-curve448_X0F:SHAKE256_ELL2_R0_

msg      =
P.x      = 5ea5ff623d27c75e73717514134e73e419f831a875ca9e82915fdf
          c7069d0a9f8b532cfb32b1d8dd04ddeedbe3fa1d0d681c01e825d6
          a9ea
P.y      = afadd8de789f8f8e3516efbbe313a7eba364c939ecba00dabf4ced
          5c563b18e70a284c17d8f46b564c4e6ce11784a3825d9411166221
          28c1
u[0]     = c704c7b3d3b36614cf3eedd0324fe6fe7d1402c50efd16cff89ff6
          3f50938506280d3843478c08e24f7842f4e3ef45f6e3c4897f9d97
          6148
u[1]     = c25427dc97fff7a5ad0a78654e2c6c27b1c1127b5b53c7950cd1fd
          6edd2703646b25f341e73deedfeb022d1d3cecd02b93b4d585ead
          3ed7
Q0.x     = 3ba318806f89c19cc019f51e33eb6b8c038dab892e858ce7c7f2c2
          ac58618d06146a5fef31e49af49588d4d3db1bcf02bd4e4a733e37
          065d
Q0.y     = b30b4cfc2fd14d9d4b70456c0f5c6f6070be551788893d570e7955
          675a20f6c286d01d6e90d2fb500d2efb8f4e18db7f8268bb9b7fbc
          5975
Q1.x     = f03a48cf003f63be61ca055fec87c750434da07a15f8aa6210389f
          f85943b5166484339c8bea1af9fc571313d35ed2fbb779408b760c
          4cbd
Q1.y     = 23943a33b2954dc54b76a8222faf5b7e18405a41f5ecc61bf1b8df
          1f9cbfad057307ed0c7b721f19c0390b8ee3a2dec223671f9ff905
          fda7

msg      = abc
P.x      = 9b2f7ce34878d7cebf34c582db14958308ea09366d1ec71f646411
          d3de0ae564d082b06f40cd30dfc08d9fb7cb21df390cf207806ad9
          d0e4
P.y      = 138a0eef0a4993ea696152ed7db61f7ddb4e8100573591e7466d61
          c0c568ecaec939e36a84d276f34c402526d8989a96e99760c4869e
          d633
u[0]     = 2dd95593dfee26fe0d218d3d9a0a23d9e1a262fd1d0b602483d084
          15213e75e2db3c69b0a5bc89e71bcefc8c723d2b6a0cf263f02ad2

```

[illegible]

```

22318256aa8693b0c85b955cf2d6540b8ed71e729af1b8d5ca3b11
6cd7
Q0.x = c2d275826d6ad55e41a22318f6b6240f1f862a2e231120ff41eadb
ec319756032e8cef2a7ac6c10214fa0608c17fcdf61ec2694a8a2b
358b
Q0.y = 93d2e092762b135509840e609d413200df800d99da91d8b8284066
6cac30e7a3520adbaa4b089b9fdc86132e42729f651d022f4782502
f12c
Q1.x = 3c0880ece7244036e9a45944a85599f9809d772f770cc237ac41b2
1aa71615e4f3bb08f64fca618896e4f6cf5bd92e16b89d2cf6e195
6bfb
Q1.y = 45cce4beb96505cac5976b3d2673641e9bcd18d3462bbb453d293e
5282740a6389cfeae610adc7bd425c728541ceec83fcc999164af4
3fb5

msg = a512_aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = ea441c10b3636ecedd5c0dfcae96384cc40de8390a0ab648765b45
08da12c586d55dc981275776507ebca0e4d1bcaa302bb69dcfa31b
3451
P.y = fee0192d49bcc0c28d954763c2cbe739b9265c4bebe3883803c649
71220cfda60b9ac99ad986cd908c0534b260b5cfca46f6c2b0f3f2
1bda
u[0] = 8cba93a007bb2c801b1769e026b1fa1640b14a34cf3029db3c7fd6
392745d6fec0f7870b5071d6da4402cedbbde28ae4e50ab30e1049
a238
u[1] = 4223746145069e4b8a981acc3404259d1a2c3ecfed5d864798a89d
45f81a2c59e2d40eb1d5f0fe11478cbb2bb30246dd388cb932ad7b
b330
Q0.x = 4321ab02a9849128691e9b80a5c5576793a218de14885fddccb91f
17ceb1646ea00a28b69ad211e1f14f17739612dbde3782319bdf00
9689
Q0.y = 1b8a7b539519eec0ea9f7a46a43822e16cba39a439733d6847ac44
a806b8adb3e1a75ea48a1228b8937ba85c6cb6ee01046e10cad895
3b1e
Q1.x = 126d744da6a14fddec0f78a9cee4571c1320ac7645b600187812e4
d7021f98fc4703732c54daec787206e1f34d9dbbf4b292c68160b8
bfbf
Q1.y = 136eebe6020f2389d448923899a1a38a4c8ad74254e0686e91c4f9
3c1f8f8e1bd619ffb7c1281467882a9c957d22d50f65c5b72b2aee
11af

```

J.6.2. curve448_X0F:SHAKE256_ELL2_NU_

```

suite = curve448_X0F:SHAKE256_ELL2_NU_
dst   = QUUX-V01-CS02-with-curve448_X0F:SHAKE256_ELL2_NU_
msg   =

```

P.x	=	b65e8dbb279fd656f926f68d463b13ca7a982b32f5da9c7cc58afc f6199e4729863fb75ca9ae3c95c6887d95a5102637a1c5c40ff0aa fadc
P.y	=	ea1ea211cf29eca11c057fe8248181591a19f6ac51d45843a65d4b b8b71bc83a64c771ed7686218a278ef1c5d620f3d26b5316218864 5453
u[0]	=	242c70f74eac8184116c71630d284cf8a742fc463e710545847ff6 4d8e9161cb9f599728a18a32dbd8b67c3bec5d64c9b1d2f2cde7b5 888d
Q.x	=	e6304424de5af3f556d3e645600530c53ad949891c3e60ba041dd5 f68a93901bfff8440164477d348c13d28e27bfcd360c44c80b4c7d 4cea
Q.y	=	4160a8f2043a347185406a6a7e50973b98b82edbdfa3209b0e1c90 118e10eeb45045b0990d4b2b0708a30eca17df40ad53c9100f20c1 0b44
msg	=	abc
P.x	=	51aceca4fa95854bbaba58d8a5e17a86c07acadef32e1188cafda2 6232131800002cc2f27c7aec454e5e0c615bddffb7df6a5f7f0f14 793f
P.y	=	c590c9246eb28b08dee816d608ef233ea5d76e305dc458774a1e1b d880387e6734219e2018e4aa50a49486dce0ba8740065da37e6cf5 212c
u[0]	=	ef6dcb75b696d325fb36d66b104700df1480c4c17ea9190d447eee 1e7e4c9b7f36bbfb8ba7ba7c4cb6b07fed16531c1ac7a26a3618b4 0b34
Q.x	=	de0dc93df9ce7953452f20e270699c1e7dacd5d571c226d77f53b7 e3053d16f8a81b1601efb362054e973c8e733b663af93f00cb81ba f130
Q.y	=	8c5bdec6fa6690905f6eff966b0f98f5a8161493bd04976684d4ec 1f4512fa8743d86860b2ff2c5d67e9c145fd906f2cb89ff812c6b9 883f
msg	=	abcdef0123456789
P.x	=	c6d65987f146b8d0cb5d2c44e1872ac3af1f458f6a8bd8c232ffe8 b9d09496229a5a27f350eb7d97305bcc4e0f38328718352e8e3129 ed71
P.y	=	4d2f901bf333fdc4135b954f20d59207e9f6a4ecf88ce5af11c892 b44f79766ec4ecc9f60d669b95ca8940f39b1b7044140ac2040c1b f659
u[0]	=	3012ba5d9b3bb648e4613833a26ecaeadb3e8c8bba07fc90ac3da0 375769289c44d3dc87474b23df7f45f9a4030892cda689e343ae a6ad
Q.x	=	dc29532761f03c24d57f530da4c24acc4c676d185becaa89fcc083 266541fb7f10ecec91dac64a34cd988274633ae25c4d784aee52de 47a8
Q.y	=	a5f6da11259c69f2e07fce6a7b6afec4c25bd2df83426765f9c070 4111da24c6a0550d5c7aac7d648d55f7640d50be99c926195e852a daac
msg	=	q128_ff ff ff
P.x	=	9b8d008863beb4a02fb9e4efefd2eba867307fb1c7ce01746115d3 2e1db551bb254e8e3e4532d5c74a83949a69a60519ecc9178083cb e943


```
P.y = 346a1fca454d1e67c628437c270ec0f0c4256bb774fe6c0e49de70  
04ff6d9199e2cd99d8f7575a96aaafc4dc8db1811ba0a44317581f4  
1371  
u[0] = fe952ac0149f92436bba12ea2e542aa226f4fc074d79ff462c41b3  
27968a649a495a8a93b6c3044af2273456abb5e166ce4fb8c9b10c  
8c2e  
Q.x = 512803d89f59c57376e6570cd54c4e901643e089cd9456f549daa4  
372b8b52679860b68aa8bedfaa88970f15ab6098d5f252083ac98a  
58c9  
Q.y = 3d9b6593c7941a20d76161c9a171f1e507495a08f03dfcae33a2ac  
3602698e46a74d1039b583c984036f590eaa43d20ba5ada3ffb55  
2f77  
  
msg = a512_aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa  
  
P.x = 8746dc34799112d1f20acda9d7f722c9abb29b1fb6b7e9e5669838  
43c20bd7c9bfad21b45c5166b808d2f5d44e188f1fdaf29cdee8a7  
2e4c  
P.y = 7c1293484c9287c298a1a0600c64347eee8530acf563cd8705e057  
28274d8cd8101835f8003b6f3b78b5beb28f5be188a3d7bce1ec5a  
36b1  
u[0] = afd3d7ad9d819be7561706e050d4f30b634b203387ab682739365f  
62cd7393ca2cf18cd07a3d3af8dd163f043ac7457c2eb145b4a561  
70a9  
Q.x = 08aed6480793218034fd3b3b0867943d7e0bd1b6f76b4929e0885b  
d082b84d4449341da6038bb08229ad9eb7d518dff2c7ea50148e70  
a4db  
Q.y = e00d32244561ebd4b5f4ef70fcac75a06416be0a1c1b304e7bd361  
a6a6586915bb902a323eaf73cf7738e70d34282f61485395ab2833  
d2c1
```

J.7. edwards448

J.7.1. edwards448 XOF:SHAKE256 ELL2 R0

```
suite      = edwards448_X0F:SHAKE256_ELL2_R0_
dst        = QUUX-V01-CS02-with-edwards448_X0F:SHAKE256_ELL2_R0_

msg        =
P.x        = 73036d4a88949c032f01507005c133884e2f0d81f9a950826245dd
           a9e844fc78186c39daaa7147ead3e462cff60e9c6340b58134480b
           4d17
P.y        = 94c1d61b43728e5d784ef4fcb1f38e1075f3aef5e99866911de5a2
           34f1aafdc26b554344742e6ba0420b71b298671bbec2b773661863
           4610
u[0]       = 0847c5ebf957d3370b1f98fde499fb3e659996d9fc9b5707176ade
           785ba72cd84b8a5597c12b1024be5f510fa5ba99642c4cec7f3f69
           d3e7
```

u[1] = f8cbd8a7ae8c8deed071f3ac4b93e7cfcb8f1eac1645d699fd6d3881cb295a5d3006d9449ed7cad412a77a1fe61e84a9e41d59ef384d6f9a

Q0.x = c08177330869db17fb81a5e6e53b36d29086d806269760f2e4caba a4015f5dbadb7ca2ba594d96a89d0ca4f0944489e1ef393d53db85096f

Q0.y = 02e894598c050eeb7195f5791f1a5f65da3776b7534be37640bcbf95d4b915bd22333c50387583507169708fbd7bea0d7aa385dcc614be9c

Q1.x = 770877fd3b6c5503398157b68a9d3609f585f40e1ebebdd69bb0e4d3d9aa811995ce75333fdadfa50db886a35959cc59cffd5c9710daca25

Q1.y = b27fef77aa6231fbbcb27538fa90eaca8abd03eb1e62fdae4ec5e828117c3b8b3ff8c34d0a6e6d79fff16d339b94ae8ede33331d5b464c792

msg = abc

P.x = 4e0158acacffa545adb818a6ed8e0b870e6abc24dfc1dc45cf9a052e98469275d9ff0c168d6a5ac7ec05b742412ee090581f12aa398f9f8c

P.y = 894d3fa437b2d2e28cdc3bfaade035430f350ec5239b6b406b5501da6f6d6210fff26719cad83b63e97ab26a12df6dec851d6bf38e294af9a

u[0] = 04d975cd938ab49be3e81703d6a57cca84ed80d2ff6d4756d3f22947fb5b70ab0231f0087cbfb4b7cae73b41b0c9396b356a4831d9a14322

u[1] = 2547ca887ac3db7b5fad3a098aa476e90078afe1358af6c63d677d6edfd2100bc004e0f5db94dd2560fc5b308e223241d00488c9ca6b0ef2

Q0.x = 7544612a97f4419c94ab0f621a1ee8ccf46c6657b8e0778ec9718bf4b41bc774487ad87d9b1e617aa49d3a4dd35a3cf57cd390ebf0429952

Q0.y = d3ab703e60267d796b485bb58a28f934bd0133a6d1bbdfeda5277fa293310be262d7f653a5adffa608c37ed45c0e6008e54a16e1a342e4df

Q1.x = 6262f18d064bc131ade1b8bbcbf1cbdf984f4f88153fcc9f94c888af35d5e41aae84c12f169a55d8abf06e6de6c5b23079e587a58cf73303e

Q1.y = 6d57589e901abe7d947c93ab02c307ad9093ed9a83eb0b6e829fb7318d590381ca25f3cc628a36a924a9ddfcf3cbcdf94edf3b338ea77403

msg = abcdef0123456789

P.x = 2c25b4503fadc94b27391933b557abdecc601c13ed51c5de68389484f93dbd6c22e5f962d9babf7a39f39f994312f8ca23344847e1fbf176

P.y = d5e6f5350f430e53a110f5ac7fcc82a96cb865aeca982029522d32601e41c042a9dfbdfbfa2b0bdc3bc58cca8a7cd546803083d3a8548

u[0] = 10659ce25588db4e4be6f7c791a79eb21a7f24aaaca76a6ca3b83b80aaf95aa328fe7d569a1ac99f9cd216edf3915d72632f1a8b990e250c

u[1] = 9243e5b6c480683fd533e81f4a778349a309ce00bd163a29eb9fa8dbc8f549242bef33e030db21cffacd408d2c4264b93e476c6a8590e7aa

Q0.x = 1457b60c12e00e47ceb3ce64b57e7c3c61636475443d704a8e2b2a

Q0.y = b0a5ac7e4b3909435416784e16e19929c653b1bdcd9478a8e5331c
a9ae
935d9f75f7a0babbc39c0a1c3b412518ed8a24bc2c4886722fb4b7
d4a747af98e4e2528c75221e2dfffd3424abb436e10539a74caaafa
3ea3
Q1.x = b44d9e34211b4028f24117e856585ed81448f3c8b934987a1c5939
c86048737a08d85934fec6b3c2ef9f09cbd365cf22744f2e4ce697
62a4
Q1.y = dc996c1736f4319868f897d9a27c45b02dd3bc6b7ca356a039606e
5406e131a0bbe8238208b327b00853e8af84b58b13443e70542556
3323
msg = q128_qqq
qq
qq
P.x = a1861a9464ae31249a0e60bf38791f3663049a3f5378998499a832
92e159a2fecff838eb9bc6939e5c6ae76eb074ad4aae39b55b72ca
0b9a
P.y = 580a2798c5b904f8adfec5bd29fb49b4633cd9f8c2935eb4a0f12e
5dfa0285680880296bb729c6405337525fb5ed3dff930c137314f6
0401
u[0] = c80390020e578f009ead417029efff6cd0926110922db63ab98395e
3bdfdd5d8a65b1a2b8d495dc8c5e59b7f3518731f7dfc0f93ace5d
ee4b
u[1] = 1c4dc6653a445bbef2add81d8e90a6c8591a788deb91d0d3f1519a
2e4a460313041b77c1b0817f2e80b388e5c3e49f37d787dc1f85e4
324a
Q0.x = 9d355251e245e4b13ed4ea3e5a3c55bf9b7211f1704771f2e1d8f1
a65610c468b1cf70c6c2ce30dcaad54ad9e5439471ec554b862ec8
875a
Q0.y = 6689ba36a242af69ac2aadb955d15e982d9b04f5d77f7609ebf742
9587feb7e5ce27490b9c72114509f89565122074e46a614d7fd7c8
00bd
Q1.x = c4b3d3ad4d2d62739a62989532992c1081e9474a201085b4616da5
706cab824693b9fb428a201bcd1639a4588cc43b9eb841dbca7421
9b1f
Q1.y = 265286f5dee8f3d894b5649da8565b58e96b4cfd44b462a2883ea6
4dbcda21a00706ea3fea53fc2d769084b0b74589e91d0384d71189
09fb
msg = a512_aaa
aa
aa
aa
aa
aa
aa
aa
aa
aa
aa
aa
aa
aa
P.x = 987c5ac19dd4b47835466a50b2d9feba7c8491b8885a04edf577e1
5a9f2c98b203ec2cd3e5390b3d20bba0fa6fc3eecefb5029a31723
4401
P.y = 5e273fcfff6b007bb6771e90509275a71ff1480c459ded26fc7b10
664db0a68aaa98bc7ecb07e49cf05b80ae5ac653fbdd14276bbd35
ccbc

```

u[0]    = 163c79ab0210a4b5e4f44fb19437ea965bf5431ab233ef16606f0b
          03c5f16a3feb7d46a5a675ce8f606e9c2bf74ee5336c54a1e54919
          f13f
u[1]    = f99666bde4995c4088333d6c2734687e815f80a99c6da02c47df4b
          51f6c9d9ed466b4fecf7d9884990a8e0d0be6907fa437e0b1a27f4
          9265
Q0.x    = d1a5eba4a332514b69760948af09ceaeddbbb9fd4cb1f19b78349c
          2ee4cf9ee86dbcf9064659a4a0566fe9c34d90aec86f0801edc131
          ad9b
Q0.y    = 5d0a75a3014c3269c33b1b5da80706a4f097893461df286353484d
          8031cd607c98edc2a846c77a841f057c7251eb45077853c7b20595
          7e52
Q1.x    = 69583b00dc6b2aced6ffa44630cc8c8cd0dd0649f57588dd0fb1da
          ad2ce132e281d01e3f25ccd3f405be759975c6484268bfe8f5e5f2
          3c30
Q1.y    = 8418484035f60bdccf48cb488634c2dfb40272123435f7e654fb6f
          254c6c42e7e38f1fa79a637a168a28de6c275232b704f9ded0ff76
          dd94

```

J.7.2. edwards448_XOF:SHAKE256_ELL2_NU_

```

suite   = edwards448_XOF:SHAKE256_ELL2_NU_
dst     = QUUX-V01-CS02-with-edwards448_XOF:SHAKE256_ELL2_NU_

msg     =
P.x     = eb5a1fc376fd73230af2de0f3374087cc7f279f0460114cf0a6c12
          d6d044c16de34ec2350c34b26bf110377655ab77936869d085406a
          f71e
P.y     = df5dcea6d42e8f494b279a500d09e895d26ac703d75ca6d118e8ca
          58bf6f608a2a383f292fce1563ff995dce75aede1fdc8e7c0c737a
          e9ad
u[0]    = 1368aefc0416867ea2cfc515416bcbeecc9ec81c4ecbd52ccdb91e
          06996b3f359bc930eef6743c7a2dd7adb785bc7093ed044efed950
          86d7
Q.x     = 4b2abf8c0fca49d027c2a81bf73bb5990e05f3e76c7ba137cc0b89
          415ccd55ce7f191cc0c11b0560c1cdc2a8085dd56996079e05a3cd
          8dde
Q.y     = 82532f5b0cb3bfb8542d3228d055bfe61129dbeae8bace80cf61f1
          7725e8ec8226a24f0e687f78f01da88e3b2715194a03dca7c0a96b
          bf04

msg     = abc
P.x     = 4623a64bceaba3202df76cd8b6e3daf70164f3fcbda6d6e340f7fa
          b5cdf89140d955f722524f5fe4d968fef6ba2853ff4ea086c2f67d
          8110
P.y     = abaac321a169761a8802ab5b5d10061fec1a83c670ac6bc9595470
          0317ee5f82870120e0e2c5a21b12a0c7ad17ebd343363604c4bcec
          afd1
u[0]    = cda3b0ecfe054c4077007d7300969ec24f4c741300b630ec9188eb
          ab31a5ae0065612ee22d9f793733179ffc2e10c53ca5b539057aaf
          dc2f
Q.x     = b1ca5bef2f157673a210f56c9b0039db8399e4749585abac64f831
          f74ed1ec5f591928976c687c06d57686bacb98440e77af878349cd
          f2d2
Q.y     = 5bbfd6a3730d517b03c3cd9e2eed94af12891334ec090e0495c2ed
          c588e9e10b6f63b03a62076808cbcd6da95adfb5af76c136b2d42e

```

0dac

```
msg = abcdef0123456789
P.x = e9eb562e76db093baa43a31b7edd04ec4aadcef3389a7b9c58a19c
      f87f8ae3d154e134b6b3ed45847a741e33df51903da681629a4b8b
      cc2e
P.y = 0cf6606927ad7eb15dbc193993bc7e4dda744b311a8ec4274c8f73
      8f74f605934582474c79260f60280fe35bd37d4347e59184cbfa12
      cbc4
u[0] = d36bae98351512c382c7a3e1eba22497574f11fef9867901b1a270
      0b39fa2cd0d38ed4380387a99162b7ba0240c743f0532ef60d577c
      413d
Q.x = 958a51e2f02e0dfd3930709010d5d16f869adb9d8a8f7c01139911
      d206c20cdb7bfb40ee33ba30536a99f49362fa7633d0f417fc3914
      fe21
Q.y = f4307a36ab6612fa97501497f01afa109733ce85875935551c3ca9
      0f0fa7e0097a8640bb7e5dbcc38ab32b23b748790f2261f2c44c3b
      f3ba
```

```
msg = q128_qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq  
      qqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqqq  
P.x   = 122a3234d34b26c69749f23356452bf9501efa2d94859d5ef741fe  
        f024156d9d191a03a2ad24c38186f93e02d05572575968b083d8a3  
        9738  
P.y   = ddf55e74eb4414c2c1fa4aa6bc37c4ab470a3fed6bb5af1e435703  
        09b162fb61879bb15f9ea49c712efd42d0a71666430f9f0d4a2050  
        5050  
u[0]  = 5945744d27122f89da3daf76ab4db9616053df64e25d30ec9a0066  
        7ee6710240579c1db8f8ef3386f3f4f413cfb325ac14094d582026  
        a971  
Q.x   = e7e1f2d13548ac2c8fc346e4c63606545bf93652011721e83ac3b  
        64226f77a8823d3881e164bc6ca45505b236e8e3721c028052fcc9  
        ade5  
Q.y   = 7e0f340501bf25f018b9d374c2acbddd43c07261d85a6ef3c855113  
        d4e023634db59a87b8fab9efe04ed1fee302c8a4994e83bdda32bd  
        9c0b
```

```
msg = a512_aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
      aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x = 221704949b1ce1ab8dd174dc9b8c56fcffa27179569ce9219c0c2f
      e183d3d23343a4c42a0e2e9d6b9d0feb1df3883ec489b6671d1fa6
      4089
P.y = ebdecfdc87142d1a919034bf22ecfad934c9a85effff14b594ae2c
      00943ca62a39d6ee3be9df0bb504ce8a9e1669bc6959c42ad6a1d3
      b686
u[0] = 1192e378043f01cedc7ea0209321519213b0184ea0d8575816bcd9
      182a367823e1eccc2faf1df8f79b24027a4b9bfa208cd320e79bef
```

```

Q.x      = 06ea
          = 0fd3bb833c1d7a5b319d1d4117406a23b9aece976186ecb18a11a6
          = 35e6fbdb920d47e04762b1f2a8c59d2f8435d0fdefe501f544cda2
          = 3dbf
Q.y      = f13b0dad4d5eeb120f2443ac4392f8096a1396f5014ec2a3506a34
          = 7fef8076a7282035cf619599b1919cf29df5ce87711c11688aab77
          = 00a6

```

J.8. secp256k1

J.8.1. secp256k1_XMD:SHA-256_SSWU_RO_

```

suite    = secp256k1_XMD:SHA-256_SSWU_RO_
dst      = QUUX-V01-CS02-with-secp256k1_XMD:SHA-256_SSWU_RO_

msg      =
P.x      = c1cae290e291aee617ebaef1be6d73861479c48b841eaba9b7b585
          = 2ddfeb1346
P.y      = 64fa678e07ae116126f08b022a94af6de15985c996c3a91b64c406
          = a960e51067
u[0]     = 6b0f9910dd2ba71c78f2ee9f04d73b5f4c5f7fc773a701abea1e57
          = 3cab002fb3
u[1]     = 1ae6c212e08fe1a5937f6202f929a2cc8ef4ee5b9782db68b0d579
          = 9fd8f09e16
Q0.x     = 74519ef88b32b425a095e4ebcc84d81b64e9e2c2675340a720bb1a
          = 1857b99f1e
Q0.y     = c174fa322ab7c192e11748beed45b508e9fdb1ce046dee9c2cd3a2
          = a86b410936
Q1.x     = 44548adb1b399263ded3510554d28b4bead34b8cf9a37b4bd0bd2b
          = a4db87ae63
Q1.y     = 96eb8e2faf05e368efe5957c6167001760233e6dd2487516b46ae7
          = 25c4cce0c6

msg      = abc
P.x      = 3377e01eab42db296b512293120c6cee72b6ecf9f9205760bd9ff1
          = 1fb3cb2c4b
P.y      = 7f95890f33efebd1044d382a01b1bee0900fb6116f94688d487c6c
          = 7b9c8371f6
u[0]     = 128aab5d3679a1f7601e3bdf94ced1f43e491f544767e18a4873f3
          = 97b08a2b61
u[1]     = 5897b65da3b595a813d0fdcc75c895dc531be76a03518b044daaa0
          = f2e4689e00
Q0.x     = 07dd9432d426845fb19857d1b3a91722436604ccbbbadad8523b8f
          = c38a5322d7
Q0.y     = 604588ef5138cffe3277bbd590b8550bcbe0e523bbaf1bed4014a4
          = 67122eb33f
Q1.x     = e9ef9794d15d4e77dde751e06c182782046b8dac05f8491eb88764
          = fc65321f78
Q1.y     = cb07ce53670d5314bf236ee2c871455c562dd76314aa41f012919f
          = e8e7f717b3

msg      = abcdef0123456789
P.x      = bac54083f293f1fe08e4a70137260aa90783a5cb84d3f35848b324
          = d0674b0e3a
P.y      = 4436476085d4c3c4508b60fcf4389c40176adce756b398bdee27bc
          = a19758d828

```



```

aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
P.x      = 082aabaeb7dedb0e78aeb619ad3bfd9277a2f77ba7fad20ef6aab
          dc6c31d19ba5a6d12283553294c1825c4b3ca2dcfe
P.y      = 05b84ae5a942248eea39e1d91030458c40153f3b654ab7872d779a
          d1e942856a20c438e8d99bc8abfbf74729ce1f7ac8
u[0]     = 0a8ffa7447f6be1c5a2ea4b959c9454b431e29ccc0802bc052413a
          9c5b4f9aac67a93431bd480d15be1e057c8a08e8c6
u[1]     = 05d487032f602c90fa7625dbafe0f4a49ef4a6b0b33d7bb349ff4c
          f5410d297fd6241876e3e77b651cfc8191e40a68b7
Q0.x     = 0cf97e6dbd0947857f3e578231d07b309c622ade08f2c08b32ff37
          2bd90db19467b2563cc997d4407968d4ac80e154f8
Q0.y     = 127f0cddf2613058101a5701f4cb9d0861fd6c2a1b8e0afe194fcc
          f586a3201a53874a2761a9ab6d7220c68661a35ab3
Q1.x     = 092f1acfa62b05f95884c6791fba989bbe58044ee6355d100973bf
          9553ade52b47929264e6ae770fb264582d8dce512a
Q1.y     = 028e6d0169a72cfedb737be45db6c401d3adfb12c58c619c82b93a
          5dfcccef12290de530b0480575ddc8397cda0bbebf

```

J.9.2. BLS12381G1_XMD:SHA-256_SSWU_NU_

```

suite    = BLS12381G1_XMD:SHA-256_SSWU_NU_
dst      = QUUX-V01-CS02-with-BLS12381G1_XMD:SHA-256_SSWU_NU_

msg      =
P.x      = 184bb665c37ff561a89ec2122dd343f20e0f4cbcaec84e3c3052ea
          81d1834e192c426074b02ed3dca4e7676ce4ce48ba
P.y      = 04407b8d35af4dacc809927071fc0405218f1401a6d15af775810e
          4e460064bcc9468beeba82fdc751be70476c888bf3
u[0]     = 156c8a6a2c184569d69a76be144b5cdc5141d2d2ca4fe341f011e2
          5e3969c55ad9e9b9ce2eb833c81a908e5fa4ac5f03
Q.x      = 11398d3b324810a1b093f8e35aa8571cced95858207e7f49c4fd74
          656096d61d8a2f9a23cdb18a4dd11cd1d66f41f709
Q.y      = 19316b6fb2ba7717355d5d66a361899057e1e84a6823039efc7bec
          cefe09d023fb2713b1c415fcf278eb0c39a89b4f72

msg      = abc
P.x      = 009769f3ab59bfd551d53a5f846b9984c59b97d6842b20a2c565ba
          a167945e3d026a3755b6345df8ec7e6acb6868ae6d
P.y      = 1532c00cf61aa3d0ce3e5aa20c3b531a2abd2c770a790a26138183
          03c6b830ffc0ecf6c357af3317b9575c567f11cd2c
u[0]     = 147e1ed29f06e4c5079b9d14fc89d2820d32419b990c1c7bb7dbea
          2a36a045124b31fffbde7c99329c05c559af1c6cc82
Q.x      = 1998321bc27ff6d71df3051b5aec12ff47363d81a5e9d2dff55f44
          4f6ca7e7d6af45c56fd029c58237c266ef5cda5254
Q.y      = 034d274476c6307ae584f951c82e7ea85b84f72d28f4d647173235
          6121af8d62a49bc263e8eb913a6cf6f125995514ee

msg      = abcdef0123456789
P.x      = 1974dbb8e6b5d20b84df7e625e2fbfecb2cdb5f77d5eae5fb2955e
          5ce7313cae8364bc2fff520a6c25619739c6bdcb6a
P.y      = 15f9897e11c6441eaa676de141c8d83c37aab8667173cbe1dfd6de
          74d11861b961dccebcd9d289ac633455dfcc7013a3
u[0]     = 04090815ad598a06897dd89bcda860f25837d54e897298ce31e694
          7378134d3761dc59a572154963e8c954919ecfa82d

```

[illegible]

J.10. BLS12-381 G2

J.10.1. BLS12381G2 XMD:SHA-256 SSWU R0

```
suite      = BLS12381G2_XMD:SHA-256_SSWU_RO
dst        = QUUX-V01-CS02-with-BLS12381G2_XMD:SHA-256_SSWU_RO_

msg        =
P.x        = 0141ebfbdca40eb85b87142e130ab689c673cf60f1a3e98d693352
           66f30d9b8d4ac44c1038e9dcdd5393faf5c41fb78a
           + I * 05cb8437535e20ecffaef7752baddf98034139c38452458baeefab
           379ba13dff5bf5dd71b72418717047f5b0f37da03d
P.y        = 0503921d7f6a12805e72940b963c0cf3471c7b2a524950ca195d11
           062ee75ec076daf2d4bc358c4b190c0c98064fdd92
           + I * 12424ac32561493f3fe3c260708a12b7c620e7be00099a974e259d
           dc7d1f6395c3c811cdd19f1e8dbf3e9ecfdcbab8d6
```

```

u[0]      = 03dbc2cce174e91ba93cbb08f26b917f98194a2ea08d1cce75b2b9
           + I * 05a2acec64114845711a54199ea339abd125ba38253b70a92c876d
u[1]      = 02f99798e8a5acdeed60d7e18e9120521ba1f47ec090984662846b
           + I * 145a81e418d4010cc027a68f14391b30074e89e60ee7a22f87217b
Q0.x      = 019ad3fc9c72425a998d7ab1ea0e646a1f6093444fc6965f1cad5a
           + I * 171c88b0b0efb5eb2b88913a9e74fe111a4f68867b59db252ce586
Q0.y      = 0ba10604e62bdd9eeeb4156652066167b72c8d743b050fb4c1016c
           + I * 0eb22c7a543d3d376e9716a49b72e79a89c9bfe9feee8533ed931c
Q1.x      = 113d2b9cd4bd98aee53470b27abc658d91b47a78a51584f3d4b950
           + I * 13855912321c5cb793e9d1e88f6f8d342d49c0b0dbac613ee9e17e
Q1.y      = 0fd3def0b7574a1d801be44fde617162aa2e89da47f464317d9bb5
           + I * 056f617902b3c0d0f78a9a8cbda43a26b65f602f8786540b9469b0
           60db7b38417915b413ca65f875c130bebfaa59790c

msg        = abc
P.x        = 02c2d18e033b960562aae3cab37a27ce00d80ccd5ba4b7fe0e7a21
           + I * 0245129dbec7780ccc7954725f4168aff2787776e6
P.y        = 1787327b68159716a37440985269cf584bcb1e621d3a7202be6ea0
           + I * 00aa65dae3c8d732d10ecd2c50f8a1baf3001578f71c694e03866e
u[0]      = 15f7c0aa8f6b296ab5ff9c2c7581ade64f4ee6f1bf18f55179ff44
           + I * 01c8067bf4c0ba709aa8b9abc3d1cef589a4758e09ef53732d670f
u[1]      = 187111d5e088b6b9acfdfad078c4dacf72dcd17ca17c82be35e79f
           + I * 08b852331c96ed983e497ebc6dee9b75e373d923b729194af8e72a
Q0.x      = 051ea586f3538a6ebb1e80881a082fa2b24df9f566
           + I * 12b2e525281b5f4d2276954e84ac4f42cf4e13b6ac4228624e1776
Q0.y      = 0faf94ce5706d53f0ca1952f1c5ef75239aeed55ad
           + I * 05d8a724db78e570e34100c0bc4a5fa84ad5839359b40398151f37
Q1.x      = cff5a51de945c563463c9efbdda569850ee5a53e77
           + I * 02eacdc556d0bdb5d18d22f23dcb086dd106cad713777c7e640794
Q1.y      = 3edbe0b3d1efe391eedf11e977fac55f9b94f2489c
           + I * 04bbe48bfd5814648d0b9e30f0717b34015d45a861425fab1ee06
           fdfce36384ae2c808185e693ae97dcde118f34de41
           + I * 19f18cc5ec0c2f055e47c802acc3b0e40c337256a208001dde14b2
           5afced146f37ea3d3ce16834c78175b3ed61f3c537
           + I * 15b0dadcd256a258b4c68ea43605dfffa6d312eef215c19e6474b3e1
           01d33b661dfec43b51abbf96fee68fc6043ac56a58
           + I * 05e47c1781286e61c7ade887512bd9c2cb9f640d3be9cf87ea0bad
           24bd0ebfe946497b48a581ab6c7d4ca74b5147287f

```

+ I * 19f98db2f4a1fcdf56a9ced7b320ea9deecf57c8e59236b0dc21f6
ee7229aa9705ce9ac7fe7a31c72edca0d92370c096

msg = abcdef0123456789

P.x = 121982811d2491fde9ba7ed31ef9ca474f0e1501297f68c298e9f4
c0028add35aea8bb83d53c08cfc007c1e005723cd0

+ I * 190d119345b94fbd15497bcba94ecf7db2cbfd1e1fe7da034d26cb
ba169fb3968288b3fafb265f9ebd380512a71c3f2c

P.y = 05571a0f8d3c08d094576981f4a3b8eda0a8e771fcdcc8ecceaf13
56a6acf17574518acb506e435b639353c2e14827c8

+ I * 0bb5e7572275c567462d91807de765611490205a941a5a6af3b169
1bfe596c31225d3aabd15faff860cb4ef17c7c3be

u[0] = 0313d9325081b415bfd4e5364efaef392ecf69b087496973b22930
3e1816d2080971470f7da112c4eb43053130b785e1

+ I * 062f84cb21ed89406890c051a0e8b9cf6c575cf6e8e18ecf63ba86
826b0ae02548d83b483b79e48512b82a6c0686df8f

u[1] = 1739123845406baa7be5c5dc74492051b6d42504de008c635f3535
bb831d478a341420e67dcc7b46b2e8cba5379cca97

+ I * 01897665d9cb5db16a27657760bbea7951f67ad68f8d55f7113f24
ba6ddd82caef240a9bfa627972279974894701d975

Q0.x = 0f48f1ea1318ddb713697708f7327781fb39718971d72a9245b973
1faaca4dbaa7cca433d6c434a820c28b18e20ea208

+ I * 06051467c8f85da5ba2540974758f7a1e0239a5981de441fdd8768
0a995649c211054869c50edbacc1f3a86c561ba3162

Q0.y = 168b3d6df80069dbbedb714d41b32961ad064c227355e1ce5fac8e
105de5e49d77f0c64867f3834848f152497eb76333

+ I * 134e0e8331cee8cb12f9c2d0742714ed9eee78a84d634c9a95f6a7
391b37125ed48bfc6e90bf3546e99930ff67cc97bc

Q1.x = 004fd03968cd1c99a0dd84551f44c206c84dcdbd78076c5bfee24e
89a92c8508b52b88b68a92258403cbe1ea2da3495f

+ I * 1674338ea298281b636b2eb0fe593008d03171195fd6dcd4531e8a
1ed1f02a72da238a17a635de307d7d24aa2d969a47

Q1.y = 0dc7fa13fff6b12558419e0a1e94bfc3cfaf67238009991c5f24ee
94b632c3d09e27eca329989aee348a67b50d5e236c

+ I * 169585e164c131103d85324f2d7747b23b91d66ae5d947c449c819
4a347969fc6bbd967729768da485ba71868df8aed2

msg = q128_qqq
qq
qq

P.x = 19a84dd7248a1066f737cc34502ee5555bd3c19f2ecdb3c7d9e24d
c65d4e25e50d83f0f77105e955d78f4762d33c17da

+ I * 0934aba516a52d8ae479939a91998299c76d39cc0c035cd18813be
c433f587e2d7a4fef038260eef0cef4d02aae3eb91

P.y = 14f81cd421617428bc3b9fe25afbb751d934a00493524bc4e06563
5b0555084dd54679df1536101b2c979c0152d09192

+ I * 09bcccfa036b4847c9950780733633f13619994394c23ff0b32fa6
b795844f4a0673e20282d07bc69641cee04f5e5662

u[0] = 025820cefc7d06fd38de7d8e370e0da8a52498be9b53cba9927b2e
f5c6de1e12e12f188bbc7bc923864883c57e49e253

+ I * 034147b77ce337a52e5948f66db0bab47a8d038e712123bb381899
b6ab5ad20f02805601e6104c29df18c254b8618c7b

u[1] = 0930315cae1f9a6017c3f0c8f2314baa130e1cf13f6532bfff0a8a1
790cd70af918088c3db94bda214e896e1543629795

+ I * 10c4df2cacf67ea3cb3108b00d4cbd0b3968031ebc8eac4b1ebcef
e84d6b715fde66bef0219951ece29d1facc8a520ef

Q0.x = 09eccbc53df677f0e5814e3f86e41e146422834854a224bf5a83a5
+ I * 0c3633943f91daee715277bd644fba585168a72f96ded64fc5a384
Q0.y = 0eb6186a0457d5b12d132902d4468bfeb7315d83320b6c32f1c875
+ I * 119e3cf167e69eb16c1c7830e8df88856d48be12e3ff0a40791a5c
Q1.x = 0eb3aabc1ddfce17ff18455fcc7167d15ce6b60ddc9eb9b59f8d40
+ I * 198bdfb19d7441ebcca61e8ff774b29d17da16547d2c10c273227a
Q1.y = 0aaf1dee3adf3ed4c80e481c09b57ea4c705e1b8d25b897f0ceec
+ I * 3990748716575f92abff22a1c8f4582aff7b872d52
msg = a512_aaa
P.x = 01a6ba2f9a11fa5598b2d8ace0fbc0a0eachb65deceb476fbbcb64f
+ I * 11fca2ff525572795a801eed17eb12785887c7b63fb77a42be46ce
P.y = 0b6798718c8aed24bc19cb27f866f1c9effcddf92397ad6448b5c9
+ I * 03a47f8e6d1763ba0cad63d6114c0accbef65707825a511b251a66
u[0] = 190b513da3e66fc9a3587b78c76d1d132b1152174d0b83e3c11140
+ I * 12ab625b0fe0ebd1367fe9fac57bb1168891846039b4216b9d9400
u[1] = 0e6a42010cf435fb5bacc156a585e1ea3294cc81d0ceb81924d950
+ I * 117d9a0defc57a33ed208428cb84e54c85a6840e7648480ae42883
Q0.x = 17cadf8d04a1a170f8347d42856526a24cc466cb2ddfd506cff011
+ I * 0d13ba91f2a8b0051cf3279ea0ee63a9f19bc9cb8bfcc7d78b3cbd
Q0.y = 17ef19497d6d9246fa94d35575c0f8d06ee02f21a284dbeaa78768
+ I * 12c3c913ba4ed03c24f0721a81a6be7430f2971ffca8fd1729aafe
Q1.x = 16ec57b7fe04c71dfe34fb5ad84dbce5a2dbbd6ee085f1d8cd17f4
+ I * 13103f7aace1ae1420d208a537f7d3a9679c287208026e4e3439ab
Q1.y = 0958b2c4c2c10fcef5a6c59b9e92c4a67b0fae3e2e0f1b6b5edad9

+ I * c940b8f3524ba9ebbc3f2ceb3cfe377655b3163bd7
0ccb594ed8bd14ca64ed9cb4e0aba221be540f25dd0d6ba15a4a4b
e5d67bcf35df7853b2d8dad3ba245f1ea3697f66aa

J.10.2. BLS12381G2_XMD:SHA-256_SSWU_NU_

suite = BLS12381G2_XMD:SHA-256_SSWU_NU_
dst = QUUX-V01-CS02-with-BLS12381G2_XMD:SHA-256_SSWU_NU_

msg =
P.x = 00e7f4568a82b4b7dc1f14c6aaa055edf51502319c723c4dc2688c
7fe5944c213f510328082396515734b6612c4e7bb7
+ I * 126b855e9e69b1f691f816e48ac6977664d24d99f8724868a18418
6469ddfd4617367e94527d4b74fc86413483afb35b
P.y = 0caeada0fd7b6176c01436833c79d305c78be307da5f6af6c133c47
311def6ff1e0babf57a0fb5539fce7ee12407b0a42
+ I * 1498aadcf7ae2b345243e281ae076df6de84455d766ab6fcdaad71
fab60abb2e8b980a440043cd305db09d283c895e3d
u[0] = 07355d25caf6e7f2f0cb2812ca0e513bd026ed09dda65b177500fa
31714e09ea0ded3a078b526bed3307f804d4b93b04
+ I * 02829ce3c021339ccb5caf3e187f6370e1e2a311dec9b753631170
63ab2015603ff52c3d3b98f19c2f65575e99e8b78c
Q.x = 18ed3794ad43c781816c523776188deafba67ab773189b8f18c49b
c7aa841cd81525171f7a5203b2a340579192403bef
+ I * 0727d90785d179e7b5732c8a34b660335fed03b913710b60903cf4
954b651ed3466dc3728e21855ae822d4a0f1d06587
Q.y = 00764a5cf6c5f61c52c838523460eb2168b5a5b43705e19cb612e0
06f29b717897facfd15dd1c8874c915f6d53d0342d
+ I * 19290bb9797c12c1d275817aa2605ebe42275b66860f0e4d04487e
bc2e47c50b36edd86c685a60c20a2bd584a82b011a

msg = abc
P.x = 108ed59fd9fae381abfd1d6bce2fd2fa220990f0f837fa30e0f279
14ed6e1454db0d1ee957b219f61da6ff8be0d6441f
+ I * 0296238ea82c6d4adb3c838ee3cb2346049c90b96d602d7bb1b469
b905c9228be25c627bffee872def773d5b2a2eb57d
P.y = 033f90f6057aadacae7963b0a0b379dd46750c1c94a6357c99b65f
63b79e321ff50fe3053330911c56b6ceea08fee656
+ I * 153606c417e59fb331b7ae6bce4fbf7c5190c33ce9402b5ebe2b70
e44fca614f3f1382a3625ed5493843d0b0a652fc3f
u[0] = 138879a9559e24cecee8697b8b4ad32cced053138ab913b9987277
2dc753a2967ed50aabc907937aefb2439ba06cc50c
+ I * 0a1ae7999ea9bab1dcc9ef8887a6cb6e8f1e22566015428d220b7e
ec90ffa70ad1f624018a9ad11e78d588bd3617f9f2
Q.x = 0f40e1d5025ecef0d850aa0bb7bbeceab21a3d4e85e6bee857805b
09693051f5b25428c6be343edba5f14317fcc30143
+ I * 02e0d261f2b9fee88b82804ec83db330caa75fbb12719cfa71ccce
1c532dc4e1e79b0a6a281ed8d3817524286c8bc04c
Q.y = 0cf4a4adc5c66da0bca4caddc6a57ecd97c8252d7526a8ff478e0d
fed816c4d321b5c3039c6683ae9b1e6a3a38c9c0ae
+ I * 11cad1646bb3768c04be2ab2bbe1f80263b7ff6f8f9488f5bc3b68
50e5a3e97e20acc583613c69cf3d2bfe8489744ebb

msg = abcdef0123456789
P.x = 038af300ef34c7759a6caaa4e69363cafeed218a1f207e93b2c70d
91a1263d375d6730bd6b6509dcac3ba5b567e85bf3

[illegible]

[illegible]

[illegible]

[illegible]

c6aeaea0d62d5d62

```
msg =
len_in_bytes = 0x80
DST_prime = acb9736c0867fdfbd6385519b90fc8c034b5af04a95897321295
0132d035792f20
msg_prime = 0080acb9736c0867fdfbd6385519b90fc8c034b5af04a9589732
12950132d035792f20
uniform_bytes = 3890dbab00a2830be398524b71c2713bbef5f4884ac2e6f0
70b092effdb19208c7df943dc5dcbaee3094a78c267ef276632ee2
c8ea0c05363c94b6348500fae4208345dd3475fe0c834c2beac7fa
7bc181692fb728c0a53d809fc8111495222ce0f38468b11becb15b
32060218e285c57a60162c2c8bb5b6bde13973cd41819
```

```
msg         = abc
len_in_bytes = 0x80
DST_prime   = acb9736c0867fdfdbd6385519b90fc8c034b5af04a95897321295
            0132d035792f20
msg_prime    = 6162630080acb9736c0867fdfdbd6385519b90fc8c034b5af04a9
            58973212950132d035792f20
uniform_bytes = 41b7ffa7a301b5c1441495ebb9774e2a53dbbf4e54b9a1af
            6a20fd41eafd69ef7b9418599c5545b1ee422f363642b01d4a5344
            9313f68da3e49dddb9cd25b97465170537d45dcdbdf92391b5bdfbf3
            44db4bd06311a05bca7dcd360b6caec849c299133e5c9194f4e15e
            3e23cfaab4003fab776f6ac0bfae9144c6e2e1c62e7d57
```

```
msg          = abcdef0123456789
len_in_bytes = 0x80
DST_prime    = acb9736c0867fdfeb6385519b90fc8c034b5af04a95897321295
              0132d035792f20
msg_prime    = 616263646566303132333435363738390080acb9736c0867fdfeb
              d6385519b90fc8c034b5af04a958973212950132d035792f20
uniform_bytes = 55317e4a21318472cd2290c3082957e1242241d9e0d04f47
              026f03401643131401071f01aa03038b2783e795bdfa8a3541c194
              ad5de7cb9c225133e24af6c86e748deb52e560569bd54ef4dac034
              65111a3a44b0ea490fb36777ff8ea9f1a8a3e8e0de3cf0880b4b2f
              8dd37d3a85a8b82375aee4fa0e909f9763319b55778e71
```

[illegible]

[illegible]

[illegible]

Acknowledgements

The authors would like to thank Adam Langley for his detailed writeup of Elligator 2 with Curve25519 [L13]; Dan Boneh, Benjamin Lipp, Christopher Patton, and Leonid Reyzin for educational discussions; and David Benjamin, Daniel Bourdrez, Frank Denis, Sean Devlin, Justin Drake, Bjoern Haase, Mike Hamburg, Dan Harkins, Daira Hopwood, Thomas Icart, Andy Polyakov, Thomas Pornin, Mamy Ratsimbazafy, Michael Scott, Filippo Valsorda, and Mathy Vanhoef for helpful reviews and feedback.

Contributors

Sharon Goldberg
Boston University
Email: goldbe@cs.bu.edu

Ela Lee
Royal Holloway, University of London
Email: Ela.Lee.2010@live.rhul.ac.uk

Michele Orru
Email: michele.orrु@ens.fr

Authors' Addresses

Armando Faz-Hernandez
Cloudflare, Inc.
101 Townsend St
San Francisco, CA 94107
United States of America
Email: armfazh@cloudflare.com

Sam Scott
Oso Security, Inc.
335 Madison Ave
New York, NY 10017
United States of America
Email: sam.scott89@gmail.com

Nick Sullivan
Cloudflare, Inc.
101 Townsend St
San Francisco, CA 94107
United States of America
Email: nicholas.sullivan@gmail.com

Riad S. Wahby
Stanford University

Email: rsw@cs.stanford.edu

Christopher A. Wood
Cloudflare, Inc.
101 Townsend St
San Francisco, CA 94107
United States of America
Email: caw@heapingbits.net