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DECoupled Application Data Enroute (DECADE)

Abstract

Content distribution applications, such as those employing peer-to-peer (P2P) technologies, are widely used on the Internet and make up a large portion of the traffic in many networks. Often, however, content distribution applications use network resources inefficiently. One way to improve efficiency is to introduce storage capabilities within the network and enable cooperation between end-host and in-network content distribution mechanisms. This is the capability provided by a DECoupled Application Data Enroute (DECADE) system, which is introduced in this document. DECADE enables applications to take advantage of in-network storage when distributing data objects as opposed to using solely end-to-end resources. This document presents the underlying principles and key functionalities of such a system and illustrates operation through a set of examples.

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1. Introduction

Content distribution applications, such as peer-to-peer (P2P) applications, are widely used on the Internet to distribute data objects and make up a large portion of the traffic in many networks. Said applications can often introduce performance bottlenecks in otherwise well-provisioned networks. In some cases, operators are forced to invest substantially in infrastructure to accommodate the use of such applications. For instance, in many subscriber networks, it can be expensive to upgrade network equipment in the "last mile", because it can involve replacing equipment and upgrading wiring and devices at individual homes, businesses, DSLAMs (Digital Subscriber Line Access Multiplexers), and CMTSs (Cable Modem Termination Systems) in remote locations. It may be more practical and economical to upgrade the core infrastructure, instead of the "last mile" of the network, as this involves fewer components that are shared by many subscribers. See [RFC6646] and [RFC6392] for a more complete discussion of the problem domain and general discussions of the capabilities envisioned for a DECADE system. As a historical point, it should be noted that [RFC6646] and [RFC6392] came out of the now closed DECADE Working Group. This document aims to advance some of the valuable concepts from that now closed Working Group.

This document presents mechanisms for providing in-network storage that can be integrated into content distribution applications. The primary focus is P2P-based content distribution, but DECADE may be useful to other applications with similar characteristics and requirements (e.g., Content Distribution Networks (CDNs) or hybrid P2P/CDNs). The approach we adopt in this document is to define the core functionalities and protocol functions that are needed to support a DECADE system. This document provides illustrative examples so that implementers can understand the main concepts in DECADE, but it is generally assumed that readers are also familiar with the terms and concepts used in [RFC6646] and [RFC6392].

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

2. Terminology

This document uses the following terminology.

Application Endpoint

A host that includes a DECADE client along with other application functionalities (e.g., peer-to-peer (P2P) client, video streaming client).

Content Distribution Application

A specific type of application that may exist in an Application Endpoint. A content distribution application is an application (e.g., P2P) designed for dissemination of large amounts of content (e.g., files or video streams) to multiple peers. Content distribution applications may divide content into smaller blocks for dissemination.

Data Object

A data object is the unit of data stored and retrieved from a DECADE server. The data object is a sequence of raw bytes. The server maintains metadata associated with each data object, but the metadata is physically and logically separate from the data object.

DECADE Client

A DECADE client uploads and/or retrieves data from a DECADE server.

DECADE Resource Protocol (DRP)

A logical protocol for communication of access control and resource-scheduling policies from a DECADE client to a DECADE server, or between DECADE servers. In practice, the functionality of the DRP may be distributed over one or more actual protocols.

DECADE Server

A DECADE server stores data inside the network for a DECADE client or another DECADE server, and thereafter it manages both the stored data and access to that data by other DECADE clients.

DECADE Storage Provider

A DECADE storage provider deploys and/or manages DECADE servers within a network.

DECADE System

An in-network storage system that is composed of DECADE clients and DECADE servers. The DECADE servers may be deployed by one or more DECADE storage providers.

In-Network Storage

A service inside a network that provides storage to applications. In-network storage may reduce upload/transit/backbone traffic and improve application performance. In-network storage may, for example, be co-located with the border router (network-attached storage) or inside a data center. A DECADE system is an example of an in-network storage system.

Standard Data Transfer (SDT) Protocol
A logical protocol used to transfer data objects between a DECADE client and DECADE server, or between DECADE servers. The intent is that in practice the SDT should map to an existing, well-known protocol already in use over the Internet for transporting data.

Overview

A DECADE system provides a distributed storage service for content distribution applications (e.g., P2P). The system consists of clients and servers. A client first uploads data objects to one or more selected servers and optionally requests distribution of these data objects to other servers. The client then selectively authorizes other clients to download these data objects. Such a system is employed in an overall application context (e.g., P2P file sharing), and it is expected that DECADE clients take part in application-specific communication sessions.

Figure 1 is a schematic of a simple DECADE system with two DECADE clients and two DECADE servers. As illustrated, a DECADE client, which is part of an Application Endpoint, uses the DECADE Resource Protocol (DRP) to convey to a server information related to access control and resource-scheduling policies. DRP can also be used between servers for exchanging this type of information. A DECADE system employs the Standard Data Transfer (SDT) protocol to transfer data objects to and from a server, as we will explain later.

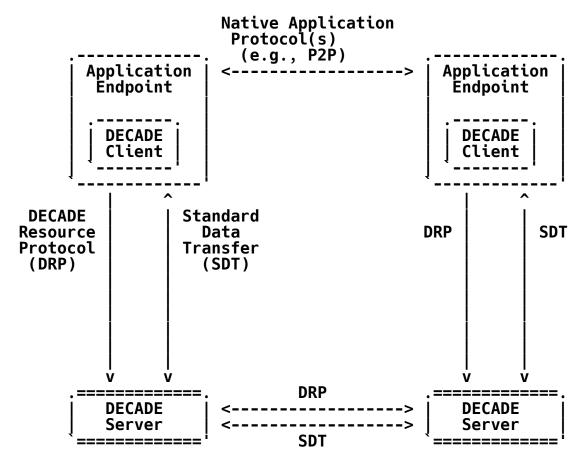


Figure 1: DECADE Overview

With Figure 1 at hand, assume that Application Endpoint B requests a data object from Application Endpoint A using their native application protocols (e.g., P2P protocol) as in Figure 2. In this case, Endpoint A will act as the sender, and Endpoint B as the receiver for said data object. S(A) is the DECADE storage server which is access controlled. This means, first, that Endpoint A has a right to store the data object in S(A). Secondly, Endpoint B needs to obtain authorization before being able to retrieve the data object from S(A).

The four steps involved in a DECADE session are illustrated in Figure 2. The sequence starts with the initial contact between Endpoint B and Endpoint A, where Endpoint B requests a data object using their native application protocol (e.g., P2P). Next, Endpoint A uses DRP to obtain a token corresponding to the data object that was requested by Endpoint B. There may be several ways for Endpoint A to obtain such a token, e.g., compute it locally or request one from its DECADE storage server, S(A). Once obtained, Endpoint A then

provides the token to Endpoint B (again, using their native application protocol). Finally, Endpoint B provides the received token to S(A) via DRP, and subsequently requests and downloads the data object via SDT. Again, it is assumed that DECADE is employed in an overall application context (e.g., P2P file-sharing session).

For completeness, note that there is an important prerequisite step (not shown) to Figure 2, where Endpoint A first discovers and then stores the data object(s) of interest in S(A).

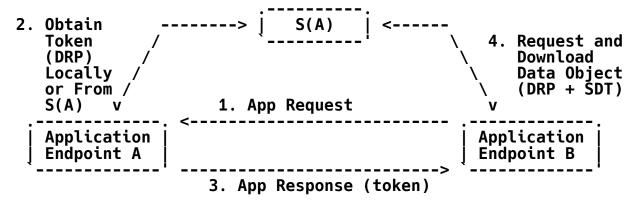


Figure 2: Download from Storage Server

4. Architectural Principles

This section presents the key principles followed by any DECADE system.

4.1. Data- and Control-Plane Decoupling

DECADE SDT and DRP can be classified as belonging to data-plane functionality. The algorithms and signaling for a P2P application, for example, would belong to control-plane functionality.

A DECADE system aims to be application independent and should support multiple content distribution applications. Typically, a complete content distribution application implements a set of control-plane functions including content search, indexing and collection, access control, replication, request routing, and QoS scheduling. Implementers of different content distribution applications may have unique considerations when designing the control-plane functions. For example, with respect to the metadata management scheme, traditional file systems provide a standard metadata abstraction: a recursive structure of directories to offer namespace management where each file is an opaque byte stream. Content distribution applications may use different metadata management schemes. For

instance, one application might use a sequence of blocks (e.g., for file sharing), while another application might use a sequence of frames (with different sizes) indexed by time.

With respect to resource-scheduling algorithms, a major advantage of many successful P2P systems is their substantial expertise in achieving efficient utilization of peer resources. For instance, many streaming P2P systems include optimization algorithms for constructing overlay topologies that can support low-latency, high-bandwidth streaming. The research community as well as implementers of such systems continuously fine-tune existing algorithms and invent new ones. A DECADE system should be able to accommodate and benefit from all new developments.

In short, given the diversity of control-plane functions, a DECADE system should allow for as much flexibility as possible to the control plane to implement specific policies (and be decoupled from data-plane DRP/SDT). Decoupling the control plane from the data plane is not new, of course. For example, OpenFlow [OpenFlow] is an implementation of this principle for Internet routing, where the computation of the forwarding table and the application of the forwarding table are separated. The Google File System [GoogleFileSystem] applies the same principle to file system design by utilizing a Master to handle metadata management and several Chunk servers to handle data-plane functions (i.e., read and write of chunks of data). Finally, NFSv4.1's parallel NFS (pNFS) extension [RFC5661] also adheres to this principle.

4.2. Immutable Data Objects

A common property of bulk content to be broadly distributed is that it is immutable -- once content is generated, it is typically not modified. For example, once a movie has been edited and released for distribution, it is very uncommon that the corresponding video frames and images need to be modified. The same applies to document distribution, such as RFCs; audio files, such as podcasts; and program patches. Focusing on immutable data can substantially simplify data-plane design, since consistency requirements can be relaxed. It also simplifies data reuse and the removal of duplicates.

Depending on its specific requirements, an application may store immutable data objects in DECADE servers such that each data object is completely self-contained (e.g., a complete, independently decodable video segment). An application may also divide data into data objects that require application-level assembly. Many content distribution applications divide bulk content into data objects for multiple reasons, including (a) fetching different data objects from

different sources in parallel and (b) faster recovery and verification as individual data objects might be recovered and verified. Typically, applications use a data object size larger than a single packet in order to reduce control overhead.

A DECADE system should be agnostic to the nature of the data objects and should not specify a fixed size for them. A protocol specification based on this architecture may prescribe requirements on minimum and maximum sizes for compliant implementations.

Note that immutable data objects can still be deleted. Applications can support modification of existing data stored at a DECADE server through a combination of storing new data objects and deleting existing data objects. For example, a metadata management function of the control plane might associate a name with a sequence of immutable data objects. If one of the data objects is modified, the meta-data management function changes the mapping of the name to a new sequence of immutable data objects.

4.3. Data Object Identifiers

A data object stored in a DECADE server shall be accessed by DECADE clients via a data object identifier. Each DECADE client may be able to access more than one storage server. A data object that is replicated across different storage servers managed by a storage provider may be accessed through a single identifier. Since data objects are immutable, it shall be possible to support persistent identifiers for data objects.

Data object identifiers should be created by DECADE clients when uploading the corresponding objects to a DECADE server. The scheme for the assignment/derivation of the data object identifier to a data object depends as the data object naming scheme and is out of scope of this document. One possibility is to name data objects using hashes as described in [RFC6920]. Note that [RFC6920] describes naming schemes on a semantic level only, but specific SDTs and DRPs use specific representations.

In particular, for some applications, it is important that clients and servers be able to validate the name-object binding, i.e., by verifying that a received object really corresponds to the name (identifier) that was used for requesting it (or that was provided by a sender). If a specific application requires name-object binding validation, the data object identifiers can support it by providing message digests or so-called self-certifying naming information.

Different name-object binding validation mechanisms may be supported in a single DECADE system. Content distribution applications can decide what mechanism to use, or to not provide name-object validation (e.g., if authenticity and integrity can by ascertained by alternative means). We expect that applications may be able to construct unique names (with high probability) without requiring a registry or other forms of coordination. Names may be self-describing so that a receiving DECADE client understands, for example, which hash function to use for validating name-object binding.

Some content distribution applications will derive the name of a data object from the hash over the data object; this is made possible by the fact that DECADE objects are immutable. But there may be other applications such as live streaming where object names will not based on hashes but rather on an enumeration scheme. The naming scheme will also enable those applications to construct unique names.

In order to enable the uniqueness, flexibility and self-describing properties, the naming scheme used in a DECADE system should provide a "type" field that indicates the name-object validation function type (for example, "sha-256" [RFC5754]) and the cryptographic data (such as an object hash) that corresponds to the type information. Moreover, the naming scheme may additionally provide application or publisher information.

4.4. Explicit Control

To support the functions of an application's control plane, applications should be able to keep track and coordinate which data is stored at particular servers. Thus, in contrast with traditional caches, applications are given explicit control over the placement (selection of a DECADE server), deletion (or expiration policy), and access control for stored data objects. Consider deletion/expiration policy as a simple example. An application might require that a DECADE server stores data objects for a relatively short period of time (e.g., for live-streaming data). Another application might need to store data objects for a longer duration (e.g., for video on demand), and so on.

4.5. Resource and Data Access Control through Delegation

A DECADE system provides a shared infrastructure to be used by multiple Application Endpoints. Thus, it needs to provide both resource and data access control, as discussed in the following subsections.

4.5.1. Resource Allocation

There are two primary interacting entities in a DECADE system. First, storage providers coordinate DECADE server provisioning, including their total available resources. Second, applications coordinate data transfers amongst available DECADE servers and between servers and clients. A form of isolation is required to enable each of the concurrently running applications to explicitly manage its own data objects and share of resources at the available servers. Therefore, a storage provider should delegate resource management on a DECADE server to uploading DECADE clients, enabling them to explicitly and independently manage their own share of resources on a server.

4.5.2. User Delegation

DECADE storage providers will have the ability to explicitly manage the entities allowed to utilize the resources available on a DECADE server. This is needed for reasons such as capacity-planning and legal considerations in certain deployment scenarios. The DECADE server should grant a share of the resources to a DECADE client. The client can in turn share the granted resources amongst its (possibly) multiple applications. The share of resources granted by a server is called a User Delegation. As a simple example, a DECADE server operated by an ISP might be configured to grant each ISP subscriber 1.5 Mbit/s of network capacity and 1 GB of memory. The ISP subscriber might in turn divide this share of resources amongst a video-streaming application and file-sharing application that are running concurrently.

5. System Components

As noted earlier, the primary focus of this document is the architectural principles and the system components that implement them. While specific system components might differ between implementations, this document details the major components and their overall roles in the architecture. To keep the scope narrow, we only discuss the primary components related to protocol development. Particular deployments will require additional components (e.g., monitoring and accounting at a server), but they are intentionally omitted from this document.

5.1. Application Endpoint

Content distribution applications have many functional components. For example, many P2P applications have components and algorithms to manage overlay topology, rate allocation, piece selection, and so on. In this document, we focus on the components directly engaged in a DECADE system. Figure 3 illustrates the components discussed in this section from the perspective of a single Application Endpoint.

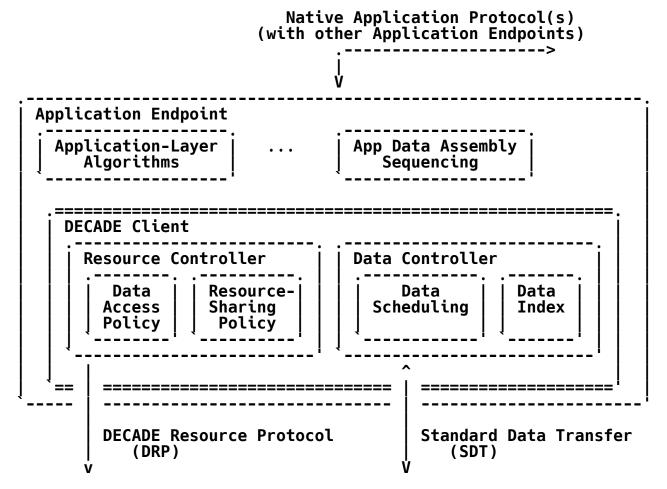


Figure 3: Application and DECADE Client Components

A DECADE system is geared towards supporting applications that can distribute content using data objects (e.g., P2P). To accomplish this, applications can include a component responsible for creating the individual data objects before distribution and for reassembling them later. We call this component Application Data Assembly. In producing and assembling data objects, two important considerations are sequencing and naming. A DECADE system assumes that applications

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implement this functionality themselves. In addition to DECADE DRP/SDT, applications will most likely also support other, native application protocols (e.g., P2P control and data transfer protocols).

5.2. DECADE Client

The DECADE client provides the local support to an application, and it can be implemented standalone, embedded into the application, or integrated in other software entities within network devices (i.e., hosts). In general, applications may have different resource-sharing policies and data access policies with regard to DECADE servers. These policies may be existing policies of applications or custom policies. The specific implementation is decided by the application.

Recall that DECADE decouples the control and the data transfer of applications. A data-scheduling component schedules data transfers according to network conditions, available servers, and/or available server resources. The Data Index indicates data available at remote servers. The Data Index (or a subset of it) can be advertised to other clients. A common use case for this is to provide the ability to locate data amongst distributed Application Endpoints (i.e., a data search mechanism such as a Distributed Hash Table (DHT)).

5.3. DECADE Server

Figure 4 illustrates the primary components of a DECADE server. Note that the description below does not assume a single-host or centralized implementation -- a DECADE server is not necessarily a single physical machine; it can also be implemented in a distributed manner on a cluster of machines.

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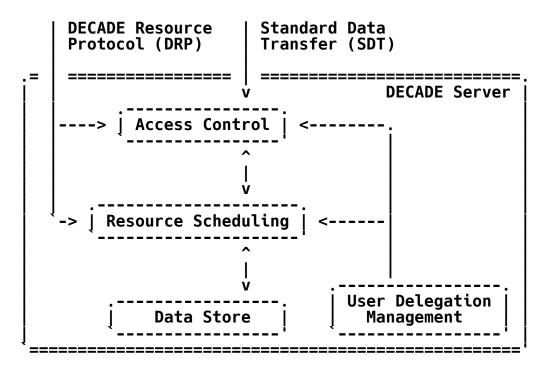


Figure 4: DECADE Server Components

Provided sufficient authorization, a client shall be able to access its own data or other client's data in a DECADE server. Clients may also authorize other clients to store data. If access is authorized by a client, the server should provide access. Applications may apply resource-sharing policies or use a custom policy. DECADE servers will then perform resource scheduling according to the resource-sharing policies indicated by the client as well as any other previously configured User Delegations. Data from applications will be stored at a DECADE server. Data may be deleted from storage either explicitly or automatically (e.g., after a Time To Live (TTL) expiration).

5.4. Data Sequencing and Naming

The DECADE naming scheme implies no sequencing or grouping of objects, even if this is done at the application layer. To illustrate these properties, this section presents several examples of use.

5.4.1. Application with Fixed-Size Chunks

Consider an application in which each individual application-layer segment of data is called a "chunk" and has a name of the form: "CONTENT_ID:SEQUENCE_NUMBER". Furthermore, assume that the application's native protocol uses chunks of size 16 KB. Now, assume that this application wishes to store data in a DECADE server in data objects of size 64 KB. To accomplish this, it can map a sequence of 4 chunks into a single data object, as shown in Figure 5.

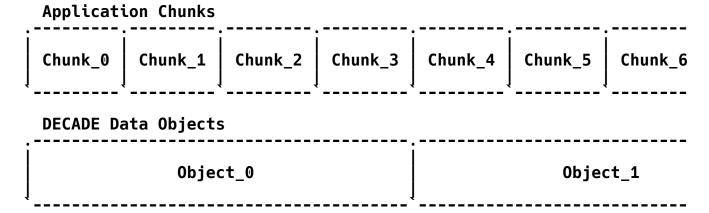


Figure 5: Mapping Application Chunks to DECADE Data Objects

In this example, the application maintains a logical mapping that is able to determine the name of a DECADE data object given the chunks contained within that data object. The name may be conveyed from either the original uploading DECADE client, another Endpoint with which the application is communicating, etc. As long as the data contained within each sequence of chunks is globally unique, the corresponding data objects have globally unique names.

5.4.2. Application with Continuous Streaming Data

Consider an application whose native protocol retrieves a continuous data stream (e.g., an MPEG2 stream) instead of downloading and redistributing chunks of data. Such an application could segment the continuous data stream to produce either fixed-sized or variable-sized data objects. Figure 6 depicts how a video streaming application might produce variable-sized data objects such that each data object contains 10 seconds of video data. In a manner similar to the previous example, the application may maintain a mapping that is able to determine the name of a data object given the time offset of the video chunk.

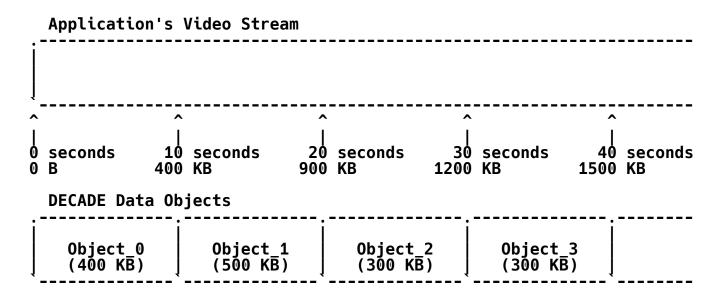


Figure 6: Mapping a Continuous Data Stream to DECADE Data Objects

5.5. Token-Based Authorization and Resource Control

A key feature of a DECADE system is that an Application Endpoint can authorize other Application Endpoints to store or retrieve data objects from its in-network storage via tokens. The peer client then uses the token when sending requests to the DECADE server. Upon receiving a token, the server validates the signature and the operation being performed.

This is a simple scheme, but has some important advantages over an alternative approach, for example, in which a client explicitly manipulates an Access Control List (ACL) associated with each data object. In particular, it has the following advantages when applied to DECADE systems. First, authorization policies are implemented within the application, thus the Application Endpoint explicitly controls when tokens are generated, to whom they are distributed, and for how long they will be valid. Second, fine-grained access and resource control can be applied to data objects. Third, there is no messaging between a client and server to manipulate data object permissions. This can simplify, in particular, applications that share data objects with many dynamic peers and need to frequently adjust access control policies attached to data objects. Finally, tokens can provide anonymous access, in which a server does not need to know the identity of each client that accesses it. This enables a client to send tokens to clients belonging to other storage providers, and to allow them to read or write data objects from the storage of its own storage provider. In addition to clients' ability to apply access control policies to data objects, the server may be

configured to apply additional policies based on user, object properties, geographic location, etc. A client might thus be denied access even though it possesses a valid token.

5.6. Discovery

A DECADE system should include a discovery mechanism through which DECADE clients locate an appropriate DECADE server. A discovery mechanism should allow a client to determine an IP address or some other identifier that can be resolved to locate the server for which the client will be authorized to generate tokens (via DRP). (The discovery mechanism might also result in an error if no such servers can be located.) After discovering one or more servers, a DECADE client can distribute load and requests across them (subject to resource limitations and policies of the servers themselves) according to the policies of the Application Endpoint in which it is embedded. The discovery mechanism outlined here does not provide the ability to locate arbitrary DECADE servers to which a client might obtain tokens from others. To do so will require application-level knowledge, and it is assumed that this functionality is implemented in the content distribution application.

As noted above, the discovered DECADE server should be authorized to allow the client to store data objects and then generate tokens to allow other clients to retrieve these data objects. This authorization may be:

- a result of off-line administrative procedures;
- access network dependent (e.g., all the subscribers to a particular ISP may be allowed by the ISP);
- due to a prior subscription;
- etc.

The particular protocol used for discovery is out of scope of this document, but any specification should reuse well-known protocols wherever possible.

6. DECADE Protocol Considerations

This section presents the DRP and the SDT protocol in terms of abstract protocol interactions that are intended to be mapped to specific protocols in an implementation. In general, the DRP/SDT functionality for DECADE client-server interaction is very similar to that for server-server interaction. Any differences are highlighted below. DRP is used by a DECADE client to configure the resources and authorization used to satisfy requests (reading, writing, and management operations concerning data objects) at a server. SDT will be used to transport data between a client and a server, as illustrated in Figure 1.

6.1. Naming

A DECADE system SHOULD use [RFC6920] as the recommended and default naming scheme. Other naming schemes that meet the guidelines in Section 4.3 MAY alternatively be used. In order to provide a simple and generic interface, the DECADE server will be responsible only for storing and retrieving individual data objects.

The DECADE naming format SHOULD NOT attempt to replace any naming or sequencing of data objects already performed by an application. Instead, naming is intended to apply only to data objects referenced by DECADE-specific purposes. An application using a DECADE client may use a naming and sequencing scheme independent of DECADE names. The DECADE client SHOULD maintain a mapping from its own data objects and their names to the DECADE-specific data objects and names. Furthermore, the DECADE naming scheme implies no sequencing or grouping of objects, even if this is done at the application layer.

6.2. Resource Protocol

DRP will provide configuration of access control and resource-sharing policies on DECADE servers. A content distribution application (e.g., a live P2P streaming session) can have permission to manage data at several servers, for instance, servers belonging to different storage providers. DRP allows one instance of such an application, i.e., an Application Endpoint, to apply access control and resource-sharing policies on each of them.

On a single DECADE server, the following resources SHOULD be managed: a) communication resources in terms of bandwidth (upload/download) and also in terms of number of active clients (simultaneous connections); and b) storage resources.

6.2.1. Access and Resource Control Token

The tokens SHOULD be generated by an entity trusted by both the DECADE client and the server at the request of a DECADE client. For example, this entity could be the client, a server trusted by the client, or another server managed by a storage provider and trusted by the client. It is important for a server to trust the entity generating the tokens since each token may incur a resource cost on the server when used. Likewise, it is important for a client to trust the entity generating the tokens since the tokens grant access to the data stored at the server.

The token does not normally include information about the identity of the authorized client (i.e., it is typically an anonymous token). However, it is not prohibited to have a binding of the token to an identity if desired (e.g., binding of the token to the IP address of the authorized party).

Upon generating a token, a DECADE client can distribute it to another client. Token confidentiality SHOULD be provided by whatever protocol it is carried in (i.e., Application Protocol, DRP, or SDT). The receiving client can then connect to the server specified in the token and perform any operation permitted by the token. The token SHOULD be sent along with the operation. The server SHOULD validate the token to identify the client that issued it and whether the requested operation is permitted by the contents of the token. If the token is successfully validated, the server SHOULD apply the resource control policies indicated in the token while performing the operation.

Tokens SHOULD include a unique identifier to allow a server to detect when a token is used multiple times and reject the additional usage attempts. Since usage of a token incurs resource costs to a server (e.g., bandwidth and storage) and an uploading DECADE client may have a limited budget, the uploading DECADE client should be able to indicate if a token may be used multiple times.

It SHOULD be possible to revoke tokens after they are generated. This could be accomplished by supplying the server the unique identifiers of the tokens that are to be revoked.

6.2.2. Status Information

DRP SHOULD provide a status request service that clients can use to request status information of a server. Access to such status information SHOULD require client authorization; that is, clients need to be authorized to access the requested status information. This authorization is based on the user delegation concept as

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described in Section 4.5. The following status information elements SHOULD be obtained: a) list of associated data objects (with properties); and b) resources used/available. In addition, the following information elements MAY be available: c) list of servers to which data objects have been distributed (in a certain time frame); and d) list of clients to which data objects have been distributed (in a certain time frame).

For the list of servers/clients to which data objects have been distributed to, the server SHOULD be able to decide on time bounds for which this information is stored and specify the corresponding time frame in the response to such requests. Some of this information may be used for accounting purposes, e.g., the list of clients to which data objects have been distributed.

Access information MAY be provided for accounting purposes, for example, when uploading DECADE clients are interested in access statistics for resources and/or to perform accounting per user. Again, access to such information requires client authorization and SHOULD be based on the delegation concept as described in Section 4.5. The following type of access information elements MAY be requested: a) what data objects have been accessed by whom and how many times; and b) access tokens that a server has seen for a given data object.

The server SHOULD decide on time bounds for which this information is stored and specify the corresponding time frame in the response to such requests.

6.2.3. Data Object Attributes

Data objects that are stored on a DECADE server SHOULD have associated attributes (in addition to the object identifier) that relate to the data storage and its management. These attributes may be used by the server (and possibly the underlying storage system) to perform specialized processing or handling for the data object, or to attach related server or storage-layer properties to the data object. These attributes have a scope local to a server. In particular, these attributes SHOULD NOT be applied to a server or client to which a data object is copied.

Depending on authorization, clients SHOULD be permitted to get or set such attributes. This authorization is based on the delegation as per Section 4.5. DECADE does not limit the set of permissible attributes, but rather specifies a set of baseline attributes that SHOULD be supported:

Expiration Time: time at which the data object can be deleted

Data Object size: in bytes

Media type: labeling of type as per [RFC6838]

Access statistics: how often the data object has been accessed (and what tokens have been used)

The data object attributes defined here are distinct from application metadata. Application metadata is custom information that an application might wish to associate with a data object to understand its semantic meaning (e.g., whether it is video and/or audio, its playback length in time, or its index in a stream). If an application wishes to store such metadata persistently, it can be stored within data objects themselves.

6.3. Data Transfer

A DECADE server will provide a data access interface, and SDT will be used to write data objects to a server and to read (download) data objects from a server. Semantically, SDT is a client-server protocol; that is, the server always responds to client requests.

To write a data object, a client first generates the object's name (see Section 6.1), and then uploads the object to a server and supplies the generated name. The name can be used to access (download) the object later; for example, the client can pass the name as a reference to other clients that can then refer to the object. Data objects can be self-contained objects such as multimedia resources, files, etc., but also chunks, such as chunks of a P2P distribution protocol that can be part of a containing object or a stream. If supported, a server can verify the integrity and other security properties of uploaded objects.

A client can request named data objects from a server. In a corresponding request message, a client specifies the object name and a suitable access and resource control token. The server checks the validity of the received token and its associated properties related to resource usage. If the named data object exists on the server and the token can be validated, the server delivers the requested object in a response message. If the data object cannot be delivered, the server provides a corresponding status/reason information in a response message. Specifics regarding error handling, including additional error conditions (e.g., overload), precedence for returned errors and its relation with server policy, are deferred to eventual protocol specification.

6.4. Server-Server Protocols

An important feature of a DECADE system is the capability for one server to directly download data objects from another server. This capability allows applications to directly replicate data objects between servers without requiring end-hosts to use uplink capacity to upload data objects to a different server.

DRP and SDT SHOULD support operations directly between servers. Servers are not assumed to trust each other nor are they configured to do so. All data operations are performed on behalf of clients via explicit instruction. However, the objects being processed do not necessarily have to originate or terminate at the client (i.e., the data object might be limited to being exchanged between servers even if the instruction is triggered by the client). Clients thus will be able to indicate to a server which remote server(s) to access, what operation is to be performed, or in which server the object is to be stored, and the credentials indicating access and resource control to perform the operation at the remote server.

Server-server support is focused on reading and writing data objects between servers. The data object referred to at the remote server is the same as the original data object requested by the client. Object attributes might also be specified in the request to the remote server. In this way, a server acts as a proxy for a client, and a client can instantiate requests via that proxy. The operations will be performed as if the original requester had its own client colocated with the server. When a client sends a request to a server with these additional parameters, it is giving the server permission to act (proxy) on its behalf. Thus, it would be prudent for the supplied token to have narrow privileges (e.g., limited to only the necessary data objects) or validity time (e.g., a small expiration time).

In the case of a retrieval operation, the server is to retrieve the data object from the remote server using the specified credentials, and then optionally return the object to a client. In the case of a storage operation, the server is to store the object to the remote server using the specified credentials. The object might optionally be uploaded from the client or might already exist at the server.

6.5. Potential DRP/SDT Candidates

Having covered the key DRP/SDT functionalities above, it is useful to consider some potential DRP/SDT candidates as guidance for future DECADE protocol implementations. To recap, the DRP is a protocol for communication of access control and resource-scheduling policies from a DECADE client to a DECADE server, or between DECADE servers. The

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SDT is a protocol used to transfer data objects between a DECADE client and DECADE server, or between DECADE servers. An evaluation of existing protocols for their suitability for DRP and SDT is given in Appendix A. Also, [INTEGRATION-EX] provides some experimental examples of how to integrate DECADE-like in-network storage infrastructure into P2P applications.

7. How In-Network Storage Components Map to DECADE

This section evaluates how the basic components of an in-network storage system (see Section 3 of [RFC6392]) map into a DECADE system.

With respect to the data access interface, DECADE clients can read and write objects of arbitrary size through the client's Data Controller, making use of standard data transfer (SDT). With respect to data management operations, clients can move or delete previously stored objects via the client's Data Controller, making use of SDT. Clients can enumerate or search contents of servers to find objects matching desired criteria through services provided by the content distribution application (e.g., buffer-map exchanges, a DHT, or peer exchange). In doing so, Application Endpoints might consult their local Data Index in the client's Data Controller (Data Search Capability).

With respect to access control authorization, all methods of access control are supported: public-unrestricted, public-restricted, and private. Access control policies are generated by a content distribution application and provided to the client's Resource Controller. The server is responsible for implementing the access control checks. Clients can manage the resources (e.g., bandwidth) on the DECADE server that can be used by other Application Endpoints (Resource Control Interface). Resource-sharing policies are generated by a content distribution application and provided to the client's Resource Controller. The server is responsible for implementing the resource-sharing policies.

Although the particular protocol used for discovery is outside the scope of this document, different options and considerations have been discussed in Section 5.6. Finally, with respect to the storage mode, DECADE servers provide an object-based storage mode. Immutable data objects might be stored at a server. Applications might consider existing blocks as data objects, or they might adjust block sizes before storing in a server.

8. Security Considerations

In general, the security considerations mentioned in [RFC6646] apply to this document as well. A DECADE system provides a distributed storage service for content distribution and similar applications. The system consists of servers and clients that use these servers to upload data objects, to request distribution of data objects, and to download data objects. Such a system is employed in an overall application context (for example, in a P2P application), and it is expected that DECADE clients take part in application-specific communication sessions. The security considerations here focus on threats related to the DECADE system and its communication services, i.e., the DRP/SDT protocols that have been described in an abstract fashion in this document.

8.1. Threat: System Denial-of-Service Attacks

A DECADE network might be used to distribute data objects from one client to a set of servers using the server-server communication feature that a client can request when uploading an object. Multiple clients uploading many objects at different servers at the same time and requesting server-server distribution for them could thus mount massive distributed denial-of-service (DDOS) attacks, overloading a network of servers. This threat is addressed by the server's access control and resource control framework. Servers can require Application Endpoints to be authorized to store and to download objects, and Application Endpoints can delegate authorization to other Application Endpoints using the token mechanism. Of course the effective security of this approach depends on the strength of the token mechanism. See below for a discussion of this and related communication security threats.

Denial-of-service attacks against a single server (directing many requests to that server) might still lead to considerable load for processing requests and invalidating tokens. SDT therefore MUST provide a redirection mechanism to allow requests to other servers. Analogous to how an HTTP reverse proxy can redirect and load balance across multiple HTTP origin servers [RFC2616].

8.2. Threat: Authorization Mechanisms Compromised

A DECADE system does not require Application Endpoints to authenticate in order to access a server for downloading objects, since authorization is not based on Endpoint or user identities but on a delegation-based authorization mechanism. Hence, most protocol security threats are related to the authorization scheme. The security of the token mechanism depends on the strength of the token mechanism and on the secrecy of the tokens. A token can represent

authorization to store a certain amount of data, to download certain objects, to download a certain amount of data per time, etc. If it is possible for an attacker to guess, construct, or simply obtain tokens, the integrity of the data maintained by the servers is compromised.

This is a general security threat that applies to authorization delegation schemes. Specifications of existing delegation schemes such as [RFC6749] discuss these general threats in detail. We can say that the DRP has to specify appropriate algorithms for token generation. Moreover, authorization tokens should have a limited validity period that should be specified by the application. Token confidentiality should be provided by application protocols that carry tokens, and the SDT and DRP should provide secure (confidential) communication modes.

8.3. Threat: Spoofing of Data Objects

In a DECADE system, an Application Endpoint is referring other Application Endpoints to servers to download a specified data object. An attacker could "inject" a faked version of the object into this process, so that the downloading Endpoint effectively receives a different object (compared to what the uploading Endpoint provided). As a result, the downloading Endpoint believes that is has received an object that corresponds to the name it was provided earlier, whereas in fact it is a faked object. Corresponding attacks could be mounted against the application protocol (that is used for referring other Endpoints to servers), servers themselves (and their storage subsystems), and the SDT by which the object is uploaded, distributed, and downloaded.

A DECADE systems fundamental mechanism against object spoofing is name-object binding validation, i.e., the ability of a receiver to check whether the name it was provided and that it used to request an object actually corresponds to the bits it received. As described above, this allows for different forms of name-object binding, for example, using hashes of data objects, with different hash functions (different algorithms, different digest lengths). For those application scenarios where hashes of data objects are not applicable (for example, live streaming), other forms of name-object binding can be used. This flexibility also addresses cryptographic algorithm evolution: hash functions might get deprecated, better alternatives might be invented, etc., so that applications can choose appropriate mechanisms that meet their security requirements.

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DECADE servers MAY perform name-object binding validation on stored objects, but Application Endpoints MUST NOT rely on that. In other words, Application Endpoints SHOULD perform name-object binding validation on received objects.

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Appendix A. Evaluation of Candidate Protocols for DECADE DRP/SDT

In this section we evaluate how well the abstract protocol interactions specified in this document for DECADE DRP and SDT can be fulfilled by the existing protocols of HTTP, CDMI, and OAuth.

A.1. HTTP

HTTP [RFC2616] is a key protocol for the Internet in general and especially for the World Wide Web. HTTP is a request-response protocol. A typical transaction involves a client (e.g., web browser) requesting content (resources) from a web server. Another example is when a client stores or deletes content from a server.

A.1.1. HTTP Support for DRP Primitives

DRP provides configuration of access control and resource-sharing policies on DECADE servers.

A.1.1.1. Access Control Primitives

Access control requires mechanisms for defining the access policies for the server and then checking the authorization of a user before it stores or retrieves content. HTTP supports a rudimentary access control via "HTTP Secure" (HTTPS). HTTPS is a combination of HTTP with SSL/TLS. The main use of HTTPS is to authenticate the server and encrypt all traffic between the client and the server. There is also a mode to support client authentication, though this is less frequently used.

A.1.1.2. Resource Control Primitives for Communication

Communication resources include bandwidth (upload/download) and the number of simultaneously connected clients (connections). HTTP supports bandwidth control indirectly through "persistent" HTTP connections. Persistent HTTP connections allows a client to keep open the underlying TCP connection to the server to allow streaming and pipelining (multiple simultaneous requests for a given client).

HTTP does not have direct support for controlling the communication resources for a given client. However, servers typically perform this function via implementation algorithms.

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A.1.1.3. Resource Control Primitives for Storage

Storage resources include the amount of memory and lifetime of storage. HTTP does not allow direct control of storage at the server endpoint. However, HTTP supports caching at intermediate points such as a web proxy. For this purpose, HTTP defines cache control mechanisms that define how long and in what situations the intermediate point may store and use the content.

A.1.2. HTTP Support for SDT Primitives

SDT is used to write objects and read (download) objects from a DECADE server. The object can be either a self-contained object such as a multimedia file or a chunk from a P2P system.

A.1.2.1. Writing Primitives

Writing involves uploading objects to the server. HTTP supports two methods of writing called PUT and POST. In HTTP, the object is called a resource and is identified by a URI. PUT uploads a resource to a specific location on the server. POST, on the other hand, submits the object to the server, and the server decides whether to update an existing resource or to create a new resource.

For DECADE, the choice of whether to use PUT or POST will be influenced by which entity is responsible for the naming. If the client performs the naming, then PUT is appropriate. If the server performs the naming, then POST should be used (to allow the server to define the URI).

A.1.2.2. Downloading Primitives

Downloading involves fetching of an object from the server. HTTP supports downloading through the GET and HEAD methods. GET fetches a specific resource as identified by the URL. HEAD is similar but only fetches the metadata ("header") associated with the resource, not the resource itself.

A.1.3. Primitives for Removing Duplicate Traffic

To challenge a remote entity for an object, the DECADE server should provide a seed number, which is generated by the server randomly, and ask the remote entity to return a hash calculated from the seed number and the content of the object. The server may also specify the hash function that the remote entity should use. HTTP supports the challenge message through the GET methods. The message type

("challenge"), the seed number, and the hash function name are put in a URL. In the reply, the hash is sent in an Entity Tag (ETag) header.

A.1.4. Other Operations

HTTP supports deleting of content on the server through the DELETE method.

A.1.5. Conclusions

HTTP can provide a rudimentary DRP and SDT for some aspects of DECADE, but it will not be able to satisfy all the DECADE requirements. For example, HTTP does not provide a complete access control mechanism nor does it support storage resource controls at the endpoint server.

It is possible, however, to envision combining HTTP with a custom suite of other protocols to fulfill most of the DECADE requirements for DRP and SDT. For example, Google Storage for Developers is built using HTTP (with extensive proprietary extensions such as custom HTTP headers). Google Storage also uses OAuth [RFC6749] (for access control) in combination with HTTP [GoogleStorageDevGuide]. An example of using OAuth for DRP is given in Appendix A.3.

A.2. CDMI

The Cloud Data Management Interface (CDMI) specification defines a functional interface through which applications can store and manage data objects in a cloud storage environment. The CDMI interface for reading/writing data is based on standard HTTP requests, with CDMI-specific encodings using JavaScript Object Notation (JSON). CDMI is specified by the Storage Networking Industry Association (SNIA) [CDMI].

A.2.1. CDMI Support for DRP Primitives

DRP provides configuration of access control and resource-sharing policies on DECADE servers.

A.2.1.1. Access Control Primitives

Access control includes mechanisms for defining the access policies for the server and then checking the authorization of a user before allowing content storage or retrieval. CDMI defines an Access Control List (ACL) per data object and thus supports access control (read and/or write) at the granularity of data objects. An ACL

contains a set of Access Control Entries (ACEs), where each ACE specifies a principal (i.e., user or group of users) and a set of privileges that are granted to that principal.

CDMI requires that an HTTP authentication mechanism be available for the server to validate the identity of a principal (client). Specifically, CDMI requires that either HTTP Basic Authentication or HTTP Digest Authentication be supported. CDMI recommends that HTTP over TLS (HTTPS) is supported to encrypt the data sent over the network.

A.2.1.2. Resource Control Primitives for Communication

Communication resources include bandwidth (upload/download) and the number of simultaneously connected clients (connections). CDMI supports two key data attributes that provide control over the communication resources to a client: "cdmi_max_throughput" and "cdmi_max_latency". These attributes are defined in the metadata for data objects and indicate the desired bandwidth or delay for transmission of the data object from the cloud server to the client.

A.2.1.3. Resource Control Primitives for Storage

Storage resources include amount of quantity and lifetime of storage. CDMI defines metadata for individual data objects and general storage system configuration that can be used for storage resource control. In particular, CDMI defines the following metadata fields:

- -cdmi_data_redundancy: desired number of copies to be maintained
- -cdmi_geographic_placement: region where object is permitted to be stored
- -cdmi_retention_period: time interval object is to be retained
- -cdmi_retention_autodelete: whether object should be automatically deleted after retention period

A.2.2. CDMI Support for SDT Primitives

SDT is used to write objects and read (download) objects from a DECADE server. The object can be either a self-contained object such as a multimedia file or a chunk from a P2P system.

A.2.2.1. Writing Primitives

Writing involves uploading objects to the server. CDMI supports standard HTTP methods for PUT and POST as described in Appendix A.1.2.1.

A.2.2.2. Downloading Primitives

Downloading involves fetching of an object from the server. CDMI supports the standard HTTP GET method as described in Appendix A.1.2.2.

A.2.3. Other Operations

CDMI supports DELETE as described in Appendix A.1.4. CDMI also supports COPY and MOVE operations.

CDMI supports the concept of containers of data objects to support joint operations on related objects. For example, GET may be done on a single data object or an entire container.

CDMI supports a global naming scheme. Every object stored within a CDMI system will have a globally unique object string identifier (ObjectID) assigned at creation time.

A.2.4. Conclusions

CDMI has a rich array of features that can provide a good base for DRP and SDT for DECADE. An initial analysis finds that the following CDMI features may be useful for DECADE:

- access control
- storage resource control
- communication resource control
- COPY/MOVE operations
- data containers
- naming scheme

A.3. OAuth

As mentioned in Appendix A.1, OAuth [RFC6749] may be used as part of the access and resource control of a DECADE system. In this section, we provide an example of how to configure OAuth requests and responses for DRP.

An OAuth request to access DECADE data objects should include the following fields:

response_type: Value should be set to "token".

client_id: The client_id indicates either the application that is using the DECADE service or the end user who is using the DECADE service from a DECADE storage service provider. DECADE storage service providers should provide the ID distribution and management function.

scope: Data object names that are requested.

An OAuth response should include the following information:

token type: "Bearer"

expires_in: The lifetime in seconds of the access token.

access_token: A token denotes the following information.

service_uri: The server address or URI which is providing the service;

permitted_operations (e.g., read, write) and objects (e.g., names
of data objects that might be read or written);

priority: Value should be set to be either "Urgent", "High",
"Normal" or "Low".

bandwidth: Given to requested operation, a weight value used in a weighted bandwidth sharing scheme, or an integer in number of bits per second;

amount: Data size in number of bytes that might be read or written.

token_signature: The signature of the access token.

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