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## A Description of the ARIA Encryption Algorithm

### Abstract

This document describes the ARIA encryption algorithm. ARIA is a 128-bit block cipher with 128-, 192-, and 256-bit keys. The algorithm consists of a key scheduling part and data randomizing part.

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## 1. Introduction

### 1.1. ARIA Overview

ARIA is a general-purpose block cipher algorithm developed by Korean cryptographers in 2003. It is an iterated block cipher with 128-, 192-, and 256-bit keys and encrypts 128-bit blocks in 12, 14, and 16 rounds, depending on the key size. It is secure and suitable for most software and hardware implementations on 32-bit and 8-bit processors. It was established as a Korean standard block cipher algorithm in 2004 [ARIAKS] and has been widely used in Korea, especially for government-to-public services. It was included in PKCS #11 in 2007 [ARIAPKCS].

## 2. Algorithm Description

The algorithm consists of a key scheduling part and data randomizing part.

### 2.1. Notations

The following notations are used in this document to describe the algorithm.

- $\wedge$  bitwise XOR operation
- $\lll$  left circular rotation
- $\ggg$  right circular rotation
- $||$  concatenation of bit strings
- 0x hexadecimal representation

### 2.2. Key Scheduling Part

Let  $K$  denote a master key of 128, 192, or 256 bits. Given the master key  $K$ , we first define 128-bit values  $K_L$  and  $K_R$  as follows.

$$K_L || K_R = K || 0 \dots 0,$$

where the number of zeros is 128, 64, or 0, depending on the size of  $K$ . That is,  $K_L$  is set to the leftmost 128 bits of  $K$  and  $K_R$  is set to the remaining bits of  $K$  (if any), right-padded with zeros to a 128-bit value. Then, we define four 128-bit values ( $W_0$ ,  $W_1$ ,  $W_2$ , and  $W_3$ ) as the intermediate round values appearing in the encryption of  $K_L || K_R$  by a 3-round, 256-bit Feistel cipher.

$$\begin{aligned} W_0 &= K_L, \\ W_1 &= F_0(W_0, CK_1) \wedge K_R, \\ W_2 &= F_E(W_1, CK_2) \wedge W_0, \\ W_3 &= F_0(W_2, CK_3) \wedge W_1. \end{aligned}$$

Here,  $F_0$  and  $F_E$ , respectively called odd and even round functions, are defined in Section 2.4.1.  $CK_1$ ,  $CK_2$ , and  $CK_3$  are 128-bit constants, taking one of the following values.

```
C1 = 0x517cc1b727220a94fe13abe8fa9a6ee0
C2 = 0x6db14acc9e21c820ff28b1d5ef5de2b0
C3 = 0xdb92371d2126e9700324977504e8c90e
```

These values are obtained from the first 128\*3 bits of the fractional part of  $1/\pi$ , where  $\pi$  is the circle ratio. Now the constants  $CK_1$ ,  $CK_2$ , and  $CK_3$  are defined by the following table.

Key size	$CK_1$	$CK_2$	$CK_3$
128	$C_1$	$C_2$	$C_3$
192	$C_2$	$C_3$	$C_1$
256	$C_3$	$C_1$	$C_2$

For example, if the key size is 192 bits,  $CK_1 = C_2$ ,  $CK_2 = C_3$ , and  $CK_3 = C_1$ .

Once  $W_0$ ,  $W_1$ ,  $W_2$ , and  $W_3$  are determined, we compute encryption round keys  $ek_1$ , ...,  $ek_{17}$  as follows.

```
ek1  = W0 ^ (W1 >>> 19),
ek2  = W1 ^ (W2 >>> 19),
ek3  = W2 ^ (W3 >>> 19),
ek4  = (W0 >>> 19) ^ W3,
ek5  = W0 ^ (W1 >>> 31),
ek6  = W1 ^ (W2 >>> 31),
ek7  = W2 ^ (W3 >>> 31),
ek8  = (W0 >>> 31) ^ W3,
ek9  = W0 ^ (W1 <<< 61),
ek10 = W1 ^ (W2 <<< 61),
ek11 = W2 ^ (W3 <<< 61),
ek12 = (W0 <<< 61) ^ W3,
ek13 = W0 ^ (W1 <<< 31),
ek14 = W1 ^ (W2 <<< 31),
ek15 = W2 ^ (W3 <<< 31),
ek16 = (W0 <<< 31) ^ W3,
ek17 = W0 ^ (W1 <<< 19).
```

The number of rounds depends on the size of the master key as follows.

Key size	Number of Rounds
128	12
192	14
256	16

Due to an extra key addition layer in the last round, 12-, 14-, and 16-round algorithms require 13, 15, and 17 round keys, respectively.

Decryption round keys are derived from the encryption round keys.

$$\begin{aligned} dk_1 &= ek_{n+1}, \\ dk_2 &= A(ek_n), \\ dk_3 &= A(ek_{n-1}), \\ &\vdots \\ dk_n &= A(ek_2), \\ dk_{n+1} &= ek_1. \end{aligned}$$

Here,  $A$  and  $n$  denote the diffusion layer of ARIA and the number of rounds, respectively. The diffusion layer  $A$  is defined in Section 2.4.3.

## 2.3. Data Randomizing Part

The data randomizing part of the ARIA algorithm consists of the encryption and decryption processes. The encryption and decryption processes use functions  $F_0$ ,  $FE$ ,  $A$ ,  $SL_1$ , and  $SL_2$ . These functions are defined in Section 2.4.

### 2.3.1. Encryption Process

#### 2.3.1.1. Encryption for 128-Bit Keys

Let  $P$  be a 128-bit plaintext and  $K$  be a 128-bit master key. Let  $ek_1, \dots, ek_{13}$  be the encryption round keys defined by  $K$ . Then the ciphertext  $C$  is computed by the following algorithm.

$P_1$	$= F_0(P, ek_1);$	// Round 1
$P_2$	$= FE(P_1, ek_2);$	// Round 2
$P_3$	$= F_0(P_2, ek_3);$	// Round 3
$P_4$	$= FE(P_3, ek_4);$	// Round 4
$P_5$	$= F_0(P_4, ek_5);$	// Round 5
$P_6$	$= FE(P_5, ek_6);$	// Round 6
$P_7$	$= F_0(P_6, ek_7);$	// Round 7
$P_8$	$= FE(P_7, ek_8);$	// Round 8
$P_9$	$= F_0(P_8, ek_9);$	// Round 9
$P_{10}$	$= FE(P_9, ek_{10});$	// Round 10
$P_{11}$	$= F_0(P_{10}, ek_{11});$	// Round 11
$C$	$= SL_2(P_{11} \wedge ek_{12}) \wedge ek_{13};$	// Round 12

### 2.3.1.2. Encryption for 192-Bit Keys

Let  $P$  be a 128-bit plaintext and  $K$  be a 192-bit master key. Let  $ek_1, \dots, ek_{15}$  be the encryption round keys defined by  $K$ . Then the ciphertext  $C$  is computed by the following algorithm.

```
P1 = F0(P , ek1 );           // Round 1
P2 = FE(P1 , ek2 );          // Round 2
P3 = F0(P2 , ek3 );           // Round 3
P4 = FE(P3 , ek4 );           // Round 4
P5 = F0(P4 , ek5 );           // Round 5
P6 = FE(P5 , ek6 );           // Round 6
P7 = F0(P6 , ek7 );           // Round 7
P8 = FE(P7 , ek8 );           // Round 8
P9 = F0(P8 , ek9 );           // Round 9
P10 = FE(P9 , ek10);          // Round 10
P11 = F0(P10, ek11);          // Round 11
P12 = FE(P11, ek12);          // Round 12
P13 = F0(P12, ek13);          // Round 13
C   = SL2(P13 ^ ek14) ^ ek15; // Round 14
```

### 2.3.1.3. Encryption for 256-Bit Keys

Let  $P$  be a 128-bit plaintext and  $K$  be a 256-bit master key. Let  $ek_1, \dots, ek_{17}$  be the encryption round keys defined by  $K$ . Then the ciphertext  $C$  is computed by the following algorithm.

```
P1 = F0(P , ek1 );           // Round 1
P2 = FE(P1 , ek2 );          // Round 2
P3 = F0(P2 , ek3 );           // Round 3
P4 = FE(P3 , ek4 );           // Round 4
P5 = F0(P4 , ek5 );           // Round 5
P6 = FE(P5 , ek6 );           // Round 6
P7 = F0(P6 , ek7 );           // Round 7
P8 = FE(P7 , ek8 );           // Round 8
P9 = F0(P8 , ek9 );           // Round 9
P10 = FE(P9 , ek10);          // Round 10
P11 = F0(P10, ek11);          // Round 11
P12 = FE(P11, ek12);          // Round 12
P13 = F0(P12, ek13);          // Round 13
P14 = FE(P13, ek14);          // Round 14
P15 = F0(P14, ek15);          // Round 15
C   = SL2(P15 ^ ek16) ^ ek17; // Round 16
```

### 2.3.2. Decryption Process

The decryption process of ARIA is the same as the encryption process except that encryption round keys are replaced by decryption round keys. For example, encryption round keys  $ek_1, \dots, ek_{13}$  of the 12-round ARIA algorithm are replaced by decryption round keys  $dk_1, \dots, dk_{13}$ , respectively.

## 2.4. Components of ARIA

### 2.4.1. Round Functions

There are two types of round functions for ARIA. One is called an odd round function and is denoted by  $F_0$ . It takes as input a pair  $(D, RK)$  of two 128-bit strings and outputs

$$F_0(D, RK) = A(SL_1(D \wedge RK)).$$

The other is called an even round function and is denoted by  $F_E$ . It takes as input a pair  $(D, RK)$  of two 128-bit strings and outputs

$$F_E(D, RK) = A(SL_2(D \wedge RK)).$$

Functions  $SL_1$  and  $SL_2$ , called substitution layers, are described in Section 2.4.2. Function  $A$ , called a diffusion layer, is described in Section 2.4.3.

### 2.4.2. Substitution Layers

ARIA has two types of substitution layers that alternate between rounds. Type 1 is used in the odd rounds, and type 2 is used in the even rounds.

Type 1 substitution layer  $SL_1$  is an algorithm that takes a 16-byte string  $x_0 || x_1 || \dots || x_{15}$  as input and outputs a 16-byte string  $y_0 || y_1 || \dots || y_{15}$  as follows.

$$\begin{aligned} y_0 &= SB_1(x_0), & y_1 &= SB_2(x_1), & y_2 &= SB_3(x_2), & y_3 &= SB_4(x_3), \\ y_4 &= SB_1(x_4), & y_5 &= SB_2(x_5), & y_6 &= SB_3(x_6), & y_7 &= SB_4(x_7), \\ y_8 &= SB_1(x_8), & y_9 &= SB_2(x_9), & y_{10} &= SB_3(x_{10}), & y_{11} &= SB_4(x_{11}), \\ y_{12} &= SB_1(x_{12}), & y_{13} &= SB_2(x_{13}), & y_{14} &= SB_3(x_{14}), & y_{15} &= SB_4(x_{15}). \end{aligned}$$

Type 2 substitution layer  $SL_2$  is an algorithm that takes a 16-byte string  $x_0 || x_1 || \dots || x_{15}$  as input and outputs a 16-byte string  $y_0 || y_1 || \dots || y_{15}$  as follows.

$y_0 = \text{SB3}(x_0), \quad y_1 = \text{SB4}(x_1), \quad y_2 = \text{SB1}(x_2), \quad y_3 = \text{SB2}(x_3),$   
 $y_4 = \text{SB3}(x_4), \quad y_5 = \text{SB4}(x_5), \quad y_6 = \text{SB1}(x_6), \quad y_7 = \text{SB2}(x_7),$   
 $y_8 = \text{SB3}(x_8), \quad y_9 = \text{SB4}(x_9), \quad y_{10} = \text{SB1}(x_{10}), \quad y_{11} = \text{SB2}(x_{11}),$   
 $y_{12} = \text{SB3}(x_{12}), \quad y_{13} = \text{SB4}(x_{13}), \quad y_{14} = \text{SB1}(x_{14}), \quad y_{15} = \text{SB2}(x_{15}).$

Here, SB1, SB2, SB3, and SB4 are S-boxes that take an 8-bit string as input and output an 8-bit string. These S-boxes are defined by the following look-up tables.

SB1:

	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
00	63	7c	77	7b	f2	6b	6f	c5	30	01	67	2b	fe	d7	ab	76
10	ca	82	c9	7d	fa	59	47	f0	ad	d4	a2	af	9c	a4	72	c0
20	b7	fd	93	26	36	3f	f7	cc	34	a5	e5	f1	71	d8	31	15
30	04	c7	23	c3	18	96	05	9a	07	12	80	e2	eb	27	b2	75
40	09	83	2c	1a	1b	6e	5a	a0	52	3b	d6	b3	29	e3	2f	84
50	53	d1	00	ed	20	fc	b1	5b	6a	cb	be	39	4a	4c	58	cf
60	d0	ef	aa	fb	43	4d	33	85	45	f9	02	7f	50	3c	9f	a8
70	51	a3	40	8f	92	9d	38	f5	bc	b6	da	21	10	ff	f3	d2
80	cd	0c	13	ec	5f	97	44	17	c4	a7	7e	3d	64	5d	19	73
90	60	81	4f	dc	22	2a	90	88	46	ee	b8	14	de	5e	0b	db
a0	e0	32	3a	0a	49	06	24	5c	c2	d3	ac	62	91	95	e4	79
b0	e7	c8	37	6d	8d	d5	4e	a9	6c	56	f4	ea	65	7a	ae	08
c0	ba	78	25	2e	1c	a6	b4	c6	e8	dd	74	1f	4b	bd	8b	8a
d0	70	3e	b5	66	48	03	f6	0e	61	35	57	b9	86	c1	1d	9e
e0	e1	f8	98	11	69	d9	8e	94	9b	1e	87	e9	ce	55	28	df
f0	8c	a1	89	0d	bf	e6	42	68	41	99	2d	0f	b0	54	bb	16

SB2:

	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
00	e2	4e	54	fc	94	c2	4a	cc	62	0d	6a	46	3c	4d	8b	d1
10	5e	fa	64	cb	b4	97	be	2b	bc	77	2e	03	d3	19	59	c1
20	1d	06	41	6b	55	f0	99	69	ea	9c	18	ae	63	df	e7	bb
30	00	73	66	fb	96	4c	85	e4	3a	09	45	aa	0f	ee	10	eb
40	2d	7f	f4	29	ac	cf	ad	91	8d	78	c8	95	f9	2f	ce	cd
50	08	7a	88	38	5c	83	2a	28	47	db	b8	c7	93	a4	12	53
60	ff	87	0e	31	36	21	58	48	01	8e	37	74	32	ca	e9	b1
70	b7	ab	0c	d7	c4	56	42	26	07	98	60	d9	b6	b9	11	40
80	ec	20	8c	bd	a0	c9	84	04	49	23	f1	4f	50	1f	13	dc
90	d8	c0	9e	57	e3	c3	7b	65	3b	02	8f	3e	e8	25	92	e5
a0	15	dd	fd	17	a9	bf	d4	9a	7e	c5	39	67	fe	76	9d	43
b0	a7	e1	d0	f5	68	f2	1b	34	70	05	a3	8a	d5	79	86	a8
c0	30	c6	51	4b	1e	a6	27	f6	35	d2	6e	24	16	82	5f	da
d0	e6	75	a2	ef	2c	b2	1c	9f	5d	6f	80	0a	72	44	9b	6c
e0	90	0b	5b	33	7d	5a	52	f3	61	a1	f7	b0	d6	3f	7c	6d
f0	ed	14	e0	a5	3d	22	b3	f8	89	de	71	1a	af	ba	b5	81

SB3:

	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
00	52	09	6a	d5	30	36	a5	38	bf	40	a3	9e	81	f3	d7	fb
10	7c	e3	39	82	9b	2f	ff	87	34	8e	43	44	c4	de	e9	cb
20	54	7b	94	32	a6	c2	23	3d	ee	4c	95	0b	42	fa	c3	4e
30	08	2e	a1	66	28	d9	24	b2	76	5b	a2	49	6d	8b	d1	25
40	72	f8	f6	64	86	68	98	16	d4	a4	5c	cc	5d	65	b6	92
50	6c	70	48	50	fd	ed	b9	da	5e	15	46	57	a7	8d	9d	84
60	90	d8	ab	00	8c	bc	d3	0a	f7	e4	58	05	b8	b3	45	06
70	d0	2c	1e	8f	ca	3f	0f	02	c1	af	bd	03	01	13	8a	6b
80	3a	91	11	41	4f	67	dc	ea	97	f2	cf	ce	f0	b4	e6	73
90	96	ac	74	22	e7	ad	35	85	e2	f9	37	e8	1c	75	df	6e
a0	47	f1	1a	71	1d	29	c5	89	6f	b7	62	0e	aa	18	be	1b
b0	fc	56	3e	4b	c6	d2	79	20	9a	db	c0	fe	78	cd	5a	f4
c0	1f	dd	a8	33	88	07	c7	31	b1	12	10	59	27	80	ec	5f
d0	60	51	7f	a9	19	b5	4a	0d	2d	e5	7a	9f	93	c9	9c	ef
e0	a0	e0	3b	4d	ae	2a	f5	b0	c8	eb	bb	3c	83	53	99	61
f0	17	2b	04	7e	ba	77	d6	26	e1	69	14	63	55	21	0c	7d

SB4:

	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
00	30	68	99	1b	87	b9	21	78	50	39	db	e1	72	9	62	3c
10	3e	7e	5e	8e	f1	a0	cc	a3	2a	1d	fb	b6	d6	20	c4	8d
20	81	65	f5	89	cb	9d	77	c6	57	43	56	17	d4	40	1a	4d
30	c0	63	6c	e3	b7	c8	64	6a	53	aa	38	98	0c	f4	9b	ed
40	7f	22	76	af	dd	3a	0b	58	67	88	06	c3	35	0d	01	8b
50	8c	c2	e6	5f	02	24	75	93	66	1e	e5	e2	54	d8	10	ce
60	7a	e8	08	2c	12	97	32	ab	b4	27	0a	23	df	ef	ca	d9
70	b8	fa	dc	31	6b	d1	ad	19	49	bd	51	96	ee	e4	a8	41
80	da	ff	cd	55	86	36	be	61	52	f8	bb	0e	82	48	69	9a
90	e0	47	9e	5c	04	4b	34	15	79	26	a7	de	29	ae	92	d7
a0	84	e9	d2	ba	5d	f3	c5	b0	bf	a4	3b	71	44	46	2b	fc
b0	eb	6f	d5	f6	14	fe	7c	70	5a	7d	fd	2f	18	83	16	a5
c0	91	1f	05	95	74	a9	c1	5b	4a	85	6d	13	07	4f	4e	45
d0	b2	0f	c9	1c	a6	bc	ec	73	90	7b	cf	59	8f	a1	f9	2d
e0	f2	b1	00	94	37	9f	d0	2e	9c	6e	28	3f	80	f0	3d	d3
f0	25	8a	b5	e7	42	b3	c7	ea	f7	4c	11	33	03	a2	ac	60

For example,  $SB1(0x23) = 0x26$  and  $SB4(0xef) = 0xd3$ . Note that SB3 and SB4 are the inverse functions of SB1 and SB2, respectively, and accordingly SL2 is the inverse of SL1.

#### 2.4.3. Diffusion Layer

Diffusion layer A is an algorithm that takes a 16-byte string  $x0 || x1 || \dots || x15$  as input and outputs a 16-byte string  $y0 || y1 || \dots || y15$  by the following equations.



```

y0 = x3 ^ x4 ^ x6 ^ x8 ^ x9 ^ x13 ^ x14,
y1 = x2 ^ x5 ^ x7 ^ x8 ^ x9 ^ x12 ^ x15,
y2 = x1 ^ x4 ^ x6 ^ x10 ^ x11 ^ x12 ^ x15,
y3 = x0 ^ x5 ^ x7 ^ x10 ^ x11 ^ x13 ^ x14,
y4 = x0 ^ x2 ^ x5 ^ x8 ^ x11 ^ x14 ^ x15,
y5 = x1 ^ x3 ^ x4 ^ x9 ^ x10 ^ x14 ^ x15,
y6 = x0 ^ x2 ^ x7 ^ x9 ^ x10 ^ x12 ^ x13,
y7 = x1 ^ x3 ^ x6 ^ x8 ^ x11 ^ x12 ^ x13,
y8 = x0 ^ x1 ^ x4 ^ x7 ^ x10 ^ x13 ^ x15,
y9 = x0 ^ x1 ^ x5 ^ x6 ^ x11 ^ x12 ^ x14,
y10 = x2 ^ x3 ^ x5 ^ x6 ^ x8 ^ x13 ^ x15,
y11 = x2 ^ x3 ^ x4 ^ x7 ^ x9 ^ x12 ^ x14,
y12 = x1 ^ x2 ^ x6 ^ x7 ^ x9 ^ x11 ^ x12,
y13 = x0 ^ x3 ^ x6 ^ x7 ^ x8 ^ x10 ^ x13,
y14 = x0 ^ x3 ^ x4 ^ x5 ^ x9 ^ x11 ^ x14,
y15 = x1 ^ x2 ^ x4 ^ x5 ^ x8 ^ x10 ^ x15.

```

Note that  $A$  is an involution. That is, for any 16-byte input string  $x$ ,  $x = A(A(x))$  holds.

### 3. Security Considerations

ARIA is designed to be resistant to all known attacks on block ciphers [ARIA03]. Its security was analyzed by the COSIC group of K.U.Leuven in Belgium [ARIAEVAL] and no security flaw has been found.

### 4. Informative References

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## Appendix A. Example Data of ARIA

Here are test data for ARIA in hexadecimal form.

### A.1. 128-Bit Key

- Key : 000102030405060708090a0b0c0d0e0f
- Plaintext : 00112233445566778899aabbccddeeff
- Ciphertext: d718fbd6ab644c739da95f3be6451778
  
- Round key generators
  - W0: 000102030405060708090a0b0c0d0e0f
  - W1: 2afbea741e1746dd55c63ba1afcea0a5
  - W2: 7c8578018bb127e02dfe4e78c288e33c
  - W3: 6785b52b74da46bf181054082763ff6d
  
- Encryption round keys
  - e1: d415a75c794b85c5e0d2a0b3cb793bf6
  - e2: 369c65e4b11777ab713a3e1e6601b8f4
  - e3: 0368d4f13d14497b6529ad7ac809e7d0
  - e4: c644552b549a263fb8d0b50906229eec
  - e5: 5f9c434951f2d2ef342787b1a781794c
  - e6: afea2c0ce71db6de42a47461f4323c54
  - e7: 324286db44ba4db6c44ac306f2a84b2c
  - e8: 7f9fa93574d842b9101a58063771eb7b
  - e9: aab9c57731fcd213ad5677458fcfe6d4
  - e10: 2f4423bb06465abada5694a19eb88459
  - e11: 9f8772808f5d580d810ef8ddac13abeb
  - e12: 8684946a155be77ef810744847e35fad
  - e13: 0f0aa16daee61bd7dfee5a599970fb35
  
- Intermediate round values
  - P1: 7fc7f12befd0a0791de87fa96b469f52
  - P2: ac8de17e49f7c5117618993162b189e9
  - P3: c3e8d59ec2e62d5249ca2741653cb7dd
  - P4: 5d4aebb165e141ff759f669e1e85cc45
  - P5: 7806e469f68874c5004b5f4a046bbcfa
  - P6: 110f93c9a630cdd51f97d2202413345a
  - P7: e054428ef088fef97928241cd3be499e
  - P8: 5734f38ea1ca3ddd102e71f95e1d5f97
  - P9: 4903325be3e500cccd52fba4354a39ae
  - P10: cb8c508e2c4f87880639dc896d25ec9d
  - P11: e7e0d2457ed73d23d481424095afdca0

## A.2. 192-Bit Key

Key : 000102030405060708090a0b0c0d0e0f  
 1011121314151617  
 Plaintext : 00112233445566778899aabbccddeeff  
 Ciphertext: 26449c1805dbe7aa25a468ce263a9e79

## A.3. 256-Bit Key

Key : 000102030405060708090a0b0c0d0e0f  
 101112131415161718191a1b1c1d1e1f  
 Plaintext : 00112233445566778899aabbccddeeff  
 Ciphertext: f92bd7c79fb72e2f2b8f80c1972d24fc

## Appendix B. OIDs

Here is an ASN.1 module conforming to the 2002 version of ASN.1  
 [X.680][X.681][X.682][X.683].

```
AriaModesOfOperation {
iso(1) member-body(2) korea(400) nsri(200046) algorithm (1)
symmetric-encryption-algorithm(1) asn1-module(0) alg-oids(0) }
```

```
DEFINITIONS IMPLICIT TAGS ::=
BEGIN
```

```
OID ::= OBJECT IDENTIFIER
```

```
-- Synonyms --
```

```
id-algorithm OID ::= { iso(1) member-body(2) korea(410) nsri(200046)
algorithm(1)}
```

```
id-sea OID ::= { id-algorithm symmetric-encryption-algorithm(1)}
id-pad OID ::= { id-algorithm pad(2)}
```

```
id-pad-null RELATIVE-OID ::= {0} -- no padding algorithms identified
id-pad-1 RELATIVE-OID ::= {1}
-- padding method 2 of ISO/IEC 9797-1:1999
```

```
-- confidentiality modes:
-- ECB, CBC, CFB, OFB, CTR
```

```
id-aria128-ecb OID ::= { id-sea aria128-ecb(1)}
id-aria128-cbc OID ::= { id-sea aria128-cbc(2)}
id-aria128-cfb OID ::= { id-sea aria128-cfb(3)}
id-aria128-ofb OID ::= { id-sea aria128-ofb(4)}
id-aria128-ctr OID ::= { id-sea aria128-ctr(5)}
```

```
id-aria192-ecb OID ::= { id-sea aria192-ecb(6)}
id-aria192-cbc OID ::= { id-sea aria192-cbc(7)}
id-aria192-cfb OID ::= { id-sea aria192-cfb(8)}
id-aria192-ofb OID ::= { id-sea aria192-ofb(9)}
id-aria192-ctr OID ::= { id-sea aria192-ctr(10)}

id-aria256-ecb OID ::= { id-sea aria256-ecb(11)}
id-aria256-cbc OID ::= { id-sea aria256-cbc(12)}
id-aria256-cfb OID ::= { id-sea aria256-cfb(13)}
id-aria256-ofb OID ::= { id-sea aria256-ofb(14)}
id-aria256-ctr OID ::= { id-sea aria256-ctr(15)}

-- authentication modes: CMAC

id-aria128-cmac OID ::= { id-sea aria128-cmac(21)}
id-aria192-cmac OID ::= { id-sea aria192-cmac(22)}
id-aria256-cmac OID ::= { id-sea aria256-cmac(23)}

-- modes for both confidentiality and authentication
-- OCB 2.0, GCM, CCM, Key Wrap

id-aria128-ocb2 OID ::= { id-sea aria128-ocb2(31)}
id-aria192-ocb2 OID ::= { id-sea aria192-ocb2(32)}
id-aria256-ocb2 OID ::= { id-sea aria256-ocb2(33)}

id-aria128-gcm OID ::= { id-sea aria128-gcm(34)}
id-aria192-gcm OID ::= { id-sea aria192-gcm(35)}
id-aria256-gcm OID ::= { id-sea aria256-gcm(36)}

id-aria128-ccm OID ::= { id-sea aria128-ccm(37)}
id-aria192-ccm OID ::= { id-sea aria192-ccm(38)}
id-aria256-ccm OID ::= { id-sea aria256-ccm(39)}

id-aria128-kw OID ::= { id-sea aria128-kw(40)}
id-aria192-kw OID ::= { id-sea aria192-kw(41)}
id-aria256-kw OID ::= { id-sea aria256-kw(42)}

-- ARIA Key-Wrap with Padding Algorithm (AES version: RFC 5649)

id-aria128-kwp OID ::= { id-sea aria128-kwp(43)}
id-aria192-kwp OID ::= { id-sea aria192-kwp(44)}
id-aria256-kwp OID ::= { id-sea aria256-kwp(45)}
```

```
AriaModeOfOperation ::= AlgorithmIdentifier
{ {AriaModeOfOperationAlgorithms} }
```

```
AriaModeOfOperationAlgorithms ALGORITHM ::= {
aria128ecb   | aria128cbc   | aria128cfb   | aria128ofb   | aria128ctr   |
aria192ecb   | aria192cbc   | aria192cfb   | aria192ofb   | aria192ctr   |
aria256ecb   | aria256cbc   | aria256cfb   | aria256ofb   | aria256ctr   |
aria128cmac  | aria192cmac  | aria256cmac  |               |               |
aria128ocb2  | aria192ocb2  | aria256ocb2  |               |               |
aria128gcm   | aria192gcm   | aria256gcm   |               |               |
aria128ccm   | aria192ccm   | aria256ccm   |               |               |
aria128kw    | aria192kw    | aria256kw    |               |               |
aria128kwp   | aria192kwp   | aria256kwp   |               |               |
... --Extensible
}
```

```
aria128ecb ALGORITHM ::=
{ OID id-aria128-ecb PARAMS AriaEcbParameters }
aria128cbc ALGORITHM ::=
{ OID id-aria128-cbc PARAMS AriaCbcParameters }
aria128cfb ALGORITHM ::=
{ OID id-aria128-cfb PARAMS AriaCfbParameters }
aria128ofb ALGORITHM ::=
{ OID id-aria128-ofb PARAMS AriaOfbParameters }
aria128ctr ALGORITHM ::=
{ OID id-aria128-ctr PARAMS AriaCtrParameters }
```

```
aria192ecb ALGORITHM ::=
{ OID id-aria192-ecb PARAMS AriaEcbParameters }
aria192cbc ALGORITHM ::=
{ OID id-aria192-cbc PARAMS AriaCbcParameters }
aria192cfb ALGORITHM ::=
{ OID id-aria192-cfb PARAMS AriaCfbParameters }
```

```
aria192ofb ALGORITHM ::=
{ OID id-aria192-ofb PARAMS AriaOfbParameters }
aria192ctr ALGORITHM ::=
{ OID id-aria192-ctr PARAMS AriaCtrParameters }
```

```
aria256ecb ALGORITHM ::=
{ OID id-aria256-ecb PARAMS AriaEcbParameters }
aria256cbc ALGORITHM ::=
{ OID id-aria256-cbc PARAMS AriaCbcParameters }
aria256cfb ALGORITHM ::=
{ OID id-aria256-cfb PARAMS AriaCfbParameters }
aria256ofb ALGORITHM ::=
{ OID id-aria256-ofb PARAMS AriaOfbParameters }
aria256ctr ALGORITHM ::=
{ OID id-aria256-ctr PARAMS AriaCtrParameters }

aria128cmac ALGORITHM ::=
{ OID id-aria128-cmac PARAMS AriaCmacParameters }
aria192cmac ALGORITHM ::=
{ OID id-aria192-cmac PARAMS AriaCmacParameters }
aria256cmac ALGORITHM ::=
{ OID id-aria256-cmac PARAMS AriaCmacParameters }

aria128ocb2 ALGORITHM ::=
{ OID id-aria128-ocb2 PARAMS AriaOcb2Parameters }
aria192ocb2 ALGORITHM ::=
{ OID id-aria192-ocb2 PARAMS AriaOcb2Parameters }
aria256ocb2 ALGORITHM ::=
{ OID id-aria256-ocb2 PARAMS AriaOcb2Parameters }

aria128gcm ALGORITHM ::=
{ OID id-aria128-gcm PARAMS AriaGcmParameters }
aria192gcm ALGORITHM ::=
{ OID id-aria192-gcm PARAMS AriaGcmParameters }
aria256gcm ALGORITHM ::=
{ OID id-aria256-gcm PARAMS AriaGcmParameters }

aria128ccm ALGORITHM ::=
{ OID id-aria128-ccm PARAMS AriaCcmParameters }
aria192ccm ALGORITHM ::=
{ OID id-aria192-ccm PARAMS AriaCcmParameters }
aria256ccm ALGORITHM ::=
{ OID id-aria256-ccm PARAMS AriaCcmParameters }

aria128kw ALGORITHM ::= { OID id-aria128-kw }
aria192kw ALGORITHM ::= { OID id-aria192-kw }
aria256kw ALGORITHM ::= { OID id-aria256-kw }

aria128kwp ALGORITHM ::= { OID id-aria128-kwp }
aria192kwp ALGORITHM ::= { OID id-aria192-kwp }
aria256kwp ALGORITHM ::= { OID id-aria256-kwp }
```

```
AriaPadAlgo ::= CHOICE {
    specifiedPadAlgo    RELATIVE-OID,
    generalPadAlgo      OID
}

AriaEcbParameters ::= SEQUENCE {
    padAlgo    AriaPadAlgo    DEFAULT specifiedPadAlgo:id-pad-null
}

AriaCbcParameters ::= SEQUENCE {
    m          INTEGER          DEFAULT 1,
    -- number of stored ciphertext blocks
    padAlgo    AriaPadAlgo    DEFAULT specifiedPadAlgo:id-pad-1
}

AriaCfbParameters ::= SEQUENCE {
    r          INTEGER,
    -- bit-length of feedback buffer, 128<=r<=128*1024
    k          INTEGER,
    -- bit-length of feedback variable, 1<=k<=128
    j          INTEGER,
    -- bit-length of plaintext/ciphertext block, 1<=j<=k
    padAlgo    AriaPadAlgo    DEFAULT specifiedPadAlgo:id-pad-null
}

AriaOfbParameters ::= SEQUENCE {
    j          INTEGER,
    -- bit-length of plaintext/ciphertext block, 1<=j<=128
    padAlgo    AriaPadAlgo    DEFAULT specifiedPadAlgo:id-pad-null
}

AriaCtrParameters ::= SEQUENCE {
    j          INTEGER,
    -- bit-length of plaintext/ciphertext block, 1<=j<=128
    padAlgo    AriaPadAlgo    DEFAULT specifiedPadAlgo:id-pad-null
}

AriaCmacParameters ::= INTEGER -- bit-length of authentication tag

Aria0cb2Parameters ::= INTEGER -- bit-length of authentication tag

AriaGcmParameters ::= SEQUENCE {
    s          INTEGER, -- bit-length of starting variable
    t          INTEGER  -- bit-length of authentication tag
}
```



```
AriaCcmParameters ::= SEQUENCE {  
    w      INTEGER (2|3|4|5|6|7|8),  
    -- length of message length field in octets  
    t      INTEGER (32|48|64|80|96|112|128)  
    -- bit-length of authentication tag  
}  
  
ALGORITHM ::= CLASS {  
    &id      OBJECT IDENTIFIER UNIQUE,  
    &Type     OPTIONAL  
}  
WITH SYNTAX { OID &id  [PARAMS &Type] }  
  
AlgorithmIdentifier { ALGORITHM:AlgoSet } ::= SEQUENCE {  
    algorithm      ALGORITHM.&id( {AlgoSet} ),  
    parameters     ALGORITHM.&Type( {AlgoSet}{@algorithm} ) OPTIONAL  
}  
  
END
```

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