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Using RSA Algorithms with CBOR Object Signing and Encryption (COSE) Messages

#### **Abstract**

The CBOR Object Signing and Encryption (COSE) specification defines cryptographic message encodings using Concise Binary Object Representation (CBOR). This specification defines algorithm encodings and representations enabling RSA algorithms to be used for COSE messages. Encodings are specified for the use of RSA Probabilistic Signature Scheme (RSASSA-PSS) signatures, RSA Encryption Scheme - Optimal Asymmetric Encryption Padding (RSAES-OAEP) encryption, and RSA keys.

#### Status of This Memo

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#### 1. Introduction

The CBOR Object Signing and Encryption (COSE) [RFC8152] specification defines cryptographic message encodings using Concise Binary Object Representation (CBOR) [RFC7049]. This specification defines algorithm encodings and representations enabling RSA algorithms to be used for COSE messages.

## 1.1. Requirements Notation and Conventions

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

## 2. RSASSA-PSS Signature Algorithm

The RSASSA-PSS signature algorithm is defined in [RFC8017].

The RSASSA-PSS signature algorithm is parameterized with a hash function (h), a mask generation function (mgf), and a salt length (sLen). For this specification, the mask generation function is fixed to be MGF1 as defined in [RFC8017]. It has been recommended that the same hash function be used for hashing the data as well as in the mask generation function. This specification follows this recommendation. The salt length is the same length as the hash function output.

Implementations need to check that the key type is 'RSA' when creating or verifying a signature.

The RSASSA-PSS algorithms specified in this document are in the following table.

Name	Value	+   Hash	Salt Length	Description
PS256	-37	SHA-256	32	RSASSA-PSS w/ SHA-256
PS384	-38	SHA-384	48	RSASSA-PSS w/ SHA-384
PS512	-39	SHA-512	64	RSASSA-PSS w/ SHA-512

Table 1: RSASSA-PSS Algorithm Values

## 3. RSAES-OAEP Key Encryption Algorithm

RSAES-OAEP is an asymmetric key encryption algorithm. The definition of RSAEA-OAEP can be found in Section 7.1 of [RFC8017]. The algorithm is parameterized using a mask generation function (mgf), a hash function (h), and encoding parameters (P). For the algorithm identifiers defined in this section:

- o mgf is always set to MGF1 as defined in [RFC8017] and uses the same hash function as h.
- o P is always set to the empty octet string.

The following table summarizes the rest of the values.

Name	Value	Hash	Description
RSAES-OAEP w/ RFC 8017   default parameters	-40	SHA-1	RSAES-OAEP w/ SHA-1
RSAES-OAEP w/ SHA-256	-41	SHA-256	RSAES-OAEP w/ SHA-256
RSAES-OAEP w/ SHA-512	-42	SHA-512	RSAES-OAEP w/ SHA-512

Table 2: RSAES-OAEP Algorithm Values

The key type MUST be 'RSA'.

## 4. RSA Keys

Key types are identified by the 'kty' member of the COSE\_Key object. This specification defines one value for this member in the following table.

Name	Value	Description
i RSA	3	RSA Key

**Table 3: Key Type Values** 

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This document defines a key structure for both the public and private parts of RSA keys. Together, an RSA public key and an RSA private key form an RSA key pair.

The document also provides support for the so-called "multi-prime" RSA keys, in which the modulus may have more than two prime factors. The benefit of multi-prime RSA is lower computational cost for the decryption and signature primitives. For a discussion on how multi-prime affects the security of RSA cryptosystems, the reader is referred to [MultiPrimeRSA].

This document follows the naming convention of [RFC8017] for the naming of the fields of an RSA public or private key, and the corresponding fields have identical semantics. The requirements for fields for RSA keys are as follows:

- o For all keys, 'kty' MUST be present and MUST have a value of 3.
- o For public keys, the fields 'n' and 'e' MUST be present. All other fields defined in the following table below MUST be absent.
- o For private keys with two primes, the fields 'other', 'r\_i',
   'd\_i', and 't\_i' MUST be absent; all other fields MUST be present.
- o For private keys with more than two primes, all fields MUST be present. For the third to nth primes, each of the primes is represented as a map containing the fields 'r\_i', 'd\_i', and 't\_i'. The field 'other' is an array of those maps.
- o All numeric key parameters are encoded in an unsigned big-endian representation as an octet sequence using the CBOR byte string type (major type 2). The octet sequence MUST utilize the minimum number of octets needed to represent the value. For instance, the value 32,768 is represented as the CBOR byte sequence 0b010\_00010, 0x80 0x00 (major type 2, additional information 2 for the length).

The following table provides a summary of the label values and the types associated with each of those labels.

Key Type	Name	Label	CB0R   Type	Description
3   3   3   3   3   3   3   3	n e d p qP dQ qInv other r_i d_i t_i	-1 -2 -3 -4 -5 -6 -7 -8 -9 -10 -11 -12	bstr bstr bstr bstr bstr bstr bstr array bstr bstr	the RSA modulus n the RSA public exponent e the RSA private exponent d the prime factor p of n the prime factor q of n dP is d mod (p - 1) dQ is d mod (q - 1) qInv is the CRT coefficient q^(-1) mod p other prime infos, an array a prime factor r_i of n, where i >= 3 d_i = d mod (r_i - 1) the CRT coefficient t_i = (r_1 * r_2 * * r_(i-1))^(-1) mod r_i

**Table 4: RSA Key Parameters** 

## 5. IANA Considerations

## 5.1. COSE Algorithms Registrations

IANA has registered the following values in the IANA "COSE Algorithms" registry [IANA.COSE].

o Name: PS256 o Value: -37

o Description: RSASSA-PSS w/ SHA-256
o Reference: Section 2 of this document

o Recommended: Yes

o Name: PS384 o Value: -38

o Description: RSASSA-PSS w/ SHA-384
o Reference: Section 2 of this document

o Recommended: Yes

- o Name: PS512 o Value: -39
- o Description: RSASSA-PSS w/ SHA-512
  o Reference: Section 2 of this document
- o Recommended: Yes
- o Name: RSAES-OAEP w/ RFC 8017 default parameters
- o Value: -40
- o Description: RSAES-OAEP w/ SHA-1
- o Reference: Section 3 of this document
- o Recommended: Yes
- o Name: RSAES-OAEP w/ SHA-256
- o Value: -41
- o Description: RSAES-OAEP w/ SHA-256
  o Reference: Section 3 of this document
- o Recommended: Yes
- o Name: RSAES-OAEP w/ SHA-512
- o Value: -42
- o Description: RSAES-OAEP w/ SHA-512
  o Reference: Section 3 of this document
- o Recommended: Yes

### 5.2. COSE Key Type Registrations

IANA has registered the following value in the IANA "COSE Key Types" registry [IANA.COSE].

- o Name: RSA
  o Value: 3
- o Description: RSA Key
- o Reference: Section 4 of this document

#### 5.3. COSE Key Type Parameters Registrations

IANA has registered the following values in the IANA "COSE Key Type Parameters" registry [IANA.COSE].

- o Key Type: 3 o Name: n
- o Label: -1
- o CBOR Type: bstr
- o Description: the RSA modulus n
- o Reference: Section 4 of this document

```
o Key Type: 3
o Name: e
o Label: -2
o CBOR Type: bstr
  Description: the RSA public exponent e
   Reference: Section 4 of this document
o Key Type: 3
o Name: d
  Label: -3
0
o CBOR Type: bstr
o Description: the RSA private exponent d
o Reference: Section 4 of this document
0
  Key Type: 3
  Name: p
Label: -4
0
0
  CBOR Type: bstr
0
  Description: the prime factor p of n
   Reference: Section 4 of this document
o Key Type: 3
o Name: q
o Label: -5
  CBOR Type: bstr
0
  Description: the prime factor q of n
  Reference: Section 4 of this document
0
  Key Type: 3
  Name: dP
0
  Label: -6
0
0
   CBOR Type: bstr
  Description: dP is d mod (p - 1)
   Reference: Section 4 of this document
  Key Type: 3
0
o Name: dQ
o Label: -7
o CBOR Type: bstr
o Description: dQ is d mod (q - 1)
o Reference: Section 4 of this document
o Key Type: 3
o Name: qInv
  Label: -8
0
o CBOR Type: bstr
o Description: qInv is the CRT coefficient q^(-1) mod p
  Reference: Section 4 of this document
```

```
o Key Type: 3
  Name: other
0
  Label: -9
0
   CBOR Type: array
0
   Description: other prime infos, an array
   Reference: Section 4 of this document
  Key Type: 3
0
  Name: r_i
0
  Label: -10
0
   CBOR Type: bstr
0
   Description: a prime factor r_i of n, where i \ge 3
  Reference: Section 4 of this document
  Key Type: 3
0
  Name: d_i
0
   Label: -11
0
   CBOR Type: bstr
0
   Description: d_i = d \mod (r_i - 1)
   Reference: Section 4 of this document
  Key Type: 3
0
o Name: ti
  Label: -12
   CBOR Type: bstr
o Description: the CRT coefficient t_i = (r_1 * r_2 * ... *
  r_(i-1))^(-1) mod r_i
Reference: Section 4 of this document
```

#### 6. Security Considerations

## 6.1. Key Size Security Considerations

A key size of 2048 bits or larger MUST be used with these algorithms. This key size corresponds roughly to the same strength as provided by a 128-bit symmetric encryption algorithm. Implementations SHOULD be able to encrypt and decrypt with modulus between 2048 and 16K bits in length. Applications can impose additional restrictions on the length of the modulus.

In addition to needing to worry about keys that are too small to provide the required security, there are issues with keys that are too large. Denial-of-service attacks have been mounted with overly large keys or oddly sized keys. This has the potential to consume resources with these keys. It is highly recommended that checks on the key length be done before starting a cryptographic operation.

There are two reasonable ways to address this attack. First, a key should not be used for a cryptographic operation until it has been verified that it is controlled by a party trusted by the recipient. This approach means that no cryptography will be done until a trust decision about the key has been made, a process described in Appendix D, Item 4 of [RFC7515]. Second, applications can impose maximum— as well as minimum—length requirements on keys. This limits the resources that would otherwise be consumed by the use of overly large keys.

## 6.2. RSASSA-PSS Security Considerations

There is a theoretical hash substitution attack that can be mounted against RSASSA-PSS [HASHID]. However, the requirement that the same hash function be used consistently for all operations is an effective mitigation against it. Unlike an Elliptic Curve Digital Signature Algorithm (ECDSA), hash function outputs are not truncated so that the full hash value is always signed. The internal padding structure of RSASSA-PSS means that one needs to have multiple collisions between the two hash functions to be successful in producing a forgery based on changing the hash function. This is highly unlikely.

## 6.3. RSAES-OAEP Security Considerations

A version of RSAES-OAEP using the default parameters specified in Appendix A.2.1 of [RFC8017] is included because this is the most widely implemented set of OAEP parameter choices. (Those default parameters are the SHA-1 hash function and the MGF1 with SHA-1 mask generation function.)

Keys used with RSAES-OAEP MUST follow the constraints in Section 7.1 of [RFC8017]. Also, keys with a low private key exponent value, as described in Section 3 of "Twenty Years of Attacks on the RSA Cryptosystem" [Boneh99], MUST NOT be used.

#### 7. References

### 7.1. Normative References

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