

# Deviations of triangle counts in the binomial random graph

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joint work with Matan Harel & Frank Mousset

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## Our setting:

$G_{n,p}$ : the binomial random graph with  
vx set  $\{1, \dots, n\}$  and edge prob.  $p$

$$X := \#\text{triangles in } G_{n,p} \quad \mathbb{E}[X] = \binom{n}{3} p^3$$

Problem. For every  $\delta > 0$ , determine the asymptotics of

$$\underbrace{\log P(X \geq (1+\delta) \mathbb{E}[X])}_{\text{logarithmic upper tail probability}}.$$

Theorem (Chatterjee / DeMarco-Kahn 2012)

If  $p \gg \log n / n$ , then, for every  $\delta > 0$ ,

$$-\log \mathbb{P}(X \geq (1+\delta) \mathbb{E}[X]) = \Theta_{\delta}(\frac{n^2 p^2}{n} \log(\frac{1}{p})).$$

Question. What is the implicit constant?

Definition. For  $\delta > 0$ ,

$$\psi(\delta) := \min \left\{ c(G) : \underbrace{\mathbb{E}[X \mid G \subseteq G_{mp}]}_{\mathbb{E}_G[X]} \geq (1+\delta) \mathbb{E}[X] \right\}.$$

Proposition. If  $\psi(\delta) \rightarrow \infty$ , then

$$\mathbb{P}(X \geq (1+\delta)\mathbb{E}[X]) \geq p^{(1+\delta)^{\binom{n}{2}}\psi(\delta)}.$$

Proof:  $\mathbb{P}_G(\cdot) = \mathbb{P}(\cdot | G \subseteq G_{n,p})$

Pick a small  $\varepsilon > 0$  and a  $G \subseteq K_n$  with  $\psi(\delta + \varepsilon)$  edges s.t.  $\mathbb{E}_G[X] \geq (1+\delta+\varepsilon)\mathbb{E}[X]$ . Note that:

$$(i) \quad \mathbb{P}(UT_\delta) \geq \mathbb{P}(G \subseteq G_{n,p}) \cdot \mathbb{P}_G(UT_\delta) = p^{\psi(\delta+\varepsilon)} \cdot \mathbb{P}_G(UT_\delta).$$

$$(ii) \quad \mathbb{E}_G[X] \leq \mathbb{P}_G(UT_\delta^c) \cdot (1+\delta)\mathbb{E}[X] + \mathbb{P}_G(UT_\delta) \cdot \binom{n}{3}.$$

$$\text{Finally, } (ii) \Rightarrow \mathbb{P}_G(UT_\delta) \geq \varepsilon \mathbb{E}[X] / \binom{n}{3} = \varepsilon p^3 = p^{\circ(\psi(\delta))}. \quad \square$$

Theorem (...) If  $n^{-\alpha} \leq p \ll 1$ , then

$$\mathbb{P}(X \geq (1+\delta)\mathbb{E}[X]) \leq p^{(1-o(1))\psi(\delta)}.$$

- Chatterjee - Dembo (2014)  $\alpha = \frac{1}{42}$
- Eldan (2016)  $\alpha = \frac{1}{18}$
- Cook - Dembo (2018)  $\alpha = \frac{1}{3}$
- Augeri (2018)  $\alpha = \frac{1}{2}$

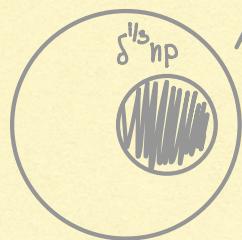
Theorem (Harel, Mousset, S. 2019) If  $n^{-1} \log n \ll p \ll 1$ , then, for every  $\delta > 0$ ,

$$\mathbb{P}(X \geq (1+\delta)\mathbb{E}[X]) \leq p^{(1-\alpha(1))\psi(\delta)}.$$

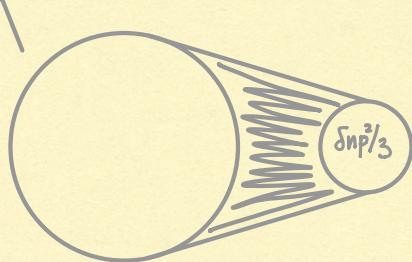
Theorem (Lubetzky, Zhao 2014) For every  $\delta > 0$ ,

$$\frac{\psi(\delta)}{n^2 p^2} \rightarrow \begin{cases} \delta^{2/3}/2 & \text{if } n^{-1} \ll p \ll n^{-1/2}, \\ \min\{\delta^{2/3}/2, \delta/3\} & \text{if } n^{-1/2} \ll p \ll 1. \end{cases}$$

clique



hub



Proof: Fix a small  $\epsilon > 0$  and a large  $C$ .

Definition. A graph  $G \subseteq K_n$  is a seed if

$$(1) \quad e(G) \leq Cn^2 p^2 \log(\frac{1}{p}) \quad \psi(\delta) = \Theta(n^2 p^2)$$

$$(2) \quad \mathbb{E}_G[X] \geq (1 + \delta - \epsilon) \mathbb{E}[X]$$

Lemma 1.  $\mathbb{P}(UT) \leq (1 + o(1)) \cdot \mathbb{P}(G_{n,p} \text{ contains a seed})$ .

Union bound over all seeds? Too naive...

$G_0$ : a seed with  $m := \psi(\delta + \epsilon)$  edges &  $e(G_0) \leq 2m$

$\Rightarrow G_0 \cup G_1$  is a seed

$$\mathbb{E}[\#\text{seeds in } G_{n,p}] \geq \binom{n^2/3}{2m} p^{3m} \geq \left( \frac{n^2 p^{3/2}}{6m} \right)^m = \Omega(p^{-1/2})^m \rightarrow \infty.$$

**Definition.** A graph  $G^*$  is a **core** if

$$(1) \quad e(G^*) \leq Cn^2p^2 \log(1/p)$$

$$(2) \quad \mathbb{E}_{G^*}[X] \geq (1 + \delta - 2\epsilon) \mathbb{E}[X]$$

$$(3) \quad \forall u, v \in G^* \quad \mathbb{E}_{G^*}[X] - \mathbb{E}_{G^* \setminus \{uv\}}[X] \geq \frac{\epsilon \mathbb{E}[X]}{Cn^2p^2 \log(1/p)}$$

**Fact.** Every seed contains a core.

**Proof:** Iteratively remove from a seed all edges failing (3).

**Lemma 2.** For every  $m$ ,

$$\#\text{cores with } m \text{ edges} \leq p^{-\epsilon m}.$$

Assuming both lemmas

$$\begin{aligned} \mathbb{P}(\text{UT}) &\stackrel{\text{L1}}{\leq} \mathbb{P}(G_{n,p} \text{ contains a seed}) \\ &\stackrel{\text{F}}{\leq} \mathbb{P}(G_{n,p} \text{ contains a core}) \\ &\leq \mathbb{E}[\#\text{cores in } G_{n,p}] \\ &\leq \sum_m \mathbb{E}[\#\text{cores with } m \text{ edges in } G_{n,p}] \\ &\stackrel{\text{L2}}{\leq} \sum_{m > \psi(\delta-2\epsilon)} p^{-\epsilon m} \cdot p^m \approx p^{(1-\epsilon)\psi(\delta-2\epsilon)} \end{aligned}$$

$\uparrow$   
 $\psi(\delta-2\epsilon) \leq \#\text{edges in a core}$

Proof of Lemma 1 : (based on Janson-Oleszkiewicz-Ruciński)

Let  $Z := \mathbb{1}[G_{n,p} \text{ does not contain a seed}]$ . As  $Z \in \{0, 1\}$ ,

$$\mathbb{P}(X \geq (1+\delta)\mathbb{E}[X]) \leq \mathbb{P}(Z=0) + \underbrace{\mathbb{P}(XZ \geq (1+\delta)\mathbb{E}[X])}_{\circledast}$$

Enough to show:  $\circledast \ll p^{2\psi(\delta)} \leq \mathbb{P}(X \geq (1+\delta)\mathbb{E}[X])$

Markov's inequality  $\Rightarrow \forall l \geq 1 \quad \circledast \leq \frac{\mathbb{E}[(XZ)^l]}{(1+\delta)^l \cdot \mathbb{E}[X]^l}$

Claim.  $3l \leq Cn^2p^2 \log(1/p) \Rightarrow \mathbb{E}[(XZ)^l] \leq (1+\delta-\varepsilon)^l \mathbb{E}[X]^l$ .

$$\text{Claim. } 3l \leq Cn^2 p^2 \log(1/p) \Rightarrow E[(XZ)^l] \leq (1+\delta-\epsilon)^l E[X]^l.$$

Proof:

$$E[(XZ)^l] = \sum_{T_1, \dots, T_l} P(T_1 \cup \dots \cup T_l \subseteq G_{n,p} \wedge Z=1)$$

$$\leq \sum_{\substack{T_1, \dots, T_l \\ T_1 \cup \dots \cup T_l \text{ not a seed}}} P(T_1 \cup \dots \cup T_l \subseteq G_{n,p})$$

$T_1 \cup \dots \cup T_l$  not a seed

$$c(T_1 \cup \dots \cup T_l) \leq 3l \leq Cn^2 p^2 \log(1/p)$$

$$E_{T_1 \cup \dots \cup T_{l-1}}[X] < (1+\delta-\epsilon) E[X]$$

$$\leq \sum_{\substack{T_1, \dots, T_{l-1} \\ T_1 \cup \dots \cup T_{l-1} \text{ not a seed}}} P(T_1 \cup \dots \cup T_{l-1} \subseteq G_{n,p}) \cdot \underbrace{\sum_{T_l} P(T_l \subseteq G_{n,p} \mid T_1 \cup \dots \cup T_{l-1} \subseteq G_{n,p})}_{E_{T_1 \cup \dots \cup T_{l-1}}[X]}$$

$$\leq \dots \leq ((1+\delta-\epsilon) E[X])^l$$

□

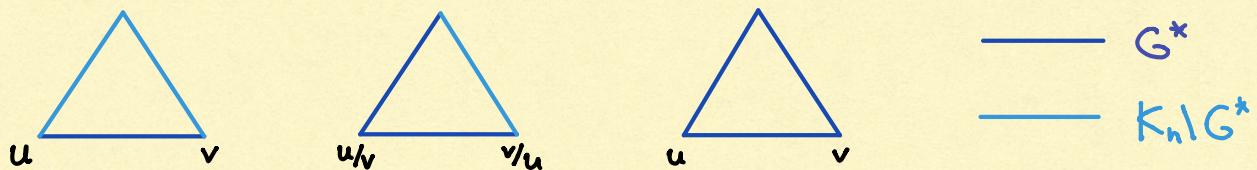
Proof of Lemma 2:

Suppose  $G^*$ : core graph with  $m$  edges

$$G^* \ni uv \mapsto \partial_{uv} := E_{G^*}[X] - E_{G^* \setminus uv}[X]$$

$$\cdot \partial_{uv} \geq \frac{\epsilon E[X]}{Cn^2 p^2 \log(\frac{1}{p})} \geq \frac{\gamma np}{\log(\frac{1}{p})}$$

$$\cdot \partial_{uv} \leq n(p^2 - p^3) + (d_u + d_v)(p - p^2) + d_{u,v}(1-p)$$



$\Rightarrow$  either  $d_u + d_v \geq \gamma'n / \log(\frac{1}{p})$  or  $d_{u,v} \geq \gamma'np / \log(\frac{1}{p})$ .

Cor.  $\forall u \in G^*$   $d_u + d_v \geq \gamma'n / \log(\frac{n}{p})$  or  $d_{u,v} \geq \gamma np / \log(\frac{n}{p})$ .

Define  $A := \{w \in V(G^*) : d_w \geq \gamma'n / (2 \log(\frac{n}{p}))\}$ .

$B := \{w \in V(G^*) : d_w \geq \gamma np / \log(\frac{n}{p})\}$ .

Obs.  $2m = 2e(G^*) = \sum_{w \in V(G^*)} d_w \geq \max\{|A| \cdot \frac{\gamma'n}{2 \log(\frac{n}{p})}, |B| \cdot \frac{\gamma np}{\log(\frac{n}{p})}\}$ .

$$\Rightarrow |A| \leq a_{\max} := \frac{4m \log(\frac{n}{p})}{\gamma'n} \quad \& \quad |B| \leq b_{\max} := \frac{2m \log(\frac{n}{p})}{\gamma np}$$

every edge has 2 endpts in B  
or  $\geq 1$  endpt in A

$$\#A \quad \#B$$

$$\#\text{cores w/ } m \text{ edges} \leq \binom{n}{a_{\max}} \cdot \binom{n}{b_{\max}} \cdot \binom{b_{\max}^2 + n \cdot a_{\max}}{m} \leq \dots \leq p^{-\varepsilon m}$$

□

Problem. Calculate

$$-\log \mathbb{P}(X_H \geq (1+\delta)\mathbb{E}[X_H]),$$

where  $X_H := \#\text{copies of } H \text{ in } G_{n,p}$ .

State-of-the-art:

- $H = K_r$  Harel-Mousset-S.
- $H$ : connected & regular HMS / Basak-Basu
- $H$ : irregular Cook-Dembo