

# Programming Languages: Imperative Program Construction

## Practicals 3. Quantifications

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1. An integer array  $X[0..N)$  is given, where  $N \geq 1$ . Explain, in words, what each of the following expressions mean.

1.  $b \equiv \langle \forall i : 0 \leq i < N : X[i] \geq 0 \rangle$ .
2.  $r = \langle \#k : 0 \leq k < N : \langle \forall i : 0 \leq i < k : X[i] < X[k] \rangle \rangle$ .
3.  $r = \langle \uparrow p, q : 0 \leq p \leq q \leq N \wedge \langle \forall i : p \leq i < q : X[i] > 0 \rangle : q - p \rangle$ .
4.  $r = \langle \#p, q : 0 \leq p < q < N : X[p] = 0 \wedge X[q] = 1 \rangle$ .
5.  $s = \langle \uparrow p, q : 0 \leq p < q < N : X[p] + X[q] \rangle$ .
6.  $b \equiv \langle \forall p, q : 0 \leq p \wedge 0 \leq q \wedge p + q = N - 1 : X[p] = X[q] \rangle$ .

2. An integer array  $X[0..N)$  is given, where  $N \geq 1$ . Express the following sentences in a formal way:

1.  $r$  is the sum of the elements of  $X$ .
2.  $X$  is increasing.
3. all values of  $X$  are distinct.
4.  $r$  is the length of a longest constant segment of  $X$ .
5.  $r$  is the maximum of the sums of the segments of  $X$ .

3. Expand the following textual substitutions. If necessary, change the dummy, according to Dummy Renaming (8.21).

1.  $\langle \star x : 0 \leq x + r < n : x + v \rangle [v \setminus 3]$
2.  $\langle \star x : 0 \leq x + r < n : x + v \rangle [x \setminus 3]$
3.  $\langle \star x : 0 \leq x + r < n : x + v \rangle [n \setminus n + x]$
4.  $\langle \star x : 0 \leq x < r : \langle \star y : 0 \leq y : x + y + n \rangle \rangle [n \setminus x + y]$
5.  $\langle \star x : 0 \leq x < r : \langle \star y : 0 \leq y : x + y + n \rangle \rangle [r \setminus y]$

4. Prove the following theorems. Provided  $0 \leq n$ ,

- (a)  $\langle \Sigma i : 0 \leq i < n + 1 : b[i] \rangle = b[0] + \langle \Sigma i : 1 \leq i < n + 1 : b[i] \rangle$
- (b)  $\langle \exists i : 0 \leq i < n + 1 : b[i] = 0 \rangle = \langle \exists i : 0 \leq i < n : b[i] = 0 \rangle \vee b[n] = 0$

5. Prove that  $\langle \forall x : R : P \rangle \equiv P \vee \langle \forall x : \neg R \rangle$ , provided  $\neg \text{occurs}(x, P)$ .

6. Prove the *range weakening* rule:  $\langle \forall x : Q \vee R : P \rangle \Rightarrow \langle \forall x : Q : P \rangle$ .

7. Prove the *body weakening* rule:  $\langle \forall x : R : P \wedge Q \rangle \Rightarrow \langle \forall x : R : P \rangle$ .