ML-specific optimizations

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What do we have today?

Objectives

- Named records
- Data types
- Pattern matching
- Equality
- Exceptions
- Module system

Tools

- Type system
- Records (tuples)
- Functions
- our invention

Named records

Examples

```
type t = { name: string, number: int }
val t_ins = { name="Sam", number=5 }
type 'a r = { name: string, number: 'a }
```

Named records

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type t = { name: string, number: int }
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Implementation

Named records

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```

Implementation

Simple records

Modules and Functors

```
Example
structure Stack1 =
    struct type 'a stack = 'a list
    fun push(a,s) = a::s
    fun top(a::rest) = a | top(nil) = raise Empty
    fun pop(a::rest) = rest | pop(nil) = raise Empty
    val empty = nil
end
```

Signatures

Example signature STACK = sig type 'a stack val empty : 'a stack val push: 'a * 'a stack -> 'a stack val top : 'a stack -> 'a val pop : 'a stack -> 'a val pop : 'a stack -> 'a

Signatures

```
Example
```

```
signature STACK =
    sig type 'a stack
    val empty : 'a stack
    val push: 'a * 'a stack -> 'a stack
    val top : 'a stack -> 'a
    val pop : 'a stack -> 'a stack
end
```

Example

Functors

Example

```
functor F(S : STACK) = struct
    val em = S.empty
end
structure T = F(Stack2)
```

Implementation

Implementation

- \bullet Modules \rightarrow Records
- ullet Functions o Functions

Linker and runtime system have no idea of module system

Data types

Examples

```
type posint = int (* positive integers *)
datatype money = COIN of posint | BILL of posint
     | CHECK of {amount:real, from: string}
datatype color = RED | BLUE | GREEN | YELLOW
datatype 'a list = nil | :: of 'a * 'a list
datatype register = REG of int
datatype tree = LEAF of int | TREE of tree * tree
datatype xxx = M | N | P of int list
datatype xxxp = \underline{M} \mid \underline{N} \mid \underline{P} of tree
datatype foo = F of int | O of tree
datatype yyy = W of int * int | X of real * real * real
datatype gen = \underline{A} \mid \underline{B} \mid \underline{C} \mid \underline{D} of int \mid \underline{E} of real
     | F of gen * gen | G of int * int * gen
```

8 / 19

Pointers indistinguishable from integers

Tagged

Pointers indistinguishable from integers

Tagged

Transparent

```
datatype register = REG of int
```

Pointers only distinguishable from small integers

Pointers only distinguishable from small integers

Constant

datatype color = RED | BLUE | GREEN | YELLOW

Pointers only distinguishable from small integers

Constant

```
datatype color = RED | BLUE | GREEN | YELLOW
```

Better transparent constructors

TransB

```
datatype xxx = \underline{M} \mid \underline{N} \mid \underline{P} of int list
```

TransU

 $\label{eq:All-pointers} \mbox{ All pointers distinguishable from all integers}$

11 / 19

All pointers distinguishable from all integers

Optimization

datatype tree = <u>LEAF</u> of int | <u>TREE</u> of tree * tree

Pointers to records of different length distinguishable from each other

Pointers to records of different length distinguishable from each other

Optimization

datatype yyy = \underline{W} of int * int | \underline{X} of real * real * real

Problems

Polymorphic datatypes

```
datatype 'a t = A of 'a | B of (real*real)
type u = int t
```

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Polymorphic datatypes

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datatype 'a t = A of 'a | B of (real*real)
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Functors

```
functor F(S: sig type 'a t
datatype 'a list = nil | :: of 'a t
end
) = struct ... end
structure A = struct
datatype 'a list = nil | :: of 'a * 'a list
type 'a t = 'a * 'a list
end
```

structure FA = F(A) (* Where is your runtime system now? *)

Conclusion

Choosing best solution

- Use only assumption 1
- Constructors
 - Tagged
 - Constant
 - Transparent
 - TransB
- Exceptions handled separately
 - Variable
 - VariableC
- Functor mismatch errors at functor-application time

Exceptions

Open type (unbounded number of constructors)

```
Example
Module A
exception C
exception D = J
Module B
exception E of int
```

Exceptions

Open type (unbounded number of constructors)

Example

Module A

exception \underline{C} exception $\underline{D} = J$

Module B

exception \underline{E} of int

Implementation

Tagged type

Pattern matching

Match between value of expression and rule of pattern-expression list

Example

```
case a
  of (false, nil) => nil
  | (true, w) => w
  | (false, x::nil) => x::x::nil
  | (false, y::z) => z
```

Pattern matching

Match between value of expression and rule of pattern-expression list

Example

```
case a
  of (false, nil) => nil
  | (true, w) => w
  | (false, x::nil) => x::x::nil
  | (false, y::z) => z
```

Implementation

Use special instruction: decon. Apply one level of deconstruction per switch.

Equality

Structural equality

```
[1,2,3] = [1,2,3] (* true *)
ref 5 = ref 5 (* false *)
```

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Enter polymorphism

Equality

Structural equality

```
[1,2,3] = [1,2,3] (* true *)
ref 5 = ref 5 (* false *)
```

Enter polymorphism

Implementation

Use runtime tags to compare records field-for-field. Recall that's assumption 3!

Unboxed updates

Example

```
datatype color = Red | Green | Blue (* unboxed *)
val x = ref Red
val _ = x := Green
```

Unboxed updates

Example

```
datatype color = Red | Green | Blue (* unboxed *)
val x = ref Red
val _ = x := Green
```

Implementation

Use types to conservatively mark safe operations to use unboxed instructions.

Our approach and problems

... Look at the code ...