

Higher order quadratures and Gaussian Quadrature

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Recall: Last time, we came up with one way to create quadratures with a higher power of h (a higher order of accuracy) from a standard Newton-Cotes composite quadrature like Trapezoidal. We went through the general method (Richardson Extrapolation), which uses the fact that we have an asymptotic expansion of the error in powers of h . We then applied (and slightly modified) this procedure to the Trapezoidal rule, to find the set of Romberg quadratures; each application of Romberg gains 2 powers of h , so:

- Romberg 1, or $R_1(h)$, is Composite Trapezoidal and has an $O(h^2)$ error.
- Romberg 2, or $R_2(h)$, was later revealed to be Simpson's rule (for $h/2$), and has an $O(h^4)$ error.
- Romberg 3, or $R_3(h)$, was later revealed to be Boole (NC $p=4$), and has an $O(h^6)$ error.
- Romberg k for $k \geq 3$ onwards gives us non interpolation-based rules with $O(h^{2k})$ error.

We discussed two implementations of Romberg: one where we compute nodes and weights based on the rule, and one where we use a "Newton tableau" approach to compute $R_1(h)$ for $h=h_0, h_0/2, \dots, h_0/2^L$, and then use the formulas to compute R_2, R_3, \dots, R_L .

Today, we will discuss a more general approach to produce rules of higher order, given $n+1$ quadrature nodes. This will eventually take us to ask: what is the best set of $n+1$ points and nodes one can pick, and what is the best error we can get out of that? The answer is Gaussian quadrature, and it turns out the best way to compute its nodes and weights involves orthogonal polynomials.

$NC \rightarrow n+1$ quad nodes (interpolation)

$n+1$ EVEN (Trapez) \rightarrow exact for $p \in \mathcal{P}_n$

- $f^{(n+1)}(n)$
- $O(h^{n+2})$ error (one int)
- $O(h^{n+1})$ error (composite)

$n+1$ ODD (Simpson) \rightarrow exact $p \in \mathcal{P}_{n+1}$

- $f^{(n+2)}(n)$
- $O(h^{n+3})$ one int
- $O(h^{n+2})$ composite.

In general, a quad rule w/ $n+1$ nodes

$$Q[f] = \sum_{j=0}^n \underbrace{\omega_j}_{n+1 \text{ weights}} \underbrace{f(x_j)}_{n+1 \text{ nodes}} = \frac{2n+2}{\text{DOFs}}$$

What if I do not constrain x_j 's?

\hookrightarrow Try make exact for $p \in \mathcal{P}_{2n+1}$

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Interval: $[-1, 1]$

$n=1$ (2 pts) $\rightarrow x_0, x_1, w_0, w_1$

$$\left\{ \begin{array}{l} 2 = \int_{-1}^1 1 dx = w_0 + w_1 \\ 0 = \int_{-1}^1 x dx = w_0 x_0 + w_1 x_1 \\ \frac{2}{3} = \int_{-1}^1 x^2 dx = w_0 x_0^2 + w_1 x_1^2 \\ 0 = \int_{-1}^1 x^3 dx = w_0 x_0^3 + w_1 x_1^3 \end{array} \right.$$

$$\left\{ w_0 = w_1 = 1, x_0 = -\frac{1}{\sqrt{3}}, x_1 = \frac{1}{\sqrt{3}} \right\}$$

2 pts, error one int $\sim h^5$
comp $\sim h^4$

NOTE: $n=0$ 1pt $\rightarrow x_0, w_0$

$$[-1, 1] \quad 2 = w_0$$

$$0 = 2x_0 \rightarrow x_0 = 0$$

Error $\sim h^3$, Comp $\sim h^2$.

□ If we set nodes $\{x_j\}_{j=0}^n$, first $n+1$ eqs \rightarrow linear system of eqs.

□ Gen. of quad rules

$$Q[f] = \sum_{j=0}^n \omega_j f(x_j) + \sum_{k=0}^m M_k f'(x_k)$$

$[-1, 1]$,

$$Q[f] = \omega_0 f(x_0) + \omega_1 f'(x_1)$$

$$\left\{ \begin{array}{l} 2 = \int_{-1}^1 1 dx = \omega_0 \\ 0 = \int_{-1}^1 x dx = \omega_0 x_0 + \omega_1 \\ 2/3 = \omega_0 x_0^2 + 2\omega_1 x_1 \\ 0 = \omega_0 x_0^3 + 3\omega_1 x_1^2 \end{array} \right.$$

GAUSS-LEGENDRE ($n+1$ pt)

$\rightarrow [-1, 1], \int_{-1}^1 f(x) dx$

$$n=0 \rightarrow x_0 = 0$$

$$n=1 \rightarrow x_0 = -1/\sqrt{3}, x_1 = 1/\sqrt{3}$$

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$$P_0(x) = 1, \quad P_1(x) = x, \quad P_2(x) = x^2 - \frac{1}{3}, \dots$$

Claim: Gauss-Legendre nodes $\{x_j\}_{j=0}^n$ are the zeroes of $P_{n+1}(x)$.

Let $p \in \mathcal{P}_{2n+1}$.

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Divide p by P_{n+1} :

$$P_{n+1} \overline{) \begin{array}{l} p(x) \\ \dots \\ r(x) \end{array}}$$
$$p(x) = q(x)P_{n+1}(x) + r(x)$$

$$p(x) = \underbrace{q(x)}_{\in \mathcal{P}_n} + \underbrace{p_{n+1}(x)}_{=0} + \underbrace{r(x)}_{\in \mathcal{P}_n}$$

$$\int_{-1}^1 p(x) dx = \underbrace{\int_{-1}^1 \cancel{q(x)} \cancel{p_{n+1}(x)} dx}_{\langle q, p_{n+1} \rangle = 0} + \int_{-1}^1 r(x) dx$$

- Given $\{x_j\}_{j=0}^n \rightarrow w_j$ s.t. they integrate poly up to degree n .

Apply quad rule to p:

$$\underbrace{\sum_{j=0}^n \omega_j p(x_j)}_{D^0} = \underbrace{\sum_{j=0}^n \omega_j q(x_j) \boxed{p_{n+1}(x_j)}}_{D^1} + \underbrace{\sum_{j=0}^n \omega_j v(x_j)}_{D^1}$$

$$\underbrace{\quad} \quad \underbrace{\quad} \quad \underbrace{\quad} \quad \int_{-1}^1 r(x) dx$$

$= 0$
for all $q \in \mathcal{P}_n$.

$$P_{n+1}(x_j) = 0$$

Formula for w_j :

$$w_j = \frac{2}{(1-x_j^2)[P'_{n+1}(x_j)]^2}$$

Gaussian Quadratures ($w(x) \geq 0$)

$$\int_a^b f(x) w(x) dx = \sum_{j=0}^n w_j f(x_j)$$

- $w(x)$ is the density function for a prob. dist. $\rightarrow E[f(x)]$
- $w(x)$ is singular but integrable (e.g. $1/\sqrt{x}$)