

November 11, 2024

Recall: Last class we wrapped up our discussion on Gaussian quadratures for a given interval $[a, b]$ and weight function $w(x)$. The main results for these quadratures are as follows:

- We construct or obtain the **3-term recursion** for the family of orthogonal polynomials $Q_j(x)$.
- The nodes $\{x_j\}_{j=0}^n$ for the $n+1$ point quadrature are the **unique $n+1$ zeroes** of $Q_{n+1}(x)$. There is a theorem that shows Q_{n+1} always has $n+1$ unique zeroes, and that all of them are in $[a, b]$.
- These zeroes can be obtained via **rootfinding** (good first guess using **interlacing + Newton**) or the **Golub-Welch** algorithm.
- The **weights w_j** can be obtained via the integrals of Lagrange polynomials or via the linear system that results from asking the rule be exact for polynomials of degree $\leq n$. However, the most efficient way to compute them is via formulas involving Q_n , Q_{n+1} and Q'_{n+1} (all which can be computed using the 3-term recursion and formulas obtained from it).
- **Weights are positive and small**; they add up to the integral of $w(x)$ from a to b . This makes Gaussian quadratures very stable.
- We can either use one interval and let n go to infinity OR do a composite quadrature for a given $n+1$ point rule. In the first case, order of accuracy goes faster to 0 than any power of $h = (b-a)/N$. In the second case, we get an $O(h^{2n+2})$ composite quadrature (the best for an $n+1$ pt quadrature).

We then introduced the concept of Adaptive Quadrature, what it is and how the algorithm to subdivide an interval adaptively works. We finished our discussion describing a recursive and non-recursive algorithm to perform Adaptive Quadrature.

Today, we start our section on numerical methods for linear algebra by going over a review of important linear algebra concepts.

LOOKING FORWARD

- ① Gauss Elimination - LU decomp
- ② Gram Schmidt - QR decomp
- ③ Eigenvalues, eigenvectors $\begin{matrix} \leftarrow UDU^* \\ \leftarrow U\Sigma V^* \end{matrix}$

KEY CONCEPTS:

- Vector space / subspace - $(V, +, *)$
 \hookrightarrow Scalar field $\mathbb{F} \rightarrow \mathbb{R} / \mathbb{C}$.
- Linear combination $\alpha_j \in \mathbb{F}$
 $\alpha_1 v_1 + \alpha_2 v_2 + \dots + \alpha_n v_n$ $v_j \in V$
- Linear span $(\{v_1, \dots, v_n\})$
 \hookrightarrow Set of all possible (finite) linear comb.
- Linear dependence / Linear independence
 $\alpha_1 v_1 + \alpha_2 v_2 + \dots + \alpha_n v_n = 0$, α_j not all 0.
 \downarrow

$$x_1 v_1 + x_2 v_2 + \dots + x_n v_n = 0, \text{ not all } 0.$$

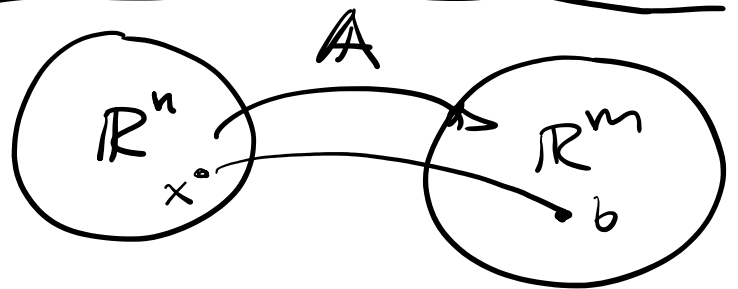
Can I write v_j as a LC of the rest?

• Basis & Dimension.

$\hookrightarrow B \subseteq V$ is a basis of V if $\text{span}(B) = V$ and B is LI set.

$$\rightarrow |B| = \dim V$$

$$\underset{m \times n}{A} \underset{n \times 1}{\vec{x}} = \underset{m \times 1}{\vec{b}}$$



1 $R(A)$, Range / Column space of A

$$R(A) = \{ b \in \mathbb{R}^m \mid \text{There is } x \text{ s.t. } Ax = b \}$$

$$Ax = x_1 v_1 + x_2 v_2 + \dots + x_n v_n \quad A = [v_1 \ v_2 \ \dots \ v_n]$$

$$\dim(R(A)) = r(A) \text{ "rank of } A\text{"}$$

2 $N(A)$ Nullspace / Kernel

$$N(A) = \{ x \in \mathbb{R}^n \mid A\vec{x} = \vec{0} \}$$

$$Ax = x_1 v_1 + x_2 v_2 + \dots + x_n v_n = 0$$

Uniqueness: If $Ax_p = b$, then

$$A(x_p + w) = b \text{ for all } w \in N(A).$$

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$N(A) = \{\vec{0}\}$, then if $Ax = b$ has a sol, it is unique.

$N(A) \neq \{\vec{0}\}$, then if $Ax = b$ has a sol, it is not unique, and is $x_p + N(A)$

$$\dim(N(A)) = \nu(A) \text{ "nullity of } A"$$

(o) Rank & Nullity thm: A $m \times n$,
 $r(A) + \nu(A) = n$ (# of cols)

$Ax = b$, A $m \times n$ matrix.

A .

Case I: Underdetermined ($m < n$)



$$r(A) \leq m \rightarrow \nu(A) > 0 (> n-m)$$

Best case: $r(A) = m$, $\nu(A) = n-m$.

Case II: Overdetermined ($m > n$)



$$r(A) \leq n, \nu(A) = 0 \text{ but } \underline{r(A) < m}.$$

\rightarrow there are b for which $Ax = b$ has no solution.

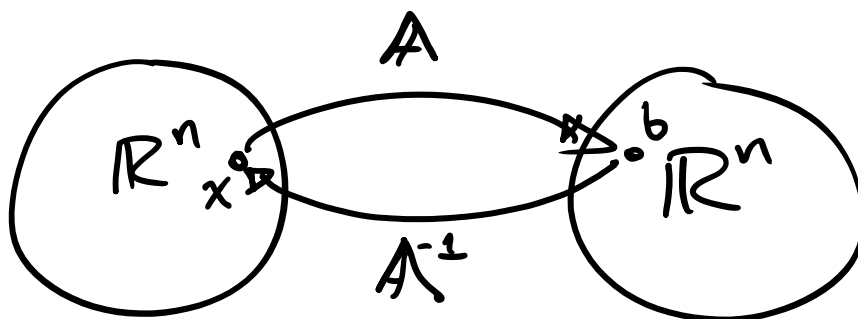
Case III: ($m = n$) (Determined)



Case III: ($m=n$) (Determined) □

\hookrightarrow I can have $v(A) = 0 \iff r(A) = n$
 Existence (for any b) and
 Uniqueness?

$A \rightarrow$ invertible (non-singular)



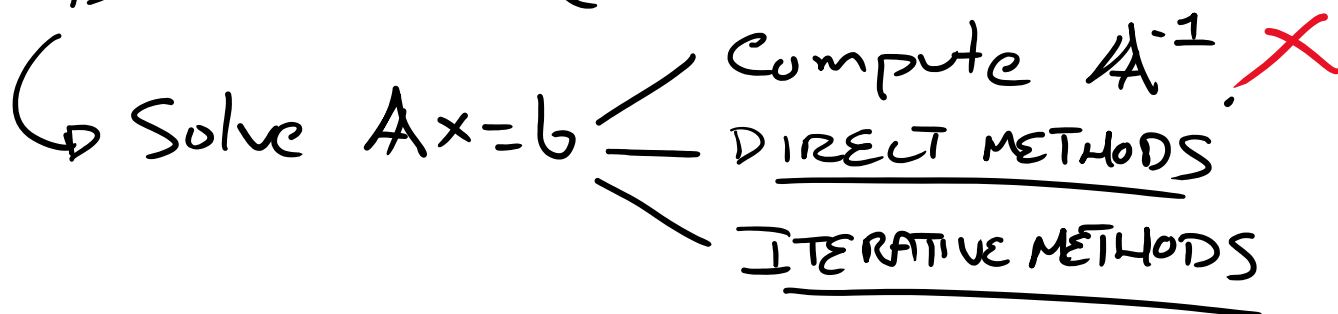
$$| \quad A A^{-1} = A^{-1} A = I$$

Equivalent conditions:

- $Ax=b$ is uniquely solvable for every b .
- A is invertible
- $r(A) = n$ (cols of A are LI)
- $\nu(A) = 0$
- A has no zero eigenvalues
- $\det(A) \neq 0$.

Best case (invertible A)

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DIRECT METHODS - Gauss Elim. (LU)

DO BEST WHEN n is not too big
OR

$\left\{ \begin{array}{l} \text{QR-based} \\ \text{UDU}^* / \text{SVD} \end{array} \right.$

I want to solve many $Ax = b_n$.

- Less dependent on cond #.

Iterative Methods (CG, GMRES, Multigrid, Fixed pt method)

Producing seq of $x_n \xrightarrow{n \rightarrow \infty} x$.

Benefit: Cost p/iter is low
(usually $A \cdot x$) Mat Vec

Cons: Cost = (# iters) (Cost p iter)

A not invertible!

$\left(\min_x \|Ax - b\|^2 \right)$

$\hookrightarrow \min \|Ax - b\|^2$



\hookrightarrow IF NOT
 \rightarrow min norm

