Hacking LLVM for Grad Students

Scott A. Carr Purdue University



ACM Software System Award 2012

- For research, LLVM replaced GCC
- Used for:
 - Just-in-time compilers
 - Secure browser extensions
 - Language virtual machines
 - Static analysis
 - Automatic vectorization
 - GPU programming
 - Software verification
 - Embedded code generators
 - Language implementations

Contents

- LLVM Basics
- Your First Pass
- Your Second Pass
- Your Nth Pass
- Your Runtime Support Library
- Misc. Tips
- Conclusion

LLVM Basics

Read (some of) the docs

- 1. Language Reference
 - Describes IR instructions
- 2. Programmer's Manual
 - How to use LLVM API
- 3. Writing an LLVM Pass
- 4. Community
 - Brandon Holt
 - Adrian Sampson
 - Mailing lists
- 5. Doxygen

LLVM IR's 3 Forms

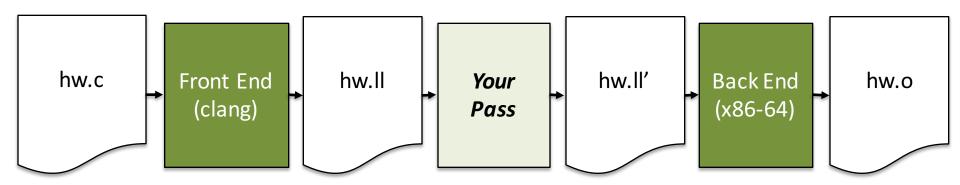
- LLVM is a library for manipulating IR
- Static Single Assignment (SSA)
- Three forms:
 - 1. In-memory
 - 2. On-disk bitcode
 - 3. Human readable assembly

```
%x = add i32 1, %y
```

What Passes Do

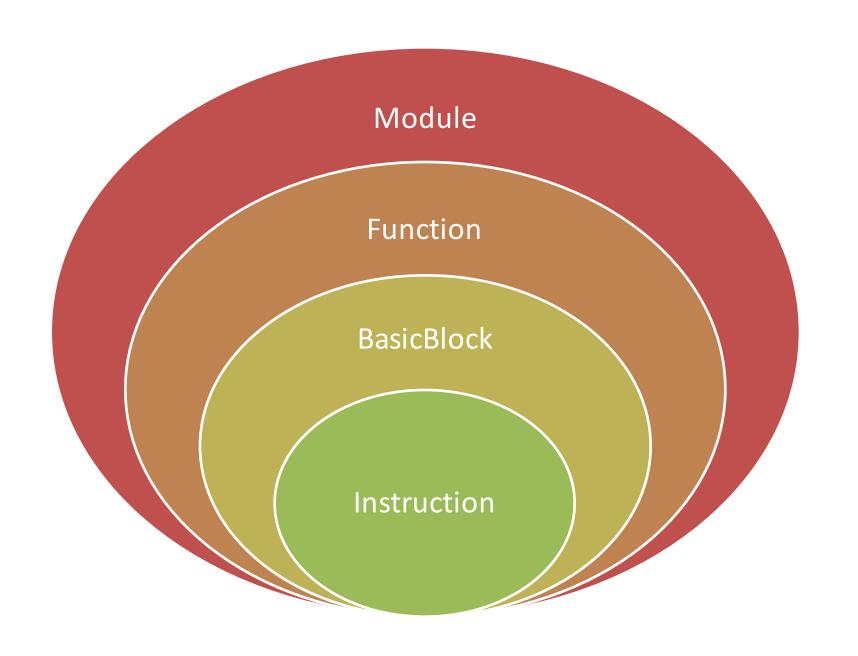
Pass: a class that,

- modifies or analyzes IR and
- is invoked by PassManager



The Important Instructions

LLVM IR	C
LoadInst	x = *y;
StoreInst	*x = y;
GetElementPtrInst	obj.x // obj+offset(x)
CallInst	Foo()
BitCastInst	(void*)ptr

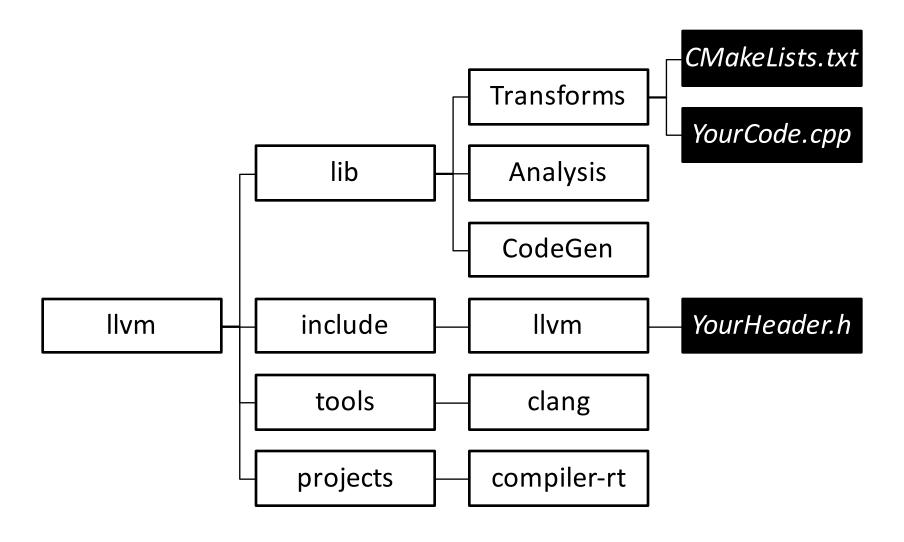


Your First Pass

Your First Pass

- For each function foo
 - Print "Hello: foo"
- From Writing an LLVM Pass

Directory Structure



```
#include "llvm/Pass.h"
#include "llvm/IR/Function.h"
#include "llvm/Support/raw_ostream.h"
using namespace llvm;
namespace {
   struct Hello : public FunctionPass
   {
       static char ID;
       Hello() : FunctionPass(ID) {}
       bool runOnFunction(Function &F) override {
           errs() << "Hello: ";
           errs().write escaped(F.getName()) << '\n';
           return false;
   };
char Hello::ID = 0;
static RegisterPass<Hello> X("hello", "Hello World Pass", false,
   false);
```

```
#include "llvm/Pass.h"
#include "llvm/IR/Function.h"
#include "llvm/Support/raw_ostream.h"
using namespace llvm;
namespace {
   struct Hello : public FunctionPass
   {
       static char ID;
       Hello() : FunctionPass(ID) {}
```

```
char Hello::ID = 0;
static RegisterPass<Hello> X("hello", "Hello World Pass", false,
   false);
```

```
bool runOnFunction(Function &F) override {
   errs() << "Hello: ";
   errs().write escaped(F.getName()) << '\n';
   return false;
```

Building

- Use CMake
- BUILD_SHARED_LIBS=ON
- CMAKE_BUILD_TYPE=Debug
- CMAKE_C_FLAGS=-fstandalone-debug
- Use clang (top tree) as your C/C++ compiler

Running Your Pass

- \$ clang -S -emit-llvm pro.c
 Compile a bitcode version of pro.c, pro.ll
- \$ opt -load Hello.so -hello < pro.ll > pro2.ll
 Run your pass on pro.ll and produce pro2.ll
- \$ clang pro2.11
 Compile pro2.ll to a binary

```
; Module p.ll
define i32 @foo(i32 %x) {
define i32 @bar(i32 %x) {
define i32 @bar(i32 %x) {
```

```
$ opt -load Hello.so -hello < p.ll
Hello: foo
Hello: bar
Hello: baz</pre>
```

Your Second Pass

Your Second Pass

- For each StoreInst,
 - Get its value and pointer operands
 - Call a function, passing these two operands to the function

```
bool runOnFunction(Function& F) {
   for (auto& basicBlock : F) {
      vector<StoreInst> stores;
      for (auto& I : basicBlock) {
         if (auto store = dyn_cast<StoreInst>(&I)) {
             stores.push_back(store);
      for (auto store : stores) {
         // do something
```

bool runOnFunction(Function& F) {

```
for (auto& basicBlock : F) {
```

```
vector<StoreInst> stores;
for (auto& I : basicBlock) {
   if (auto store = dyn_cast<StoreInst>(&I)) {
```

```
if (auto store = dyn_cast<StoreInst>(&I)) {
   stores.push_back(store);
}
```

```
bool runOnFunction(Function& F) {
      for (auto store : stores) {
         // do something
```

Doing Something

```
Create:
  IRBuilder::CreateXXX(...)
  Needs an insertion point
Replace:
  Instruction::replaceAllUsesWith(Instruction *)
  Uses and Users lists
Delete:
  Instruction::eraseFromParent()
```

```
bool runOnFunction(Function& F) {
   for (auto store : stores) {
      // do something
      IRBuilder<> IRB(store);
      auto val = store->getValueOperand();
      auto ptr = store->getPointerOperand();
      auto val2 = IRB.CreateBitCast(val, myValType);
      auto ptr2 = IRB.CreateBitCast(ptr, myPtrType);
      auto fnCall = IRB.CreateCall(myStoreFn, {val2, ptr2});
```

```
IRBuilder<> IRB(store);
auto val = store->getValueOperand();
```

```
auto val = store->getValueOperand();
auto ptr = store->getPointerOperand();
```

```
auto val2 = IRB.CreateBitCast(val, myValType);
auto ptr2 = IRB.CreateBitCast(ptr, myPtrType);
```

```
auto fnCall = IRB.CreateCall(myStoreFn, {val2, ptr2});
```

```
; Module p.ll
...
store %x, %y
```

```
$ opt -load Hello.so -hello < p.ll > p2.ll
$ cat p2.ll
```

```
; Module p2.11
...
%x2 = bitcast i8* %x
%y2 = bitcast i8** %y
call void @myStoreFn(%x2, %y2)
store i32* %x, i32** %y
```

Your Nth Pass

Your Nth Pass

- Use an analysis
 - CFG, CallGraph, AliasAnalysis, etc.
- Get the size of variables with DataLayout
- Examine library functions with TargetLibraryInfo
- Add a command line option
- Use DebugInfo or Metadata

Pass Manager and Analysis

- BLACK MAGIC!
- Copy another pass
- In getAnalysisUsage, AnalysisUsage.addRequired(...)
- INITIALIZE_PASS_DEPENDENCY(...)
- Module Passes can use Function Analysis
 - getAnalysis<FnAnalysis>(llvm::Function *)
- Lots of things are already implemented for you!
- Many analyses are "best effort!"
 - False negatives, not false positives

DataLayout

- Module::getDataLayout()
- Query the size of a type
 - Like sizeof in C
- Get the offset of a struct field
 - For GetElementPtr

TargetLibraryInfo

- getAnalysis<TargetLibraryInfoWrapperPass>().getTLI()
- What library functions are available?
- isMallocLikeFn
- isFreeCall
- isOperatorNewLikeFn

Add a Command Line Option

```
    cl::opt <T> VarName(...);
    -Arguments:

            "flagname"
            cl::desc("my description")
            cl::init(default_val)
```

- Now, just use VarName
- To set flag:

```
-mllvm -flagname <value>
```

See: CommandLine 2.0 Library Manual

DebugInfo

- Defines relationship between source code and generated code (IR)
- Many-to-one mapping
- API just changed
- Get the source line # that generated an IR instruction

Metadata

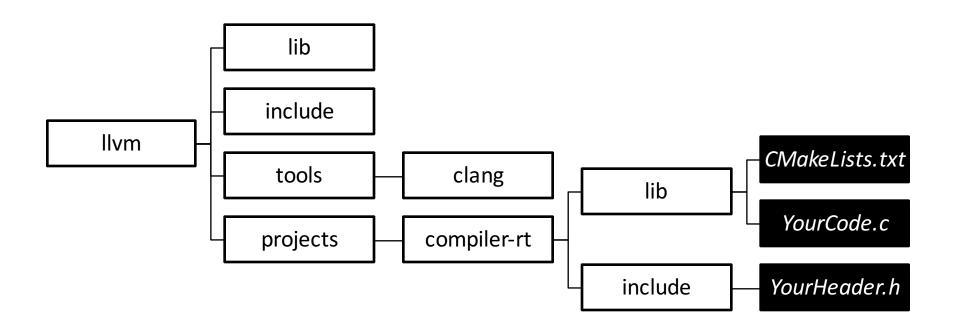
- Information about the IR
- __attribute__((annotate(...))) is metadata
- Not available at runtime
- Ex: Type Based Alias Analysis

Your Runtime Support Library

Your Runtime Support Library

- Add new built-in functions
- Insert calls to your functions with IRBuilder
- Written in C (for clang)

Runtime Directory Structure



Using Your Runtime

```
    Module::getOrInsertFunction(...)
    IRB.CreateCall(
        yourFn,
        {arg1, ... argN});
```

Misc. Tips

Misc. Tips

- Save the IR before and after your pass with raw_fd_ostream
- Eventually, run your pass from within clang/LLVM
 - Find a pass that runs at the same stage and copy it
- Debug with LLDB
- Use LTO for whole program analysis

LLDB

- clang -v to get the command
- 11db -- <the command>
- X->dump() while single stepping
- See Holt's website for more

LTO

- Link Time Optimization
- Why we care: whole program analysis
- Building complex software with LTO requires hacking the build
 - Ex: Chrome

Conclusion

Conclusion

- Use LLVM to "do stuff with code"
- Find a pass that does something similar and tweak it
- Community is one of LLVM's strengths
 - i.e., use Google

Questions?



Links

- Adrian Sampson
 - http://adriansampson.net/blog/llvm.html
- Brandon Holt
 - http://homes.cs.washington.edu/~bholt/posts/llvm-debugging.html
- LLVM Weekly (Alex Bradbury)
 - http://llvmweekly.org/
- Tillman Sheller
 - http://blogs.s-osg.org/an-introduction-to-accelerating-your-build-with-clang/
 - http://blogs.s-osg.org/a-conclusion-to-accelerating-your-build-with-clang/
- Dev Meetings
 - http://llvm.org/devmtg/
- Google for the LLVM Docs ;)
- Me: http://scottandrewcarr.com or @ScottCarr

EXTRA SLIDES

Example: Code Pointer Integrity

- Published in OSDI '14
- Prevents control-flow hijack attacks by ensuring the integrity of function pointers, return addresses, etc.
- Protects a complete OS (FreeBSD)
- Implemented as LLVM pass and runtime

ACM Software System Award 2012

Due to its clean and flexible design and easy to use programming interfaces, LLVM has quickly replaced GCC as the infrastructure of choice for doing research on program translation, optimization, and analysis. Researchers routinely use it for projects as diverse as building link-time interprocedural optimizers, just-in-time compilers, secure browser extensions, language virtual machines, static analysis tools, automatic vectorization, GPU programming, software verification, hardware synthesis tools, embedded code generators, and numerous language implementations.

Example: Data Confidentiality & Integrity

- Provide strong (high overhead) protection for some data and weaker (lower overhead) protection for other data
- Pass
 - Identifies sensitive variables
 - Inserts a bounds check before each load/store
- Runtime
 - Bounds metadata

PHINode

```
%i = icmp eq i32 %x, %y
                br i1 %i, label %BB1, label %BB2
; label %BB1
                                       ; label %BB2
%p1 = add i32 %a, %b
                                      %p2 = add i32 %c, %d
br label %end
                                       br label %end
            ; label %end
           %p = phi i32 [ %p1, %BB1], [%p2, %BB2]
```