

If  $G$  is a connected looped multigraph the *Feynmann loop number*  $L(G)$  is the number of 1-vertices plus the number 2-vertices plus the number of holes (i.e. rank of the fundamental group =  $|E| - |V| + 1 = \chi + 1$ ).

$$L(G) = |V_1| + |V_2| + |E| - |V| + 1$$

Observe that if an edge  $e = \{vw\}$  with is added to  $G$  then

$$L(G \cup \{e\}) = \begin{cases} L(G) - 1 & \text{if } \deg(v) = \deg(w) = 2 \text{ and } v \neq w \\ L(G) & \text{if } v = w \text{ and } \deg(v) \in \{1, 2\} \\ & \text{or if } \deg(v) = 2 \text{ and } \deg(w) \neq 2 \\ L(G) + 1 & \text{else} \end{cases}$$

If  $G$  is a looped multigraph its simplification  $S(G)$  is the simple graph obtained from  $G$  by deleting all loops and collapsing multiedges. Observe that  $G$  can be obtained from  $S(G)$  by adding loops and multiedges and if  $L(G) \leq n$  then  $S(G)$  has at most  $n$  holes and  $2(n - 1)$  vertices as

$$\begin{aligned} n &\geq |V_1| + |V_2| + \frac{1}{2} \sum_k k|V_k| - |V| + 1 \\ &= \frac{1}{2}|V_1| + |V_2| + \sum_{k \geq 3} \left( \frac{k}{2} - 1 \right) |V_k| + 1 \\ &\geq \frac{1}{2}|V| + \frac{1}{2}|V_2| + 1 \\ &\geq \frac{1}{2}|V| + 1 \end{aligned}$$

These bound is tight. If  $n > 2$  then the looped graph with a single vertex and  $n$  loops achieves the bound on the number of holes. The trivalent graph with  $2(n - 1)$  vertices arranged in a cycle with every other edge a 2-multiedge has achieves a loop number  $n$ .

All Feynmann diagrams with Feynmann loop number at most  $n$  can thus be generated by first generating all trees with at most  $2(n - 1)$  vertices then iteratively adding edges between pairs of vertices. Every edge added also adds an additional hole. If looped multigraph  $G$  has more than  $n - |V_1| - |V_2|$  holes it cannot be the subgraph of a Feynmann diagram with Feynmann loop number at most  $n$ . Thus contributing diagrams can be computed by brute force exhaustion via iteratively adding edges.