CPE 323: The MSP430 Assembly Language Programming

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Outline

- Introduction
- Assembly language directives
- SUMI/SUMD
  - Adding two 32-bit numbers (decimal, integers)
- CountEs: Counting characters ‘E’
- Subroutines
  - CALL&RETURN
  - Subroutine Nesting
  - Passing parameters
  - Stack and Local Variables
- Performance

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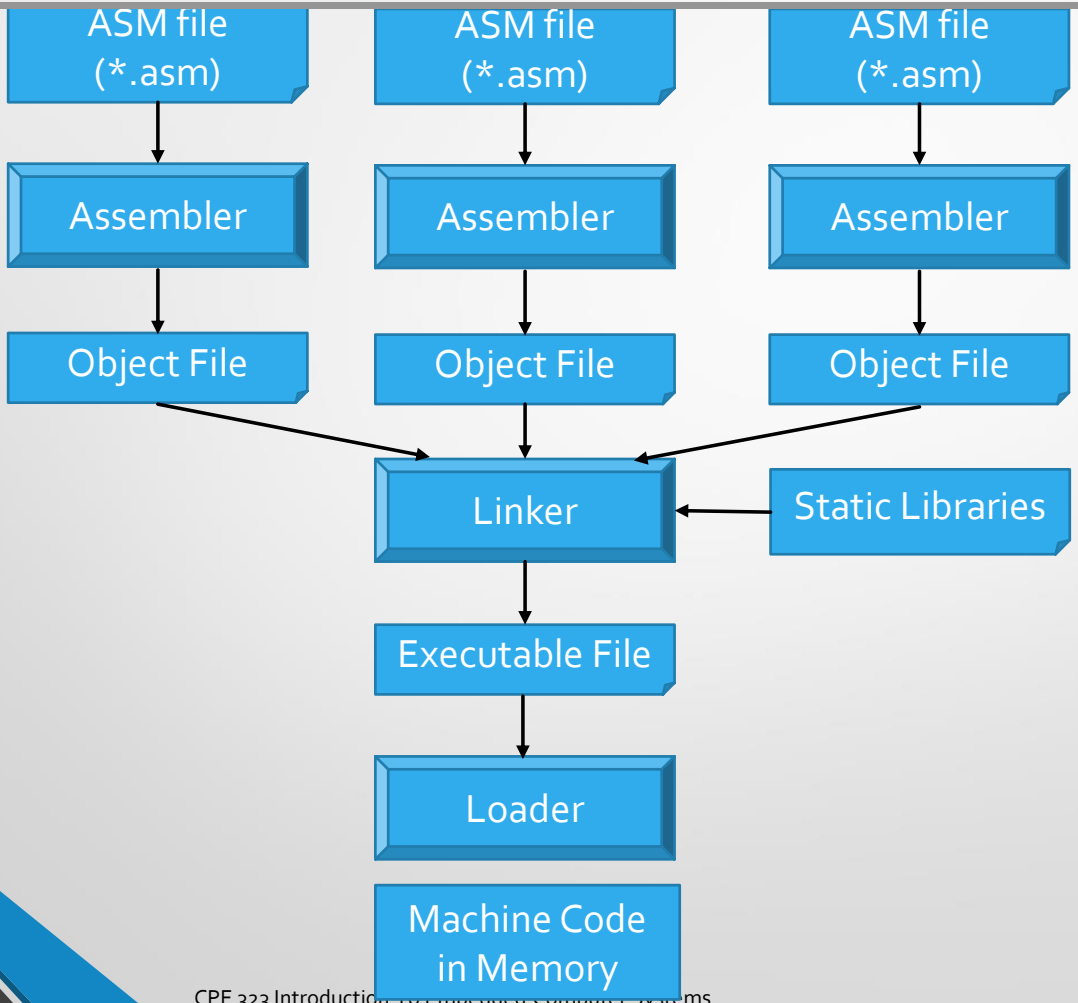
Assembly Programming Flow



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Assembly Directives

- Assembly language directives tell the assembler to
  - Set the data and program at particular addresses in address pace
  - Allocate space for constants and variables
  - Define synonyms
  - Include additional files
  - …
- Typical directives
  - Equate: assign a value to a symbol
  - Origin: set the current location pointer
  - Define space: allocate space in memory
  - Define constant: allocate space for and initialize constants
  - Include: loads another source file

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ASM430 Section Control Directives

- CCStudio ASM430 has three predefined sections into which various parts of a program are assembled
  - .bss: Uninitialized data section
  - .data: Initialized data section
  - .text: Executable code section

Description	ASM430 (CCS)	A430 (IAR)
Reserve size bytes in the uninitialized sect.	.bss	-
Assemble into the initialized data section	.data	RSEG const
Assemble into a named initialized data sect.	.sect	RSEG
Assemble into the executable code	.text	RSEG code
Reserve space in a named (uninitialized) section	.usect	-
Align on byte boundary	.align 1	-
Align on word boundary	.align 2	EVEN

Examples

```
; IAR
        RSEG DAT16_N      ; switch to DATA segment
        EVEN              ; make sure it starts at even address
MyWord: DS 2              ; allocate 2 bytes / 1 word
MyByte: DS 1              ; allocate 1 byte

; CCS Assembler (Example #1)
MyWord: .usect ".bss", 2, 2   ; allocate 1 word
MyByte: .usect ".bss", 1      ; allocate 1 byte

; CCS Assembler (Example #2)
        .bss  MyWord,2,2  ; allocate 1 word
        .bss  MyByte,1    ; allocate 1 byte
```

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Constant Initialization Directives

Description	ASM430 (CCS)	A430 (IAR)
Initialize one or more successive bytes or text strings	.byte or .string	DB
Initialize 32-bit IEEE floating-point	.float	DF
Initialize a variable-length field	.field	-
Reserve size bytes in the current location	.space	DS
Initialize one or more 16-bit integers	.word	DW
Initialize one or more 32-bit integers	.long	DL

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Directives: Dealing with Constants

```
b1:         .byte   5           ; allocates a byte in memory and initialize it with 5
b2:         .byte   -122        ; allocates a byte with constant -122
b3:         .byte   10110111b   ; binary value of a constant
b4:         .byte   0xA0        ; hexadecimal value of a constant
b5:         .byte   123q        ; octal value of a constant
tf:         .equ 25
```

Word view of Memory

Label	Address	Memory[15:8]	Memory[7:0]
b1	0x3100	0x86	0x05
b3	0x3102	0xA0	0xB7
b5	0x3104	--	0x53

Byte view of Memory

Label	Address	Memory[7:0]
b1	0x3100	0x05
b2	0x3101	0x86
b3	0x3102	0xB7
b4	0x3103	0xA0
b5	0x3104	0x53

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Directives: Dealing with Constants

```
...
w1:         .word   21          ; allocates a word constant in memory;

w2:         .word  -21
w3:         .word tf
dw1:        .long  100000       ; allocates a long word size constant in memory;
                                ; 100000 (0x0001_86A0)
dw2:        .long 0xFFFFFFEA
```

Label	Address	Memory[15:8]	Memory[7:0]
w1	0x3106	0x00	0x15
w2	0x3108	0xFF	0xEB
w3	0x310A	0x00	0x19
dw1	0x310C	0x86	0xA0
	0x310E	0x00	0x01
dw2	0x3110	0xFF	0xEA
	0x3112	0xFF	0xFF

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Directives: Dealing with Constants

```
s1:     .byte 'A', 'B', 'C', 'D‘ ; allocates 4 bytes in memory with string ABCD
s2:     .byte "ABCD", ' ' ; allocates 5 bytes in memory with string ABCD + NULL
```

Label	Address	Memory[15:8]	Memory[7:0]
s1	0x3114	0x42	0x41
	0x3116	0x44	0x43
s2	0x3118	0x42	0x41
	0x311A	0x44	0x43
	0x311C	--	0x00
	0x311E		

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Table of Symbols

- Created by the assembler (think about this as a table of synonyms)

Symbol	Value [hex]
b1	0x3100
b2	0x3101
b3	0x3102
b4	0x3103
b5	0x3104
tf	0x0019
w1	0x3106
w2	0x3108
w3	0x310A
dw1	0x310C
dw2	0x3110
s1	0x3114
s2	0x3118

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Directives: Variables in RAM

```
.bss v1b,1,1        ; allocates a byte in memory, equivalent to DS 1
.bss v2b,1,1        ; allocates a byte in memory
.bss v3w,2,2        ; allocates a word of 2 bytes in memory
.bss v4b,8,2        ; allocates a buffer of 2 long words (8 bytes)
.bss vx,1,1
```

Label	Address	Memory[15:8]	Memory[7:0]
v1b	0x1100	--	--
v3w	0x1102	--	--
v4b	0x1104	--	--
	0x1106	--	--
	0x1108	--	--
	0x110A	--	--
vx	0x110C		

Symbol	Value [hex]
v1b	0x1100
v2b	0x1101
v3w	0x1102
v4b	0x1104
vx	0x110C

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Decimal/Integer Addition of 32-bit Numbers

- Problem
  - Write an assembly program that finds a sum of two 32-bit numbers
    - Input numbers are decimal numbers (8-digit in length)
    - Input numbers are signed integers in two's complement
- Data:
  - lint1: `DC32 0x45678923`
  - lint2: `DC32 0x23456789`
  - Decimal sum: `0x69135712`
  - Integer sum: `0x68ac31ac`
- Approach
  - Input numbers: storage, placement in memory
  - Results: storage (ABSOLUTE ASSEMBLER)
  - Main program: initialization, program loops
  - Decimal addition, integer addition

Decimal/Integer Addition of 32-bit Numbers

```
;-------------------------------------------------------------------------------
; File       : LongIntAddition.asm
; Function   : Sums up two long integers represented in binary and BCD
; Description: Program demonstrates addition of two operands lint1 and lint2.
;              Operands are first interpreted as 32-bit decimal numbers and
;              and their sum is stored into lsumd;
;              Next, the operands are interpreted as 32-bit signed integers
;              in two's complement and their sum is stored into lsumi.
; Input      : Input integers are lint1 and lint2 (constants in flash)
; Output     : Results are stored in lsumd (decimal sum) and lsumi (int sum)
; Author     : A. Milenkovic, milenkovic@computer.org
; Date       : August 24, 2018
;-------------------------------------------------------------------------------
            .cdecls C,LIST,"msp430.h"       ; Include device header file

;-------------------------------------------------------------------------------
            .def    RESET                   ; Export program entry-point to
                                            ; make it known to linker.
;-------------------------------------------------------------------------------
            .text                           ; Assemble into program memory.
            .retain                         ; Override ELF conditional linking
                                            ; and retain current section.
            .retainrefs                     ; And retain any sections that have
                                            ; references to current section.
;-------------------------------------------------------------------------------
```

Decimal/Integer Addition of 32-bit Numbers (cont'd)

```
lint1:.long 0x45678923
lint2:.long 0x23456789
;-------------------------------------------------------------------------------
;-------------------------------------------------------------------------------
lsumd:.usect ".bss", 4,2 ; allocate 4 bytes for decimal result
lsumi:.usect ".bss", 4,2 ; allocate 4 bytes for integer result
;-------------------------------------------------------------------------------
RESET:      mov.w   #__STACK_END,SP         ; Initialize stack pointer
StopWDT:    mov.w   #WDTPW|WDTHOLD,&WDTCTL  ; Stop watchdog timer
;-------------------------------------------------------------------------------
```

Decimal/Integer Addition of 32-bit Numbers (cont'd)

```
;-------------------------------------------------------------------------
; Main code here
;-------------------------------------------------------------------------
            clr.w   R2                  ; clear status register
            mov.w   lint1, R8           ; get lower 16 bits from lint1 to R8
            dadd.w  lint2, R8           ; decimal addition, R8 + lower 16-bit of lint2
            mov.w   R8, lsumd           ; store the result (lower 16-bit)
            mov.w   lint1+2, R8         ; get upper 16 bits of lint1 to R8
            dadd.w  lint2+2, R8         ; decimal addition
            mov.w   R8, lsumd+2         ; store the result (upper 16-bit)
            mov.w   lint1, R8           ; get lower 16 bite from lint1 to R8
            add.w   lint2, R8           ; integer addition
            mov.w   R8, lsumi           ; store the result (lower 16 bits)
            mov.w   lint1+2, R8         ; get upper 16 bits from lint1 to R8
            addc.w  lint2+2, R8         ; add upper words, plus carry
            mov.w   R8, lsumi+2         ; store upper 16 bits of the result

           jmp $                        ; jump to current location '$'
                                        ; (endless loop)
```

Decimal/Integer Addition of 32-bit Numbers (cont'd)

```
;-------------------------------------------------------------------------------
; Stack Pointer definition
;-------------------------------------------------------------------------------
            .global __STACK_END
            .sect   .stack


;-------------------------------------------------------------------------------
; Interrupt Vectors
;-------------------------------------------------------------------------------
            .sect   ".reset"                ; MSP430 RESET Vector
            .short  RESET
```

Version 2: Decimal/Integer Addition of 32-bit Numbers (cont'd)

```
; Decimal addition
            mov.w   #lint1, R4          ; pointer to lint1
            mov.w   #lsumd, R8           ; pointer to lsumd
            mov.w   #2, R5               ; R5 is a counter (32-bit=2x16-bit)
            clr.w   R10                  ; clear R10
lda:        mov.w   4(R4), R7            ; load lint2
            mov.w   R10, R2              ; load original SR
            dadd.w  @R4+, R7             ; decimal add lint1 (with carry)
            mov.w   R2, R10              ; backup R2 in R10
            mov.w   R7, 0(R8)            ; store result (@R8+0)
            add.w   #2, R8               ; update R8
            dec.w   R5                   ; decrement R5
            jnz     lda                  ; jump if not zero to lda
```

Version 2: Decimal/Integer Addition of 32-bit Numbers (cont'd)

```
; Integer addition
            mov.w   #lint1, R4          ; pointer to lint1
            mov.w   #lsumi, R8          ; pointer to lsumi
            mov.w   #2, R5              ; R5 is a counter (32-bit=2x16-bit)
            clr.w   R10                 ; clear R10
lia:        mov.w   4(R4), R7           ; load lint2
            mov.w   R10, R2             ; load original SR
            addc.w  @R4+, R7            ; decimal add lint1 (with carry)
            mov.w   R2, R10             ; backup R2 in R10
            mov.w   R7, 0(R8)           ; store result (@R8+0)
            add.w   #2, R8              ; update R8
            dec.w   R5                  ; decrement R5
            jnz     lia                 ; jump if not zero to lia

            jmp     $                   ; jump to current location '$'
                                        ; (endless loop)
```

Count Characters 'E'

- Problem
  - Write an assembly program that processes an input string to find the number of characters 'E' in the string
  - The number of characters is "displayed" on the port 1 of the MSP430
- Example:
  - mystr="HELLO WORLD, I AM THE MSP430!", ''
  - P1OUT=0x02
- Approach
  - Input string: storage, placement in memory
  - Main program: initialization, main program loop
  - Program loop: iterations, counter, loop exit
  - Output: control of ports

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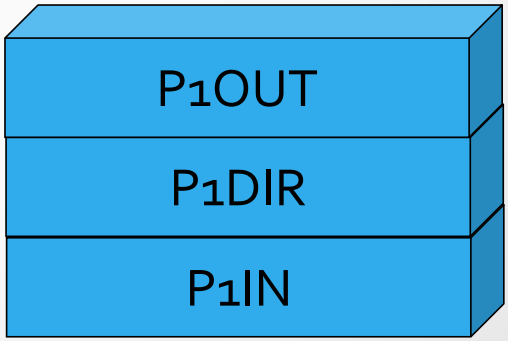
Programmer's View of Parallel Ports

- Parallel ports: x=1,2,3,4,5, …
- Each can be configured as:
  - Input: PxDIR=0x00 (default)
  - Output: PxDIR=0xFF
- Writing to an output port:
  - PxOUT=x02
- Reading from an input port:
  - My_port=P1IN

Port Registers



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Count Characters 'E'

```
;-------------------------------------------------------------------------------
; File       : Lab4_D1.asm (CPE 325 Lab4 Demo code)
; Function   : Counts the number of characters E in a given string
; Description: Program traverses an input array of characters
;              to detect a character 'E'; exits when a NULL is detected
; Input      : The input string is specified in myStr
; Output     : The port P1OUT displays the number of E's in the string
; Author     : A. Milenkovic, milenkovic@computer.org
; Date       : August 14, 2008
;-------------------------------------------------------------------------------
        .cdecls C,LIST,"msp430.h"       ; Include device header file

;-------------------------------------------------------------------------------
        .def    RESET                   ; Export program entry-point to
                                        ; make it known to linker.
myStr:  .string "HELLO WORLD, I AM THE MSP430!", ''
;-------------------------------------------------------------------------------
        .text                           ; Assemble into program memory.
        .retain                         ; Override ELF conditional linking
                                        ; and retain current section.
        .retainrefs                     ; And retain any sections that have
                                        ; references to current section.

;-------------------------------------------------------------------------------
RESET:  mov.w   #__STACK_END,SP         ; Initialize stack pointer
        mov.w   #WDTPW|WDTHOLD,&WDTCTL  ; Stop watchdog timer
```


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Count Characters 'E' (cont'd)

```
;-------------------------------------------------------------------------------
; Main loop here
;-------------------------------------------------------------------------------
main:   bis.b   #0FFh,&P1DIR            ; configure P1.x output
        mov.w   #myStr, R4              ; load the starting address of the string into R4
        clr.b   R5                      ; register R5 will serve as a counter
gnext:  mov.b   @R4+, R6                ; get a new character
        cmp     #0,R6                   ; is it a null character
        jeq     lend                    ; if yes, go to the end
        cmp.b   #'E',R6                 ; is it an 'E' character
        jne     gnext                   ; if not, go to the next
        inc.w   R5                      ; if yes, increment counter
        jmp     gnext                   ; go to the next character

lend:   mov.b   R5,&P1OUT               ; set all P1 pins (output)
        bis.w   #LPM4,SR                ; LPM4
        nop                             ; required only for Debugger

;-------------------------------------------------------------------------------
; Stack Pointer definition
;-------------------------------------------------------------------------------
        .global __STACK_END
        .sect   .stack
;-------------------------------------------------------------------------------
; Interrupt Vectors
;-------------------------------------------------------------------------------
         .sect   ".reset"               ; MSP430 RESET Vector
         .short  RESET
         .end
```

The Case for Subroutines: An Example

- Problem
  - Sum up elements of two integer arrays
  - Display results on P2OUT&P1OUT and P4OUT&P3OUT
- Example
  - arr1 .int 1, 2, 3, 4, 1, 2, 3, 4 ; the first array
  - arr2 .int 1, 1, 1, 1, -1, -1, -1 ; the second array
  - Results
    - P2OUT&P1OUT=0x000A, P4OUT&P3OUT=0x0001
- Approach
  - Input numbers: arrays
  - Main program (no subroutines): initialization, program loops

Sum Up Two Integer Arrays (ver1)

```
;-------------------------------------------------------------------------------
; File       : Lab5_D1.asm (CPE 325 Lab5 Demo code)
; Function   : Finds a sum of two integer arrays
; Description: The program initializes ports,
;              sums up elements of two integer arrays and
;              display sums on parallel ports
; Input      : The input arrays are signed 16-bit integers in arr1 and arr2
; Output     : P1OUT&P2OU displays sum of arr1, P3OUT&P4OUT displays sum of arr2
; Author     : A. Milenkovic, milenkovic@computer.org
; Date       : September 14, 2008
;-------------------------------------------------------------------------------
            .cdecls C,LIST,"msp430.h"       ; Include device header file

;-------------------------------------------------------------------------------
            .def    RESET                   ; Export program entry-point to
                                            ; make it known to linker.
;-------------------------------------------------------------------------------
            .text                           ; Assemble into program memory.
            .retain                         ; Override ELF conditional linking
                                            ; and retain current section.
            .retainrefs                     ; And retain any sections that have
                                            ; references to current section.

;-------------------------------------------------------------------------------
RESET:      mov.w   #__STACK_END,SP         ; Initialize stack pointer
StopWDT:    mov.w   #WDTPW|WDTHOLD,&WDTCTL  ; Stop watchdog timer
```

Sum up two integer arrays (ver1)

```
;-------------------------------------------------------------------------------
; Main code here
;-------------------------------------------------------------------------------
main:       bis.b   #0xFF,&P1DIR            ; configure P1.x as output
            bis.b   #0xFF,&P2DIR            ; configure P2.x as output
            bis.b   #0xFF,&P3DIR            ; configure P3.x as output
            bis.b   #0xFF,&P4DIR            ; configure P4.x as output
            ; load the starting address of the array1 into the register R4
            mov.w   #arr1, R4
            ; load the starting address of the array1 into the register R4
            mov.w   #arr2, R5
            ; Sum arr1 and display
            clr.w   R7                      ; Holds the sum
            mov.w       #8, R10                 ; number of elements in arr1
lnext1:     add.w       @R4+, R7                ; get next element
            dec.w   R10
            jnz     lnext1
            mov.b   R7, P1OUT               ; display sum of arr1
            swpb    R7
            mov.b   R7, P2OUT
```

Sum up two integer arrays (ver1)

```
 ; Sum arr2 and display
            clr.w   R7                  ; Holds the sum
            mov.w   #7, R10             ; number of elements in arr2
lnext2:     add.w   @R5+, R7            ; get next element
            dec.w   R10
            jnz     lnext2
            mov.b   R7, P3OUT           ; display sum of arr1
            swpb    R7
            mov.b   R7, P4OUT
            jmp     $

arr1:        .int     1, 2, 3, 4, 1, 2, 3, 4     ; the first array
arr2:        .int     1, 1, 1, 1, -1, -1, -1     ; the second array


;-------------------------------------------------------------------------------
; Stack Pointer definition
;-------------------------------------------------------------------------------
            .global __STACK_END
            .sect   .stack

;-------------------------------------------------------------------------------
; Interrupt Vectors
;-------------------------------------------------------------------------------
            .sect   ".reset"            ; MSP430 RESET Vector
            .short  RESET
            .end
```

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Subroutines

- A particular sub-task is performed many times on different data values
- Frequently used subtasks are known as subroutines
- Subroutines: How do they work?
  - Only one copy of the instructions that constitute the subroutine is placed in memory
  - Any program that requires the use of the subroutine simply branches to its starting location in memory
  - Upon completion of the task in the subroutine, the execution continues at the next instruction in the calling program

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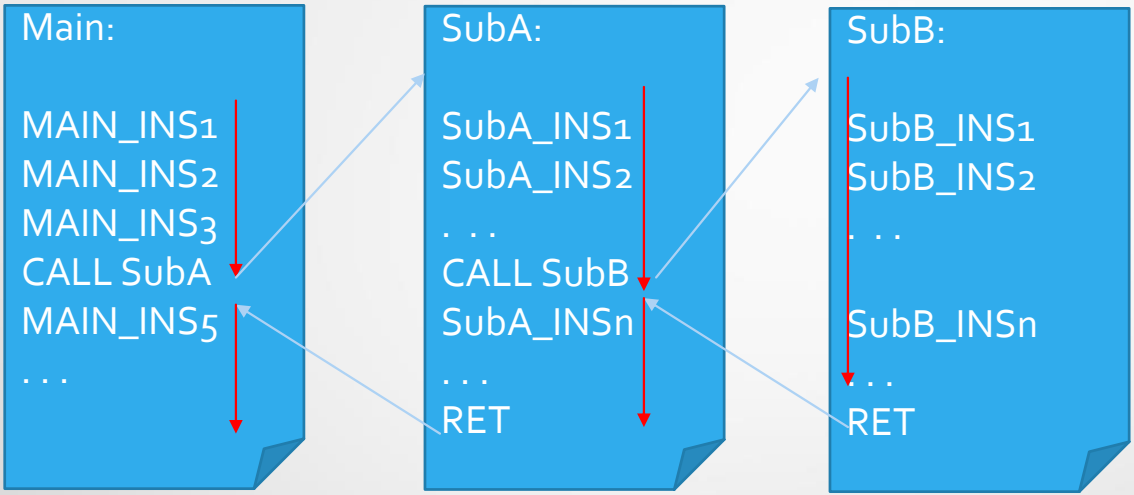
Subroutines (cont'd)

- CALL instruction:
  perform the branch to subroutines
  - SP <= SP – 2 ; allocate a word on the stack for return address
  - M[SP] <= PC ; push the return address (current PC) onto the stack
  - PC <= TargetAddress ; the starting address of the subroutine is moved into PC
- RET instruction:
  the last instruction in the subroutine
  - PC <= M[SP] ; pop the return address from the stack
  - SP <= SP + 2 ; release the stack space

Subroutine Nesting



Mechanisms for Passing Parameters

- Through registers
- Through stack
  - By value
    - Actual parameter is transferred
    - If the parameter is modified by the subroutine, the “new value” does not affect the “old value”
  - By reference
    - The address of the parameter is passed
    - There is only one copy of parameter
    - If parameter is modified, it is modified globally

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Subroutine: SUMA_RP

- Subroutine for summing up elements of an integer array
- Passing parameters through registers
  - `R12 - starting address of the array`
  - `R13 - array length`
  - `R14 - display id`
    `(0 for P2&P1, 1 for P4&P3)`

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Subroutine: SUMA_RP

```
;-------------------------------------------------------------------------------
; File       : Lab5_D2_RP.asm (CPE 325 Lab5 Demo code)
; Function   : Finds a sum of an input integer array
; Description: suma_rp is a subroutine that sums elements of an integer array
; Input      : The input parameters are:
;                   R12 -- array starting address
;                   R13 -- the number of elements (>= 1)
;                   R14 -- display ID (0 for P1&P2 and 1 for P3&P4)
; Output     : No output
; Author     : A. Milenkovic, milenkovic@computer.org
; Date       : September 14, 2008
;-------------------------------------------------------------------------------
            .cdecls C,LIST,"msp430.h"       ; Include device header file

            .def suma_rp

            .text
```


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Subroutine: SUMA_RP

```
suma_rp:
            push.w  R7              ; save the register R7 on the stack
            clr.w   R7              ; clear register R7 (keeps the sum)
lnext:      add.w   @R12+, R7       ; add a new element
            dec.w   R13             ; decrement step counter
            jnz     lnext           ; jump if not finished
            bit.w   #1, R14         ; test display ID
            jnz     lp34            ; jump on lp34 if display ID=1
            mov.b   R7, P1OUT       ; display lower 8-bits of the sum on P1OUT
            swpb    R7              ; swap bytes
            mov.b   R7, P2OUT       ; display upper 8-bits of the sum on P2OUT
            jmp     lend            ; skip to end
lp34:       mov.b   R7, P3OUT       ; display lower 8-bits of the sum on P3OUT
            swpb    R7              ; swap bytes
            mov.b   R7, P4OUT       ; display upper 8-bits of the sum on P4OUT
lend:       pop     R7              ; restore R7
            ret                     ; return from subroutine

            .end
```


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Main (ver2): Call suma_rp

```
;-------------------------------------------------------------------------------
; Main code here
;-------------------------------------------------------------------------------
main:       bis.b   #0xFF,&P1DIR            ; configure P1.x as output
            bis.b   #0xFF,&P2DIR            ; configure P2.x as output
            bis.b   #0xFF,&P3DIR            ; configure P3.x as output
            bis.b   #0xFF,&P4DIR            ; configure P4.x as output

            mov.w   #arr1, R12              ; put address into R12
            mov.w   #8, R13                 ; put array length into R13
            mov.w   #0, R14                 ; display #0 (P1&P2)
            call    #suma_rp

            mov.w   #arr2, R12              ; put address into R12
            mov.w   #7, R13                 ; put array length into R13
            mov.w   #1, R14                 ; display #0 (P3&P4)
            call    #suma_rp
            jmp     $

arr1:        .int     1, 2, 3, 4, 1, 2, 3, 4     ; the first array
arr2:        .int     1, 1, 1, 1, -1, -1, -1     ; the second array
```


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Subroutine: SUMA_SP

- Subroutine for summing up elements of an integer array
- Passing parameters through the stack
  - The calling program prepares input parameters on the stack

Main (ver3): Call suma_sp (Pass Through Stack)

```
;-------------------------------------------------------------------------------
; Main code here
;-------------------------------------------------------------------------------
main:       bis.b   #0xFF,&P1DIR                ; configure P1.x as output
            bis.b   #0xFF,&P2DIR                ; configure P2.x as output
            bis.b   #0xFF,&P3DIR                ; configure P3.x as output
            bis.b   #0xFF,&P4DIR                ; configure P4.x as output

            push    #arr1                       ; push the address of arr1
            push    #8                          ; push the number of elements
            push    #0                          ; push display id
            call    #suma_sp
            add.w   #6,SP                       ; collapse the stack
            push    #arr2                       ; push the address of arr1
            push    #7                          ; push the number of elements
            push    #1                          ; push display id
            call    #suma_sp
            add.w   #6,SP                       ; collapse the stack

            jmp     $

arr1:       .int     1, 2, 3, 4, 1, 2, 3, 4     ; the first array
arr2:       .int     1, 1, 1, 1, -1, -1, -1     ; the second array
```

Address	Stack
0x0800	OTOS
0x07FE	#arr1
0x07FC	0008
0x07FA	0000
0x07F8	Ret. Addr.

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Subroutine: SUMA_SP

```
;-------------------------------------------------------------------------------
; File       : Lab5_D3_SP.asm (CPE 325 Lab5 Demo code)
; Function   : Finds a sum of an input integer array
; Description: suma_sp is a subroutine that sums elements of an integer array
; Input      : The input parameters are on the stack pushed as follows:
;                starting addrress of the array
;                array length
;                display id
; Output     : No output
; Author     : A. Milenkovic, milenkovic@computer.org
; Date       : September 14, 2008
;-------------------------------------------------------------------------------
            .cdecls C,LIST,"msp430.h"       ; Include device header file

            .def    suma_sp

            .text
```


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Subroutine: SUMA_SP (cont'd)

```
suma_sp:
                                        ; save the registers on the stack
            push    R7                  ; save R7, temporal sum
            push    R6                  ; save R6, array length
            push    R4                  ; save R5, pointer to array
            clr.w   R7                  ; clear R7
            mov.w   10(SP), R6          ; retrieve array length
            mov.w   12(SP), R4          ; retrieve starting address
lnext:      add.w   @R4+, R7            ; add next element
            dec.w   R6                  ; decrement array length
            jnz     lnext               ; repeat if not done
            mov.w   8(SP), R4           ; get id from the stack
            bit.w   #1, R4              ; test display id
            jnz     lp34                ; jump to lp34 display id = 1
            mov.b   R7, P1OUT             ; lower 8 bits of the sum to
P1OUT
            swpb    R7                  ; swap bytes
            mov.b   R7, P2OUT           ; upper 8 bits of the sum to P2OUT
            jmp     lend                ; jump to lend
lp34:       mov.b   R7, P3OUT           ; lower 8 bits of ths sum to P3OUT
            swpb    R7                  ; swap bytes
            mov.b   R7, P4OUT           ; upper 8 bits of the sum to P4OUT
lend:       pop     R4                  ; restore R4
            pop     R6                  ; restore R6
            pop     R7                  ; restore R7
            ret                         ; return

            .end
```

Address	Stack
0x0800	OTOS
0x07FE	#arr1
0x07FC	0008
0x07FA	0000
0x07F8	Ret. Addr.
0x07F6	(R7)
0x07F4	(R6)
0x07F2	(R4)

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The Stack and Local Variables

- Subroutines often need local workspace
- We can use a fixed block of memory space – static allocation – but:
  - The code will not be relocatable
  - The code will not be reentrant
  - The code will not be able to be called recursively
- Better solution: dynamic allocation
  - Allocate all local variables on the stack
  - STACK FRAME = a block of memory allocated by a subroutine to be used for local variables
  - FRAME POINTER = an address register used to point to the stack frame

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Subroutine: SUMA_SPSF

```
;-------------------------------------------------------------------------------
; File       : Lab5_D4_SPSF.asm (CPE 325 Lab5 Demo code)
; Function   : Finds a sum of an input integer array
; Description: suma_spsf is a subroutine that sums elements of an integer array.
;              The subroutine allocates local variables on the stack:
;                  counter (SFP+2)
;                  sum (SFP+4)
; Input      : The input parameters are on the stack pushed as follows:
;                  starting address of the array
;                  array length
;                  display id
; Output     : No output
; Author     : A. Milenkovic, milenkovic@computer.org
; Date       : September 14, 2008
;-------------------------------------------------------------------------------
            .cdecls C,LIST,"msp430.h"       ; Include device header file

            .def    suma_spsf

             .text
```


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Subroutine: SUMA_SPSF (cont'd)

```
suma_spsf:
        ; save the registers on the stack
        push    R12             ; save R12 - R12 is stack frame pointer
        mov.w   SP, R12         ; R12 points on the bottom of the stack frame
        sub.w   #4, SP          ; allocate 4 bytes for local variables
        push    R4              ; pointer register
        clr.w   -4(R12)         ; clear sum, sum=0
        mov.w    6(R12), -2(R12) ; get array length
        mov.w   8(R12), R4      ; R4 points to the array starting address
lnext:  add.w   @R4+, -4(R12)   ; add next element
        dec.w   -2(R12)         ; decrement counter
        jnz     lnext           ; repeat if not done
        bit.w   #1, 4(R12)      ; test display id
        jnz     lp34            ; jump to lp34 if display id = 1
        mov.b   -4(R12), P1OUT  ; lower 8 bits of the sum to P1OUT
        mov.b   -3(R12), P2OUT  ; upper 8 bits of the sume to P2OUT
        jmp     lend            ; skip to lend
lp34:   mov.b   -4(R12), P3OUT  ; lower 8 bits of the sum to P3OUT
        mov.b   -3(R12), P4OUT  ; upper 8 bits of the sume to P4OUT
lend:   pop     R4              ; restore R4
        add.w   #4, SP          ; collapse the stack frame
        pop     R12             ; restore stack frame pointer
        ret                     ; return
        .end
```

R12 →

SP →

Address	Stack
0x0800	OTOS
0x07FE	#arr1
0x07FC	0008
0x07FA	0000
0x07F8	Ret. Addr.
0x07F6	(R12)
0x07F4	counter
0x07F2	sum
0x0731	(R4)

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Performance

- Performance: how fast a task can be completed
- Performance(X) = 1/ExecutionTime(X)
- ET: ExecutionTime

$$ET = IC \cdot CPI \cdot CCT = \frac{IC \cdot CPI}{CF}$$

  - IC: Instruction Count – the number of instructions executed in the program
  - CPI: Cycles Per Instruction – the average number of clock cycles it takes to execute an instruction
  - CCT: Clock Cycle Time – the duration of one processor clock cycle
  - CF: Clock Frequency (1/CCT)

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Performance: An Example

```
RESET:      mov.w   #__STACK_END,SP          ; 4cc
StopWDT:    mov.w   #WDTPW|WDTHOLD,&WDTCTL   ; 5cc
            push    R14                      ; 3 cc (table 3.15)
            mov.w   SP, R14                  ; 1 cc
            mov.w   #aend, R6                ; 2 cc
            mov.w   R6, R5                   ; 1 cc
            sub.w   #arr1, R5                ; 2 cc
            sub.w   R5, SP                   ; 1 cc
lnext:      dec.w   R6                       ; 1 cc  x 9
            dec.w   R14                      ; 1 cc  x 9
            mov.b   @R6, 0(R14)              ; 4 cc  x 9
            dec.w   R5                       ; 1 cc x 9
            jnz     lnext                    ; 2 cc x 9
            jmp     $

arr1     .byte   1, 2, 3, 4, 5, 6, 7, 8, 9
aend
         .end
TOTAL NUMBER OF CLOCK CYLES:                 4+5+3+1+2+1+2+1+9x(1+1+4+1+2) = 19+9x9 = 100 cc
TOTAL NUMBER OF INSTRUCITONS                  8+9x5 = 53 instructions
CPI                                          100/53 = 1.88 cc/instruction
```


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