CPE 212 - Fundamentals of Software Engineering

Recursion

Outline

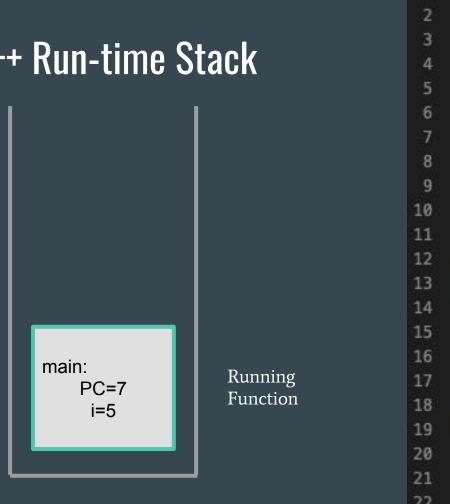
- C++ Memory Allocation
- C++ Runtime Stack Intro
- Recursion Definition
- Example
- Terminology
- Best Practices

C++ Memory Allocation

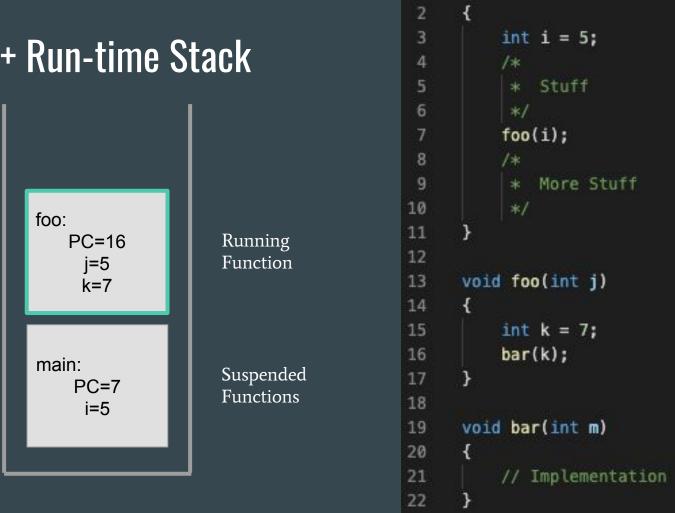
Program
CodeC++
StackFree MemoryMemory
HeapFixed
SizeGrows =><= Grows</td>

Activation Record

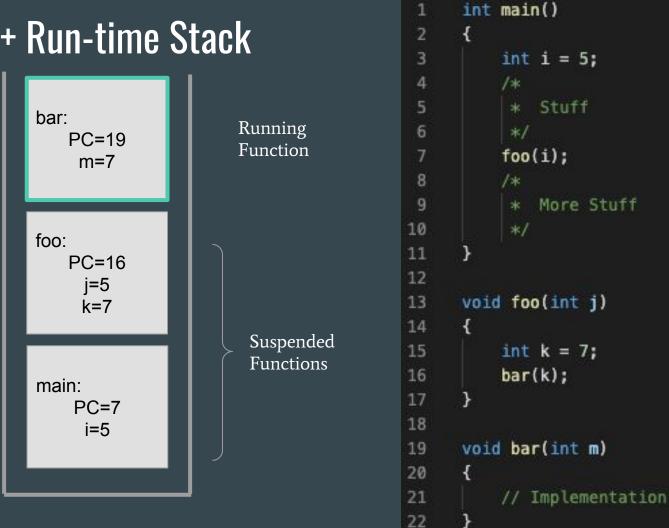
- A record used at run time to store information about a function call, including the parameters, local variables, register values, and return address
- Also called a stack frame
- Gets put on the run-time stack
 - Data structure used to keep track of activation records during the execution of a program

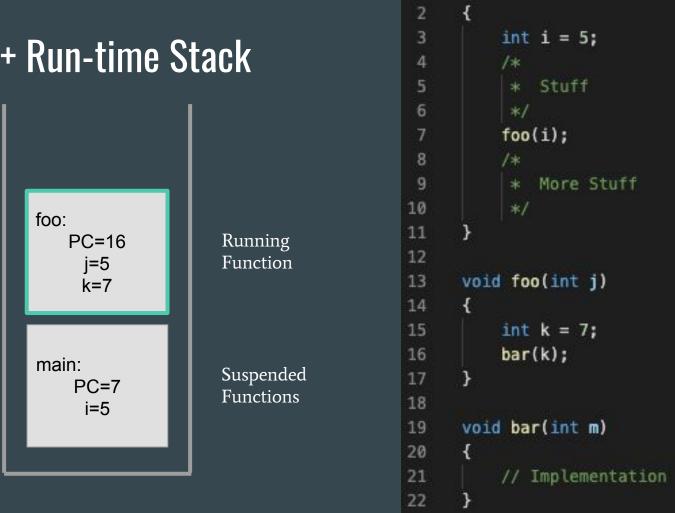


```
int main()
          int i = 5;
          /*
           * Stuff
           */
          foo(i);
          /*
              More Stuff
     void foo(int j)
          int k = 7;
          bar(k);
      void bar(int m)
          // Implementation
22
```

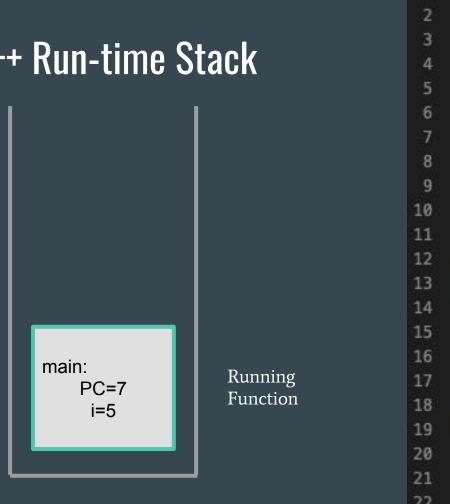


int main()





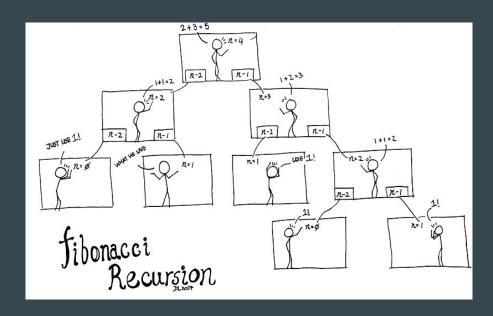
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Recursion

- Recursive Call
 - A function call in which the function being called is the same as the one making the call
- Direct Recursion
 - When a function directly calls itself
- Indirect Recursion
 - When a chain of two or more function calls returns to the function that originated the chain



https://stackoverflow.com/questions/41903017/how-to-visualize-fibonacci-recursion

Recursion Example

```
Factorial of n:
                                                              Fact(n)
n! = 1 \times 2 \times 3 \times ... \times (n-2) \times (n-1) \times n
                                                              Begin
Factorial of 3:
                                                                 if n == 0 or 1 then
3! = 1 \times 2 \times 3
                                                                    Return 1;
                                                                 else
   = 6
Factorial using recursion:
                                                                    Return n*Call Fact(n-1);
F(n) = 1 \qquad \text{when } n = 0 \text{ or } 1
                                                                 endif
      = F(n-1) when n > 1
                                                              End
```

Recursion Terminology

- Base Case
 - The case for which the solution can be stated non recursively
- General (recursive) Case
 - The case for which the solution is expressed in terms of a smaller version of itself
- Recursive Algorithm
 - A solution that is expressed in terms of (1) smaller instances of itself and (2) a base case

```
Factorial(3)

Factorial(3)

Factorial(2)
```

```
Factorial(3)

3 * Factorial(2)

Factorial(3)

Factorial(1)
```

```
Factorial(3)

3 * Factorial(2)

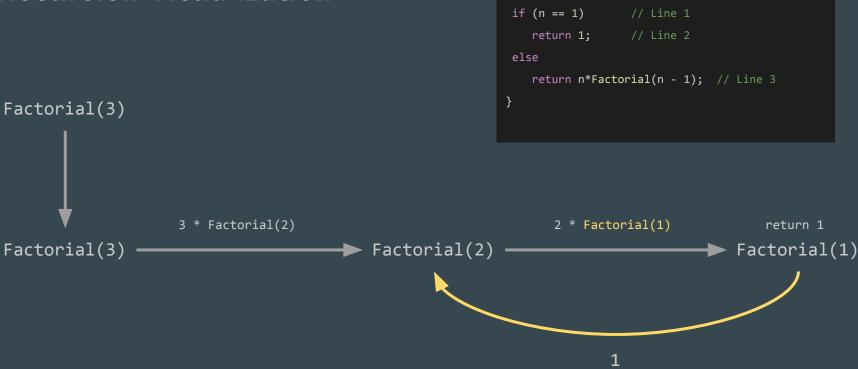
Factorial(3)

Factorial(2) -
```

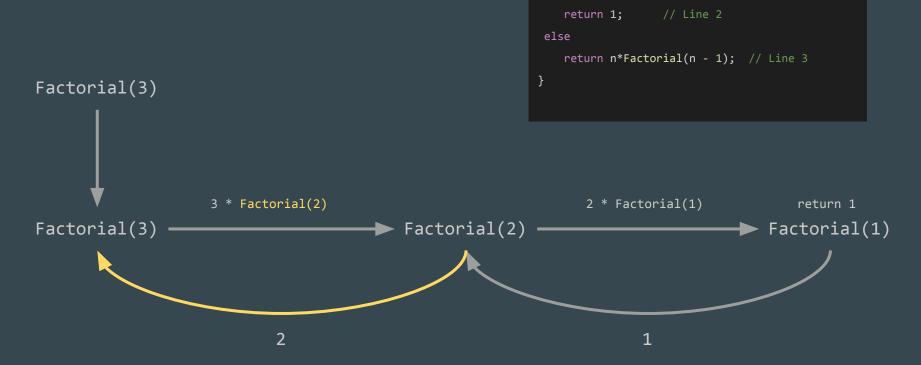
```
2 * Factorial(1)

Factorial(1)
```

return 1

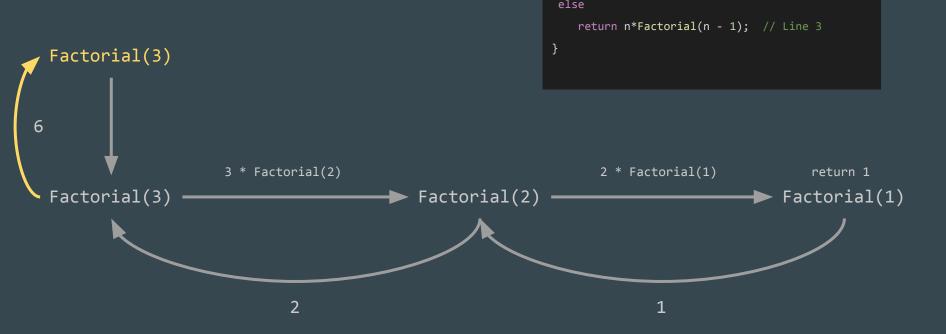


int Factorial(int n)



int Factorial(int n)

if (n == 1) // Line 1



int Factorial(int n)

if (n == 1) // Line 1

Recursion Implementation

```
int Factorial(int n)  // Non-recursive
{
  int fact = 1;
  for( int k = 2; k <= n; k++)
  {
    fact = fact*k;
  }
  return fact;
}</pre>
```

Infinite Recursion

- The situation in which a function calls itself over and over endlessly
- Consequences of Infinite Recursion
 - Run-time stack grows
 - Memory space consumed
 - Run-Time "Stack Overflow" error occurs

Verifying Recursive Functions

Three-Question Method

- 1. Base-Case Question
 - a. Is there a non-recursive way to exit the function?
 - b. Is it correct?
- 2. Smaller-Case Question
 - a. Does each recursive call to the function involve a smaller case of the original problem, leading inescapably to the base case?
- 3. General-Case Question
 - a. Assuming that the recursive calls work correctly, does the entire function work correctly?

Proof-By-Induction Procedure

- 1. Prove that f(n) is true for some value k
- 2. Assume that f(n) is true for some value n > k
- 3. Show that f(k+1) is true

Conclude that f(n) is true for all $n \ge k$

Proof-By-Induction Example

Correctness Proof:

Assume N = 1.

Does Factorial(1) equal 1!? Yes! Factorial(1) = 1 = 1!

Assume Factorial(N) is correct, i.e. Factorial(N) = N * (N-1) * ... * 2 * 1 = N!

Prove Factorial(N+1) (N+1)!

Proof-By-Induction Example

```
Mathematically: (N+1)! = (N+1) * N * (N-1) * ... * 2 * 1 = (N+1) * N!
```

According to the source code:

```
Factorial(N+1) = (N+1) * Factorial(N)
= (N+1) * N! \quad [Assuming Factorial(N) = N!]
= (N+1)!
```

Since we assumed that Factorial(N) = N!, we can rewrite Factorial(N+1) = (N+1)! =

Therefore, the function Factorial(N) will return N! for an arbitrary value of $N \ge 1$.

Recursion with Data Structures



```
void ReversePrint(NodePtr head)
       (head != NULL)
       ReversePrint(head->link);
       cout << head->component << endl;</pre>
ReversePrint(head of list ptr);
```

Tail Recursion

- A recursive function is tail recursive when recursive call is the last thing executed by the function.
- Tail recursion can be optimized by the compiler
- Is the factorial example tail recursion?

```
// An example of tail recursive function
void print(int n)
   if (n < 0) return;
   cout << " " << n;</pre>
    // The last executed statement is recursive call
   print(n-1);
```

Writing Recursive Functions

- Understand the problem first!!
- Determine the size of the problem
- Identify and solve the base case
- Identify and solve the general case using smaller instance of the general case

When should you use Recursion?

- Depth of recursion is relatively "shallow"
- Number of recursive calls grows slowly as problem size grows
- Recursive version does roughly the same amount of work as the non-recursive version
- Recursive version is shorter and simpler than the non-recursive version