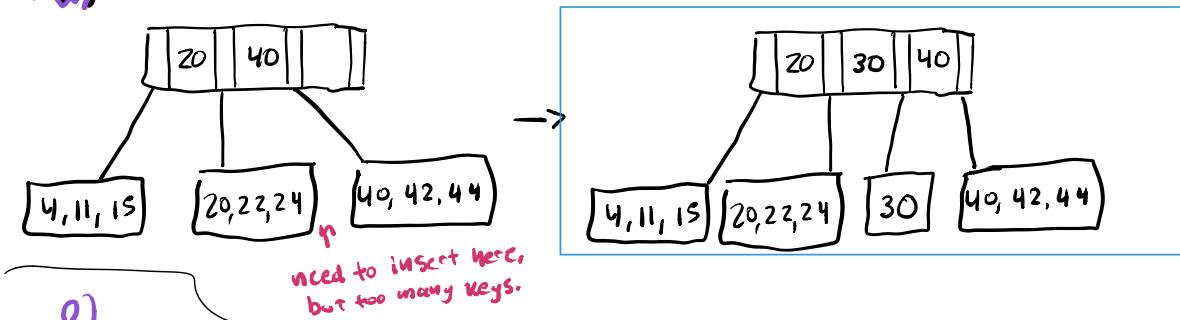
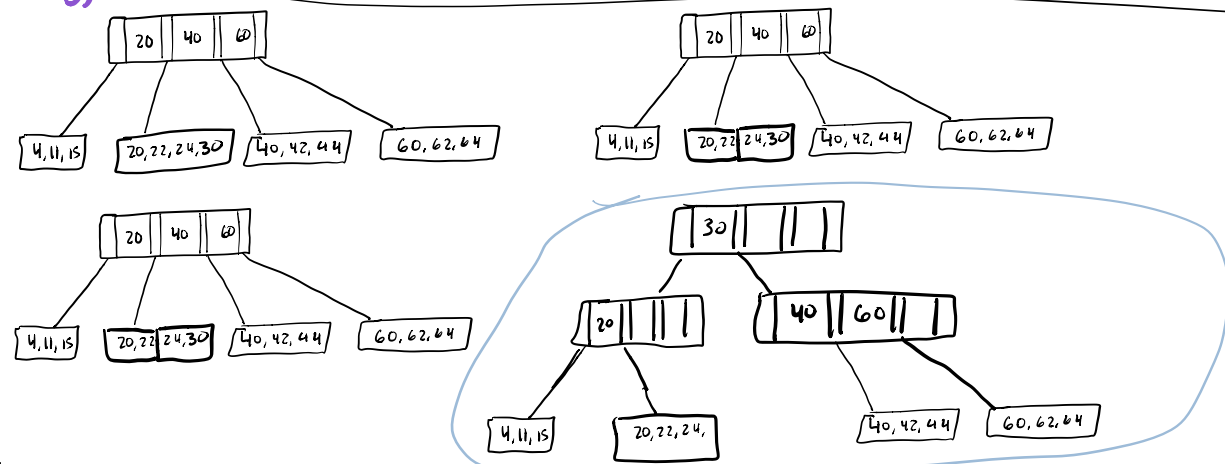


1.) a)

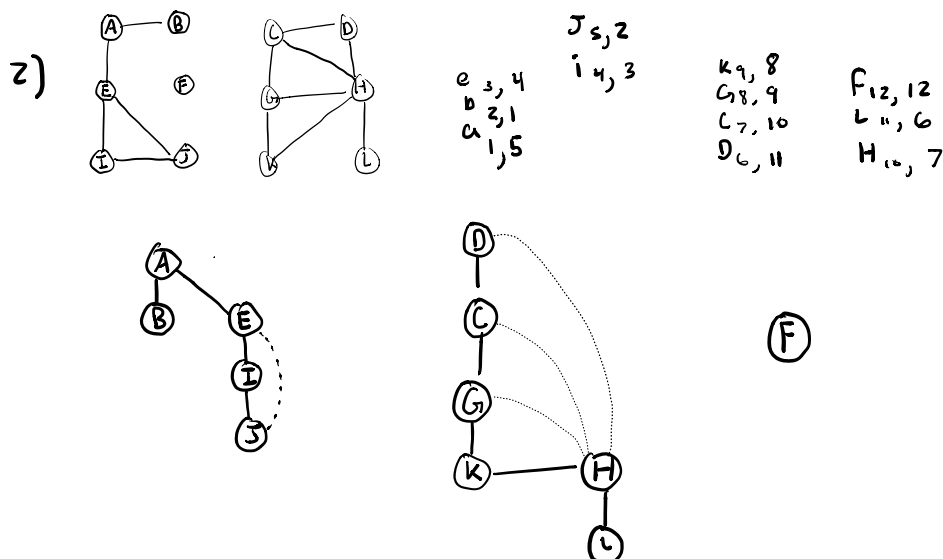


b)



c) Which space/time trade off technique...

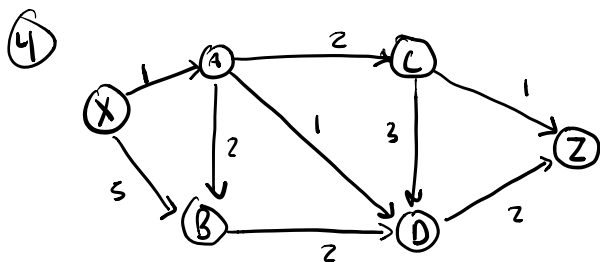
B-trees use prestructuring to use extra space but index faster.



Q3) $T(n) = 9T(\frac{n}{3}) + 3n^2 + 8n + 1$

$T(n) = 9T(\frac{n}{3}) + n(3n + 8) + 1$

$T(n) = 9(9T(\frac{n}{3})) +$ not sure... didn't study enough!



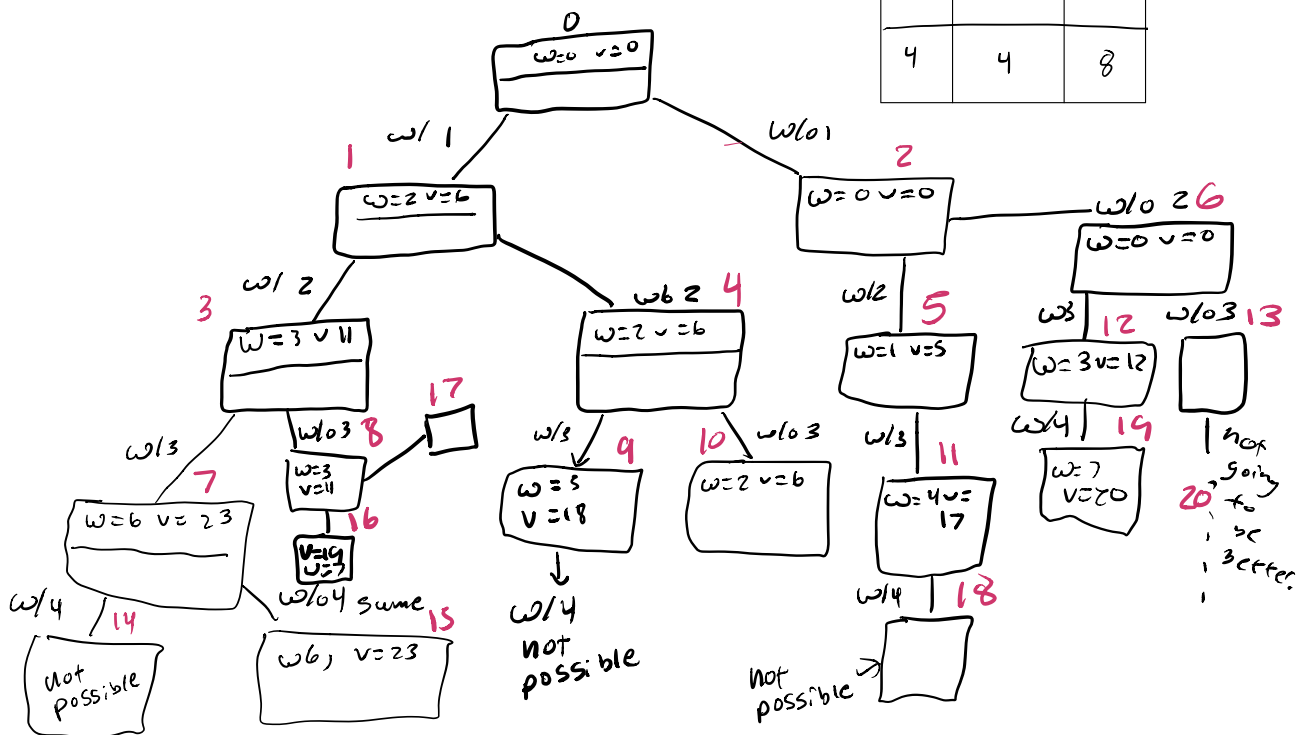
From X to A : X-A : 1
 From X to B : X-A-B : 3
 From X to C : X-A-C : 3
 From X to D : X-A-D : 2
 From X to Z : X-A-D-Z : 4

vertices	Remaining vertices
X(-, 0)	A(x, 1) B(x, 5) C(-, ∞) D(-, ∞) Z(-, ∞)
A(x, 1)	B(a, 1+2) D(a, 1+1) C(a, 1+2) Z(-, ∞)
D(a, 2)	Z(D, 2)

3.) 5, capacity is 7

The optimal solution is items
1, 2 and 3. Weight = 6
V = 23.

Item	Weight	Value
1	2	6
2	1	5
3	3	12
4	4	8



Sorry for the mess... hard to keep it clean.

(6)

- a) For a sorted array, simply compute the difference of first and last element. $O(1)$
- b) For unsorted linked list, you would need to traverse the whole array $O(n)$ to find min and max. So $O(n)$
- c) Smallest, leftmost node \rightarrow Follow left until left is null.
Largest rightmost node \rightarrow follow right until right node is Null.
 $\hookrightarrow 1 \rightarrow n$ nodes long, running time is $O(n)$

(7.)

a) I would say this is dynamic programming as you are throwing out 2/3's of places the item could be.

b) $T(n) =$ running out of time...

c)

d)

BONUS:

NP contains The knapsack problem, graph coloring, and Hamiltonian Circuits, but there are certainly more.