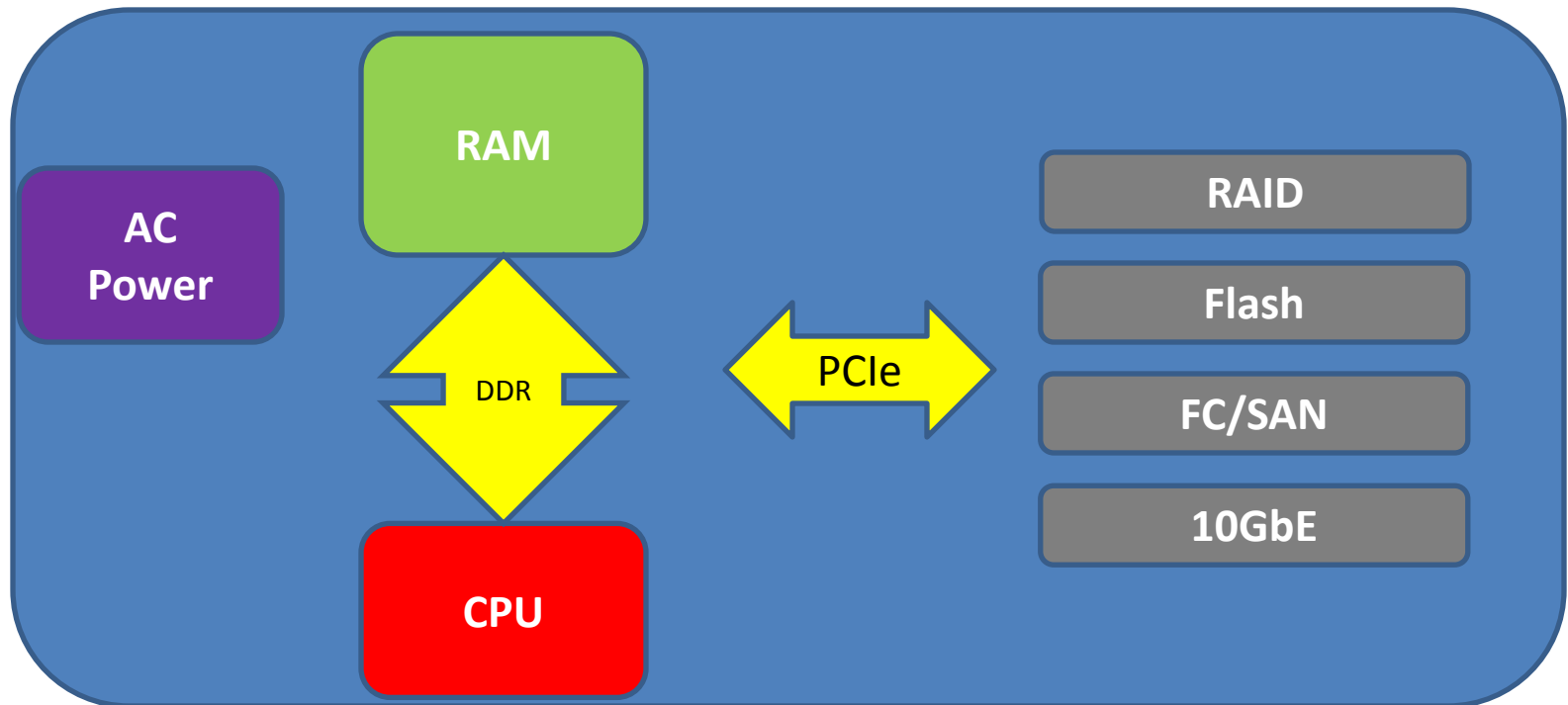


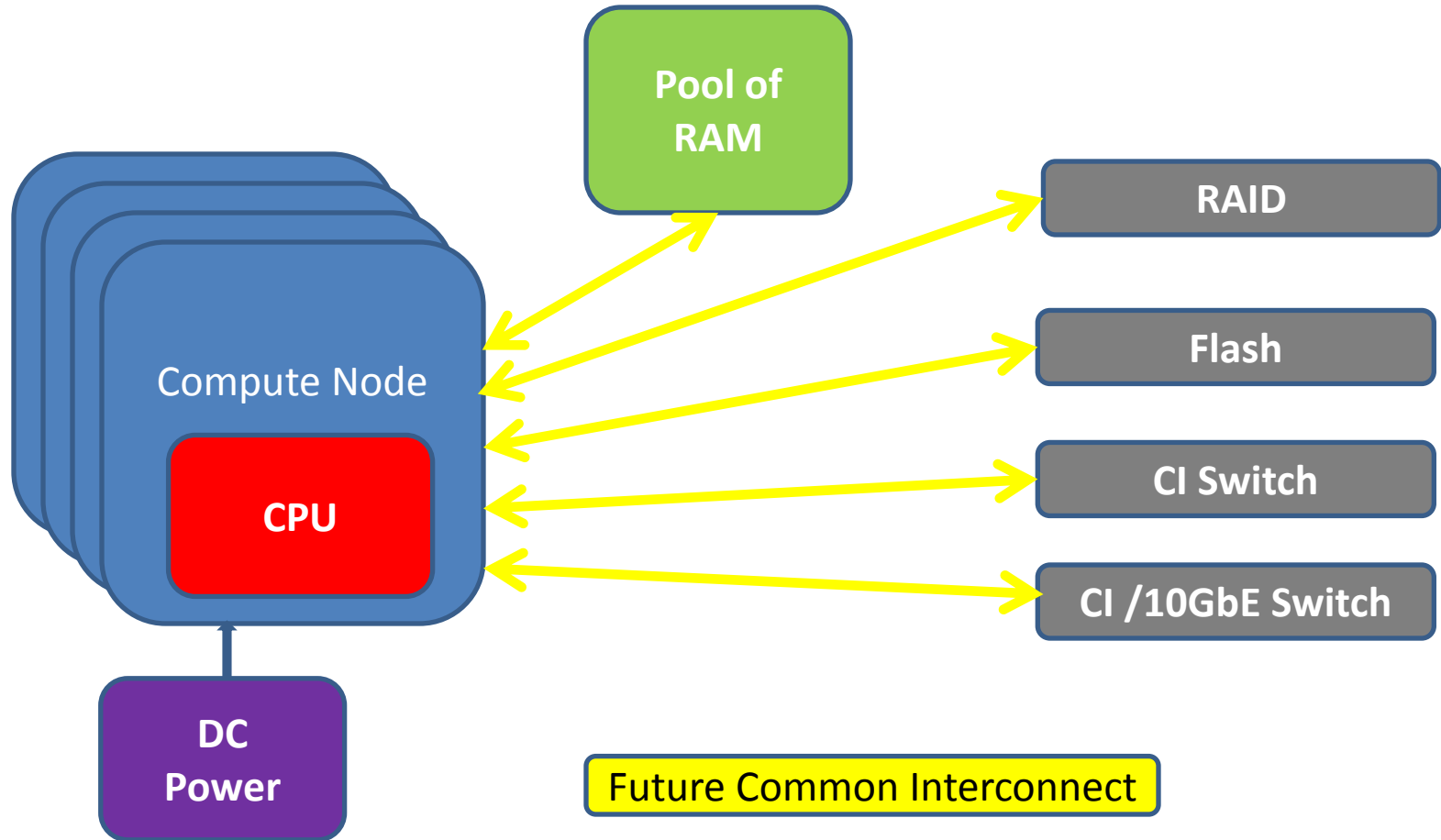
# Appmem Devices

Disaggregating RAM

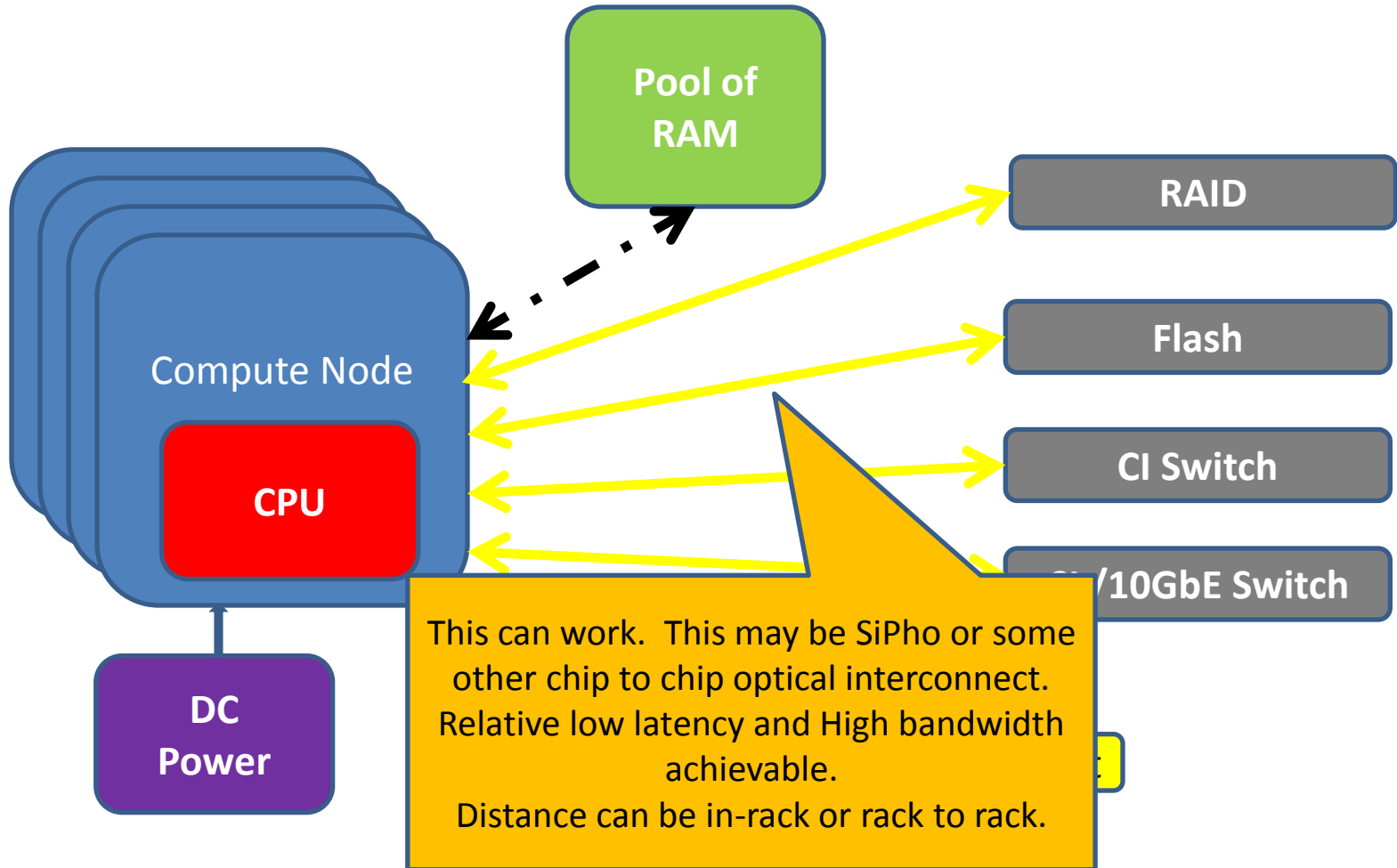
# Traditional Server



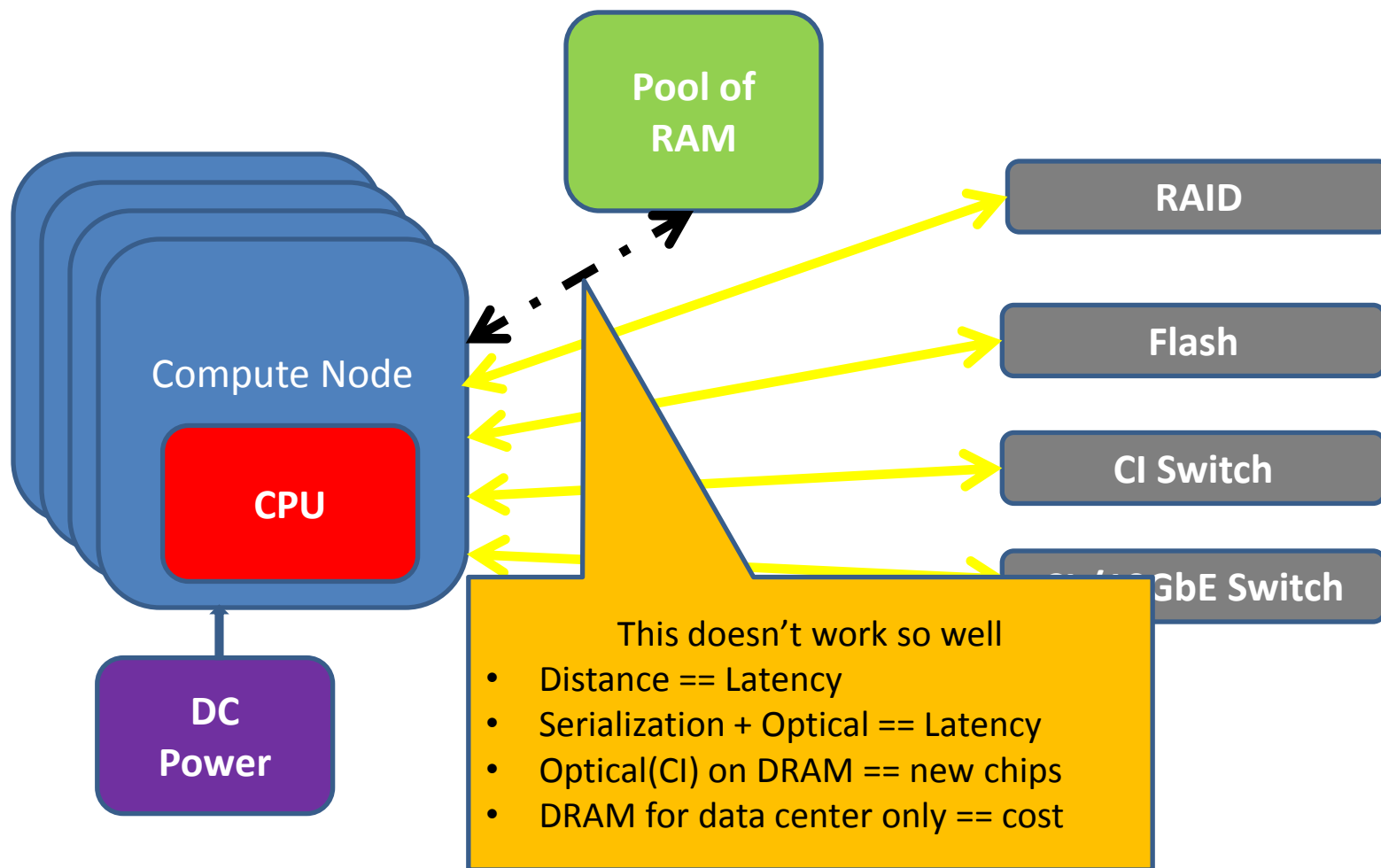
# Future Common Interconnect (CI) Disaggregated Rack



# Could work 😊



# Maybe Not ☹️



# Common interconnect

- DDRx level bandwidth achievable (40+ GB/s).
- Chip to chip optical might be technically challenging and costly but worth it in the data center.
- Lowest possible latency for the given protocol. (PCIe latencies or better + distance).
- Optical Avoids the electrical design and layout problems in todays boards (EMI, crosstalk).
- Fabrication, Marketing, and business hurdles to overcome. No fundamental technological or architectural roadblocks.

# RAM Latency is a killer

- Generations of CPU design have dealt with RAM latency, yet we still need to keep it as low as possible.
  - multi-level caches
  - Out-of-order execution
  - Hyper-threading
- Long stalls waiting on System RAM have a cascading effect on system performance.
- Architecturally to applications, all System RAM is treated the same (compilers and interpreters).
- Adding higher-latency RAM to system without application awareness will lead to inefficient programs.

# Disaggregating RAM adds latency

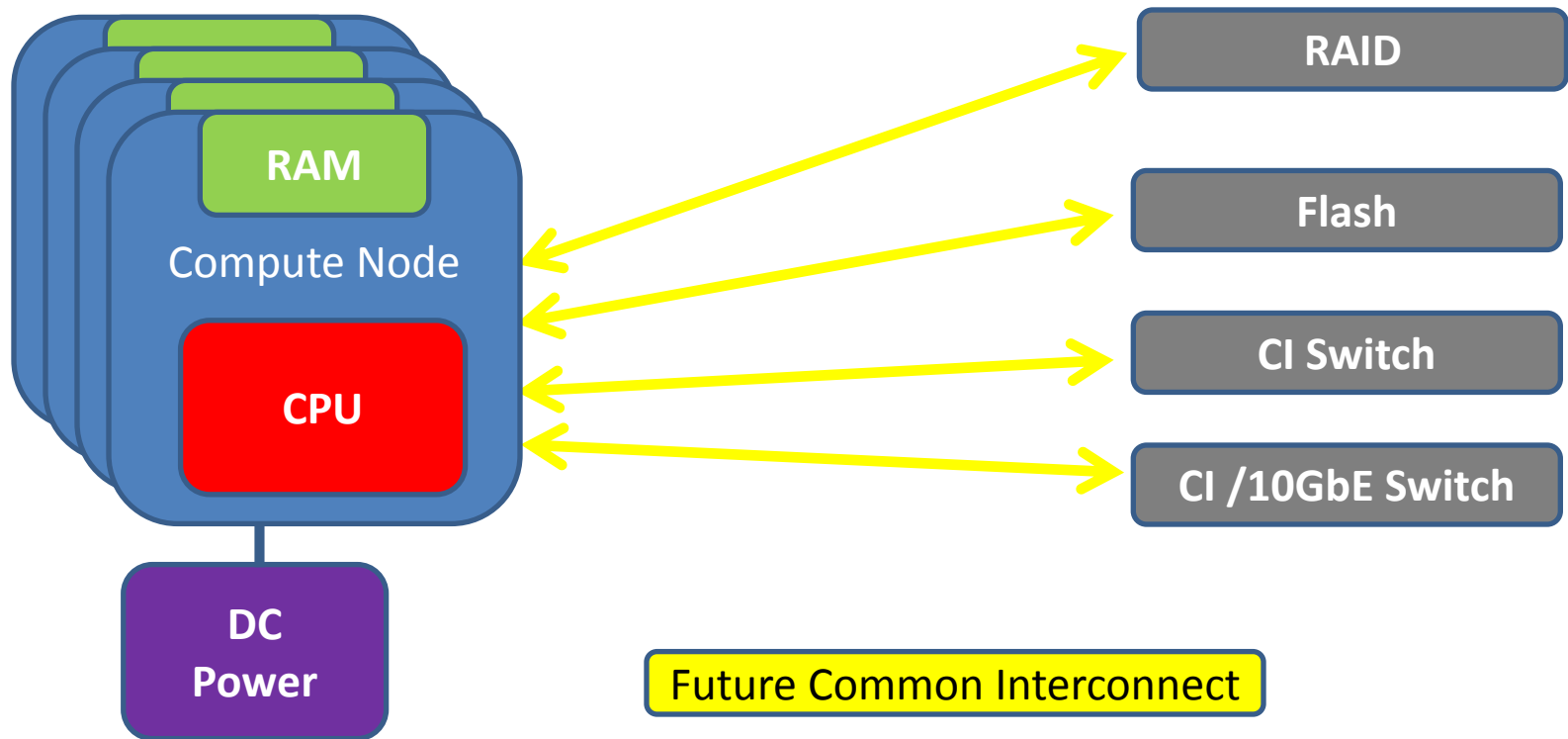
- Physical Optical latency = 5ns per meter.
  - Theoretical latency of DDR4 is around 15 ns.
  - DDR traces lengths are inches.
- Serdes of request/data stream add latency.
  - Logic at both ends needs to covert bit stream into addressable memory (dozens of clock cycles at both ends).
  - Currently DDR uses parallel bus and can transfer full address and multi-word data in single cycle.
  - Transport latency of PCIe is about 150-300ns (10-20x DDR)



# Market issues

- DRAM chips are used everywhere
  - Adding CI to DRAM chips or adapters for data center adds costs that can't easily be amortized across consumer application in the near term.
  - Specialized DRAM chips for data center only add costs reducing benefit of disaggregation.
  - DRAM chips are commodity parts – increasing complexity will increase fabrication difficulties and cost.
  - Needs of serial memory in consumer space are different than those in data-center.

# So we have to live with this?



# Aggregated CPU and RAM

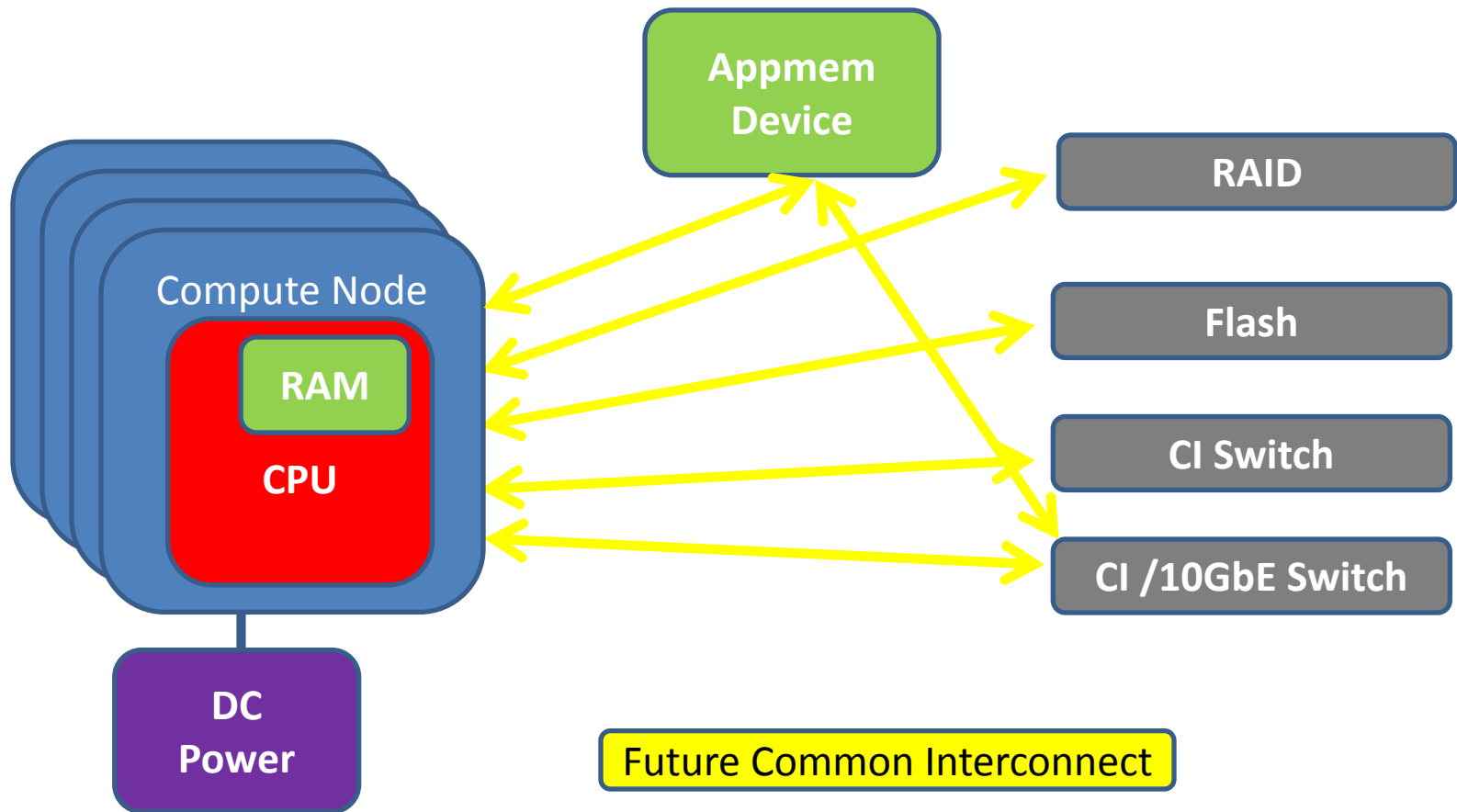
- Many applications are CPU/IO limited before they become RAM limited.
- 4-16 GB is plenty for many data center applications.
- Cost/layout reduction trend may be RAM on CPU die or package.
- Swap to flash is a solution for applications that occasionally exceed physical RAM.
- No architecture or application changes.

# Disaggregated CPU and RAM

- For some applications RAM needs aren't known or change over time.
- Larger chunks of RAM may be used for short amounts of time while data is collated and processed.
- Using secondary storage for transitory data is both inefficient and costly.
- With CPU and RAM tied, you can't really have the "lego blocks" vision of a self-configuring data center.

# Future Disaggregated Rack

Maybe it will look like this



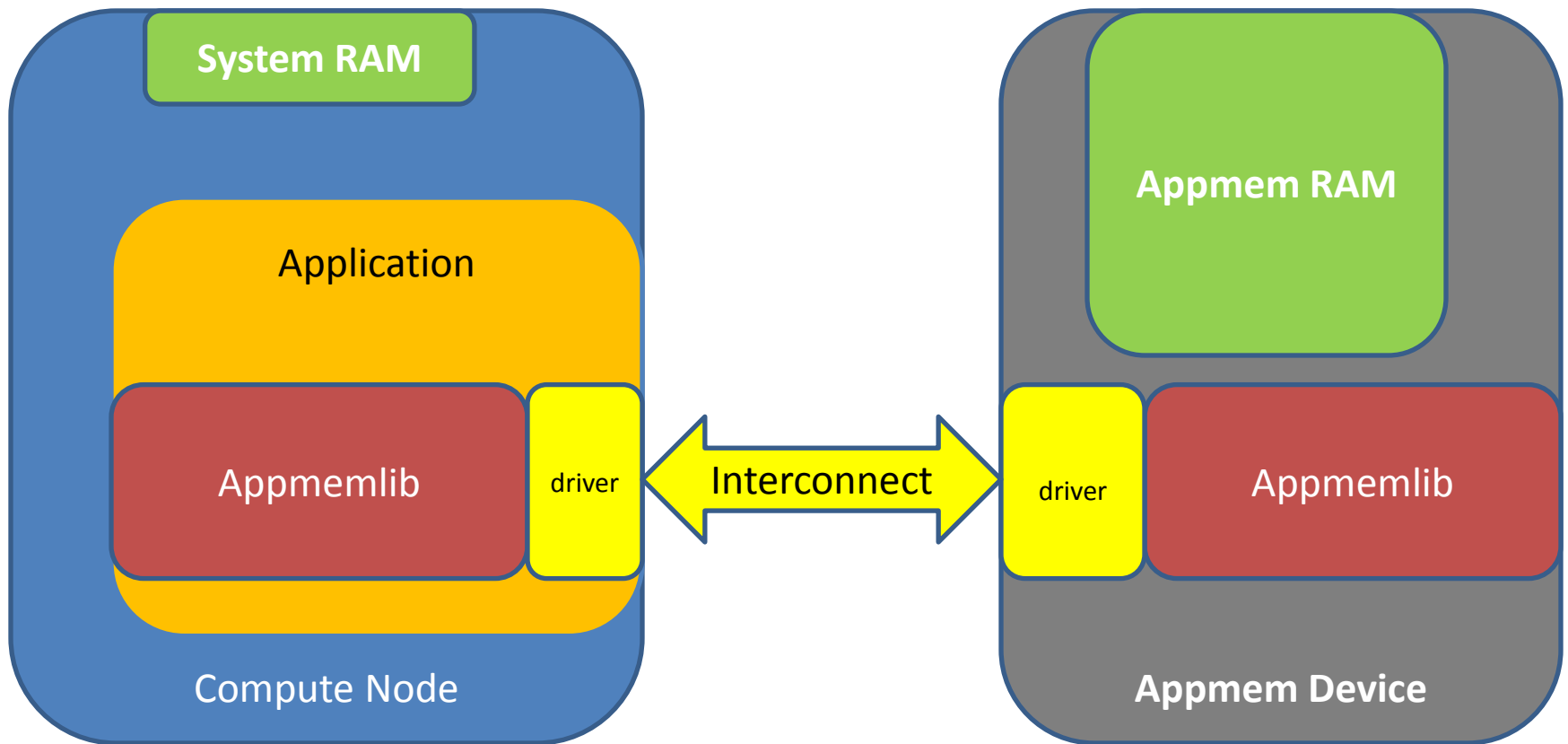
# Appmem Device

- Device that has a large amount of RAM
  - May be a traditional server with lots of RAM
  - May be an embedded device
- Available to Applications
- Has “functions” – common data structures used by applications
  - Associative and Dynamic Arrays
  - Lists, Queues, Stacks
  - Trees
  - Application defined structures

# Appmem Device Contentions

1. Memory may or may not be persistent
2. Memory may or may not be shared
3. Interconnect is irrelevant
4. Categorized by latency **at the application.**
5. Accessed using a common library.

# Software Architecture





# Appmemlib

- Uses the same library on the client and device.
- Application programmers can define their own functions and test them in System RAM, then integrate library on the device.
  - If device is Linux based, just compile appmemlib on the device.
  - If device is embedded non-Linux, vendor can provide SDK based on appmemlib.
- Interconnect specific driver either standardized or provided by the device vendor.

# Appmem Device

## Basic Advantages

- Direct access operations should have the same latency as the underlying interconnect.
  - PCIe Appmem device write/read of an object to memory should have same order of latency as mapped PCIe device memory read/write.
  - Some overhead penalty for the API.
- Complex Direct access operations are offloaded.
  - Example - In an Associative Array, hashing the key occurs on the device. Depending on speed of the device, overall latency may be greater.
  - The hashing is offloaded from the main CPU, freeing it to process other data.

# Appmem Device

## Compute Advantages

- Some Compute intensive operations can be completely offloaded.
  - Sort, Find, Copy, Delete, etc.
  - Any operation where the data elements can be described generically (offset, data type).
  - More complex and application specific functions can be created when Device needs to know structure of the data.

# Appmem Device

## Value Add

- Persistence

- Device may be have flash/block or file based persistent storage.
- Persistence may be dynamically persistent or client directed (flush or close).

- Sharing

- Single Client may own the function or it may be shared publicly or through some authorization mechanism.
- Single Client mode sharing (one at a time) or live sharing .
- Locks/atomicity handled by the Device.

# Appmem Device

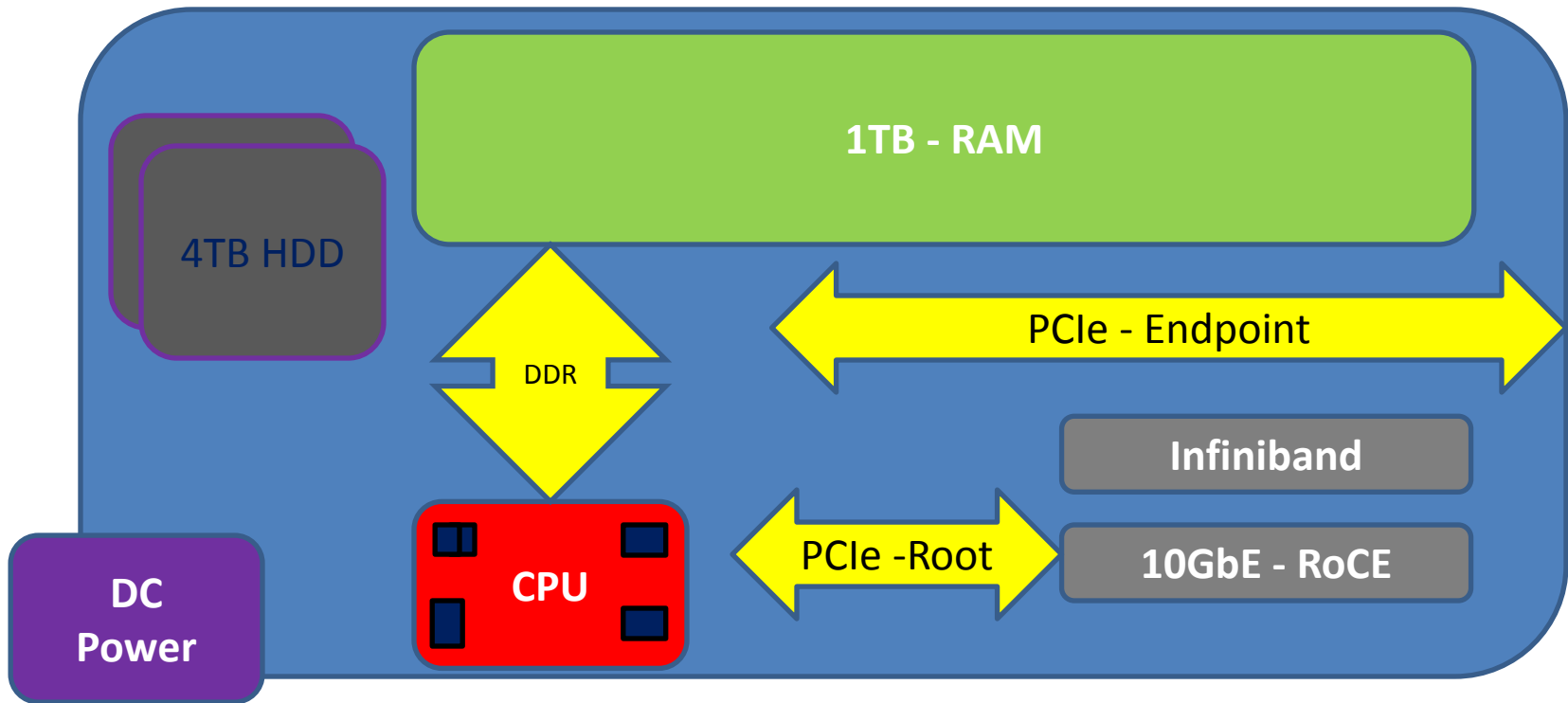
## Value Add advantages

- Data structures can be swapped out of RAM to make room for other applications when some are idle.
- One client may perform collation operation from many sources, create a huge multi-GB Tree or List, and then hand off the handle to second client who can begin some processing operations on the data. No file read/write delay between the operations.
- Many other uses cases.

# A High End Appmem Device

## Current technology

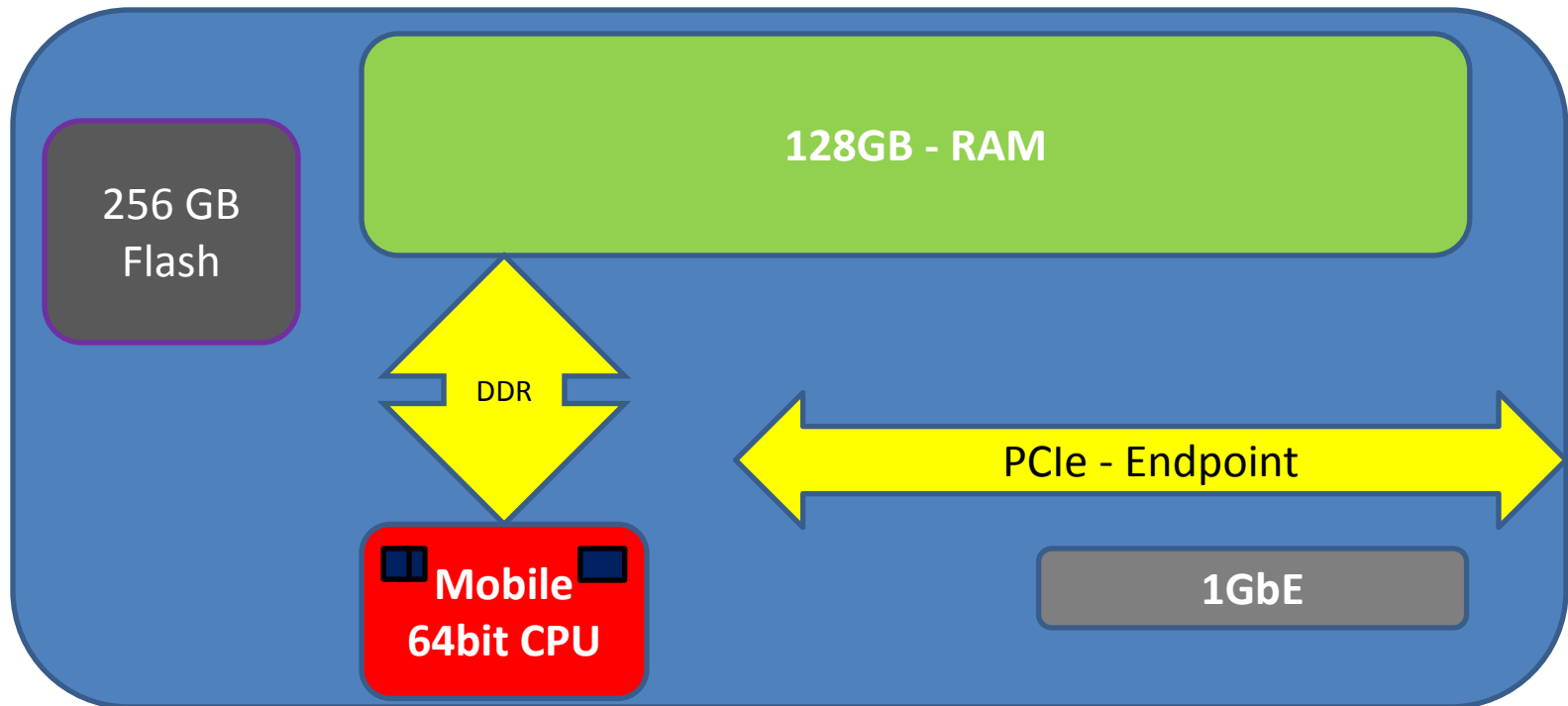
- Basically, a 4 core x86 server with a bunch of RAM, serving as a PCIe endpoint and also has Infiniband and 10GbE Adapters. May be in a rack or card slot.



# A Embedded Appmem Device

## Current technology

- A 64 bit Mobile-Class CPU, 128 GB RAM, with Flash. PCIe Endpoint with a remote access through 1GbE.



# Appmem Device classification

- Devices are categorized by latency – at the client.
- Will use multi-level cache nomenclature
  - A0 has 0-10ns latency
  - A1 100ns
  - A2 1us
  - A3 10us. ( $A5 = 10\text{ns} \wedge (5+1) = 1\text{ms}$ )
- When client discovers devices, it will perform simple latency test to determine classification.
  - High end device might be A1 when connected via PCIe.
  - Same device could be A4 when connected via Ethernet and multiple switch hops.



# Appmem project

- <https://github.com/scuzzydude/appmem>
- Very early concept code
  - Appmemtest utility. Tests both C and C++ interfaces to appmemlib.
  - Appmemlib – shared library with interfaces and basic data structures. User mode (emulation) target.
  - Appmemk – kernel driver with Memory Mapped and IOCTL interfaces and kernel emulated target.
  - Am\_targ – very basic networked target
- Supports Flat Memory, Static Arrays, and Associative Arrays.

# Next Steps

- Hardware
  - Demo real device as PCIe endpoint.
    - Server using PCIe NT as an endpoint.
    - Modify firmware of an existing PCIe card
    - FPGA based solution.
- Appmemlib
  - Real world use cases and realistic “big data” benchmark.
  - Add more data structures (Lists, Queues, Stacks, Trees, but also looking for more complex structures and algorithms to demonstrate offload.)
  - Terminology refinement (i.e. spec)
  - Refinement of library->driver layer
  - Wrappers for other languages (Java, JS, PHP).

# Reference and Contact

- [Open Compute Summit – Disaggregated Rack](#)
- [Appmem Overview](#)
- [Github Appmem Project](#)
- [Appmem API and Implementations Notes](#)
- [Appmem Architecture Notes](#)

## Contact

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