

# CSC 665: Artificial Intelligence

## Constraint Satisfaction Problems

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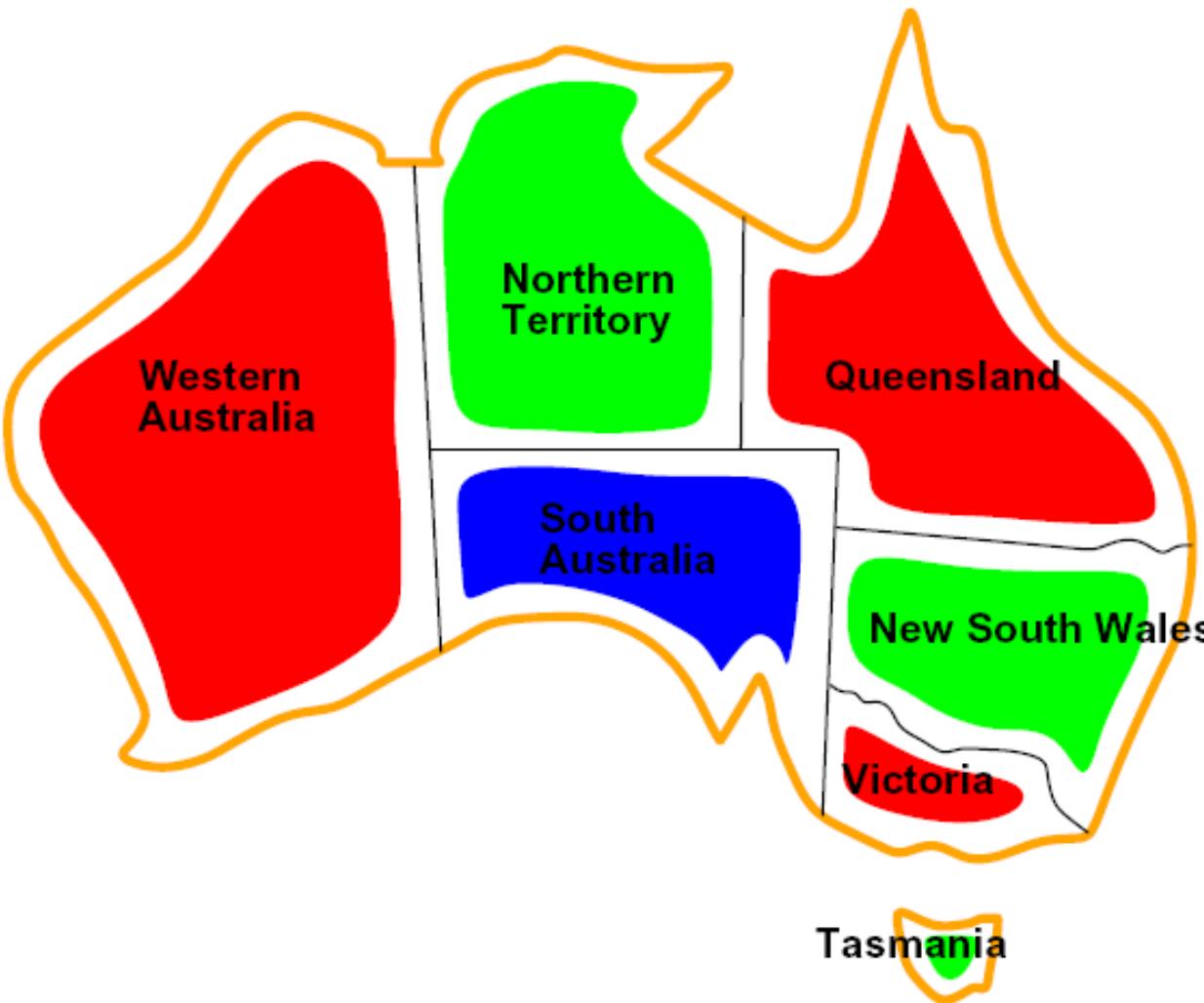
# Constraint Satisfaction Problems

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- Constraint satisfaction problems (CSPs):
  - A special subset of search problems
  - State is defined by a set of variables  $X_i$  and each variable has a domain  $D_i$
  - Goal test is a set of constraints specifying allowable combinations of values for subsets of variables

# CSP Examples

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# Example: Map Coloring

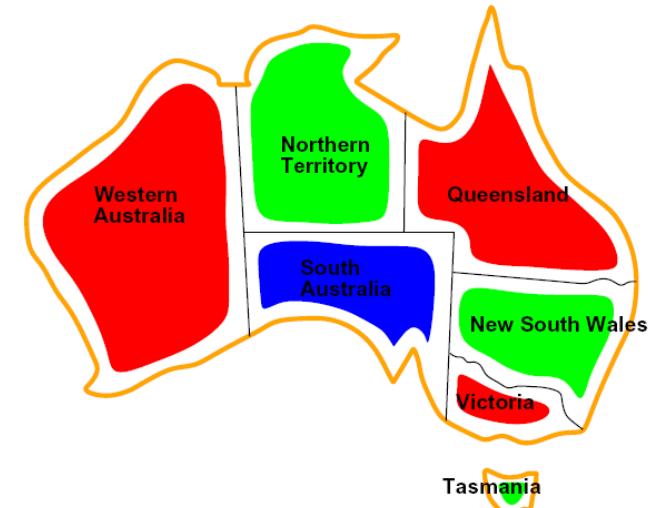
- Variables: WA, NT, Q, NSW, V, SA, T
- Domains:  $D = \{\text{red, green, blue}\}$
- Constraints: adjacent regions must have different colors

Implicit:  $\text{WA} \neq \text{NT}$

Explicit:  $(\text{WA}, \text{NT}) \in \{(\text{red, green}), (\text{red, blue}), \dots\}$

- Solutions are assignments satisfying all constraints, e.g.:

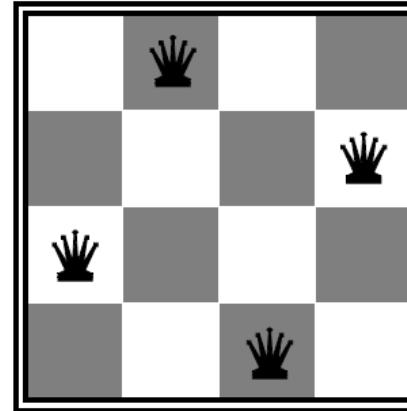
$\{\text{WA=red, NT=green, Q=red, NSW=green, V=red, SA=blue, T=green}\}$



# Example: N-Queens

## ■ Formulation 1:

- Variables:  $X_{ij}$
- Domains:  $\{0, 1\}$
- Constraints



$$\forall i, j, k \quad (X_{ij}, X_{ik}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$\forall i, j, k \quad (X_{ij}, X_{kj}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$\forall i, j, k \quad (X_{ij}, X_{i+k, j+k}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$\forall i, j, k \quad (X_{ij}, X_{i+k, j-k}) \in \{(0, 0), (0, 1), (1, 0)\}$$

$$\sum_{i,j} X_{ij} = N$$

# Example: N-Queens

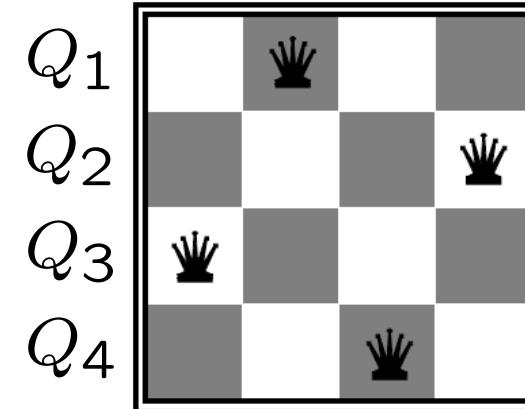
- Formulation 2:

- Variables:  $Q_k$  (Rows)
- Domains:  $\{1, 2, 3, \dots, N\}$
- Constraints:

Implicit:  $\forall i, j \text{ non-threatening}(Q_i, Q_j)$

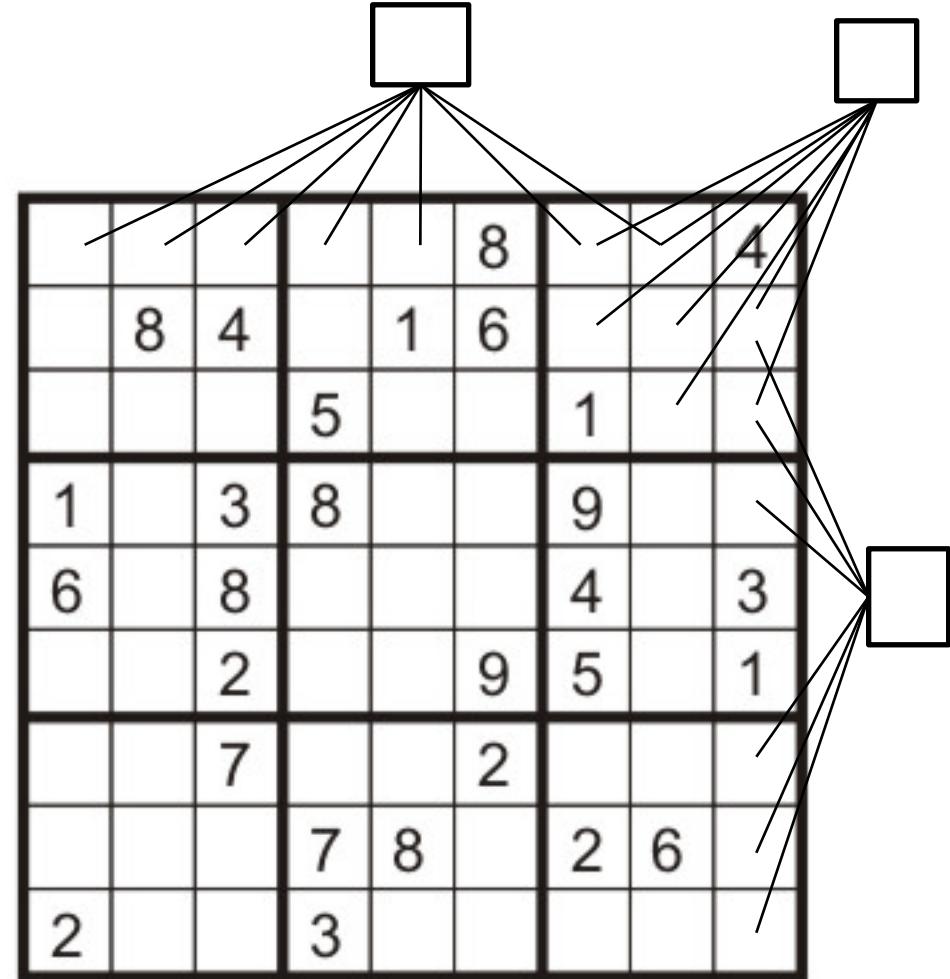
Explicit:  $(Q_1, Q_2) \in \{(1, 3), (1, 4), \dots\}$

...



# Example: Sudoku

- **Variables:**
  - Each (open) square
- **Domains:**
  - $\{1, 2, \dots, 9\}$
- **Constraints:**
  - 9-way alldiff for each column
  - 9-way alldiff for each row
  - 9-way alldiff for each region



# Constraint Satisfaction Problems: Definition

Definition:

A **Constraint Satisfaction Problem (CSP)** consists of:

- a set of variables  $V$
- a domain  $\text{dom}(V)$  for each variable  $V \in V$
- a set of constraints  $C$

Simple example:

- $V = \{V_1\}$ 
  - $\text{dom}(V_1) = \{1,2,3,4\}$
- $C = \{C_1, C_2\}$ 
  - $C_1: V_1 \neq 2$
  - $C_2: V_1 > 1$

Another example:

- $V = \{V_1, V_2\}$ 
  - $\text{dom}(V_1) = \{1,2,3\}$
  - $\text{dom}(V_2) = \{1,2\}$
- $C = \{C_1, C_2, C_3\}$ 
  - $C_1: V_2 \neq 2$
  - $C_2: V_1 + V_2 < 5$
  - $C_3: V_1 > V_2$

# Models of a CSP

Definition:

A **model** of a CSP is an assignment of values to all of its variables that **satisfies** all of its constraints.

Simple example:

- $V = \{V_1\}$ 
  - $\text{dom}(V_1) = \{1,2,3,4\}$
- $C = \{C_1, C_2\}$ 
  - $C_1: V_1 \neq 2$
  - $C_2: V_1 > 1$

All models for this CSP:

- { $V_1 = 3$ }
- { $V_1 = 4$ }

# Models of a CSP

- $\mathcal{V} = \{V_1, V_2\}$ 
  - $\text{dom}(V_1) = \{1, 2, 3\}$
  - $\text{dom}(V_2) = \{1, 2\}$
- $C = \{C_1, C_2, C_3\}$ 
  - $C_1: V_2 \neq 2$
  - $C_2: V_1 + V_2 < 5$
  - $C_3: V_1 > V_2$

Which are models for this CSP?

{ $V_1=1, V_2=1$ }

{ $V_1=2, V_2=1$ }

{ $V_1=3, V_2=1$ }

{ $V_1=3, V_2=2$ }



# Varieties of Variables

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- Discrete Variables
  - Finite domains
    - Size  $d$  means  $O(d^n)$  complete assignments
  - Infinite domains (integers, strings, etc.)
  
- Continuous variables
  - E.g., start/end times for Hubble Telescope observations

# Varieties of Constraints

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- Varieties of Constraints

- Unary constraints involve a single variable (equivalent to reducing domains), e.g.:

$SA \neq \text{green}$

- Binary constraints involve pairs of variables, e.g.:

$SA \neq WA$

- Higher-order constraints involve 3 or more variables

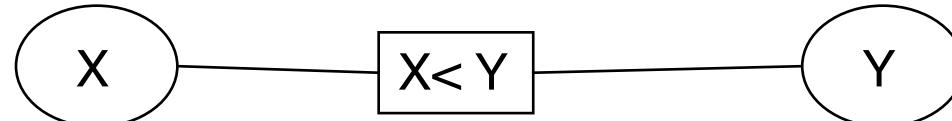
# Constraint Network: Definition

Definition: A **constraint network** is defined by a graph, with

- one **node** for every **variable** (drawn as **circle**)
- one **node** for every **constraint** (drawn as **rectangle**)
- **Edges/arcs** running between **variable nodes** and **constraint nodes** whenever a given variable is involved in a given constraint.

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- Two variables X and Y
- One constraint:  $X < Y$



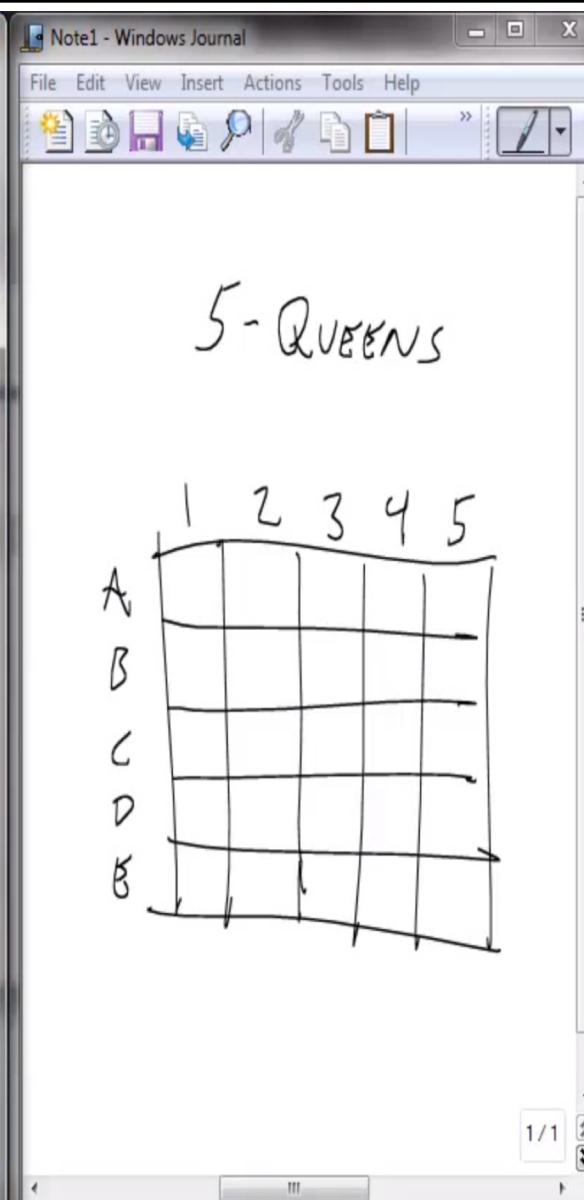
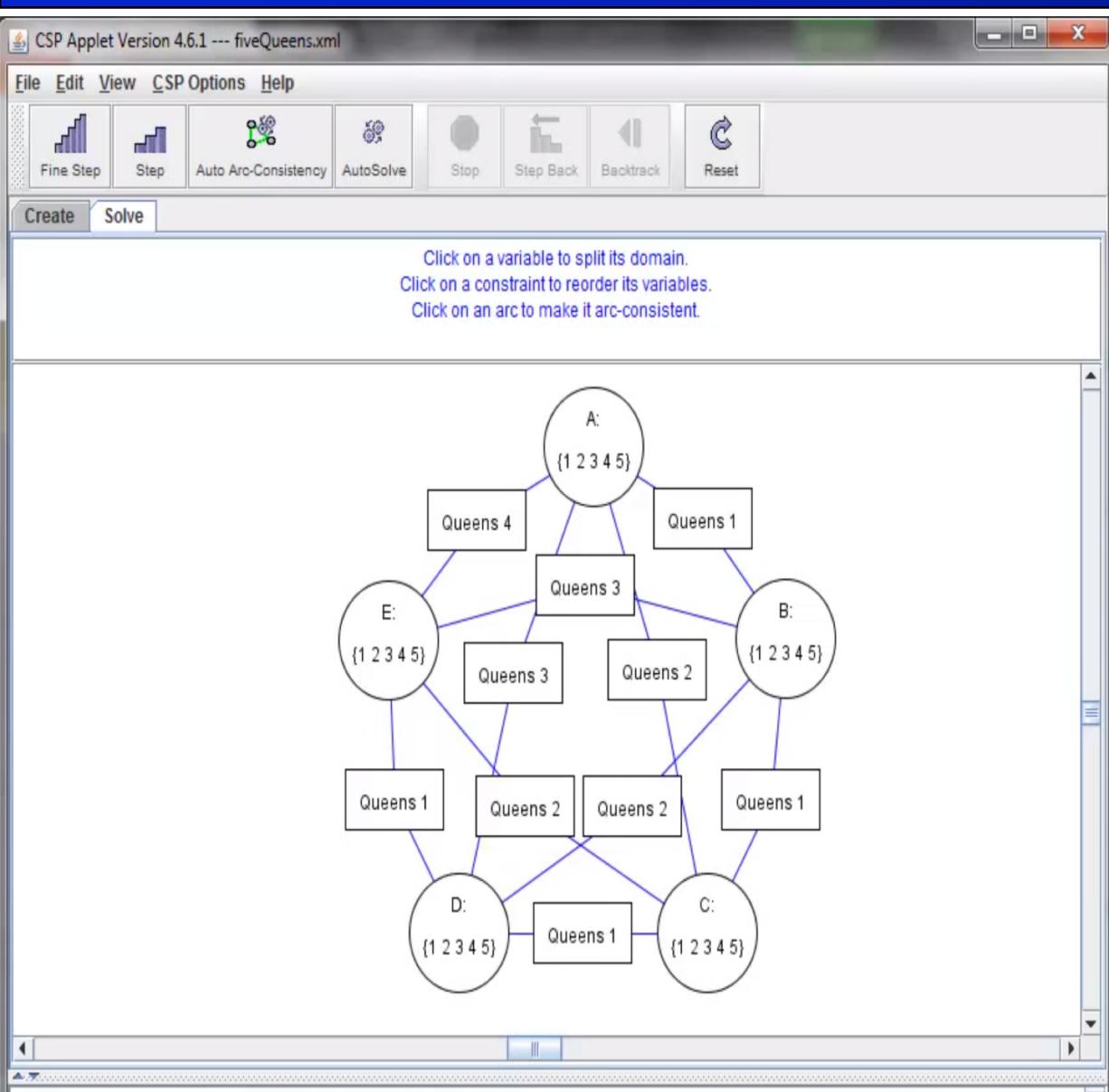
# Constraint Network: Definition

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- **Edges/arcs** running between **variable nodes** and **constraint nodes** whenever a given variable is involved in a given constraint.

- Whiteboard example: 3 Variables A,B,C
  - 3 Constraints:  $A < B$ ,  $B < C$ ,  $A + 3 = C$
  - 6 edges/arcs in the constraint network:
    - $\langle A, A < B \rangle$  ,  $\langle B, A < B \rangle$
    - $\langle B, B < C \rangle$  ,  $\langle C, B < C \rangle$
    - $\langle A, A + 3 = C \rangle$  ,  $\langle C, A + 3 = C \rangle$

# Alspace CSP Demos



<http://www.aispace.org/>

# Real-World CSPs

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- Assignment problems: e.g., who teaches what class
  - Timetabling problems: e.g., which class is offered when and where?
  - Hardware configuration
  - Transportation scheduling
  - Factory scheduling
  - Circuit layout
  - Fault diagnosis
  - ... lots more!
- 
- Many real-world problems involve real-valued variables...

# Solving CSPs

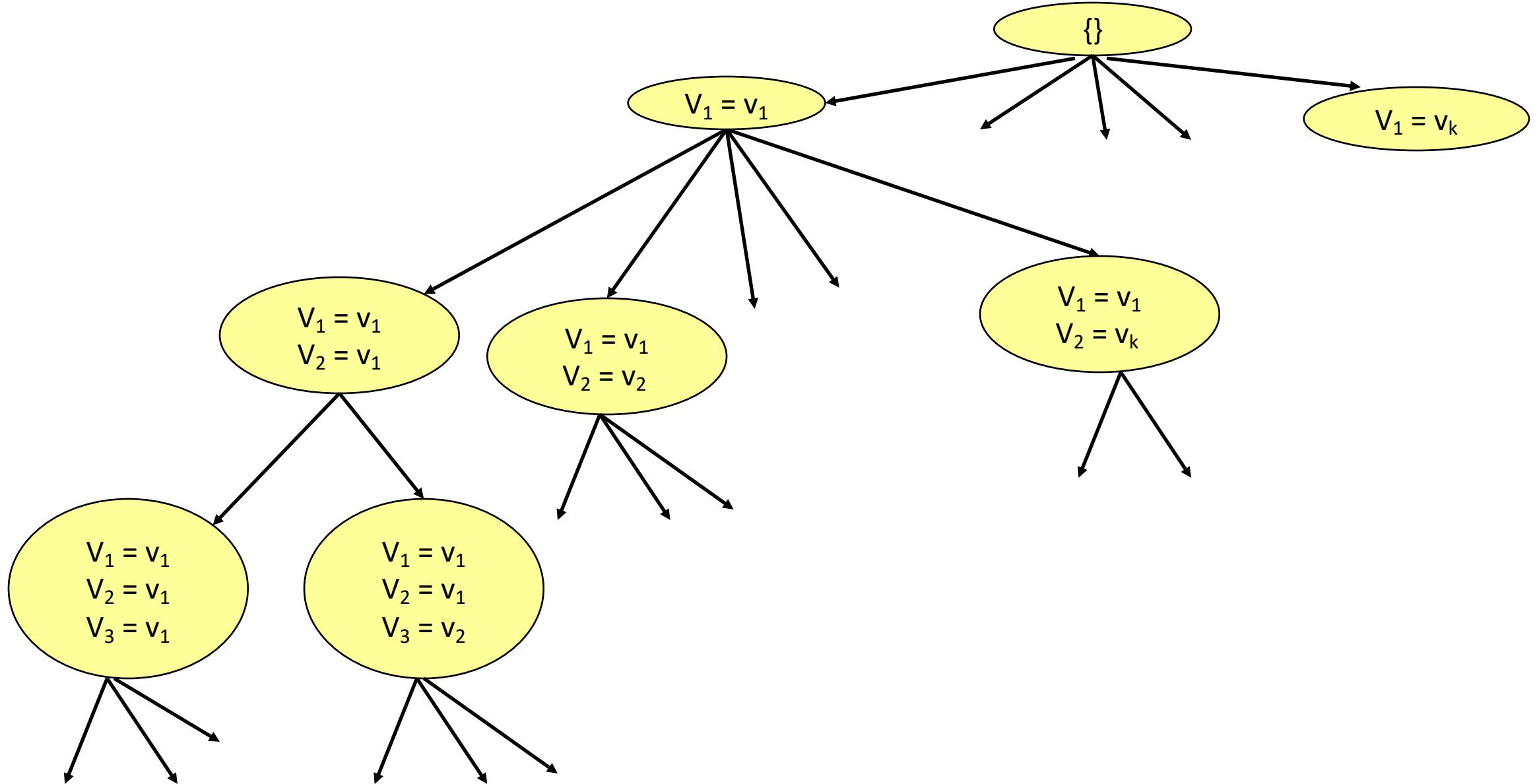
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# Standard Search Formulation

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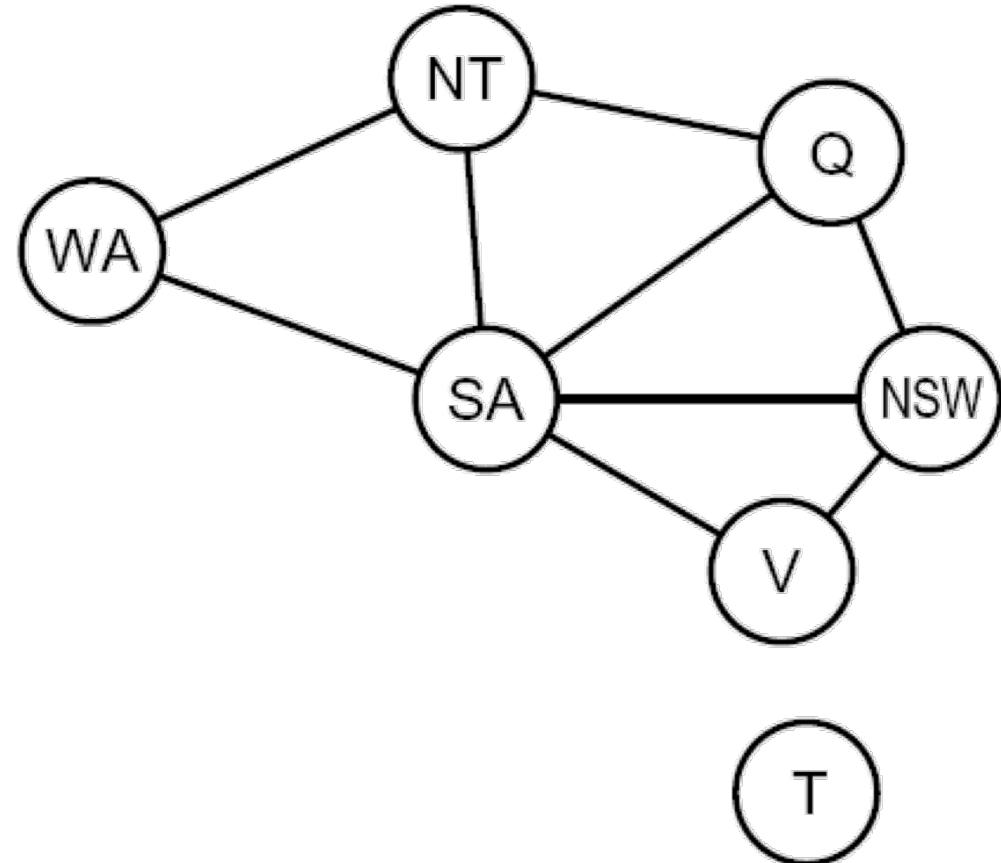
- Standard search formulation of CSPs
- States defined by the values assigned so far (partial assignments)
  - **Initial state**: the empty assignment, {}
  - **Successor function**: assign a value to an unassigned variable
  - **Goal test**: the current assignment is complete and satisfies all constraints
- We'll start with the straightforward, naïve approach, then improve it

# Standard Search Formulation: Example

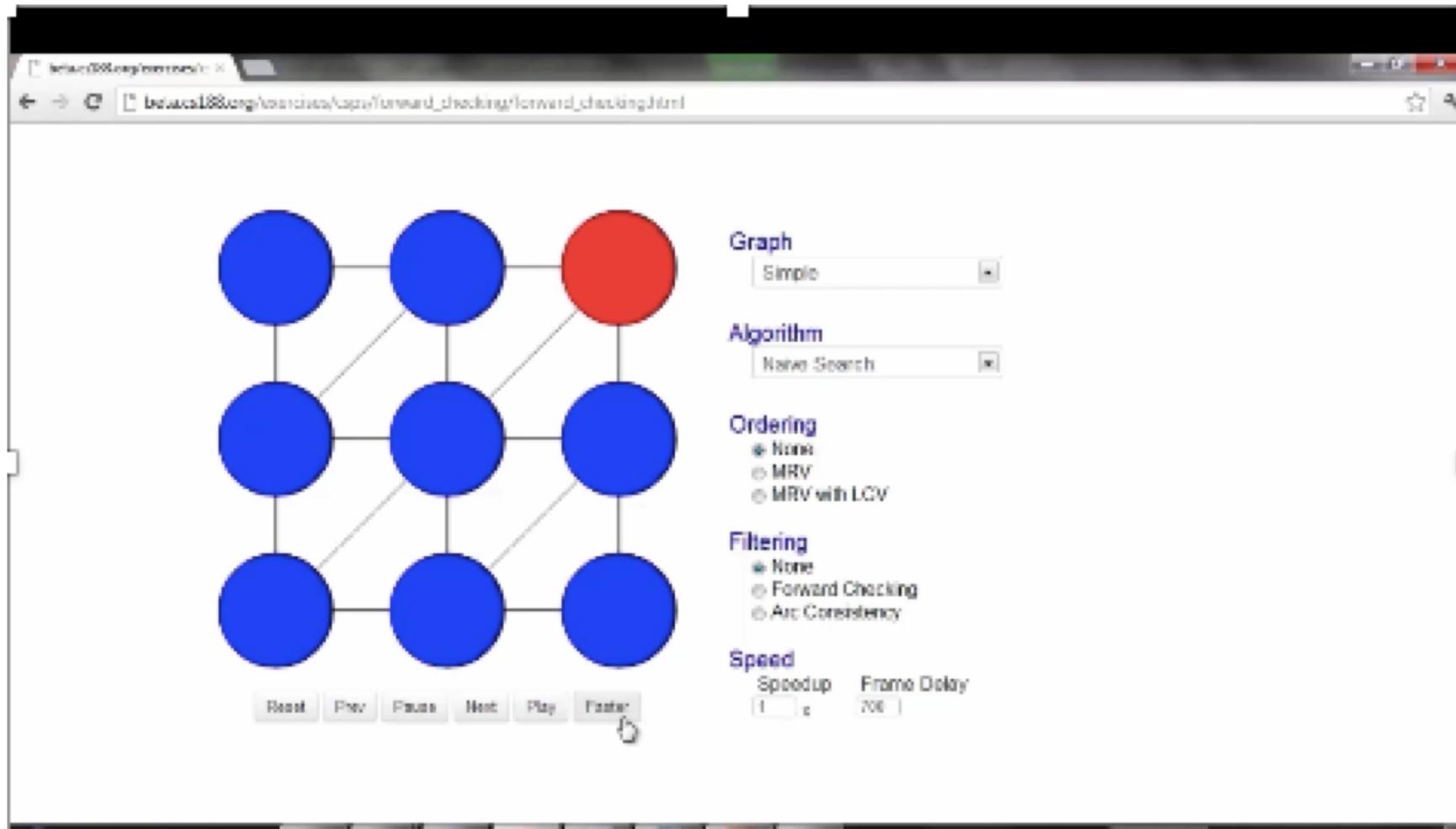


# Search Methods

- What would BFS do?
- What would DFS do?



# Video of Demo Coloring -- DFS



# Backtracking Search

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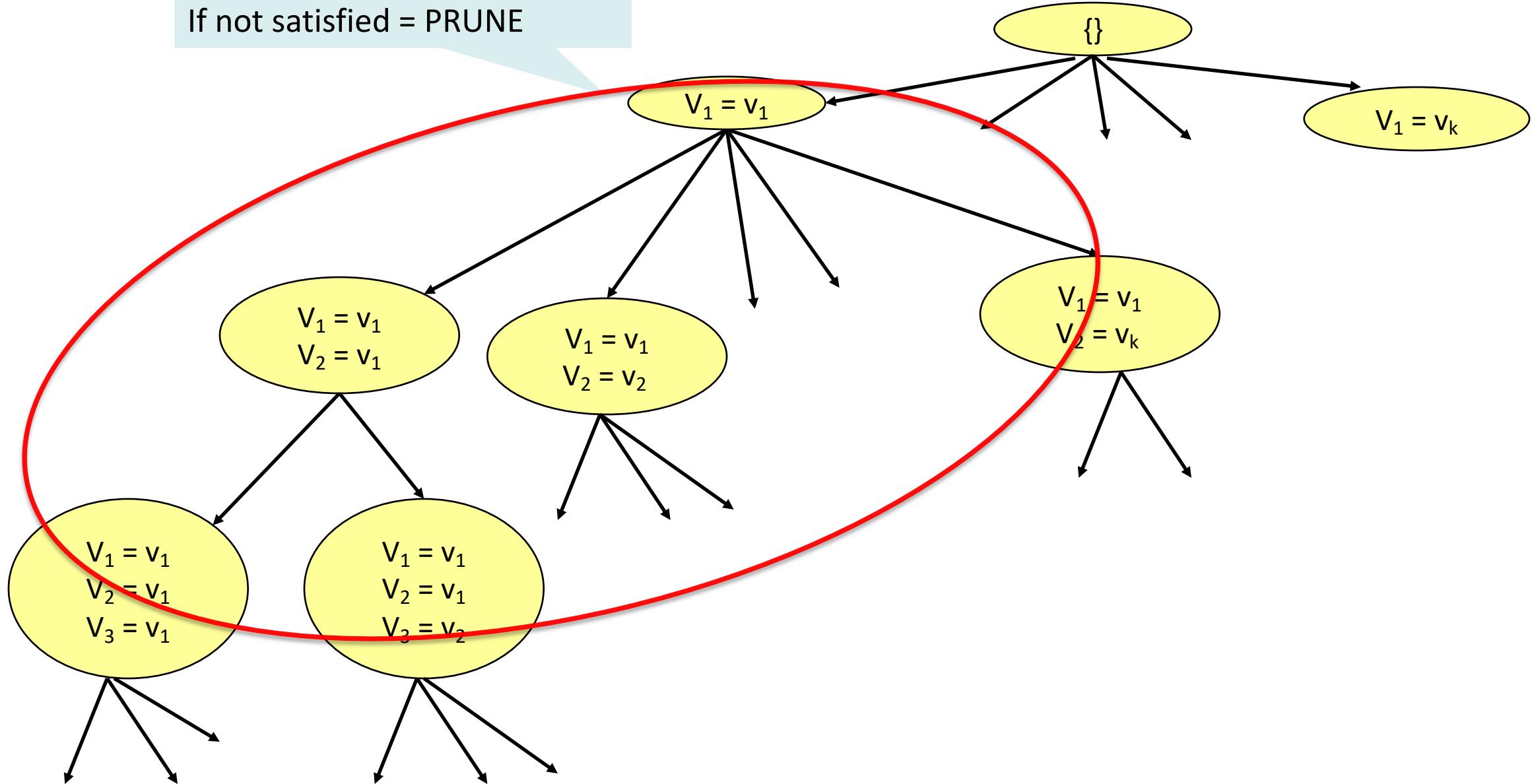
# Backtracking Search

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- Backtracking search is the basic uninformed algorithm for solving CSPs
- Idea: Check constraints as you go
  - I.e. consider only values which do not conflict previous assignments
  - Might have to do some computation to check the constraints
  - “Incremental goal test”
- Depth-first search with this improvement is called *backtracking search*
- Can solve n-queens for  $n \approx 25$

# Backtracking Example

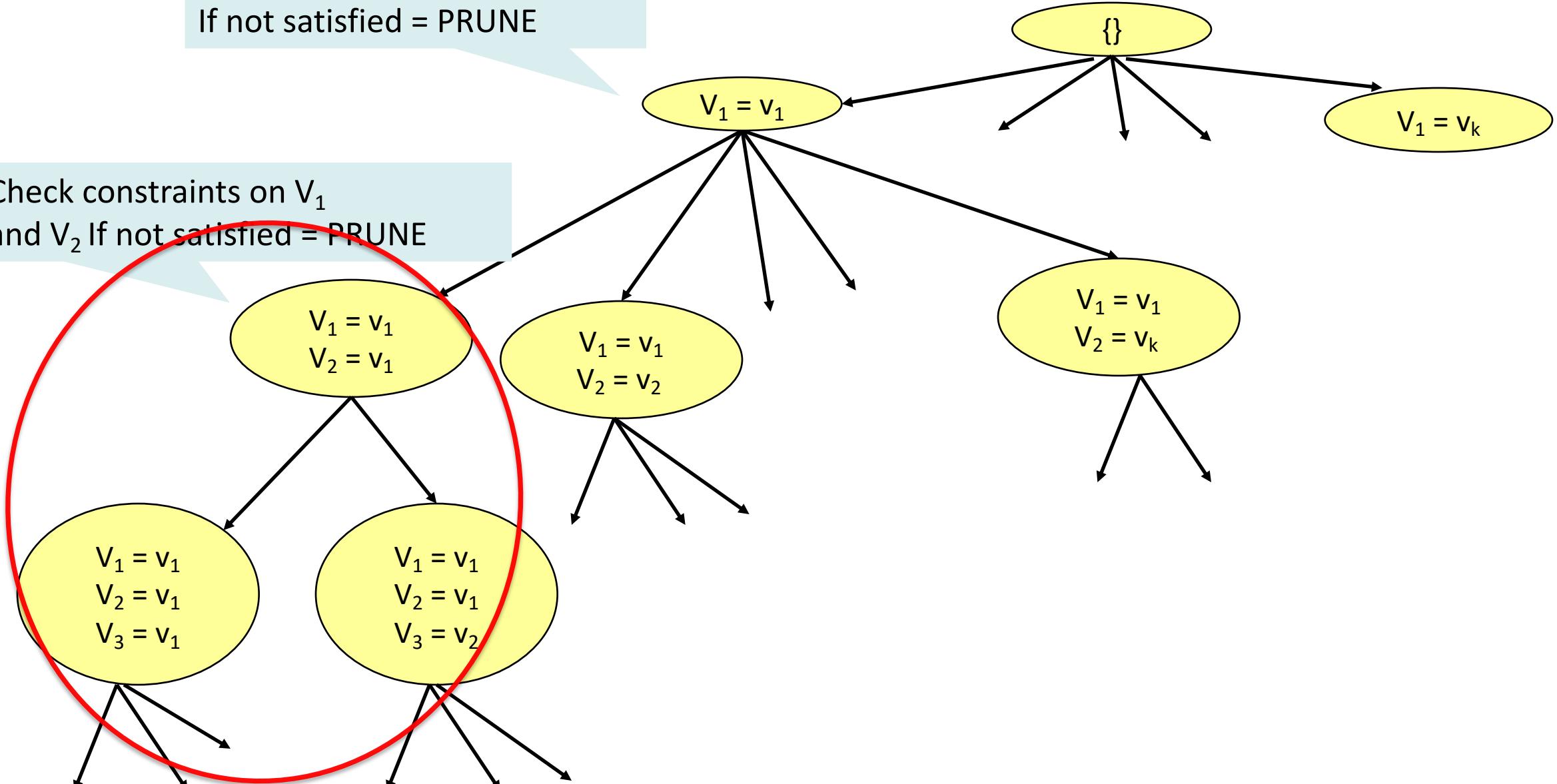
Check unary constraints on  $V_1$   
If not satisfied = PRUNE



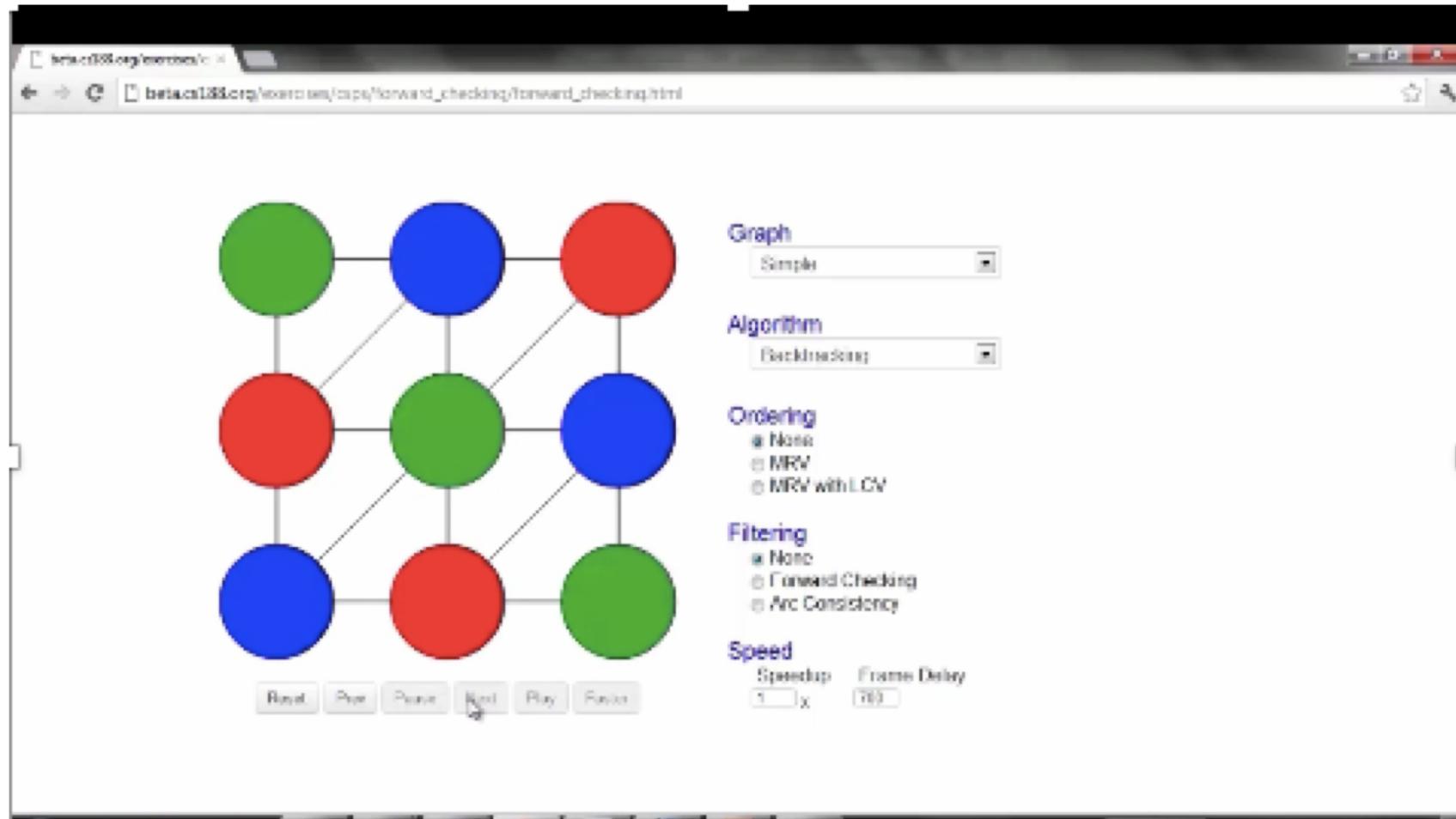
# Backtracking Example

Check unary constraints on  $V_1$   
If not satisfied = PRUNE

Check constraints on  $V_1$   
and  $V_2$  If not satisfied = PRUNE



# Video of Demo Coloring – Backtracking



# Improving Backtracking

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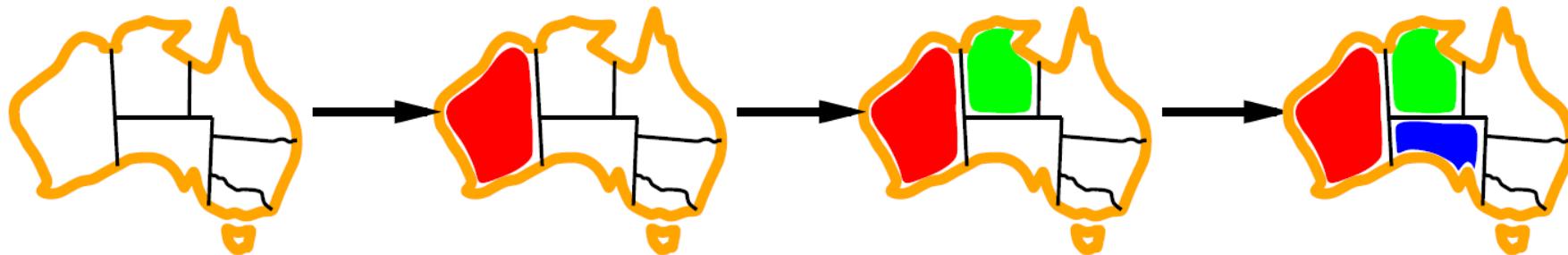
- General-purpose ideas give huge gains in speed
- Ordering:
  - Which variable should be assigned next?
  - In what order should its values be tried?
- Filtering: Can we detect inevitable failure early?

# Ordering

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# Ordering: Minimum Remaining Values

- **Variable Ordering: Minimum Remaining Values (MRV):**
  - Choose the variable with the fewest legal values left in its domain



- Why min rather than max?
- Also called “most constrained variable”
- “Fail-fast” ordering

# Ordering: Least Constraining Value

- **Value Ordering: Least Constraining Value**
  - Given a choice of variable, choose the *least constraining value*
  - I.e., the one that rules out the fewest values in the remaining variables
  - Note that it may take some computation to determine this!
- **Why least rather than most?**
  - Tries to avoid failure by assigning values that leave maximal flexibility for the remaining variables.
  - Combining these ordering ideas makes 1000 queens feasible



# Filtering

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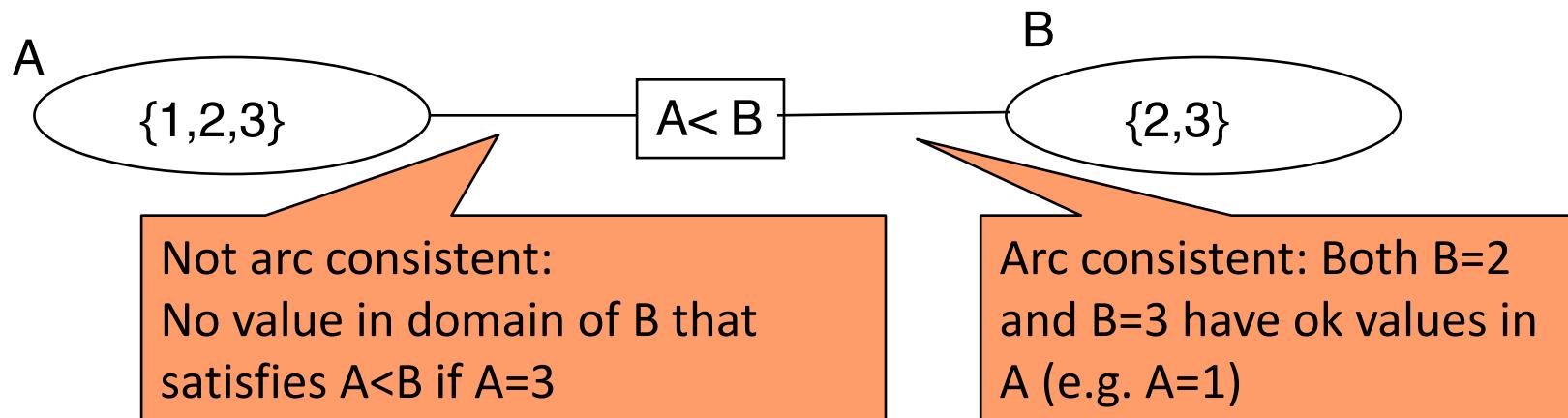
Can we detect inevitable failure early?

# Arc Consistency

Definitions:

An arc  $\langle X, c(X,Y) \rangle$  is arc consistent if for each value  $x$  in  $\text{dom}(X)$  there is some value  $y$  in  $\text{dom}(Y)$  such that  $c(X,Y)$  is satisfied.

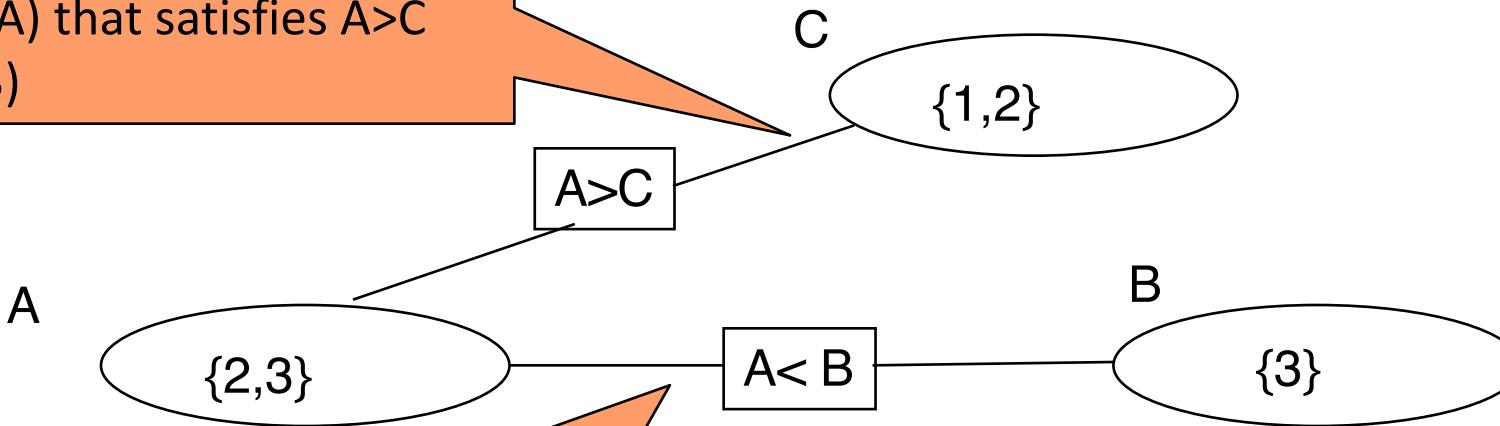
A network is arc consistent if all its arcs are arc consistent.



# Arc Consistency

Arc consistent:

For each value in  $\text{dom}(C)$ , there is one in  $\text{dom}(A)$  that satisfies  $A > C$  (namely  $A=3$ )

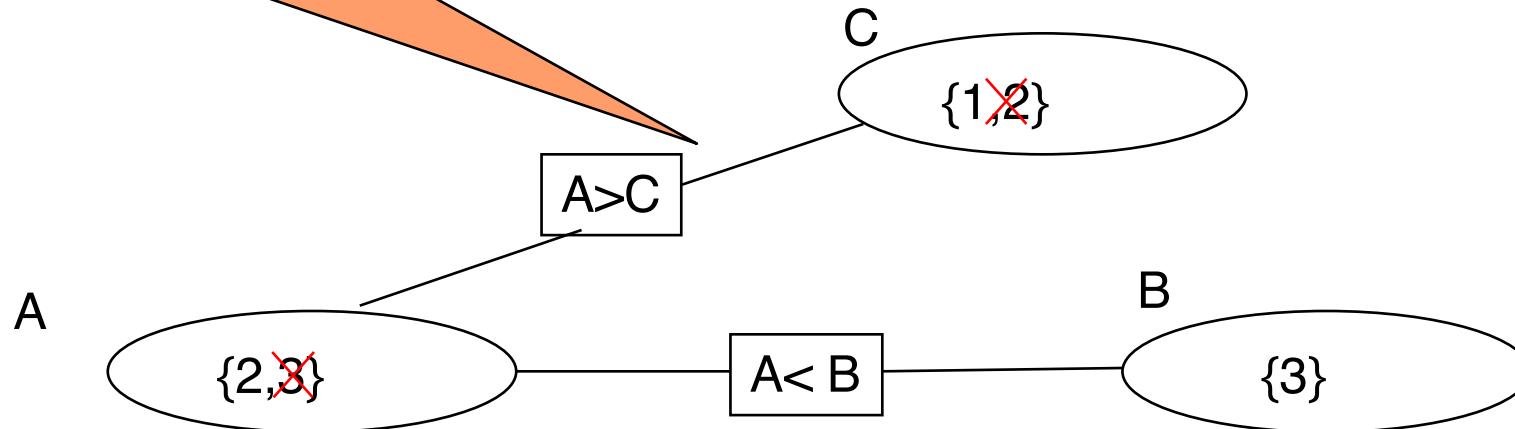


Not arc consistent:

No value in domain of B that satisfies  $A < B$  if  $A=3$

# Arc Consistency

Not arc consistent anymore:  
For C=2, there is no value in  
 $\text{dom}(A)$  that satisfies  $A > C$



# Arc Consistency: High Level Strategy

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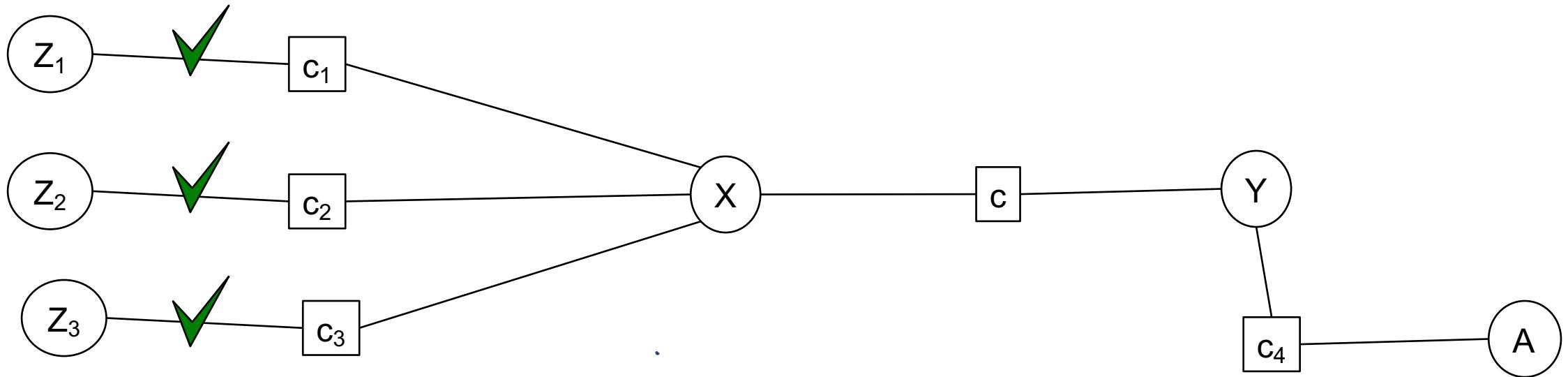
- Consider the arcs in turn, making each arc consistent
  - Reconsider arcs that could be made inconsistent again by this pruning of the domains
- Eventually reach a ‘fixed point’: all arcs consistent
- Run ‘Simple Problem 1’ in Alspace for an example:



- Arc Consistency can be run as a preprocessor or after each assignment.

# Arc Consistency Quiz

- When we reduce the domain of a variable  $X$  to make an arc  $\langle X, c \rangle$  arc consistent, which arcs do we need to reconsider?



- You do not need to reconsider other arcs
  - If arc  $\langle Y, c \rangle$  was arc consistent before, it will still be arc consistent
  - If an arc  $\langle X, c' \rangle$  was arc consistent before, it will still be arc consistent
  - Nothing changes for arcs of constraints not involving  $X$

# Arc Consistency Algorithm: Interpreting Outcomes

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- Three possible outcomes  
(when all arcs are arc consistent):
  - Each domain has a single value
    - We have a (unique) solution.
  - At least one domain is empty
    - No solution! All values are ruled out for this variable.
  - Some domains have more than one value
    - There may be a solution, multiple ones, or none
    - Need to solve this new CSP (usually simpler) problem:  
same constraints, domains have been reduced

# Arc Consistency Quiz

Can we have an arc consistent network with non-empty domains that has no solution?

YES



NO

- Example: vars A, B, C with domain {1, 2} and constraints  $A = B$ ,  $B = C$ ,  $A \neq C$
- Or see Alspace CSP applet “Simple Problem 2”

# Domain Splitting (or Case Analysis)

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- Arc consistency ends: Some domains have more than one value → may or may not have a solution
  - A. Apply Backtracking Search or
  - B. Split the problem in a number of disjoint cases:

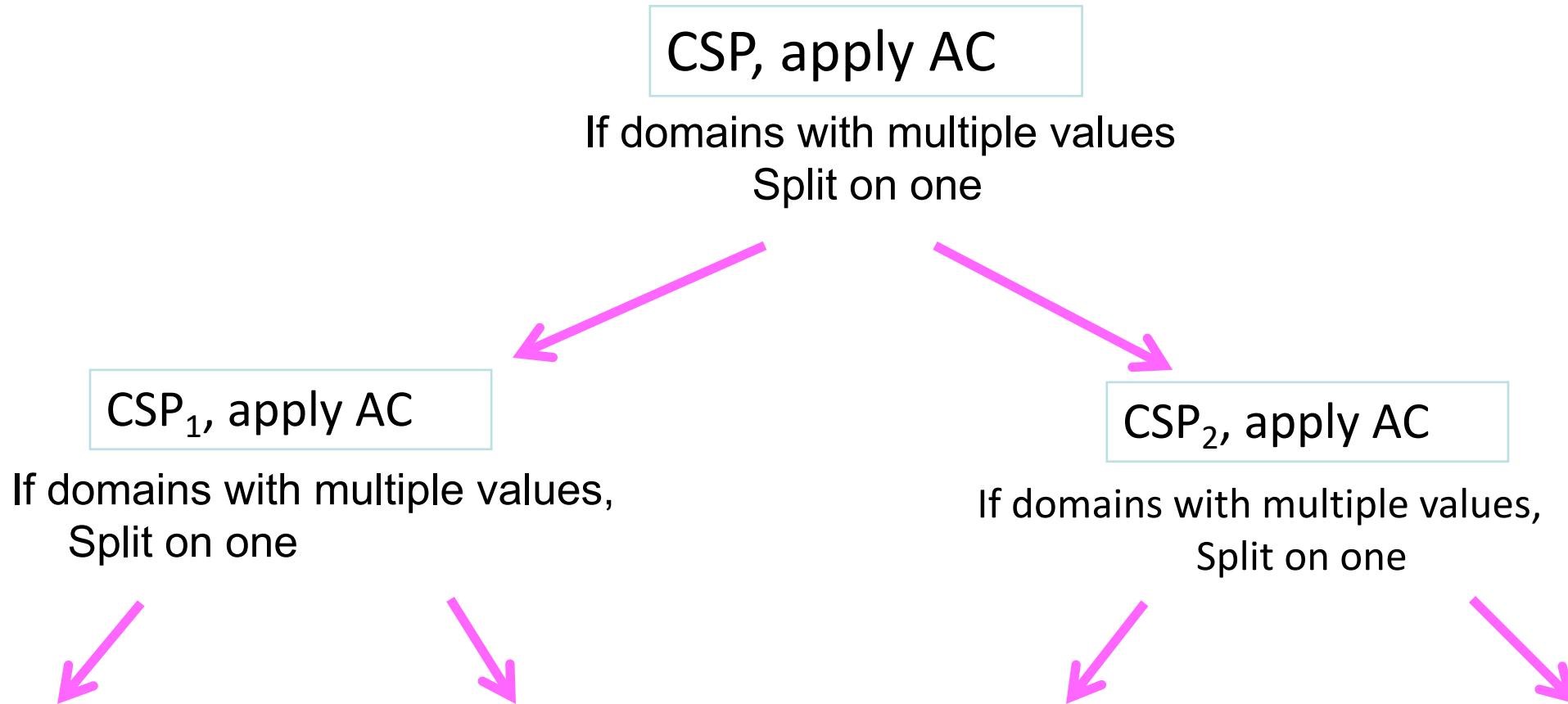
CSP with  $\text{dom}(X) = \{x_1, x_2, x_3, x_4\}$  becomes

$\text{CSP}_1$  with  $\text{dom}(X) = \{x_1, x_2\}$  and

$\text{CSP}_2$  with  $\text{dom}(X) = \{x_3, x_4\}$

- Solution to CSP is the **union** of solutions to  $\text{CSP}_i$

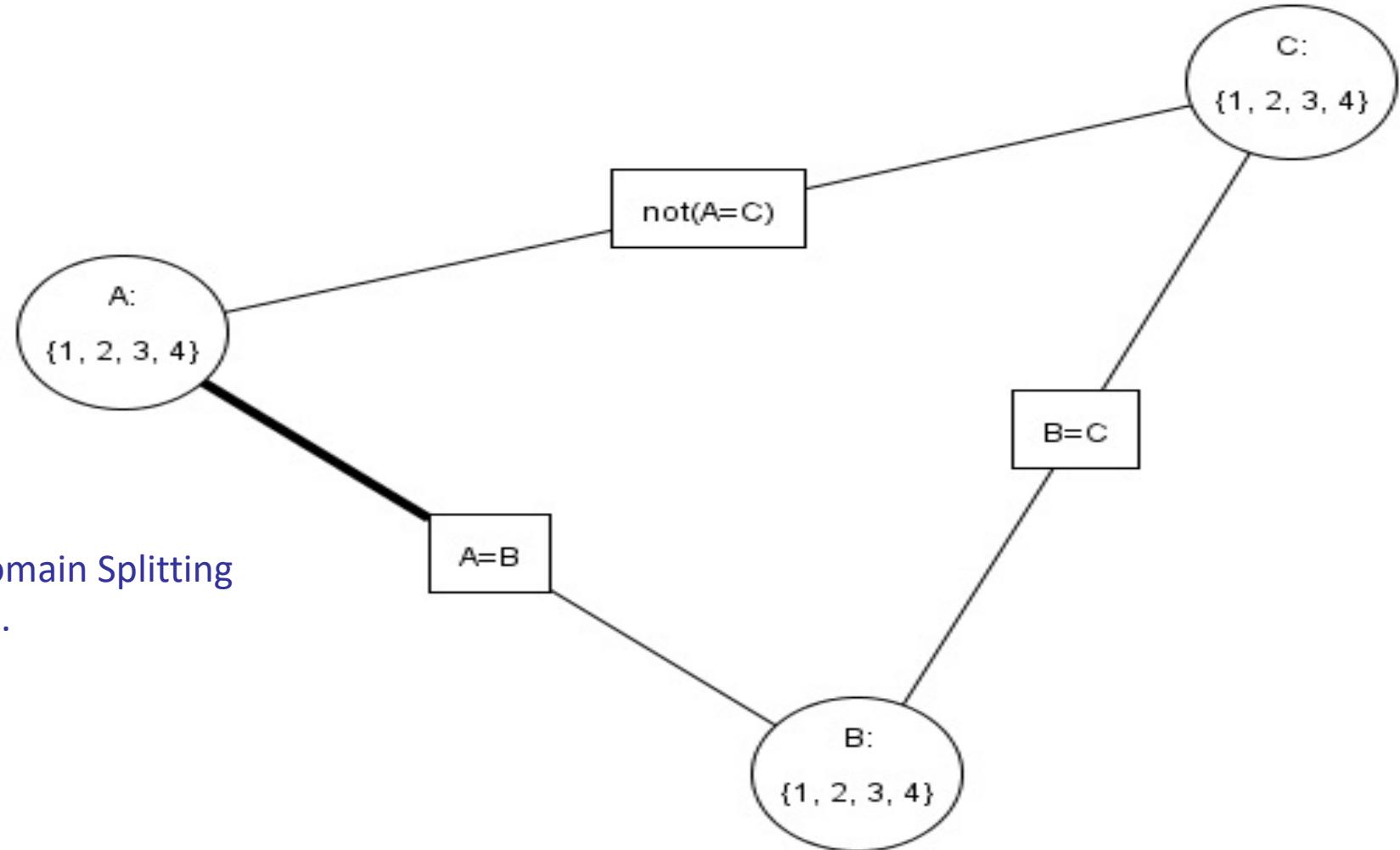
# Searching by Domain Splitting



# Example: ‘Simple Problem 2’ in Alspace

- 3 variables: A, B, C
- Domains: all {1,2,3,4}
- $A=B$ ,  $B=C$ ,  $A \neq C$

• Trace Arc Consistency + Domain Splitting  
for this network in Alspace.



# Example: ‘Simple Problem 2’ in Alspace

3 variables: A, B, C

Domains: all  $\{1,2,3,4\}$

$A=B$ ,  $B=C$ ,  $A \neq C$

$(\{1,2,3,4\}, \{1,2,3,4\}, \{1,2,3,4\})$

$\downarrow$  AC (arc consistency)

$(\{1,2,3,4\}, \{1,2,3,4\}, \{1,2,3,4\})$

$A \in \{1,3\}$

$A \in \{2,4\}$

$(\{1,3\}, \{1,2,3,4\}, \{1,2,3,4\})$

$(\{2,4\}, \{1,2,3,4\}, \{1,2,3,4\})$

$\downarrow$  AC

$\downarrow$  AC

$(\{1,3\}, \{1,3\}, \{1,3\})$

$(\{2,4\}, \{2,4\}, \{2,4\})$

$B \in \{1\}$

$B \in \{3\}$

$B \in \{2\}$

$B \in \{4\}$

$(\{1,3\}, \{1\}, \{1,3\})$

$(\{1,3\}, \{3\}, \{1,3\})$

$(\{2,4\}, \{2\}, \{2,4\})$

$(\{2,4\}, \{4\}, \{2,4\})$

$\downarrow$  AC

$\downarrow$  AC

$\downarrow$  AC

$\downarrow$  AC

$(\{\}, \{\}, \{\})$

$(\{\}, \{\}, \{\})$

$(\{\}, \{\}, \{\})$

$(\{\}, \{\}, \{\})$

No solution

No solution

No solution

No solution

# Example: ‘Simple Problem 2’ in Alspace

3 variables: A, B, C

Domains: all  $\{1,2,3,4\}$

$A=B$ ,  $B=C$ ,  $A=C$

$(\{1,2,3,4\}, \{1,2,3,4\}, \{1,2,3,4\})$

$\downarrow$  AC (arc consistency)

$(\{1,2,3,4\}, \{1,2,3,4\}, \{1,2,3,4\})$

$A \in \{1,3\}$

$A \in \{2,4\}$

$(\{1,3\}, \{1,2,3,4\}, \{1,2,3,4\})$

$(\{2,4\}, \{1,2,3,4\}, \{1,2,3,4\})$

$\downarrow$  AC

$\downarrow$  AC

$(\{1,3\}, \{1,3\}, \{1,3\})$

$(\{2,4\}, \{2,4\}, \{2,4\})$

$B \in \{1\}$

$B \in \{3\}$

$B \in \{2\}$

$B \in \{4\}$

$(\{1,3\}, \{1\}, \{1,3\})$

$(\{1,3\}, \{3\}, \{1,3\})$

$(\{2,4\}, \{2\}, \{2,4\})$

$(\{2,4\}, \{4\}, \{2,4\})$

$\downarrow$  AC

$\downarrow$  AC

$\downarrow$  AC

$\downarrow$  AC

$(\{1\}, \{1\}, \{1\})$

$(\{3\}, \{3\}, \{3\})$

$(\{2\}, \{2\}, \{2\})$

$(\{4\}, \{4\}, \{4\})$

Solution

Solution

Solution

Solution

# A quick note on local search for CSPs

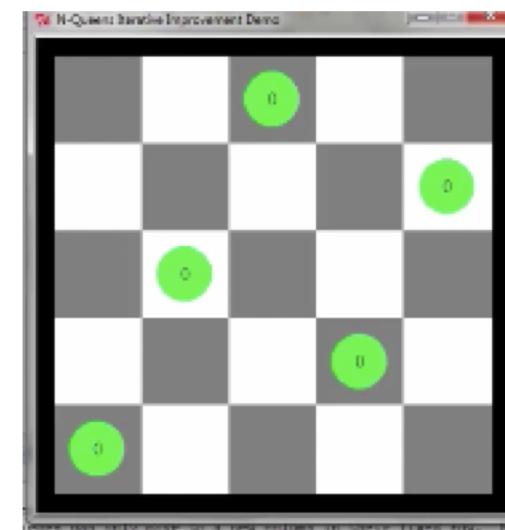
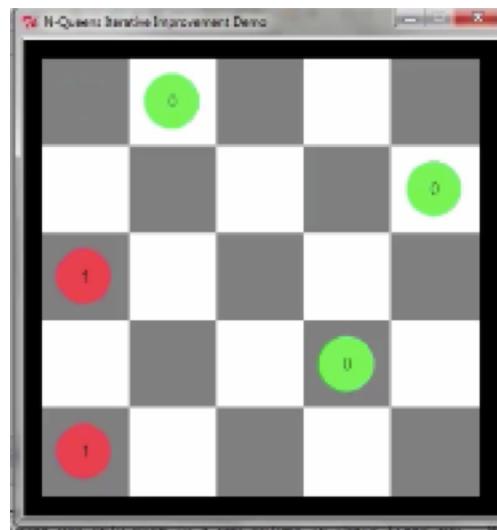
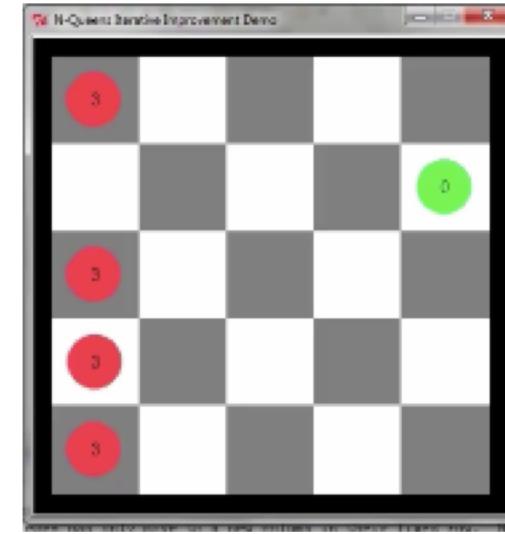
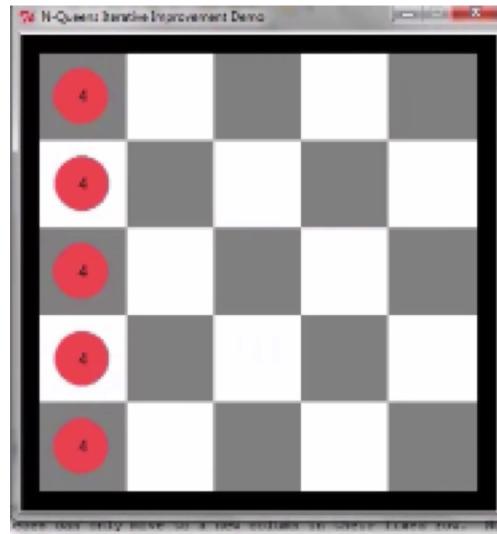
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# Greedy Descent with Min-Conflict Heuristic

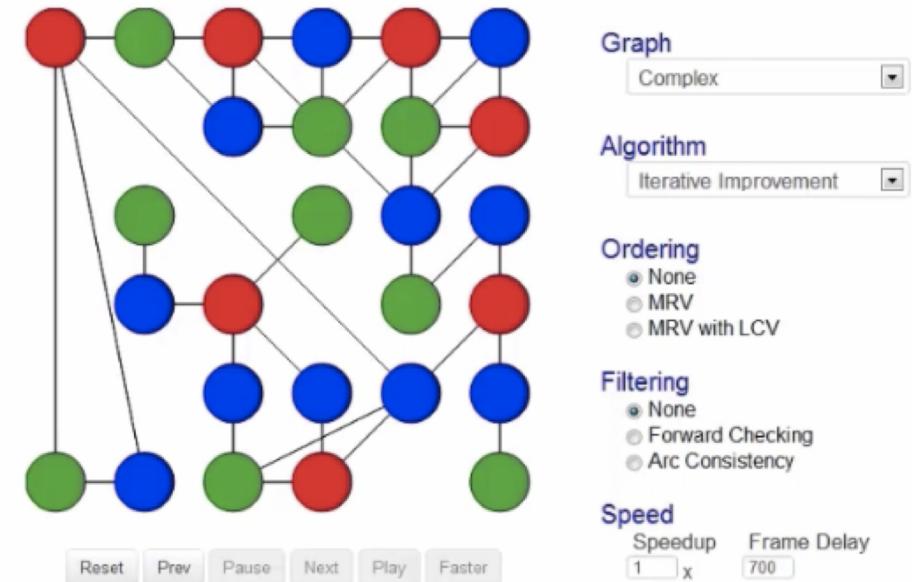
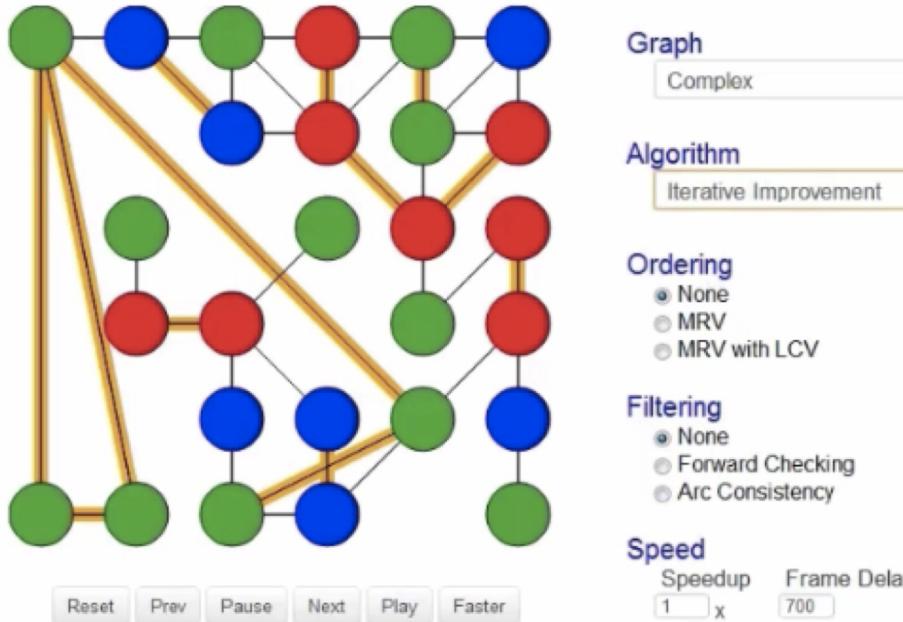
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- Algorithm: While not solved,
  - Variable selection: ***randomly select*** any conflicted variable
  - Value selection: min-conflicts heuristic:
    - Choose a value that violates the fewest constraints
    - Break ties randomly
- One of the best local search techniques for CSP solving
- Given random initial state, can solve n-queens in almost constant time for arbitrary n with high probability (e.g., n = 10,000,000)!

# Demo Min-Conflict in N-Queens



# Demo Min-Conflict – Coloring



# Constraint Optimization Problems

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- Constraint Satisfaction Problems
  - Hard constraints: need to satisfy all of them
  - All models are equally good
- Constraint Optimization Problems
  - Hard constraints: need to satisfy all of them
  - Soft constraints: need to satisfy them as well as possible
  - Can have weighted constraints
    - Minimize  $h(n) = \text{sum of weights of constraints unsatisfied in } n$
    - Hard constraints have a very large weight
    - Some soft constraints can be more important than other soft constraints → larger weight
  - All local search methods work just as well for constraint optimization
    - all they need is an evaluation function  $h$

# Example for Constraint Optimization Problem

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## Exam scheduling

- Hard constraints:
  - Cannot have an exam in too small a room
  - Cannot have multiple exams in the same room in the same time slot
  - ...
- Soft constraints
  - Students should not have multiple exams on the same day
  - It would be nice if students had their exams spread out
  - ...

# Local Search Generality: Dynamically Changing CSPs

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- The problem may change over time
  - Particularly important in scheduling
  - E.g., schedule for airline:
    - Thousands of flights and thousands of personnel assignments
    - A storm can render the schedule infeasible
- Goal: Repair the schedule with minimum number of changes
  - Often easy for Local Search starting from the current schedule
  - Other techniques usually:
    - Require more time
    - Might find solutions requiring many more changes

# Summary: CSPs

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- CSPs are a special kind of search problem:
  - States are partial assignments
  - Goal test defined by constraints
- Basic solution: backtracking search
- Speed-ups:
  - Ordering
  - Filtering
- Greedy descent with min-conflict heuristic is often effective in practice

# Reading

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- Read Chapters 6, 6.1, 6.2, 6.2.1, 6.2.2, 6.2.6, 6.3, 6.3.1, 6.4 in the AIMA textbook
- Read Chapters 4.6 in the P&M textbook (Domain Splitting)