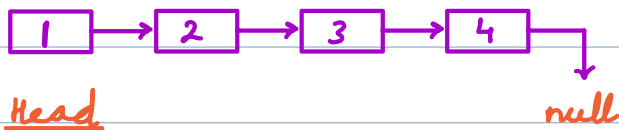


Arrays / Dynamic Array



Linked List → 1) Linear data structure.

2) Need not have continuous memory allocation.

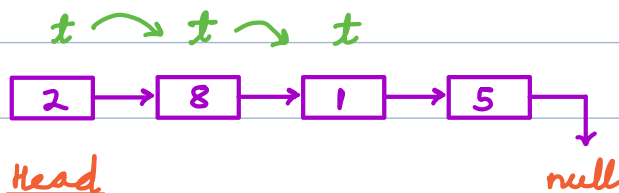


```
class Node {  
    int data; // any datatype  
    Node next;  
    Node (int x) {  
        data = x;  
        next = null;  
    }  
}
```

Operations on Linked List

1) Access K^{th} element ($K=0 \rightarrow$ first data)

Array → $A[K]$ $TC = O(1)$



Node temp = Head

for $i \rightarrow 1$ to K {

temp = temp.next

}

return temp.data

TC = $O(K)$

2) Check if x is present (Searching)

Array \rightarrow Linear Search TC = $O(N)$
sorted \rightarrow Binary Search TC = $O(\log N)$

Node temp = Head

while (temp != null) {

if (temp.data == x) return true

temp = temp.next

}

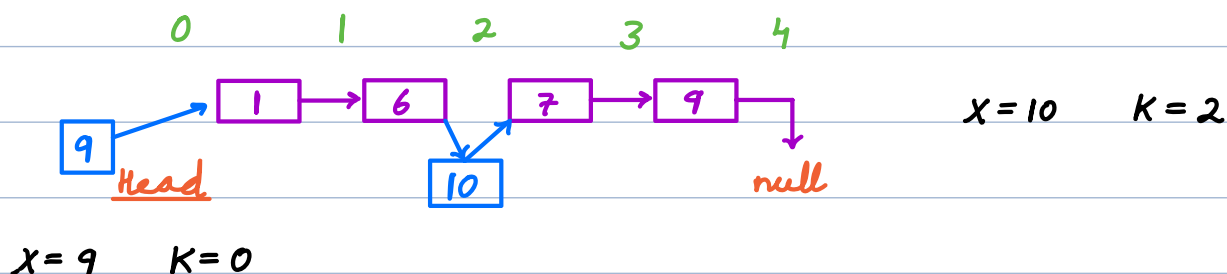
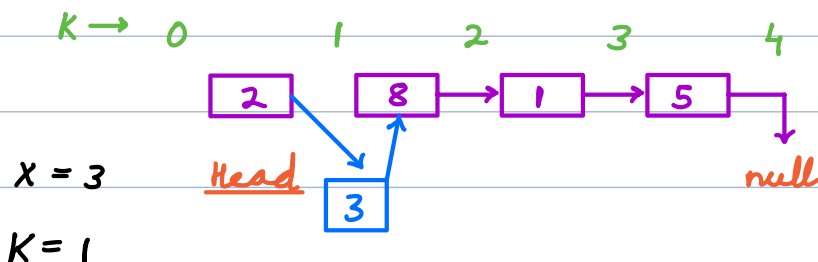
return false

TC = $O(N)$

Cannot use BS $\rightarrow \because$ we cannot jump to middle element.

3) Insert a node with data x at position K (0 \rightarrow start)

$[0 \leq K \leq N]$



$xr = \text{new Node}(x)$

if $(K == 0)$ {

$xr.\text{next} = \text{Head}$

$\text{Head} = xr$

return Head

}

$\text{temp} = \text{Head}$

for $i \rightarrow 1$ to $(K-1)$ {

$\text{temp} = \text{temp}.\text{next}$

}

$xr.\text{next} = \text{temp}.\text{next}$

$\text{temp}.\text{next} = xr$

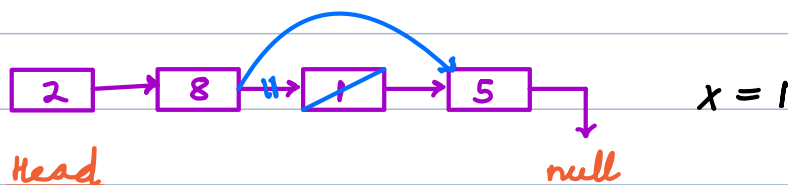
return Head



$TC = \underline{O(K)}$

4) Delete the first occurrence of x .

(If not present \rightarrow leave as it is)



Cases \rightarrow 1) Empty list ($\text{Head} == \text{null}$)

nothing to do.

2) $(\text{Head}.\text{data} == x)$

$\text{Head} = \text{Head}.\text{next}$

$5 \rightarrow 4 \rightarrow 7 \xrightarrow{t} \cancel{1} \rightarrow \text{null}$
 t (t.next) \Rightarrow error

```

if (Head == null) return Head
if (Head.data == x) {
    Head = Head.next
    return Head
}

```

```

temp = Head
while (temp.next != null) {
    if (temp.next.data == x) {
        temp.next = temp.next.next
        break
    }
    temp = temp.next
}
return Head

```

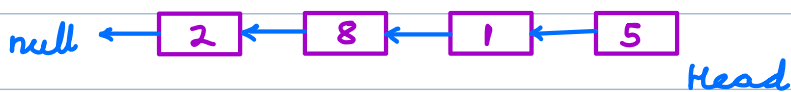
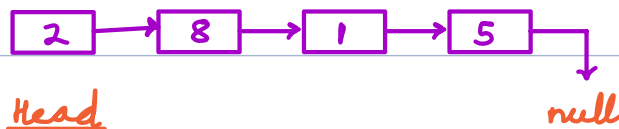
TC = $O(N)$

H.W \rightarrow Delete all occurrences of x .

K = 2

$1 \rightarrow \cancel{2} \rightarrow \cancel{2} \rightarrow \cancel{2} \rightarrow 5 \rightarrow 1 \rightarrow \text{null}$

Q \rightarrow Reverse the given linked list.

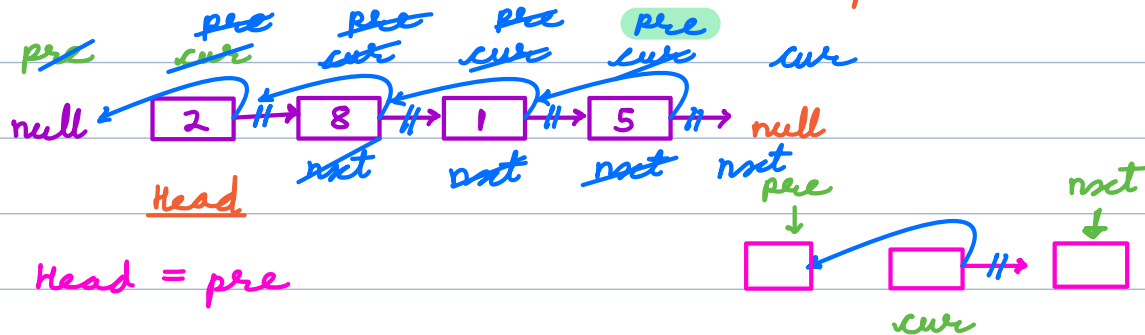


Case 2

1) Empty list \rightarrow if (Head == null) return Head

2) Single Node \rightarrow if (Head.next == null) return Head

3) Else \rightarrow Reverse the links & update Head



Head = pre

cur = Head

pre = null

while (cur != null) {

next = cur.next

cur.next = pre

pre = cur

cur = next

}

TC = $O(N)$

SC = $O(1)$

Head = pre

Q \rightarrow Check if the given linked list is palindrome.

2 \rightarrow 5 \rightarrow 8 \rightarrow 7 \rightarrow 2 \rightarrow null
Head

Ans = false

2 \rightarrow 5 \rightarrow 8 \rightarrow 5 \rightarrow 2 \rightarrow null
Head

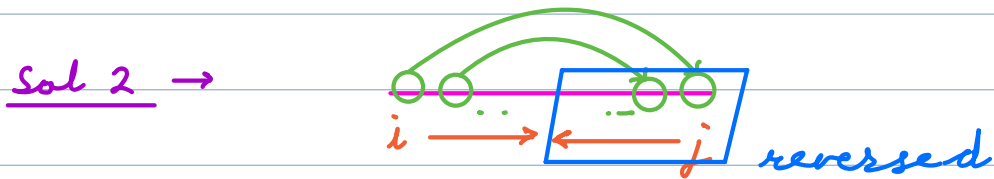
Ans = true

1 → null Ans = true
Head

Sol 1 →

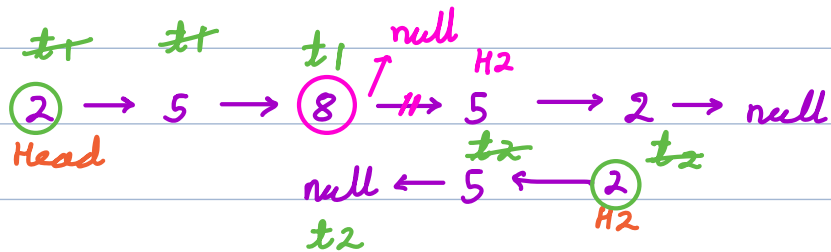
- 1) create a copy list
- 2) Reverse the copy
- 3) compare copy with original.

} $TC = O(N)$
 $SC = O(N)$



- 1) Find middle element
- 2) Reverse second half
- 3) compare reverse with first half

} $TC = O(N)$
 $SC = O(1)$



Ans = true

len = 0 temp = Head

while (temp != null) {

len ++

temp = temp.next

}