

Quene - f_ r

initially, x = -1 f = 0erqueue (10) erqueue (5) void enqueue (x) & int dequeue () & erqueue (3) if (is Empty ()) dequeue () return -1 is Empty () x = A43 f++ boolear is Empty () {
return f > r
} return x Underflow → Empty DS Overflow - completely full DS > Do not allow more elements to insert. → Use dyramic array V 2) Using Kinked List (effective use of memory writ arrays) erqueue (10) Head Tail erqueve (5) erqueue (3) dequeue () initially, Head = Tail = null is Empty () 1) erqueue (x) - insert at tail node 2> dequeue () → delete head Scheck 3) is Empty () → return (Head == null) 4) front () -> return Kead.data underflow case 5) rear () → return tail. data

X X X 3 7 11 20

erqueue (4) / dequeve () / erqueue (9) /
erqueue (3) / erqueue (7) / erqueue (11) /
erqueue (20) / dequeve () /

D→ Implement queue using 2 stacks.

(push (x) is Empty ()
(pop () peak () / top ()

Front

(P37 7 || 20)
(P3 7 W 20)
(Park (x) is Empty ()
(Park (x) is E

erqueve (4) / dequeve () / erqueve (9) /
erqueve (3) / erqueve (7) / erqueve (11) /
erqueve (20) / dequeve () + 5 entrs steps
dequeve () * 5 in TC = O(1)

void erqueue (x) { boolear is Empty () {

| st |. push (x) | return st |. is Empty ()

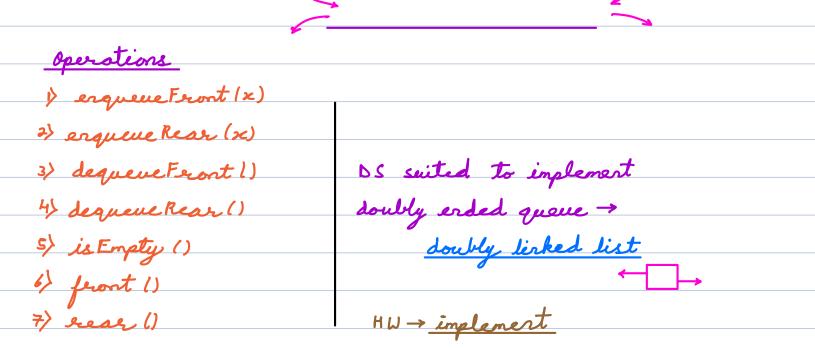
}

ll et 2. is Empty ()

int dequeve () & void more () { while (! st 1. is Empty ()) if (is Empty ()) st2. push (st1. pop()) return -1 if (st2. is Empty ()) more () return st2. pop () more K elements → TC = O(K) dequeue K items -> TC = O(K) removal -> (more + dequeue) $\Rightarrow 7c = o(2) = \underline{o(1)}$ 0 → Giver ar isteger N, fird Nth number formed by only digits 12. (Perfect number)

if (N <= 2) return N q. erqueue (1) q. enqueue (2) while (N >= 0) { while (i <= N) { d = (N% 2) + 1res = d + res // oppend z = q. dequeue () N = N/2 - 1a = 2 * 10 + 1 b = a + 1 // x * 10 + 2TC = O(log(N)) SC = O(1)if (i == N) return a if (i+1 == N) return b g. erqueue (a) TC = O(N) SC = O(N)q. erqueue (b) i += 2

Doubly Erded Queue



Q→ liver or integer array & ar integer K. Fird max element & substrays of size K.

$$A = \begin{bmatrix} 1 & 8 & 5 & 6 & 7 & 4 & 2 & 0 & 3 \end{bmatrix} \quad K = 4$$

$$Ans \rightarrow 8 \quad 8 \quad 7 \quad 7 \quad 4$$

$$A = \begin{bmatrix} 1 & 4 & 3 & 2 & 5 & 1 \\ 4 & 3 & 2 & 5 & 1 & K = 3 \\ 4 & 4 & 5 & 5 & 6 & 6 \\ 4 & 4 & 5 & 6 & 6 & 6 \\ 4 & 4 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 & 6 & 6 \\ 4 & 6 & 6 & 6 \\ 4$$

Fixed size substray > sliding wirdow

$$A = \begin{bmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 1 & 8 & 5 & 6 & 7 & 4 & 2 & 0 & 3 \end{bmatrix} \quad K = 4$$

/ oulside window → remove

✓ if greater on right > current cannot be arswer.

for i \rightarrow 0 to (K-1) {

while (! q. is Empty 22 A[q. rest ()] < A[i])

q. dequeneRest ()

}

q. erqueneRest (i)
}

```
print (A[q. front ()])

for i → k to (N-1) {

while (!q. is Empty 22 A[q. rear ()] < A[i7)

q. dequeueRear ()

}

q. erqueve Rear (i)

if (q. front () == i-k)

q. dequeueFront ()

print (A[q. front ()])

}
```

TC = O(N) SC = O(N)