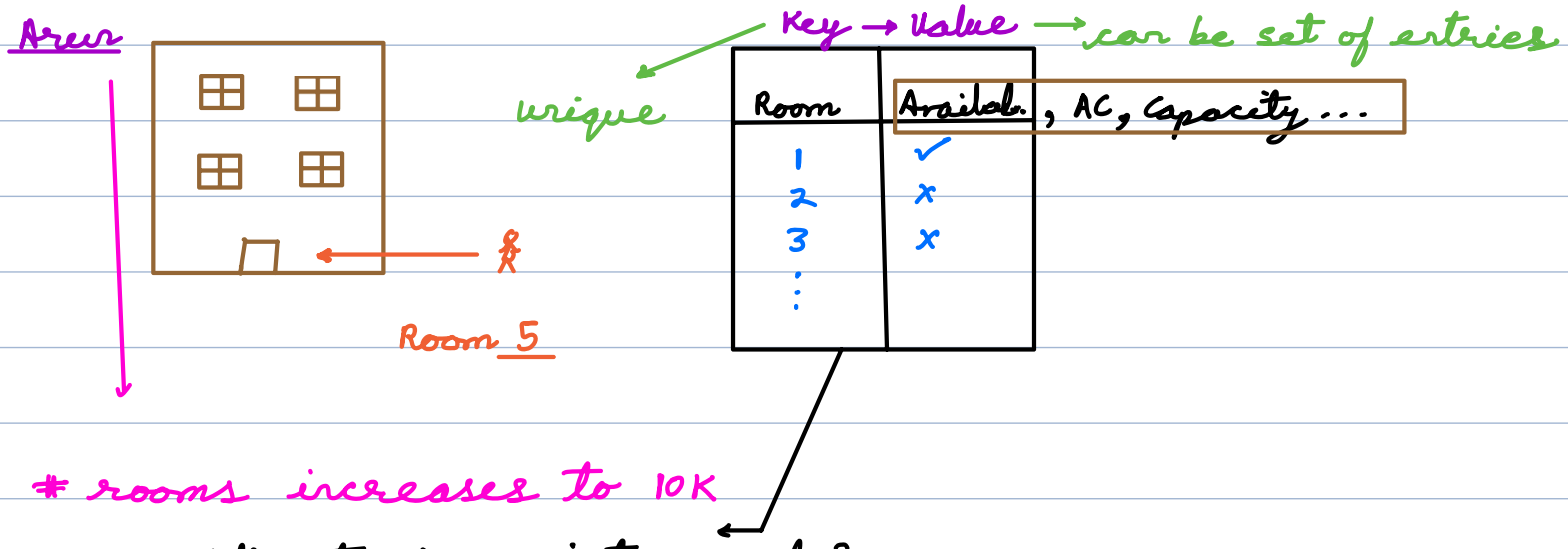

rooms increases to 10K

Will data in register work ?

No → We need to use tech.

HashMap stores data → <Key, Value>

(internal implementation → later)

Store population of every country

Hashmap < String, Long >

< Country, Population >

Store # states of every country

Hashmap < String, Int >

Country, #states

Store the name of all states of every country.

Hashmap < String, List < String >>

country list of all states

Population of each state of every country

Hashmap < String, Hashmap < String, Larg >>

country state, population

Hashset <Key> → Used to store unique data.

Functionalities

HashMap

1) insert (key, value) → put(K, V)
2) size()
3) delete (key)

} TC = O(1)

4) update (key, value) → similar to second insertion
∵ for same key value will be overridden.

5) search (key)

Hashset

1) insert (Key)
2) size()
3) delete (key)
4) search (key)

} TC = O(1)

Java	C++	Python	JS	C#
Hashmap	unordered_map	dictionary	map	dictionary
Hashset	unordered_set	set	set	Hashset

H.W → check syntax in your language.

Hashmap/Hashset → Order of insertion is not maintained.

Treeset / Treemap → Order is sorted wrt keys.

Linked Hashmap / Linked Hashset → Order of insertion is maintained.

insert → 1 5 8 stored → 1 5 8

Q → Given an integer array & Q queries where we have to find frequency of the given element.

0	1	2	3	4	5	6
2	6	3	8	2	8	2

Query → 2 → 3
6 → 1
5 → 0

Bruteforce → ∀ query, travel & check the frequency. TC = O(Q * N)

Frequency array F[i] = freq. of i

```
for i → 0 to (N-1) {
    F[A[i]] ++
}
```

SC = O(|A[i]|) 1 <= A[i] <= 10^9 (MLE)

Sol → Use hashmap to store data.

```
// hm<int, int>
for i → 0 to (N-1) {
    if (hm.containsKey(A[i]))
        hm.put(A[i], hm.get(A[i]) + 1)
    else
        hm.put(A[i], 1)
```

```
}
for i → 0 to (Q-1) {
    if (hm.containsKey(A[i]))
        print(hm.get(A[i]))
    else
        print(0)
}
```

TC = O(N+Q) SC = O(N)

Q → List of learner id wrt each attempt.
Find learners with least attempts.

id → [101 102 103 101 102 101 105]

least participation → 1
id → {103, 105}

1) Find freq ∀ ids
2) Find least value
3) Keys with least values give list of learners.

Q → Find the first non-repeating element.
(non-repeating → only present once)

N = 6 A = [1 2 (3) 2 5 1]
(Ans)

A = [4 3 3 (2) 5 6 4 5]
(Ans)

idea → use & store frequency

1 ⇒ non-repeating

Sol → ⇒ Store freq ∀ elements

```
// hm <int, int>
for i → 0 to (N-1) {
    if (hm.containsKey(A[i]))
        hm.put(A[i], hm.get(A[i]) + 1)
    else
        hm.put(A[i], 1)
}
```

⇒

```
for i → 0 to (N-1) {  // index traversal ✔
    if (hm.get(A[i]) == 1)
        return A[i]
}
return -1
```

TC = O(N) SC = O(N)

Q → Given an integer array, find the count of distinct elements.

```
   0  1  2  3  4  5  6
A = [3  5  6  5  4  6  5]
```

{3 5 6 4} Ans = 4

A = [3 3 3 3] Ans = 1

```
for i → 0 to (N-1) {
    hs.add(A[i])
}
return hs.size()
```

TC = O(N) SC = O(N)

Q → Given an integer array, check if there is a subarray with sum 0.

	0	1	2	3	4	5	6	7	8	9
A = [	2	2	1	-3	4	3	1	-2	-3	2]

Ans = true

Bruteforce → ∀ subarrays, calculate sum & check.

TC = O(N³) → O(N²) SC = O(1)

sum of subarray → prefix sum

i — j P[j] − P[i−1] = 0

⇒ P[j] = P[i−1]

case 1 → P[i] = 0 ⇒ Ans = true ✓

2 → P[j] = P[i−1] repeating value in P[] ✓

	0	i=1	2	j=3	4	5
A = [	2	2	1	-3	4	3]
P =	(2)	4	5	(2)	6	9
	0			3		

P[0] = A[0]

if (P[0] == 0) return true ✓

hs.add(P[0])

for i → 1 to (N-1) {

A = [2 3 -5 ...]

P → 2 5 0

```
        P[i] = P[i-1] + A[i]
        if (P[i] == 0) return true
        if (hs.contains(P[i])) return true
        hs.add(P[i])
    }
    return false
```

TC = O(N) SC = O(N)

Q → Find the count of subarrays with sum 0.

```
    cnt = 0
    P[0] = A[0]
    if (P[0] == 0) cnt++
    hs.add(P[0])
for i → 1 to (N-1) {
        P[i] = P[i-1] + A[i]
        if (P[i] == 0) cnt++
        if (hs.contains(P[i])) cnt++
        hs.add(P[i])
    }
    return cnt
```

TC = O(N) SC = O(N)

H.W → Find count of subarrays with sum K.
