Heap Sort

0	1	2	3	4	5	6	7
7	3	5	1	9	6	7	1

A = [7 3 5 1 9 6 7 1]

Sol 1 → 1) Create min-heap / max-heap.

2) getMin() / getMax() / element N times & store in arr.

TC = O(N + N log (N)) = O(N log (N))

SC = O(N) → O(1) Inplace Heap Build

i/p A[] ⟶ Heap ⟶ Sorted A[]

Min / Max ✓

0	1	2	3	4	5	6	7
~~7~~ 9	~~3~~ 7	~~5~~ 7	1	~~9~~ 3	6	~~7~~ 5	1

A = [9 7 7 1 3 6 5 1]

lc = A[1] | rc = A[2] | lc = A[3] | rc = A[4] | lc = A[5] | rc = A[6] | lc = A[7] | rc = A[8] x

Structure → Complete Binary Tree

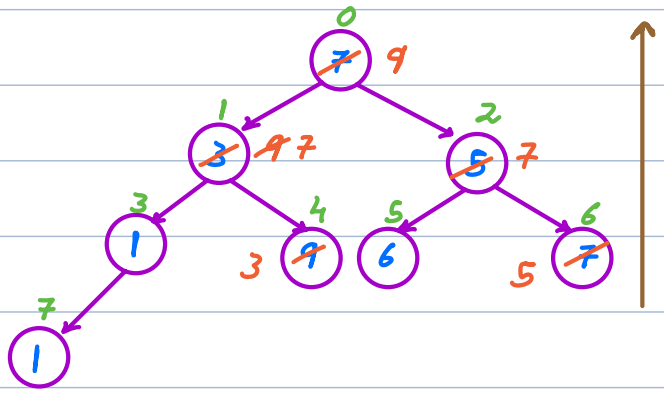
N nodes → # leaves = (N+1)/2

8 → # leaves = 4

i → lc = 2i + 1

rc = 2i + 2

parent = (i-1)/2

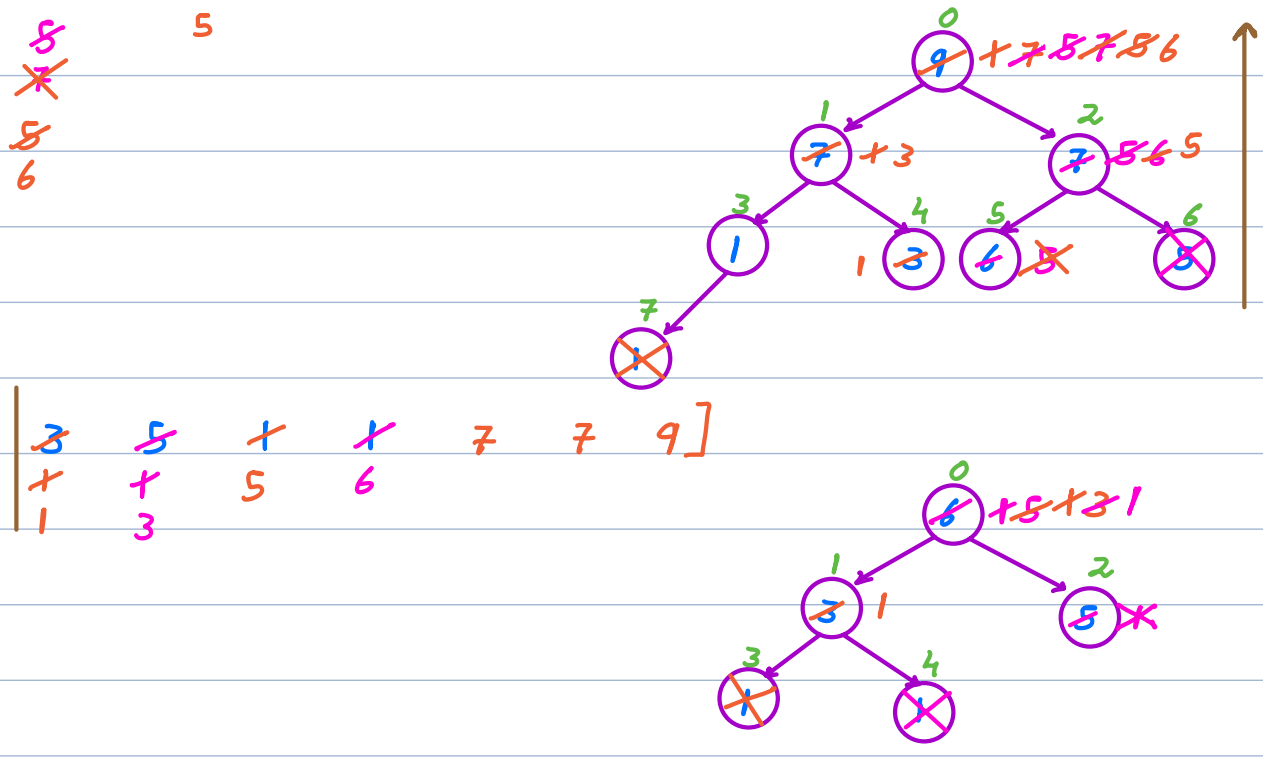


TC = O(N)

Max Heap

0	1	2	3	4	5	6	7
~~1~~ 7	~~7~~ 5	~~7~~ 6	1	~~3~~ 1	~~6~~ 5	~~5~~ 7	~~1~~ 9
7	3	6			7		

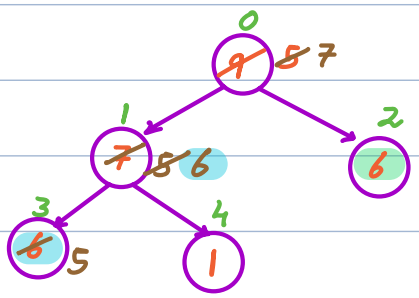
A = [7 5 6 1 1 5 7 9]



A = [1 1 3 5 6 7 7 9] SC = O(1)

TC = O(N log (N))

Unstable sorting



Q → Find K^{th} largest element.

A = [8 5 2 4 1 7] K = 3 Ans = 5

A = [1 2 3 4 5] K = 5 Ans = 1

Sol 1 → sort A[]

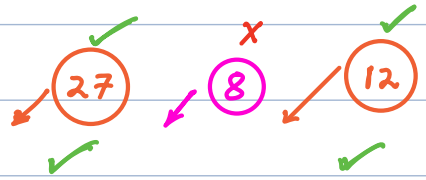
Ans = A[N-K]

TC = O(N log (N))

SC = O(1)

K	Ans
1	A[N-1]
2	A[N-2]
⋮	⋮

Team of 5 players



28 40 ~~15~~ ~~10~~ 30

A = [28 40 15 10 30 27 8 12]

10 X

Q → Find K^{th} largest element ∀ subarray from 0 to i where $i >= (K-1)$.

A = [28 40 15 10 30 27 8 18] K=5

0 1 2 3 4 5 6 7

10 15 15 18

A = [5 4 1 6 7] K=2

0 1 2 3 4

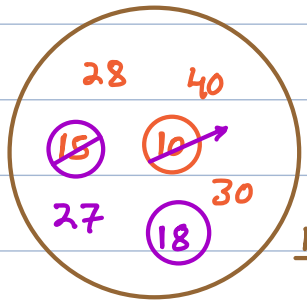
4 4 5 6

K^{th} largest → Min Heap of size K to store K largest elements.

A = [28 40 15 10 30 27 8 18] K=5

0 1 2 3 4 5 6 7

10 15 15 18



Min Heap

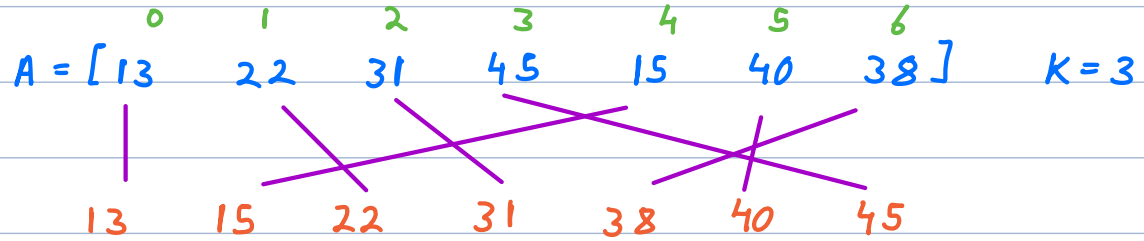
```
// h → minHeap
for i → 0 to (K-1) {
    h.insert(A[i])
}

print(h.root)          // h[0]
for i → K to (N-1) {
    if (h.root < A[i]) {
        h.getMin()     // delete root
        h.insert(A[i])
    }
    print(h.root)
}
```

SC = O(K)

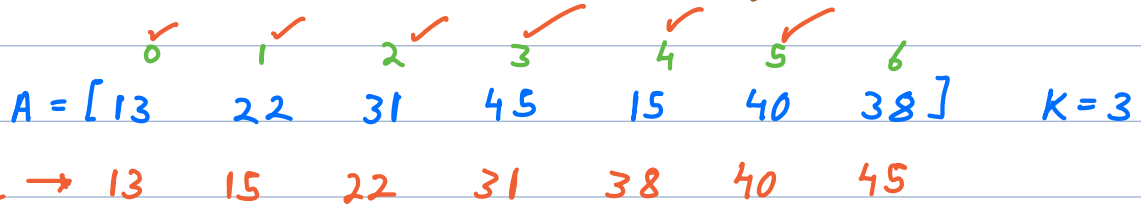
TC = O(N log(K))

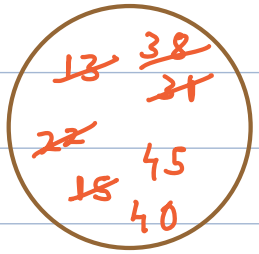
Q → Sort the given array where every element is atmost K distance away from its position in sorted order. (N >> K)



Smallest → (0 ——— K)

use minHeap of size (K+1)





SC = O(K)

TC = O(N log(K))

Q → Given a stream of integers.
Find median ∀ integer intake.

median → middle element in sorted order.

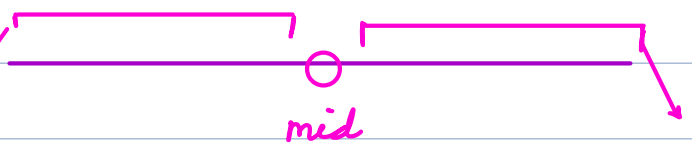
i/p → 9 10 11 ...
o/p → 9 9 10 ...

A = [1 2 4 3]
1 2 3 4
2.5

odd → middle in sorted
even → first middle in sorted. (2)

Bruteforce → ∀ intake, sort & find middle.
TC = O(N * N log(N)) SC = O(1)

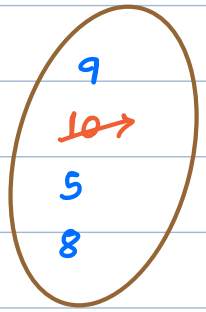
Sol →



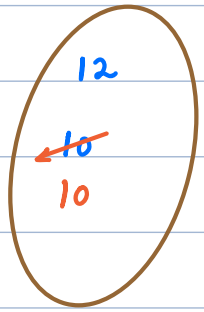
store small values in maxHeap ✓

store large values in minHeap. ✓

Ans = maxHeap root node



Max Heap (small)



Min Heap (large)

i/p → 9 12 10 5 8

o/p → 9 9 10 9 9

```
// i/p → x
if (x <= root of maxHeap) {
    insert in maxHeap
    if (size of maxHeap - size of minHeap > 1)
        move root of maxHeap to minHeap
} else {
    insert in minHeap
    if (size of minHeap - size of maxHeap > 0)
        move root of minHeap to maxHeap
}
```

TC = O(N log(N)) SC = O(N)
