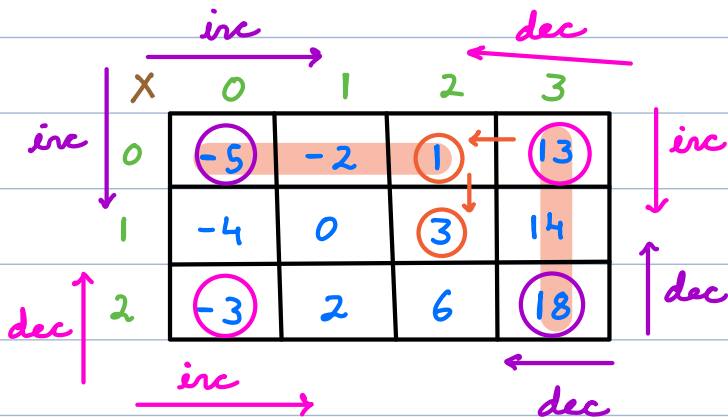


Q → Given a row wise & column wise sorted matrix.  
Find out whether element K is present or not.

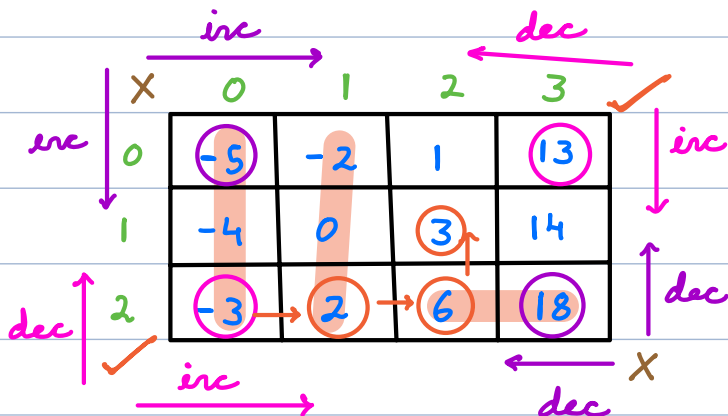
		0	1	2	3
0	-5	-2	1	13	
1	-4	0	3	14	
2	-3	2	6	18	

<u>K</u>	<u>Ans</u>
13	→ true
2	→ true
15	→ false

Brute force → check all cells.  $TC = O(N*M)$   $SC = O(1)$



$K=3 \rightarrow \text{true}$



$K=3 \rightarrow \text{true}$

// top right

$N \rightarrow \# \text{ rows}$   $M \rightarrow \# \text{ col}$

$r = 0$   $c = M - 1$

while ( $r < N$  &&  $c \geq 0$ ) {

if ( $A[r][c] == K$ ) return true

else if ( $A[r][c] < K$ )  $r++$

```

    } else c--
}

```

TC =  $O(N+M)$  SC =  $O(1)$

return false

Q → Given a binary sorted matrix of size  $N \times M$ .  
Find the row with max # 1's.

	0	1	2	3
0	0	0	1	1
1	0	0	0	1
→ 2	0	1	1	1

Ans = 2

Bruteforce →  $\forall$  rows, count the # 1's.

TC =  $O(N \times M)$  SC =  $O(1)$

TC =  $O(N)$

	0	1	2	3	4	5
0	0	0	0	0	1	1
→ 1	0	0	1	1	1	1
2	0	0	0	0	0	0
3	0	0	0	1	1	1
4	0	0	1	1	1	1

~~$r = 0 + 2 + 3 + 4 + 5$~~

~~$c = 5 + 4 + 3 + 2 + 1$~~

ans =  ~~$1 + 0 + 0 + 1$~~  ✓

ans = -1

TC =  $O(N+M)$  SC =  $O(1)$

// top right

$r = 0$   $c = M - 1$

while ( $r < N$  &&  $c \geq 0$ ) {

while ( $c \geq 0$  &&  $A[r][c] == 1$ ) {

$c--$

ans = r

```

    }
    r++
}

```

return ans

Q → Given a square matrix  $(N \times N)$ .

Print **boundary elements** in clockwise direction.



	0	1	2	3
0	1	2	3	4
1	5	6	7	8
2	9	10	11	12
3	13	14	15	16

4x4

o/p → 1 2 3 4 8 12 16  
15 14 13 9 5

	0	1	2
0	1	2	3
1	4	5	6
2	7	8	9

o/p → 1 2 3 6 9 8 7 4

$r = 0$

```

for c → 0 to (N-2) {
    print (A[r][c])
}

```

$c = N-1$

```

for r → 0 to (N-2) {
    print (A[r][c])
}

```

$r = N-1$

```

for c → (N-1) to 1 {
    print (A[r][c])
}

```

$c = 0$

```

for r → (N-1) to 1 {
    print (A[r][c])
}

```

$TC = O(4N) \rightarrow O(N)$

$SC = O(1)$

Q → Given an integer  $N$ , fill the  $N \times N$  matrix with elements **1 to  $N^2$**  in spiral order (clock wise), starting from  $(0,0)$ .

$N=4$

	0	1	2	3
0	1	2	3	4
1	12	13	14	5
2	11	16	15	6
3	10	9	8	7

4x4

$N=5$

	1	2	3	4	5
	16	17	18	19	6
	15	24	25	20	7
	14	23	22	21	8
	13	12	11	10	9

Code → HW

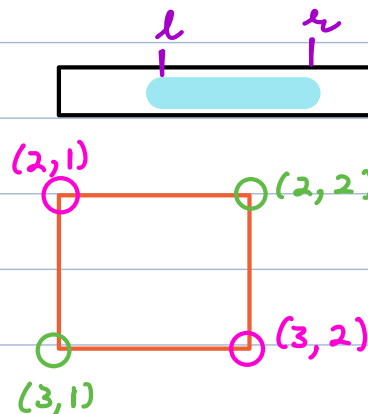
$TC = O(N^2)$

$SC = O(1)$

Submatrix

	0	1	2	3
0	1	2	3	4
1	12	13	14	5
2	11	16	15	6
3	10	9	8	7

eg.



✓ top left & bottom right OR  
top right & bottom left

Q → Given a matrix, find sum of all possible submatrix sum.

	0	1	2
0	4	9	6
1	5	-1	2

(0,0)

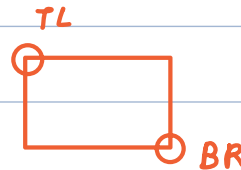
$$4 \rightarrow (0+1) * (0+1) * (2-0) * (3-0) \\ = 1 * 1 * 2 * 3 = 6$$

$[4]$ 4	$[4 \ 9]$ 13	$\begin{bmatrix} 4 \\ 5 \end{bmatrix}$ 9	$\begin{bmatrix} 4 & 9 \\ 5 & -1 \end{bmatrix}$ 17
$[9]$ 9	$[9 \ 6]$ 15	$\begin{bmatrix} 9 \\ -1 \end{bmatrix}$ 8	$\begin{bmatrix} 9 & 6 \\ -1 & 2 \end{bmatrix}$ 16
$[6]$ 6	$[5 \ -1]$ 4	$\begin{bmatrix} 6 \\ 2 \end{bmatrix}$ 8	$\begin{bmatrix} 4 & 9 & 6 \\ 5 & -1 & 2 \end{bmatrix}$ 25
$[5]$ 5	$[-1 \ 2]$ 1		
$[-1]$ -1	$[4 \ 9 \ 6]$ 19		
$[2]$ 2	$[5 \ -1 \ 2]$ 6		

Ans = 166

Contribution Technique → Ans =  $\sum_i \sum_j$  contribution of  $A[i][j]$

$A[i][j] * (\# \text{ submatrix } A[i][j] \text{ is a part of})$



0	0	1	2	3	4
0	1	2	3	4	5
1	16	17	18	19	6
2	15	24	25	20	7
3	14	23	22	21	8
4	13	12	11	10	9

# top left →  $(i+1) * (j+1)$

# bottom right →  $(N-i) * (M-j)$

$$[i \quad N-1] \rightarrow N-1-i+1 = \underline{N-i}$$

$$\# \text{ submatrices} = (i+1) * (j+1) * (N-i) * (M-j)$$

$$\text{Ans} = \sum_i \sum_j A[i][j] * (i+1) * (j+1) * (N-i) * (M-j)$$

ans = 0

for i → 0 to (N-1) {

for j → 0 to (M-1) {

ans += A[i][j] \* (i+1) \* (j+1) \* (N-i) \* (M-j)

}

} return ans

TC = O(N \* M)

SC = O(1)

	0	1	2	3	4
0	✓	✓	✓		
1	✓	✓	✓	✓	✓
2			✓	✓	✓
3			✓	✓	✓

$$6 * 9 = \underline{54}$$