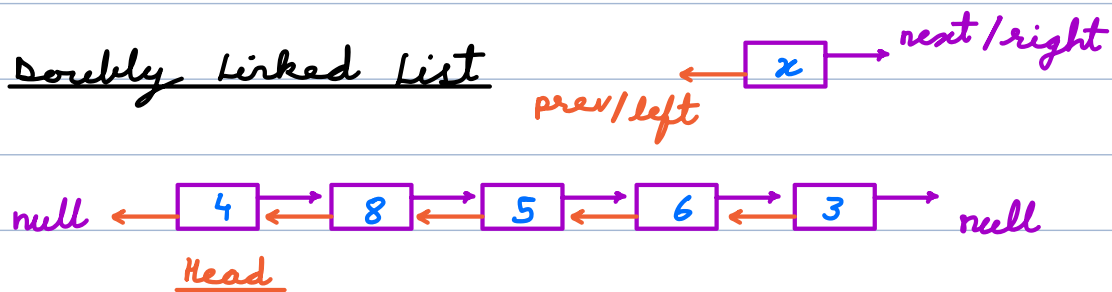
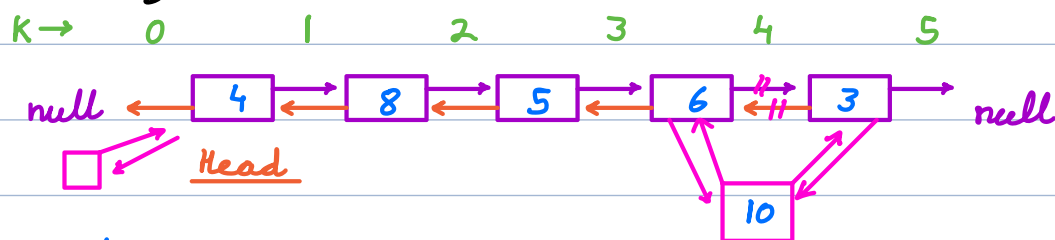


## Doubly Linked List



Q → Insert a node with data  $x$  at position  $K$  [ $0 \leq N$ ] in doubly linked list.



$K = 4$

$X = 10$

```

xr = new Node(x) // xr.next = xr.prev = null
if (Head == null) return xr
if (K == 0) {
    |   xr.next = Head
    |   Head.prev = xr
    |   return xr // updated Head
}

```

```

temp = Head
for i → 1 to (K-1) {
    |   temp = temp.next
}

```

```

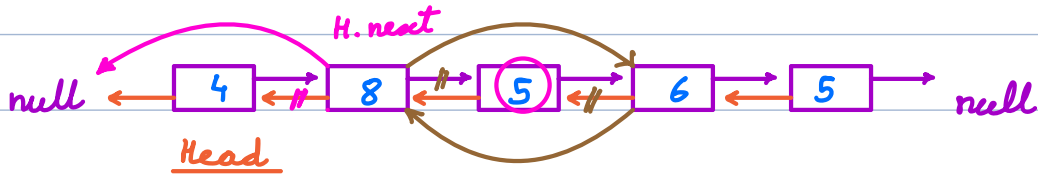
xr.prev = temp
xr.next = temp.next
if (temp.next != null)    temp.next.prev = xr
temp.next = xr

```

return Head

$TC = O(k)$      $SC = O(1)$

Q → Given a doubly linked list, delete the first occurrence of  $x$ . If not present, ignore.



$x = 5$

$\text{null} \leftarrow [x] \rightarrow \text{null}$

```
if (Head == null) return Head
```

```
if (Head.data == x) {
```

```
    if (Head.next != null)
```

```
        Head.next.pre = null
```

```
    Head = Head.next
```

```
    return Head
```

```
}
```

```
Temp = Head
```

```
while (Temp != null) {
```

```
    if (Temp.data == x) break
```

```
    Temp = Temp.next
```

```
}
```

```
if (Temp == null) return Head
```

```
// node to delete → Temp
```

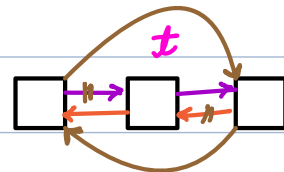
```
if (Temp.prev != null) // non-Head
```

```
    Temp.prev.next = Temp.next
```

```
if (Temp.next != null)
```

```
    Temp.next.prev = Temp.prev
```

```
return Head
```

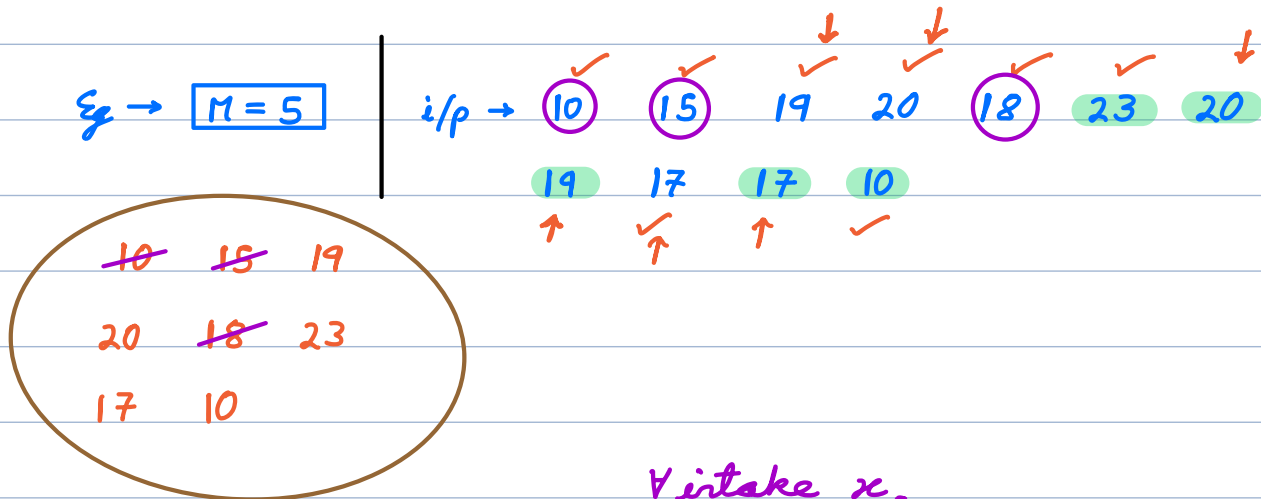


$\text{null} \leftarrow [x] \rightarrow \text{null}$

$$TC = O(N) \quad SC = O(1)$$

LRU → Least Recently Used

Q → Given a running stream of integers, & a fixed memory  $SC = O(M)$ .  $\forall$  input maintain most recent  $M$  inputs in the memory. If memory is full remove least recent data.



$\forall$  intake  $x$ ,

(once the memory is full)

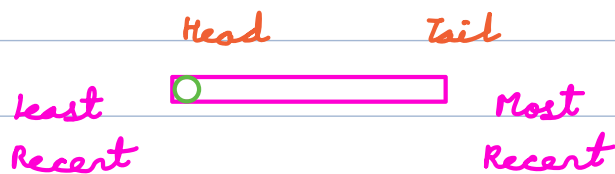
if  $x$  is not present

- 1) delete least recent item ✓
- 2) insert  $x$  as most recent item ✓

if  $x$  is already present

- 1) delete  $x$  from its position
- 2) insert  $x$  as most recent item ✓

Maintain elements in order of recency → ~~Array / Linked List /~~ Dynamic Array / DLL



if we directly reach the node to delete.

⇒ Use HashMap

$\langle x, \text{node with data } x \rangle$

```

forall x,
if (hm.containsKey(x)) {
    xr = hm.get(x)
    deleteNode(Head, xr) → TC = O(1)
    insertNode(Tail, xr)
} else if (hm.size() < M) {
    xr = new Node(x)
    insertNode(Tail, xr)
    mp.put(x, xr)
} else { // Memory full
    mp.remove(Head.data)
    deleteHead(Head)
    xr = new Node(x)
    insertNode(Tail, xr)
    mp.put(x, xr)
}

```

TC per i/p = O(1)      SC = O(M)

---

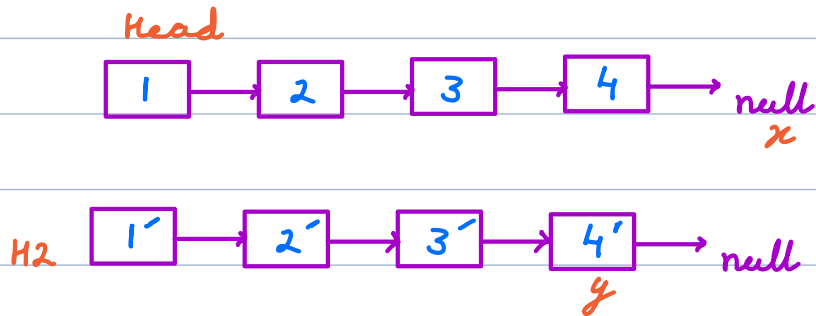
$x = \text{new Node}(5)$ 
  
 $y = x$ 
  
 $y.data = 10$ 
  
 $\text{print}(x.data) // 10$

Shallow Copy

$x = \text{new Node}(5)$ 
  
 $y = \text{new Node}(x.data)$ 
  
 $y.data = 10$ 
  
 $\text{print}(x.data) // 5$

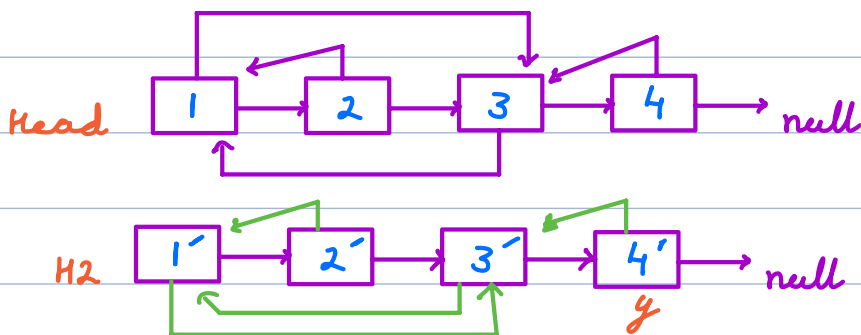
Deep Copy

Q → Create deep copy of doubly linked list with random pointers.



```

x = Head
H2 = new Node (Head.data) // Head != null
y = H2      x = x.next
while (x != null) {
    y.next = new Node (x.data)
    y = y.next
    x = x.next
}
return H2
  
```



Nodes → store reference HashMap

Nodes in copy list x,

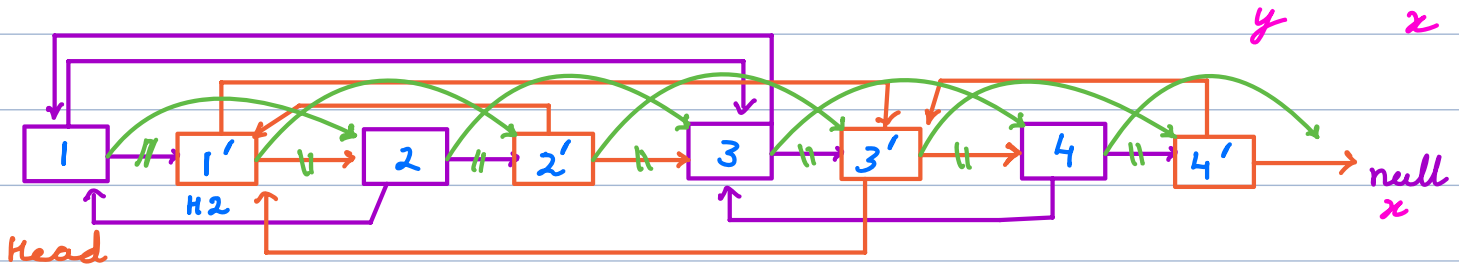
$x.random = mp.get(original-x.random)$

< node, node >  
original → copy

TC =  $O(N)$

SC =  $O(N)$

$O(1)$



①  $x = \text{Head}$

```
while (x != null) {  
    y = new Node(x.data)  
    y.next = x.next  
    x.next = y  
    x = x.next.next // y.next  
}
```

②  $x = \text{Head}$

```
while (x != null) {  
    x.next.random  
    = x.random.next  
    x = x.next.next  
}
```

③

```
H2 = Head.next  
x = Head    y = H2  
while (x != null) {  
    x.next = y.next  
    x = x.next  
    if (x != null) { y.next = x.next  
                    y = x.next }  
}
```

return H2

$TC = O(N)$

$SC = O(1)$