

Today's Agenda

Hashing - Internal Implementation.

Will start in next 3-4 minutes.

Q. Given an array of size N & Q queries. In each query, an element is given. We need to check if that element exists or not in the given array.

Arr \rightarrow [2, 4, 11, 15, 6, 8, 14, 9]

$Q=4$

$X=4$ \rightarrow true

$X=10$ \rightarrow false

$X=17$ \rightarrow false

$X=14 \rightarrow$ true.

B.f idea. ~

for every query, iterate on array elements & check if that element is present or not.

$$\left[\begin{array}{l} T.C \rightarrow O(Q * N) \\ S.C \rightarrow O(1) \end{array} \right]$$

idea-2. \rightarrow Sort the array & apply BS for every query.

$$\left[\begin{array}{l} T.C \rightarrow O((Q+N) \cdot \log N) \\ S.C \rightarrow \text{depends on inbuilt sorting algo used} \end{array} \right]$$

idea-3.

Create an array where i^{th} index denotes the presence of i .

Arr[7] = [2, 4, 11, 15, 6, 8, 14, 9]

Q=4

| | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|----|----|----|----|----|----|
| 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 1 |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |

K=4

K=10

K=17

K=14

insert $\rightarrow O(1)$

search $\rightarrow O(1)$

delete $\rightarrow O(1)$

Issues with this representation →

① Wastage of space → $[2, 10^5, 15]$

② Inability to create big arrays

arr → $[100, 3, 10^9, 16, 1]$

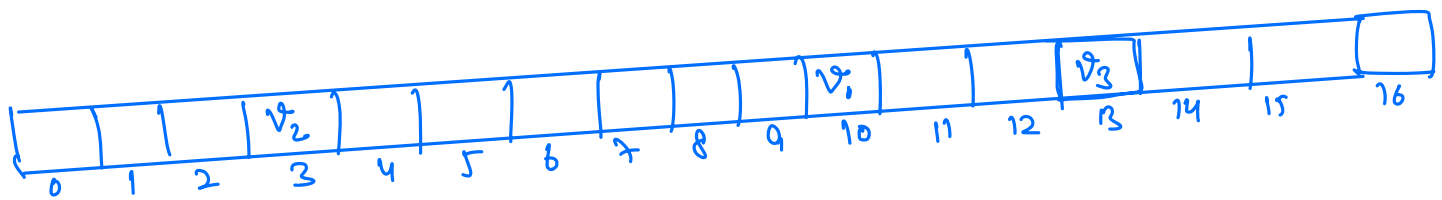
③ Adjustments for -ve values →

arr[1] → $[12, 3, -5, 10]$

| | | | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|----|----|----|----|----|----|----|----|
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 |

Assumption \rightarrow Max size of array possible $\rightarrow \underline{17}$
(0-16)

arr[7] \rightarrow [\checkmark_{10} , \checkmark_{20} , \checkmark_{30} , \checkmark_{27}]



Collision

$$10 \% 17 \rightarrow 10$$

$$20 \% 17 \rightarrow \underline{3}$$

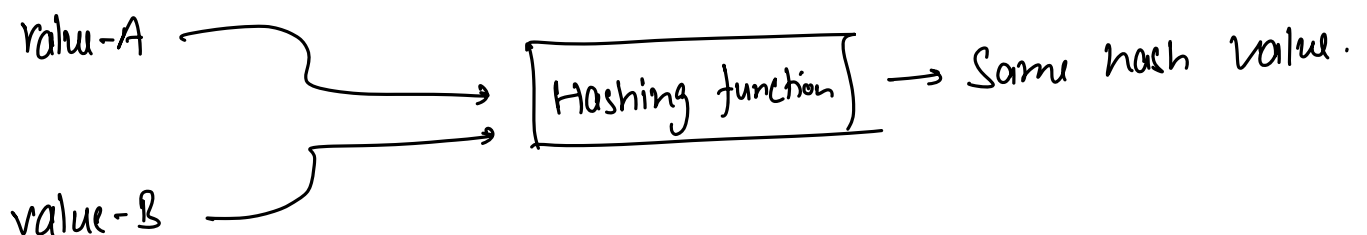
$$30 \% 17 \rightarrow 13$$

$$27 \% 17 \rightarrow 10$$

$(\text{Key} \% 17) \rightarrow$ Hashing Function.

Hashing \rightarrow Hashing is a process where we pass our data through the hash function.

\rightarrow give us the index to which data must be mapped.

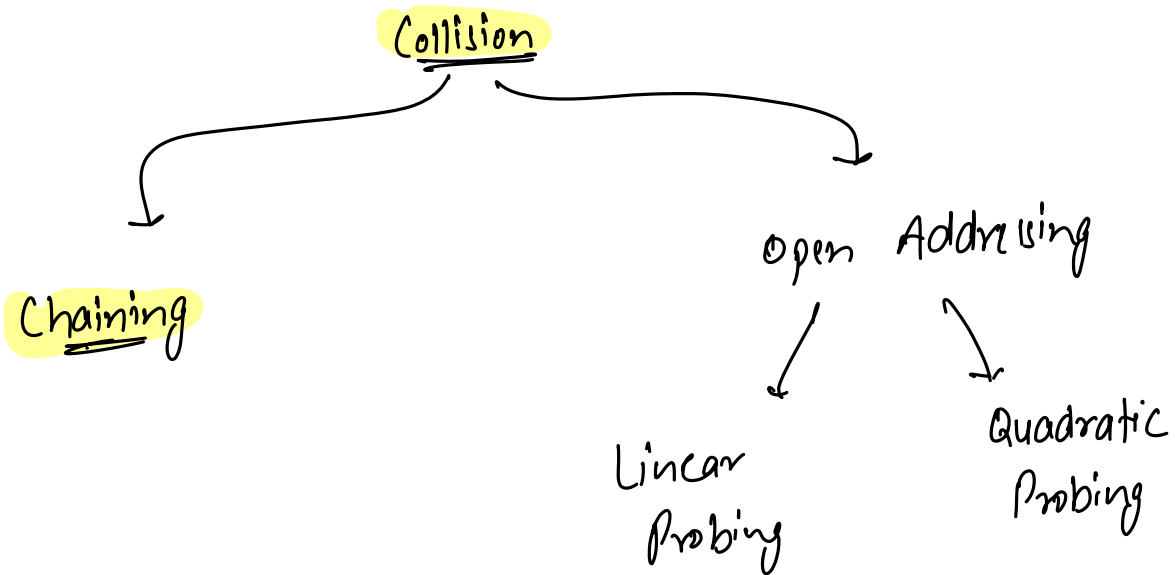
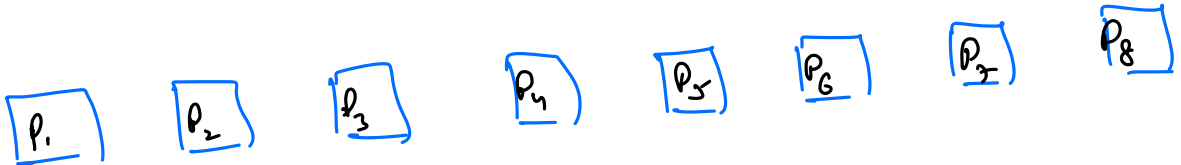


Q → Can we completely avoid collision?

Not possible.

Pigeon - Hole principle.

11 - pigeons < 8 - holes



elements →

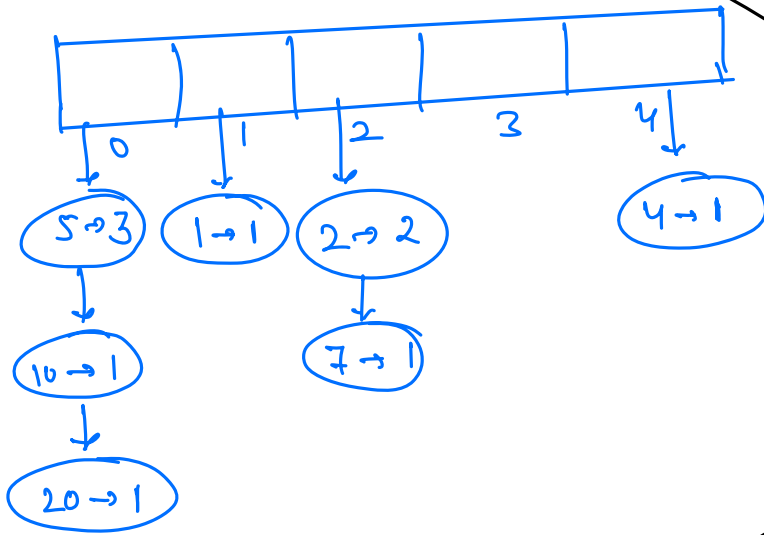
| | | | | | | | | | |
|---|---|---|---|----|---|---|----|---|---|
| 2 | 4 | 5 | 7 | 10 | 2 | 5 | 20 | 5 | 1 |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |

Assm max size of array allowed $\rightarrow 5$

→ store frequency of all the elements.

Hashing function \rightarrow element % 5

~~buckets. -~~



↓

Array of linked-lists \Rightarrow HashMap.

Average size of linked-list = $\frac{3+1+2+0+1}{5} = \frac{7}{5} = 1.4$ (load factor)

[Loading factor]
(λ)

$$\lambda = \frac{\text{total no. of nodes}}{\text{size of array}}$$

Insertion \rightarrow

- ① Calculate hash value or bucket index.
- ② Check if Key is already present in bucket
 - \rightarrow if it is already present (update value)
 - \rightarrow otherwise, add a new Node.

$$T.C \rightarrow \underline{\underline{O(\lambda)}}$$

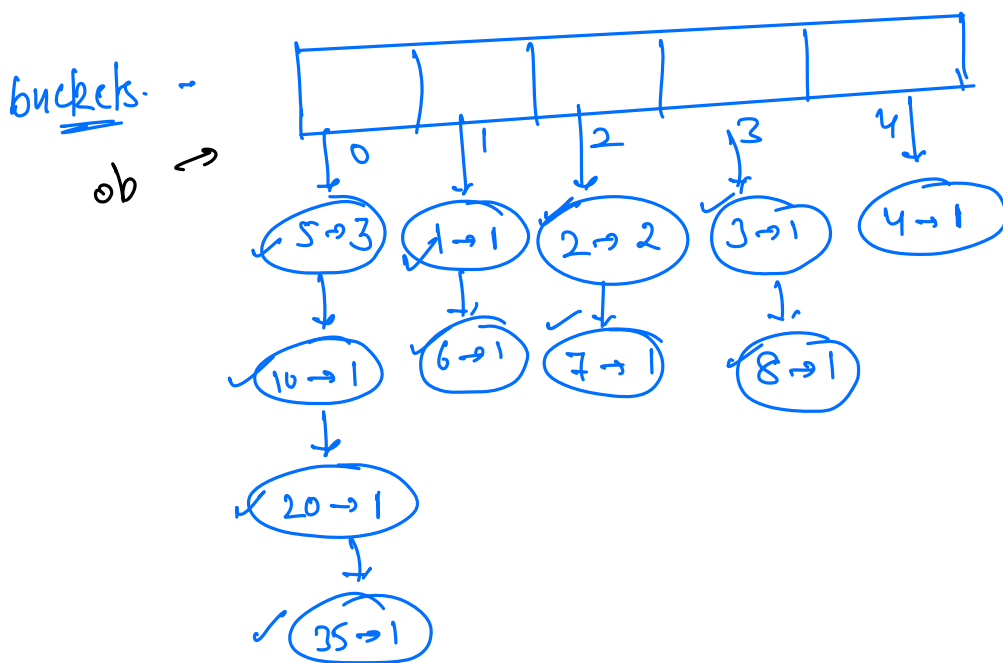
If we can somehow maintain, the average size of linked-list as constant (λ)

T.C for insertion \rightarrow Averagely constant.

Implementation of Hashmap →

https://www.scaler.com/topics/java/online-java-compiler/?snippet_slug=746dda157372a0ba6f98

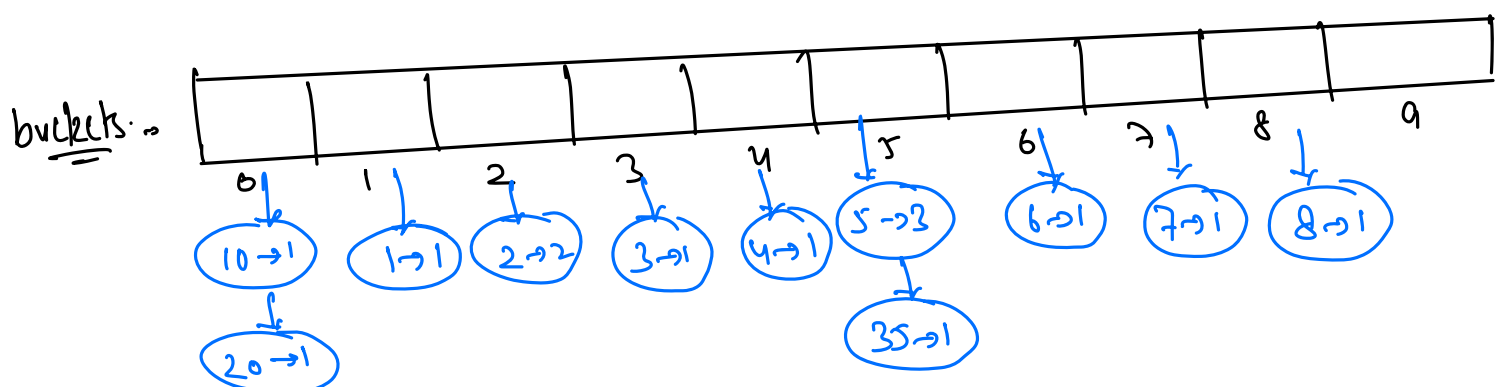
T.C for insertion, deletion, searching → $O(1)$



$$\lambda = \frac{11}{5} = 2.2$$

if $(\lambda > 2.0)$ {
 rehash()
}

↓ $\% 10$



$$\lambda = \frac{11}{10} = 1.1 < \underline{2}.$$

