

# GRI: Interpreter of a dynamic language for GGraph algorithms

Sandeep Dasgupta  
University Of Illinois at Urbana Champaign.  
sdasgup3@illinois.edu

## ABSTRACT

As graphical models are increasingly become popular in various field, the domain experts often struggle to represent and compute on such model in a convenient and efficient way. In this project we develop a dynamically typed language to represent and compute on the graphical models which provides both the desired convenience without losing much on the efficiency.

## Keywords

Graph Algorithm, dynamically typed language, interpreter

## 1. INTRODUCTION

As graphical models are increasingly being used in various fields like biochemistry (genomics), electrical engineering (communication networks and coding theory), computer science (algorithms and computation) and operations research (scheduling), organizational structures, social networking, there is a need to represent and allow computation on them in a convenient and efficient way. This involves (but not limited to)

- Designing a language which provide an convenient interface to the programmer to program those models. This is essential so that even for domain experts who are not coding experts can code and reason about their implementation. Ease of interface could be due to:
  - Expressive power of the language representing those models.
  - Intuitive extensibility of the language.
  - Ability of the language to provide exploratory programming, where the user may experiment with different ideas (without dwelling much into the language syntax) before coming to a conclusive one.
- Designed language need to be efficient in the following sense.
  - Underlying design decisions including data structures need to be carefully crafted to achieve expected run-time w.r.t the input size.

- Implementation need to be scalable w.r.t the space/time requirements. This is important because most of the graph algorithm typically work on huge input sizes.

## 2. RELATED WORK

Our work is mostly inspired by the line of work by GUESS [3] and Graphal [2]. GUESS, a novel system for graph exploration that combines an interpreted language with a graphical front end that allows researchers to rapidly prototype and deploy new visualizations. GUESS also contains a novel, interactive interpreter that connects the language and interface in a way that facilitates exploratory visualization tasks. Our language, Gython, is a domain-specific embedded language which provides all the advantages of Python with new, graph specific operators, primitives, and shortcuts.

Graphal is an interpreter of a programming language that is mainly oriented to graph algorithms. There is a command line interpreter and a graphical integrated development environment. The IDE contains text editor for programmers, compilation and script output, advanced debugger and visualization window. The progress of the interpreted and debugged graph algorithm can be displayed in 3D scene.

Our language design is very similar to above two work. But we additionally provided built-in functions for basic graph algorithms. This not only help us getting convenient short hand notations to compute those basic algorithms, but also we gain on performance due the fact that those basic algorithms are now compiled.

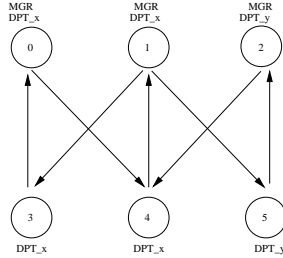
## 3. A MOTIVATING EXAMPLE

In this section we will provide some insight of designed language using some motivating examples.

**EXAMPLE 1.** *Figure1, shows a directed graph where the nodes represent the employees and the edges between them represents the email conversation from one employee to other. For example, Node 0 represents an employee in DPT\_x who is a Manager as well (this is attributed by the MGR tag). Similarly, Node 3 represents a DPT\_x employee. The edge between Node 3 and Node 0 represents the email conversation from employee Node 3 to employee node Node 0.*

*Now lets first talk about the ways to represent this graph. One way is as shown in 2.*

*At line S1, a new graph variable is created. S3 sets that the graph is directed. S5 create a node Node 0 of this graph. As nodes and edges are the basic building block of a graph we have made them the first class objects which allows users to access them directly. Lines S6 - S9 sets various properties of the vertex. For example, it says that the vertex has `__id = 0`, it belongs to `__DPT_x`, does not belongs to `__DPT_y` and its `__MGR` property is true. Similar*



**Figure 1: An example graph**

```

function main(argv) main() {
S1:  g = graph();
S2:  g.loadFromFile(argv[0]);
S3:  displayAdjMatrix(g.getAdjacencyMatrix(), g.getVertices());
S4:  ...
S5:  dpt_x_employee = g.getVertexSetWithProperty("__DPT_x", 1.0);
S6:  mgr_employee = g.getVertexSetWithProperty("__MGR", 1.0);
S7:  ...
S8:  /* Set of __DPT_x employees who are __MGR as well */
S9:  dpt_x_AND_mgr = dpt_x_employee.intersection(mgr_employee);
S10: ...
S11: println("Set of __DPT_x employees who are __MGR as well");
S12: foreach(employee; dpt_x_AND_mgr) {
S13:     println(" " + employee.__id + " ");
S14: }
S15: ...
S16: /* Set of __DPT_x employees who are NOT __MGR */
S17: dpt_x_MINUS_mgr = dpt_x_employee.difference(mgr_employee);
S18: ...
S19: it = dpt_x_MINUS_mgr.iterator();
S20: ...
S21: println("Set of __DPT_x employees who are NOT __MGR");
S22: while(it.hasNext()) {
S23:     employee = it.next();
S24:     println(" " + employee.__id + " ");
S25: }
S26: ...
S27: // Email Exchanges from MGRs to non-MGR DPT_x employee
S28: emailExchanges = mgr_employee -> dpt_x_MINUS_mgr;
S29: ...
S30: // Email Exchanges from non-MGR DPT_x employee to MGRs*/
S31: emailExchanges = mgr_employee <- dpt_x_MINUS_mgr;
S32: ...
S33: // Email Exchanges between non-MGR DPT_x employee and MGRs*/
S34: emailExchanges = mgr_employee <-> dpt_x_MINUS_mgr;
S35: ...
}

```

**Figure 3: Example code snippet to create a graph**

```

function main(argv) main() {
S1:  g = graph();
S2:  ...
S3:  g.setDirected(true);
S4:  ...
S5:  v0 = g.createVertex();
S6:  v0.__id = 0;
S7:  v0.__DPT_x = 1;
S8:  v0.__DPT_y = 0;
S9:  v0.__MGR = 1;
S10: ...
S11: v3 = g.createVertex();
S12: v3.__id = 3;
S13: v3.__DPT_x = 1;
S14: v3.__DPT_y = 0;
S15: v3.__MGR = 0;
S16: ...
S17: g.createEdge(v0, v3);
}

```

**Figure 2: Example code snippet to create a graph**

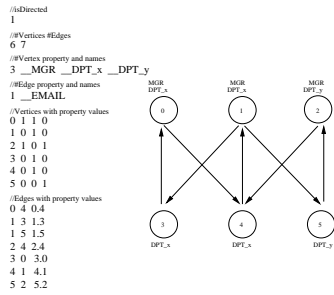
properties are set for vertex node with `__id = 3`. And finally a directed edge between them is created at line S17.

One thing to note here is that the methods like `graph`, `setDirected`, `createVertex` and `createEdge` are all built-in compiled functions.

Figure 3 shows another way to represent graph. It uses an input file to be fed to the interpreted program. The format of the input file is shown in Figure 4. Such an input file is read from the command line and used to create a graph as shown in line S2 of Figure 3. At line S3, we can see two built-in functions `getAdjacencyMatrix` & `getVertices` on graph which respectively gives an array of array representing the adjacency matrix representation of the graph and a set of vertices in the graph. `displayAdjMatrix` is just a method call, the definition of which is not shown for brevity. Now let's try to solve certain queries from the graph at Figure. In case we want to get all the nodes in the graph who are `__DPT_x` employees, then the query to get all those nodes is at line S5. Similarly, the query at line S6, gives the nodes for which the property `__MGR` is set to true. One thing to note here that the return value of both the queries are sets which are amenable to set operations.

For example, in case we want to get all the employees who are both department `__DPT_x` employee and managers, we can get that using the set intersection as shown at line S9. Line S8, shows the multilien comment we support. Line S9 shows a way to iterate over the set elements using `foreach` construct.

Again, if we want to get the set of `__DPT_x` employees who are **NOT** `__MGR`, we can write a query like at line S17. Line S19 shows



**Figure 4: The input file used for graph creation**

```

S1:  define("NUM_VERTICES", "10");
S2:  define("NOTVISITED", "0");
S3:  define("VISITED", "1");
      function dfs(v, dfsorder) {
S4:      if(v.visit == VISITED)
S5:          return;
S6:
S7:      println("vertex visited: " + v.num);
S8:      v.visit = VISITED;
S9:
S10:     dfsorder.pushBack(v);
S11:
S12:     foreach(neighbor; v.getNeighbors())
S13:         dfs(neighbor, dfsorder);
      }

      function main(argv) {
S14:         g = graph();
S15:         g.loadFromFile(argv[0]);
S16:
S17:         dfsorder = array(0);
S18:         dfs(first, dfsorder);
S19:
S20:         for(vertex: dfsorder) {
S21:             println(" " + vertex.__id + " " );
S22:         }
      }

```

**Figure 5: Example code snippet for dfs traversal on a graph**

another ways to get an iteration to iterate on composite data structures.

Now once we have two sets of nodes each with a specific property, we can query for the edges between them. For example, lines S28, 31 and 34 gives the set of edges emanating from one set of nodes to another set of nodes. Note that operators  $\leftarrow$  and  $\rightarrow$  are applicable to both directed graphs and the operator  $\leftrightarrow$  is applicable to both directed and undirected graphs. For directed graphs, it gives all the bidirectional edges between two node sets and for undirected graphs, it returns all the edges between two node sets.

□

**EXAMPLE 2.** Figure 5 represents another example implementation the depth first traversal on a graph. In this example we are trying to store the dfs traversal order of the vertices as well. For that we declared an array at line S17 and used it to store the traversal order at line S10. The method getNeighbors at line S12 is again a built-in function and returns the neighbouring vertices of the receiver vertex.

□

## 4. LANGUAGE DESIGN

The syntax of the language is an oversimplified version of C, but without mention of any types. The operations on incompatible types will be error-ed out while interpreting.

We have implemented the tokenizer and syntax analyzer using flex and bison. We are supporting syntax like `#include("filename")` and `#define("PI", "3.14")` while doing a single pass of parsing (i.e. Preprocessing of these constructs are done while parsing). This is achieved by using flex's internal stack to manage multiple buffers. The grammar rules are mostly borrowed from [1]. The rules are compiled by bison tool to generate the C parser. We are able to generate the AST corresponding to test-cases confirming to the grammar rules. Our AST is basically a list of function definitions. Each function definition object contains name of the function, a set of formal arguments and a list of body statements. These body statements could be an assignment, loop-statement, function call, etc. The leafs of the AST could be an identifier, int, float, true, false, null, string, vertex, edge or graph.

Some of the key features of the parser is as follows:

- Support of C statements like if then, if then else, while, for, foreach.
- Support of break, continue within loop-body and return in function-body. As we are representing both loop-body and function-body as compound statements (i.e. anything between "{" & "}"), so we do not have to distinguish these two cases. But we will error-out if break is used inside non-loop body. The detailed semantics of executing a break, continue and return will be discussed in the interpreter runtime section.
- Supporting graph as first class object graphnode. The syntax to declare a graph is `g = graph();` which will be represented in AST as an assignment-node with left-node containing an identifier and right node as a function call. Now this function call corresponds to a built in function that returns a graphnode (which is of one the leaf nodes of AST) on execution.
- Supporting vertices and edges as first class objects which contains a map to add properties. This feature is useful in various graph algorithms like in dfs traversal 5, we use a vertex property "visited" to keep track of vertices already explored.
- We are supporting composite data-structures like array, struct and set and iterators on them. Figure 6 represents a code snippet representing some of the operation on these data-structures.
- The language semantics will be same as that of C as we are using a subset of it.
- All variables are defined as local and are valid only in the scope of the current function (function not block).
- The language specifies no constructs for variable declaration and type specification. The interpreter uses some inner data types (like null, Bool, int, float, string, array, struct, set, graph, vertex, edge) which can be dynamically changed with assign command.

## 5. INTERPRETER RUNTIME

The following are the key features of the runtime:

- The runtime starts with searching for function definition `function main(argv) main() {` and then creates a function call out of it and execute it. While creating the function call it uses the command-line parameters as the actual parameters of the function call.
- The execution of a function call involves finding the corresponding function definition, checking if the number of formal and actuals are equal and then pushing a call stack frame ( which contains the mapping between the formal and actual values passed to them) in a global call stack. After that, the function is executed w.r.t the current context(i.e. the top of the call stack).
- The execution of the function involves executing a list of statements. The statements may add further mappings in the current call stack frame. Whenever a name (identifier) is refereed, the mapping in the current context need to be consulted to get the actual value of it.
- The semantics of `break`, `continue` or `return` is supported using the try-catch mechanism of C++.

For example, while interpreting a `loop-body`, whenever a `break` is encountered, a corresponding `break-object` is thrown, which is caught in a place outside the entire loop execution in order to implement the semantics of `break`.

Similarly, while executing a `loop-body`, whenever a `continue` is encountered, a `continue-object` is thrown, which is caught outside the `loop-body` execution so as to skip the current iteration and continue with the `loop-incr` execution (in case of for loop) and `loop-condition-expr` execution.

And finally, while executing a `function-body`, which is a list of statements, when any one of those statements is a `return`, a `return-object` is thrown, which is caught outside of the loop which is going over that list and in this way the semantics of `return` is maintained.

- Occurrence of `break` and `continue` within non-loop body triggers an error. While interpreting a node-block (which is a set statements within “{” and “}”), whenever the runtime finds a `break` or `continue` it throws a object. Now if this object is caught inside a non-loop block then error is reported.
- Division by zero and operations on incompatible types are runtime errors.

Also we are supporting a number of built-in functions 1. The advantage of using is that we can save a lot of interpretation time.

## 6. EVALUATION

The baseline of our evaluation will be the open source Graphal [2] system. Graphal [2] is an interpreter of a programming language that is mainly oriented to graph algorithms. There is a command line interpreter and a graphical integrated development environment. The IDE contains text editor for programmers, compilation and script output, advanced debugger and visualization window. The evaluation will be between command line version of Graphal with our implementation. The evaluation will cover the following two aspects of our implementation:

- The convenience and intuitive extensibility provided by our programming model.

```
function main(argv) main() {
S1:   arr = array(5);
S2:   i = 0;
S3:   ...
S4:   foreach(var ; arr)
S5:       var = i++;
S6:   ...
S7:   println("-- array items --");
S8:   foreach(var ; arr)
S9:       println(var);
S10:  ...
S11:  st = struct();
S12:  st.number = 42;
S13:  st.pi = 3.14;
S14:  st.str = "bagr";
S15:  ...
S16:  println("-- struct items --");
S17:  foreach(var ; st)
S18:      println(var);
S19:  println("-- struct items using iterator --");
S20:  it = st.iterator();
S21:  while(it.hasNext())
S22:      ` println(it.next());
S23:  ...
S24:  g = graph();
S25:  v1 = g.generateVertex();
S26:  v2 = g.generateVertex();
S27:  v3 = g.generateVertex();
S28:  e1 = g.generateEdge(v1, v2);
S29:  e2 = g.generateEdge(v2, v3);
S30:  ...
S31:  v1.color = "red";
S32:  v2.color = "green";
S33:  v3.color = "blue";
S34:  e1.value = 0.5;
S35:  e2.value = 0.4;
S36:  ...
S37:  println("-- vertex set --");
S38:  foreach(var ; g.getVertices())
S39:      println(" " + var + ": " + var.color);
S40:  ...
S41:  println("-- edge set --");
S42:  foreach(var ; g.getEdges())
S43:      println(" " + var + ": " + var.value);
}
```

**Figure 6:** Example code snippet to show operations on array, set and struct

**Table 1: List of all the built-in functions provided.**

Category	Name of the function	Synopsis
Output	<code>println</code>	Convert the object to a string representation and send the result to the standard output and append newline.
	<code>print</code>	Convert the object to a string representation and send the result to the standard output.
Container	<code>array(size)</code>	Create a new array of specified size.
	<code>struct()</code>	Return a new struct.
	<code>set()</code>	Return a new set
	<code>size(array struct set)</code>	Return the size of the argument container.
	<code>union(set, set)</code>	Return union of two sets.
	<code>intersection</code>	Return intersection of two sets.
	<code>difference</code>	Return difference of two sets.
	<code>pushFront(array, elem)</code>	push an element elem to front of array
	<code>pushBack(array, elem)</code>	push an element elem to back of array
	<code>popFront(array)</code>	Remove an elem from front of array
	<code>popBack</code>	Remove an elem from back of array
	<code>front</code>	Get the front elemnt of an array.
	<code>back</code>	Get the back element of an array
Iterator	<code>iterator(object)</code>	Return a copy of a container and set its inner iterator to the beginning.
	<code>hasNext(Object)</code>	Check if a container has a next item.
	<code>next(Object)</code>	Get the next item.
Graph	<code>graph()</code>	Returns a newly created graph
	<code>loadFromFile(graph)</code>	Load a graph from a file
	<code>saveToFile(graph, filename)</code>	Save rhe current state of the graph to a file.
	<code>setDirected(graph, bool)</code>	Set the directed flag of the graph and return the previous value.
	<code>isDirected(graph)</code>	Check if a graph is directed.
	<code>generateVertex(graph)</code>	Create a new vertex in a graph.
	<code>generateEdge(graph, vertex, vertex)</code>	Create a new edge in a graph.
	<code>deleteVertex(graph, vertex)</code>	Delete a vertex from a graph.
	<code>deleteEdge(graph, edge)</code>	Delete an edge from a graph.
	<code>getNumVertices(graph)</code>	Get count of vertices in a graph.
	<code>getNumEdges(graph)</code>	Get count of edges in a graph.
	<code>getVertices(graph)</code>	Get graph's vertices as a set object.
	<code>getEdges(graph)</code>	Get graph's edges as a set object.
	<code>getNeighbors(vertex)</code>	Get all neighbors of a vertex as a set object.
	<code>getBeginVertex(edge)</code>	Get the begin vertex of an edge. If the graph is not directed, one of the edge's vertices will be returned.
	<code>getEndVertex(edge)</code>	Get the end vertex of an edge. If the graph is not directed, one of the edge's vertices will be returned.
	<code>getAdjacencyMatrix(graph)</code>	Create adjacency matrix from a graph. An array of array will be returned.
	<code>getTransitiveClosure(graph)</code>	Returns the transitive closure of the adjascent matrix representation of the graph. Uses the Floyd Warshall Algorithm.
	<code>getShortestPath(graph, weight, start, end)</code>	Returns an array containing two elements. First, the parent of each vertex in the shortest path tree. Second, the distance of each vertex from start in that shortest path tree.  If end == null, the the algorithm will compute the shortest distance of all the vertices from start. Else it stops when the shortest distance of end vertex from start is computed.  weight (a string) signifies which of the properties of the edge need to be considered for computing min distance.  Uses Dijkstras Algorithm with min heap.
	<code>getMST(graph, weight)</code>	Return the parent of each vertex in the minimum snapping tree.  weight (a string) signifies which of the properties of the edge need to be considered for computing min distance.  Uses Prim's Algorithm.
	<code>getVertexSetWithProperty(graph, property, value)</code>	Get the set of vertices with property string equals the value.
	<code>getbfsOrdering(graph)</code>	Get the set of vertices in bfs traversal order.
	<code>getdfsOrdering(graph)</code>	Get the set of vertices in dfs traversal order.

- We will be implementing a couple of well know graph algorithms in both baseline and our implementation and use number of dynamic instructions interpreted as a measure to show the conciseness of our representation.
- Performance
  - As we are planning to keep the compiled version of frequently used graph routines (like `dfs_iterators(vertex)`, `bfs_iterators(vertex)`), we are expecting to achieve better performance in terms of runtime.

## 7. REFERENCES

- [1] <http://www.quut.com/c/ANSI-C-grammar-y-2011.html>.
- [2] Graphal: Graph algorithms interpreter.
- [3] E. Adar. Guess: A language and interface for graph exploration. In *Proceedings of the SIGCHI Conference on Human Factors in Computing Systems*, CHI '06, pages 791–800, New York, NY, USA, 2006. ACM.