allvm - Binary Decompilation

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Outline

Goal & Motivation

Possible Approaches

Our Approach

Present Status

Research Goal

- Obtain “richer” LLVM IR than native machine code.
- Enable advanced compiler techniques (e.g. pointer analysis, information flow tracking, automatic vectorization)

Why “richer” LLVM IR

- Source code analysis not possible
  - IP-protected software
  - Malicious executable
  - Legacy executable
- Source code analysis not sufficient
  - What-you-see-is-not-what-you-execute
- Platform aware and dynamic optimizations

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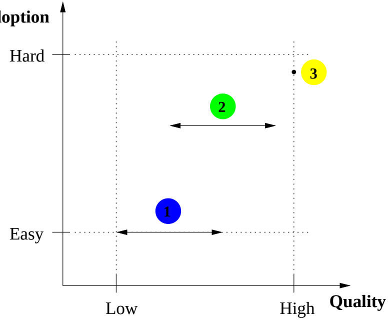
3 Possible Approaches

<u>Research Goal</u>: Obtain “richer” LLVM IR than native machine code.

<u>Possible Approaches</u>

1. Decompile Machine Code → LLVM IR
2. "Annotated" Machine Code → LLVM IR
3. Ship LLVM IR



Decompile `Machine Code` → `LLVM IR`

Benefits	Challenges
• Easy to adopt • No compiler support needed	• Reconstructing code and control flow • Variable recovery • Type recovery • Function & ABI rules recovery

- Tools Available: QEMU, BAP, Dagger, Mcsema, Fracture

“Annotated” `Machine Code` → `LLVM IR`

Benefits	Challenges

• Effective reconstruction
• Minimal compiler support needed
	• Annotations to be
	    • Minimal
	    • Compiler & IR independent
	• Adoption

- Tools Available: None

Ship LLVM IR

Benefits	Challenges
• *No loss* of information	• Adoption in Non LLVM based compilers • Code size bloat

- Tools Available: Portable Native Client, Renderscript, iOS, watchOS, tvOS apps, ThinLTO

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Our Approach

Long term goal

Minimal compiler-independent annotations to reconstruct high-quality IR

Short term goals

1. Experiment with `Machine Code` $\rightarrow$ `LLVM IR`, to **understand** the challenges better
   - To select from existing decompilation frameworks.
   - Experiment with different variable and type recovery strategies
2. Design suitable annotations for what cannot be inferred without them

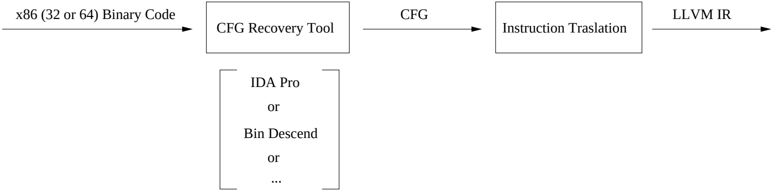
Outline

Present Status

Action Items	Status
Selected "mcsema" among the existing `Machine Code` $\rightarrow$ `LLVM IR` solutions. • Comparison of mcsema with existing tools • Evaluation of mcsema	Done
Literature survey on variable, type, function param recovery	Done
Implementing a variable and type recovery model using mcsema	Ongoing

Selecting mcsema

- Actively supported and open sourced
- Well documented
- Functional `LLVM IR`
- Separation of modules: CFG recovery and Instruction translation (`CFG` → `LLVM IR`)



Instruction Translation

- Processor state: Modeled as struct of `ints`
- Processor memory: Modeled as flat array of bytes

```
start:
    mov eax, [esp − 4h]
    add eax, 1
    ret
```

Binary Code

Translation

```
RECOVERED_FUNC ( struct RegisterContext regctx ) :
    VAR_EAX = regctx.EAX
    VAR_ESP = regctx.ESP

    VAR_EAX = [ VAR_ESP − 4 ]
    VAR_EAX += 1

    regctx.EAX = VAR_EAX
    regctx.ESP = VAR_ESP
END
```

High level view of Recovered Code

mcsema: Demo

Support & Limitations

- What Works
  - Integer Instructions
  - FPU and SSE registers
  - Callbacks, External Call, Jump tables
- In Progress
  - FPU and SSE Instructions: Not fully supported
  - Exceptions
  - Better Optimizations

Variable, Type, Function Param Recovery

- Enables
    - Fundamental analysis (Dependence, Pointer analysis)
    - Optimizations (register promotion)
- State of the art
    - Divine
        - State of the art variable recovery
    - TIE
        - Type recovery
    - Second Write
        - Heuristics for function parameter detection
        - Scalable variable and type recovery

Summary

Today: Functioning translation from `Machine Code` $\rightarrow$ executable `LLVM IR` (IR quality is poor)

Questions ?

EXTRA SLIDES: WYSNYX

The following compiler (Microsoft C++ .NET) induced vulnerability was discovered during the Windows security push in 2002

```
memset(password, '\0', len);
free(password);
```

$\longrightarrow$

```
free(password);
```