



# allvm - Binary Decompile

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March 30, 2016



**Goal & Motivation**

Possible Approaches

Our Approach

Ongoing Work

Backup



# Research Goal

- Research Goal
  - Obtain “richer” LLVM IR than native machine code.
  - Enable advanced compiler techniques ( e.g. pointer analysis, information flow tracking, automatic vectorization)



# Why “richer” LLVM IR

- Source code analysis not possible
  - IP-protected software
  - Malicious executables
  - Legacy executables
- Source code analysis not sufficient
  - What-you-see-is-not-what-you-execute
- Platform aware optimizations



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# The 3 Possible Approaches

- Decompile Machine Code  $\rightarrow$  LLVM IR
- “Annotated” Machine Code  $\rightarrow$  LLVM IR
- Ship LLVM IR



## Decompile Machine Code $\rightarrow$ LLVM IR

Benefits	Challenges
<ul style="list-style-type: none"><li>• Easy to adopt</li><li>• No compiler support needed</li></ul>	<ul style="list-style-type: none"><li>• Reconstructing code and control flow</li><li>• Variable recovery</li><li>• Function &amp; ABI rules recovery</li></ul>

- Tools Available: QEMU, BAP, Dagger, Mcsema, Fracture



## “Annotated” Machine Code $\rightarrow$ LLVM IR

Benefits	Challenges
<ul style="list-style-type: none"><li>• Effective reconstruction</li><li>• Minimal compiler support needed</li></ul>	<ul style="list-style-type: none"><li>• Annotations must be<ul style="list-style-type: none"><li>• Minimal</li><li>• Compiler &amp; IR independent</li></ul></li><li>• Adoption</li></ul>

- Tools Available: None





# Ship LLVM IR

Benefits	Challenges
<ul style="list-style-type: none"><li>• <i>No loss</i> of information</li></ul>	<ul style="list-style-type: none"><li>• Adoption in Non LLVM based compilers</li><li>• Code size bloat</li></ul>

- Tools Available: Portable Native Client, Renderscript, iOS, watchOS and tvOS apps.



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# Our Approach

## Long term goal

Minimal compiler-independent annotations to reconstruct high-quality IR

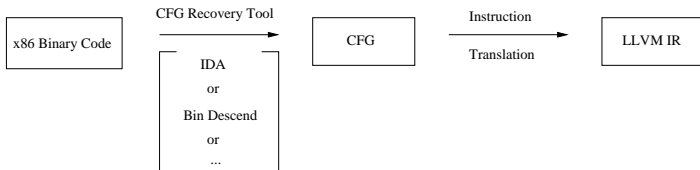
## Short term goals

- ❶ Experiment with Machine Code  $\rightarrow$  LLVM IR, to **understand** the challenges better
  - To select an existing decompilation framework.
  - Experiment with different variable and type recovery strategies
- ❷ Design suitable annotations for what cannot be inferred without them



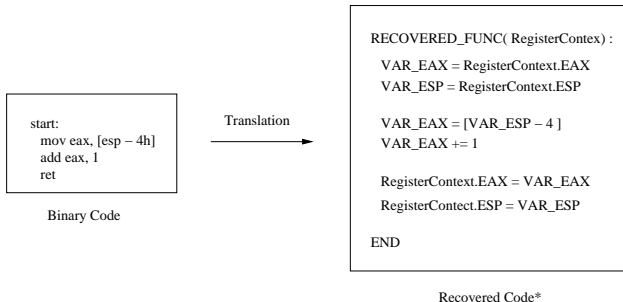
# Selecting mcsema

- Separation of modules: CFG recovery and CFG  $\rightarrow$  LLVM IR
- Actively supported and open sourced
- Well documented
- Functional LLVM IR





# Instruction Translation: Memory Model



\*: This is a functional view of the extracted code, NOT the actual recovered LLVM IR



# mcsema: Demo



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# Variable & Function Parameter Recovery

- Benefit
  - Enables many fundamental analysis (Dependence, Pointer analysis)
  - Functional IR
- State of the art
  - Divine
    - State of the art variable recovery
  - Second Write
    - Heuristics for function parameter detection
    - Scalable variable and type recovery
  - TIE
    - Type recovery





# Support & Limitations

- What Works
  - Integer Instructions
  - FPU and SSE registers
  - Callbacks, External Call, Jump tables
- In Progress
  - FPU and SSE Instructions: Not fully supported
  - Exceptions
  - Better Optimizations



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# What-you-see-is-not-what-you-execute

The following compiler (Microsoft C++ .NET) induced vulnerability was discovered during the Windows security push in 2002

```
memset(password, '\0', len);  
free(password);
```

—————>

```
free(password);
```