

QLang: Qubit Language (Reference Manual)

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Contents

1 Introduction

2 Lexical conventions

There are five kinds of tokens in the language, namely (1) identifiers, (2) keywords, (3) constants, (4) expression operators, and (5) other separators. At a given point in the parsing, the next token is chosen as to include the longest possible string of characters forming a token.

2.1 Character set

QLang supports a subset of ASCII; that is, allowed characters are `a-zA-Z0-9@#,-_ ; () [] { } < > = + / | *`, as well as tabulations `\t`, spaces, and line returns `\n` and `\r`.

2.2 Comments

Comments start with a `#` sign, which then extends until the next carriage return. Multiline comments are not supported.

2.3 Identifier (names)

An identifier is an arbitrarily long sequence of alphabetic and numeric characters, where `_` is included as “alphabetic”. It must start with a lowercase or uppercase letter, i.e. one of `a-zA-Z`.

The language is case-sensitive: `hullabaloo` and `hullABaLoo` are considered as different.

2.4 Keywords

The following identifiers are reserved for keywords: using them as function or variable name will result in an error at compilation time.

```
pi e
int float comp rvect cvect mat
true false
if elif else
def for from to by while break
or and xor
not re im norm isunit trans det adj conj sin cos tan exp
```

2.5 Constants

There are four sorts of constants in the language, namely *integer*, *float*, *complex* and *identifier* constants. The first are comprised of any sequence of integers of the form `0|[1-9][0-9]*` (recall that integers are non-negative), and have type `int`. The second are of type `float` and have the form `R`, while the third are of type `com` and have the form `R|R+Ri|Ri` where `R` consists of a (i) sign, (ii) an integer part followed by (iii) a point, (iv) a decimal part, then (v) either a `e` or a `E` followed by an exponent part, possibly signed. (i) and (v) are optional, and either (ii) or (iv) can be missing as well. In more detail, `R` is defined as `[+-]{0,1}(((A.B*|.B+)([eE][+-]?B+)?)|A[eE][+-]?B+)` and `A=0|[1-9]B*`, `B=0|[1-9]` (that is, `R` matches a real number such as `2.78e5`, `1.5E-1` or `10.25`).

check this
paragraph.

Finally, the identifier constants are a subset of the reserved keywords, and include:

e the base of natural logarithm $e = \sum_{k=0}^{\infty} \frac{1}{k!}$. Equivalent to `exp(1)`; has type `com`.

pi the constant π . Has type `com`.

true represents the Boolean value `true`. Stored internally as `int 1`.

false represents the Boolean value `false`. Stored internally as `int 0`.

2.6 Elementary operations and spacing

An operation, or language elementary unit, starts from the end of the previous one, and ends whenever a semicolon is encountered. Whitespace does not play any role, except as separators between tokens; in particular, indentation is arbitrary.

3 Objects and types

3.1 Objects and lvalues

As in C, “an object is a manipulatable region of storage; an lvalue is an expression referring to an object.”

3.2 Valid types

The language features 5 elementary types, namely `int`, `float`, `com`, `qub`, `mat`. (In particular, column and row vectors are represented respectively as $n \times 1$ or $1 \times n$ matrices.) Is also valid any type that inductively can be built from an a valid type as follows:

- elementary types are valid;
- an *array* of a valid type is valid. Arrays have fixed size (that must be declared at compilation time), and are comprised of elements of a single, fixed valid type;
- a *function* taking as input a fixed number of elements from (non-necessarily identical) valid types, and returning a valid type.

4 Conversions

Applying unary or binary operators to some values may cause an implicit conversion of their operands. In this section, we list the possible conversions, and their expected result – any conversion not listed here is impossible, and attempting to force it would generate a compilation error.

- `int` \rightarrow `float`, `float` \rightarrow `com`, `int` \rightarrow `com`.
- `com` \rightarrow `float`: the imaginary part of the complex number is dropped (will generate a warning).
- `float` \rightarrow `int`: the floating number is rounded towards zero.
- `com` \rightarrow `int`: equivalent to `com` \rightarrow `float` \rightarrow `int`.
- `com` \rightarrow `mat`: the floating number z becomes the 1×1 $\begin{bmatrix} z \end{bmatrix}$ (will generate a warning).
- `float` \rightarrow `mat`: the floating number x becomes the 1×1 matrix $\begin{bmatrix} x \end{bmatrix}$ (will generate a warning).
- `int` \rightarrow `mat`: the integer a becomes the 1×1 matrix $\begin{bmatrix} a \end{bmatrix}$ (will generate a warning).

5 Expressions

5.1 Operator Precedence

Operator Type	Operator	Associativity
Primary Expressions	() [] < >	Left
Unary	not re im norm unit trans det adj conj sin cos tan exp	Right
Binary	* / % + - @ eq lt gt leq geq or and xor ^	Left (except ^ which is Right)
Assignment	=	Left

5.2 Literals

Literals are integers, floats, complex numbers, qubits, and matrices, as well as the built-in constants of the language (e.g. `pi`). Integers are of type `int`, floats are of type `float`, complex numbers are of type `com`, qubits are of type `qub`, and matrices are of type `mat`. The built-in constants have pre-determined types described above (e.g. `pi` is of type `float`).

The remaining major subsections of this section describe the groups of *expression* operators, while the minor subsections describe the individual operators within a group.

5.3 Primary Expressions

5.3.1 identifier

Identifiers are primary expression. All identifiers have an associated type that is given to them upon declaration (e.g. `float ident` declares an identifier named `ident` that is of type `float`).

5.3.2 literals

Literals are primary expression. They are described above.

5.3.3 (*expression*)

Parenthesized expressions are primary expressions. The type and value of a parenthesized expression is the same as the type and value of the expression without parenthesis. Parentheses allow expressions to be evaluated in a desired precedence. Parenthesized expressions are evaluated relative to each other starting with the expression that is nested the most deeply and ending with the expression that is nested the least deeply (i.e. the shallowest).

5.3.4 *primary-expression(expression-list)*

Primary expressions followed by a parenthesized expression list are primary expressions. Such primary expressions can be used in the declaration of functions or function calls. The expression list must consist of one or more expressions separated by commas. If being used in function declarations, they must be preceded by the correct function declaration syntax and each *expression* in the expression list must evaluate to a type followed by an identifier. If being used in function calls each *expression* in the expression list must evaluate to an identifier.

5.3.5 *primary-expression*[*expression-list*]

Primary expressions followed by a bracketed expression list are primary expressions. Such primary expressions can be used in the declaration of matrices and arrays or to access an element of a matrix or array. The expression list must consist of one (for matrices and arrays) or two (for matrices) expressions separated by commas, and must evaluate to `int`.

5.3.6 [*expression-elementlist*]

Expression element lists in brackets are primary expressions. Such primary expressions are used to define matrices and therefore are of type `mat`. The expression element list must consist of one or more expressions separated by commas or semi-colons. Commas separate expressions into matrix columns and `colons` separate expressions into matrix rows. The expressions must evaluate to the same type and can be of type `int`, `float`, `com`, or `mat`. Additionally, the number of expressions in each row of the matrix must be the same. An example matrix is shown below.

```
1 int a = 3;
  int b = 12;
3 mat my_matrix = [ 0+1, 2, a: 5-1, 2*3-1, 12/2];
```

5.3.7 *<expression*|

Expressions with a less than sign on the left and a bar on the right are primary expression. Such expressions are used to define qubits and therefore are of type `qub`. The notation is meant to mimic the "bra-" of "bra-ket" notation and can therefore be thought of as a row vector representation of the given qubit. Following "bra-ket" notation, the expression must evaluate to an integer literal of only 0's and 1's, which represents the state of the qubit. An example "bra-" qubit is shown below.

```
1 qub b_qubit = <0100|;
```

5.3.8 |*expression*>

Expressions with a bar on the left and a greater than sign on the right are primary expression. All of the considerations are the same as for *<expression|*, except that this notation mimics the "ket" of "bra-ket" notation and can therefore be thought of as a column vector representation of the given qubit. An example "ket-" qubit is shown below.

```
1 int a = 001;
  qub k_qubit = |a>;
```

5.4 Unary Operators

5.4.1 *not expression*

The result is a Boolean indicating the logical *not* of the *expression*. The type of the expression must be `int` or `float`. In the *expressions*, 0 is considered false and all other values are considered true.

5.4.2 *re expression*

The result is the real component of the *expression*. The type of the expression must be `com`. The result has the same type as the expression (it is a complex number with 0 imaginary component).

5.4.3 *im expression*

The result is the imaginary component of the *expression*. The type of the expression must be `com`. The result has the same type as the expression (it is a complex number with 0 real component).

5.4.4 *norm expression*

The result is the norm of the *expression*. The type of the expression must be `mat`, `com`, `qub` or `float`. The result has type `float`, and corresponds to the 2-norm; in the case of `com` or `float`, this coincides with respectively the module and absolute value.

5.4.5 *isunit expression*

The result is a Boolean indicating if it is true or false that the expression is a unit matrix. The type of the expression must be `mat`.

5.4.6 *trans expression*

The result is the transpose of the *expression*. The type of the expression must be `mat`. The result has the same type as the *expression*.

5.4.7 *det expression*

The result is the determinant of the *expression*. The type of the expression must be `mat`. The result has type `float` if the expression is an integer matrix or `float` matrix and type `com` if the expression is a complex number matrix.

5.4.8 *adj expression*

The result is the `adjoint` of the *expression*. The type of the expression must be `mat`. The result has the same type as the *expression*.

5.4.9 *conj expression*

The result is the complex conjugate of the *expression*. The type of the expression must be `com` or `mat`. The result has the same type as the *expression*.

5.4.10 *sin expression*

The result is the evaluation of the trigonometric function sine on the *expression*. The type of the expression must be `int`, `float`, or `com`. The result has type `float` if the expression is of type `int` or `float` and type `com` if the expression is of type `com`.

5.4.11 *cos expression*

The result is the evaluation of the trigonometric function cosine on the *expression*. The type of the expression must be `int`, `float`, or `com`. The result has type `float` if the expression is of type `int` or `float` and type `com` if the expression is of type `com`.

5.4.12 *tan expression*

The result is the evaluation of the trigonometric function tangent on the *expression*. The type of the expression must be `int`, `float`, or `com`. The result has type `float` if the expression is of type `int` or `float` and type `com` if the expression is of type `com`. (If an error occurred because of a division by zero, a runtime exception is raised.)

5.5 Binary Operators

5.5.1 *expression ^ expression*

The result is the exponentiation of the first *expression* by the second *expression*. The types of the expression must be of type `int`, `float`, or `com`. If the expressions are of the same type, the result has the same type as the *expressions*. Otherwise, if at least one *expression* is a `com`, the result is of type `com`; if neither expressions are `com`, but at least one is `float`, the result is of type `float`.

5.5.2 *expression * expression*

The result is the product of the *expressions*. The type considerations are the same as they are for *expression ^ expression*.

5.5.3 *expression / expression*

The result is the quotient of the *expressions*, where the first *expression* is the dividend and the second is the divisor. The type considerations are the same as they are for *expression ^ expression*. Integer division is rounded towards 0 and truncated. (If an error occurred because of a division by zero, a runtime exception is raised.)

5.5.4 *expression % expression*

The result is the remainder of the division of the *expressions*, where the first *expression* is the dividend and the second is the divisor. The sign of the dividend and the divisor are ignored, so the result returned is always the remainder of the absolute value (or module) of the dividend divided by the absolute value of the divisor. The type considerations are the same as they are for *expression ^ expression*.

5.5.5 *expression + expression*

The result is the sum of the *expressions*. The types of the expressions must be of type `int`, `float`, `com`, `mat` or `qub`. If at least one *expression* is a `com`, the result is of type `com`; if neither expressions are `com`, but at least one is `float`, the result is of type `float`. Qubits and matrices are special and can only be summed with within operands of the same type (and, in the case of matrices, dimensions).

5.5.6 *expression - expression*

The result is the difference of the first and second *expression*. The type considerations are the same as they are for *expression + expression*.

5.5.7 *expression @ expression*

The result is the tensor product of the first and second *expressions*. The expressions must be of type of `mat`. The result has the same type as the *expression*.

5.5.8 *expression eq expression*

The result is a Boolean indicating if it is true or false that the two *expression* are structurally equivalent. The type of the expressions must be the same.

5.5.9 *expression lt expression*

The result is a Boolean indicating if it is true or false that the first *expression* is less than the second. The type of the expressions must be `int` or `float` and must be the same.

no ordering
for complex
numbers

5.5.10 *expression gt expression*

The result is a Boolean indicating if it is true or false that the first *expression* is greater than the second. The type of the expressions must be `int` or `float` and must be the same.

no ordering
for complex
numbers

5.5.11 *expression leq expression*

The result is a Boolean indicating if it is true or false that the first *expression* is less than or equal to the second. The type of the expressions must be `int` or `float` and must be the same.

no ordering
for complex
numbers

5.5.12 *expression geq expression*

The result is a Boolean indicating if it is true or false that the first *expression* is greater than or equal to the second. The type of the expressions must be `int` or `float` and must be the same.

no ordering
for complex
numbers

5.5.13 *expression or expression*

The result is a Boolean indicating the logical *or* of the *expressions*. The type of the expressions must be `int` or `float` and must be the same. In the *expressions*, 0 is considered `false` and all other values are considered `true`.

5.5.14 *expression and expression*

The result is a Boolean indicating the logical *and* of the *expressions*. The type considerations are the same as they are for *expression or expression*.

5.5.15 *expression xor expression*

The result is a Boolean indicating the logical *xor* of the *expressions*. The type considerations are the same as they are for *expression or expression*.

5.6 Assignment Operators

Assignment operators have left associativity

5.6.1 lvalue = *expression*

The result is the assignment of the expression to the lvalue. The lvalue must have been previously declared. The type of the expression must be of the same that the lvalue was declared as. Recall, lvalues can be declared as `int`, `float`, `comp`, `mat`, and `qubit`.

6 Declarations

Declarations are used within functions to specify how to interpret each identifier. Declarations have the form

declaration:
type-specifier declarator-list

6.1 Type Specifiers

There are five main type specifiers:

type-specifier:

`int`
`float`
`com`
`mat`
`qub`

6.2 Declarator List

The declarator-list consist of either a single declarator, or a series of declarators separated by commas.

declarator-list:
declarator
declarator , declarator-list

A declarator refers to an object with a type determined by the type-specifier in the overall declaration. Declarators can have the following form

declarator:
identifier
declarator ()
declarator [constant-expression]
(declarator)

6.3 Meaning of Declarators

Each declarator that appears in an expression is a call to create an object of the specified type. Each declarator has one identifier, and it is this identifier that is now associated with the created object.

If declarator D has the form

$$D ()$$

then the contained identifier has the type "function" that is returning an object. This object has the type which the identifier would have had if the declarator had just been D.

If a declarator has the form

$$D[\textit{constant-expression}]$$

or

$$D[]$$

then it is a declarator whose identifier is of type "array". In the first case, the constant-expression is an expression whose value can be defined at compile time. The type of that constant-expression is int. In the second case, the constant expression 1 is used.

An array may be constructed from one of the basic types, or from another array.

Parentheses in declarators do not change the type of contained identifier, but can affect the relations between the individual components of the declarator.

Not all possible combinations of the above syntax are permitted. There are certain restrictions such as how array of functions cannot be declared.

7 Statements

7.1 Expression statements

Expression statements are the building blocks of an executable program. As the name suggests, expression statements are nothing but expressions, delimited by semicolons. Expressions can actually be declarations, assignments, operations or even function calls. For example,

```
x = a + 3;
```

is a valid expression statement, and so is

```
print(a);
```

7.2 The if-elif-else statement

The `if-elif-else` statement is used for selectively executing statements based on some condition. Essentially, if the condition following the `if` keyword is satisfied, the specified statements get

executed. To specify what happens if the condition does not evaluate to true, we have the `else` keyword. In case we want to evaluate more than one condition at a time, we also have the `elif` keyword. So an `if` can be followed by any number of `elif`s, and at most one `else` block which is the end of the construct. The statements following the `else` are executed only if neither of the conditions specified before that evaluate to true.

```
if ( condition) {  
} elif (condition) {  
} else {  
}
```

Example:

```
if ( x==5) {  
    print("x is 5");  
} elif (x==3) {  
    print("x is 3");  
} else {  
    print("x is neither 5 nor 3")  
}
```

7.3 The for loop

The `for` statement is used for executing a set of statements a specified number of times. The statements within the `for` loop are executed as long as the value of the variable is within the specified range. As soon as the value goes out of range, control comes out of the `for` loop. To ensure termination, each iteration of the `for` loop increments/decrements the value of the variable, bringing it one step closer to the final value that is to be achieved.

By default, increment or decrement is by 1. However, if the desired increment is something other than one, the `by` keyword lets you specify that explicitly.

An example of `for` loop, increment by 2 is as follows:

```
for k from 1 to 10 by 2 {  
}
```

8 Scope rules

Within a program, variables may be declared and/or defined in various places. The scope of each variable is different, depending on where it is declared. There are three primary scope rules.

If a variable is defined at the outset/outer block of a program, it is visible everywhere in the program.

If a variable is defined as a parameter to a function, or inside a function/block of code, it is visible only within that function.

Declarations made after a specific declaration are not visible to it, or to any declarations before it.

For instance, consider the following snippet.

```
int x = 5;

int y = x + 10; \# this works

int z = a + 100; \# this does not

int a = 200;
```

9 Constant expressions

In order to facilitate efficiency in writing expression, the language introduces various mathematical constants such as π , e and matrices such *Pauli* matrices and *Hadamard* matrices which are frequently used in quantum computation. The keywords I , X , Y , Z , and H are reserved for this expressions.

$$I = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \quad X = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \quad Z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \quad Y = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix}.$$

The *Hadamard gate* is defined by the matrix:

$$H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}.$$

10 Examples

We present some examples that illustrates the use of Qlang in solving quantum computing problems.

10.1 Solving Quantum Computation Problem

10.1.1 Problem1

Evaluate the following expressions: a. $(H \otimes X)|00\rangle$ b. $\langle 101|000\rangle$ c. $\langle 01|H \otimes H|01\rangle$

```
2 def pseudo = evaluate () {
4   \# a qubit type declaration follows dirac notation
   qub mat0 = |00>;
6
   \# Both X and H are constant with type mat and
   \# @ corresponds to tensor product.
   mat HX = H @ X;
10
   pseudo = HX * mat0;
12 }
```

10.1.2 Problem 2

Find the matrix corresponding to the quantum circuit:

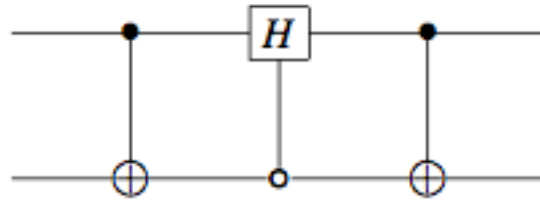


Figure 1: Quantum Circuit implementing series of control gates

```

1 def circuitMat = findMatrix () {
2
3   # all basis qubit in 2 dimension
4   qub mat0=|00>;
5   qub mat1=|01>;
6   qub mat2=|10>;
7   qub mat3=|11>;
8
9   # controlled not matrix
10  mat CNOT = [1,0,0,0;0,1,0,0;0,0,0,1;0,0,1,0]
11
12  #controlled hadmard matrix
13  mat HNOT = [1/sqrt(2),0,0,1/sqrt(2);0,1,0,0;1/sqrt(2),0,1,-1/sqrt(2)
14              ;0,0,0,0]
15  #composition of control gates
16  mat allGates = CNOT * HNOT * CNOT
17
18  # Matrix corresponding to the circuit
19  circuitMat =[allGates*mat0:allGates*mat1:allGates*mat2:allGates*mat3]
20 }

```

10.1.3 Problem 3

Consider the circuit and show the probabilities of outcome 0 where $|\Psi_{in}\rangle = |1\rangle$

```

1 def probability = outcomeZero() {
2
3   # top and bottom qubits
4   qub top = |0>;
5   qub bottom = |1>;
6
7   # Applying H on top qubit
8   mat output = (H @ I) * (top @ bottom);
9 }

```

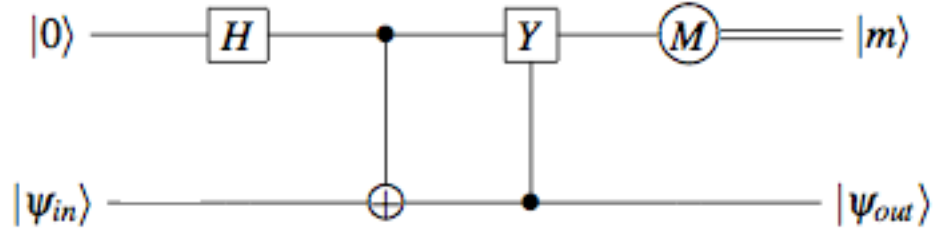


Figure 2: Quantum Circuit

```

10  # Controlled Not operator
    mat CNOT = [I, [0,0;0,0]; [0,0;0,0], X];
12
13  # Controlled Y operator
14  mat CY = [Y,[0,0;0,0];[0,0;0,0], I];
16
17  # Applying Control Operators
    output = (CY)*(CNOT)*output
18
19  # Applying measurement operator on top qubit |0> <0|
20  mat M = (|0>*<0| @ I)
22
23  # state after applying measurement operator on top qubit
    outcome = M * output;
24
25  #probability of outcome
26  probability = norm(outcome);
28 }

```

10.2 Simulation of Quantum Algorithm

10.2.1 Deutsch Jozsa Algorithm

```

2  def outcome = deutschJozsa(qub top, mat U){
4
5      # in corresponds to the qubit in top register
      # input is the tensor product of top register and bottom register
6      mat input= top @ |1>;
8
9      # application of Hadamard gate on both top and bottom inputs
      input = (H @ H)*input;
10
11     # application of U gate on the above result

```

```

12  input = U * input;

14  # application of Hadamard gate on the top register
    input = (H @ I)*input;

16

18  # application of measurement operator on the top register
    # top * Adj (top) corresponds to the Measurement operator

20  input=(top*Adj(top)@ I)*input;

22  #after the measurement is applied, check if the input is 0 or not
    if (input == 0){
24      #probability of outcome 0 is 0
        outcome = 0;
26    }
    else{
28      # probability of outcome 0 is 1
        outcome = 1;
30    }
}

```

10.2.2 Grover's Search Algorithm

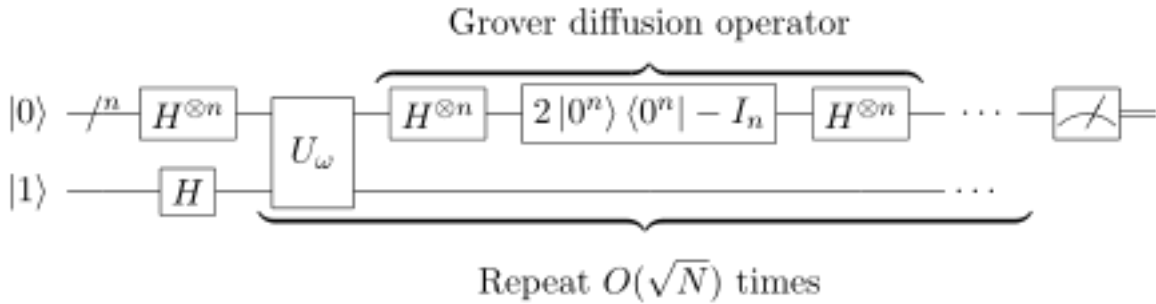


Figure 3: Grover Algorithm Circuit

```

2  def result = grover (quit top, int x0){
    # returns the probability to find x0 for a function f such that f(x0)=1
4  # x0 can be x0=0,1,?,2^n-1
    # this is a special case where n=1

6  # qubit in the bottom register
8  qub bottom = |1>;

10 # tensor product of top and bottom qubit
    mat input = top @ bottom;
12

```

```

14   #application of Hadamard
    input = (H @ H) * input;

16   #define S
    mat S = [1,0;0-1]

18

20   # k : number of time grover operator is applied
    # for n > 1 k=ceil((pi*2^(n/2))/4);
    int k = 1;

22

24   # define O operator such that  $O|x\rangle|q\rangle=|x\rangle|q \bmod f(x)\rangle$  or  $O|x\rangle=(-1)^{f(x)}|x\rangle$ 
    # for n > 1  $O = I(2^{(n+1)})$ ;
    mat O = I;
26   O(x0+1, x0+1) = -1;

28   # Grover iteration matrix
    mat GO = (G*O)^k;

30

32   # After application of Grover iteration matrix
    mat output = GO * input;

34   result = (H @ H)* output;

36 }

```